
# Advanced Library Format
# for
# ASIC Technology, Cells, & Blocks

**containing Functional, Electrical, and Physical Models for Design, Analysis, and Optimization from RTL to Layout**

**Version 1.9.2**

**September 28, 2000**

**This is a draft in progress.**
**Any section splits and x-refs will be updated.**

**Accellera**

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# Section 1
# Introduction

## 1.1 Motivation

Designing digital integrated circuits has become an increasingly complex process. More functions get integrated into a single chip, yet the cycle time of electronic products and technologies has become considerably shorter. It would be impossible to successfully design a chip of today's complexity within the time-to-market constraints without extensive use of EDA tools, which have become an integral part of the complex design flow. The efficiency of the tools and the reliability of the results for simulation, synthesis, timing and power analysis, layout and extraction rely significantly on the quality of available information about the cells in the technology library.

New challenges in the design flow, especially signal integrity, arise as the traditional tools and design flows hit their limits of capability in processing complex designs. As a result, new tools emerge, and libraries are needed in order to make them work properly. Library creation (generation) itself has become a very complex process and the choice or rejection of a particular application (tool) is often constrained or dictated by the availability of a library for that application. The library constraint can prevent designers from choosing an application program that is best suited for meeting specific design challenges. Similar considerations can inhibit the development and productization of such an application program altogether. As a result, competitiveness and innovation of the whole electronic industry can stagnate.

In order to remove these constraints, an industry-wide standard for library formats, the Advanced Library Format (ALF), is proposed. It enables the EDA industry to develop innovative products and ASIC designers to choose the best product without library format constraints. Since ASIC vendors have to support a multitude of libraries according to the preferences of their customers, a common standard library is expected to significantly reduce the library development cycle and facilitate the deployment of new technologies sooner.

## 1.2 Goals

The basic goals of the proposed library standard are:

- *simplicity* - library creation process needs to be easy to understand and not become a cumbersome process only known by a few experts.
- *generality* - tools of any level of sophistication need to be able to retrieve necessary information from the library.
- *expandability* - this needs to be done for early adoption and future enhancement possibilities.

- *flexibility* - the choice of keeping information in one library or in separate libraries needs to be in the hand of the user not the standard.
- *efficiency* - the complexity of the design information requires the process of retrieving information from the library does not become a bottleneck. The right trade-off between compactness and verbosity needs to be established.
- *ease of implementation* - backward compatibility with existing libraries shall be provided and translation to the new library needs to be an easy task.
- *conciseness* - unambiguous description and accuracy of contents shall be detailed.
- *acceptance* - there needs to be a preference for the new standard library over existing libraries.

## 1.3 Target applications

The fundamental purpose of ALF is to serve as the primary database for all third-party applications of ASIC cells. In other words, it is an elaborate and formalized version of the *databook*.

In the early days, databooks provided all the information a designer needed for choosing a cell in a particular application: Logic symbols, schematics, and a truth table provided the functional specification for simple cells. For more complex blocks, the name of the cell (e.g., asynchronous ROM, synchronous 2-port RAM, or 4-bit synchronous up-down counters) and timing diagrams conveyed the functional information. The performance characteristics of each cell were provided by the loading characteristics, delay and timing constraints, and some information about DC and AC power consumption. The designers chose the cell type according to the functionality, estimated the performance of the design, and eventually re-implemented it in an optimized way as necessary to meet performance constraints.

Design automation enabled tremendous progress in efficiency, productivity, and the ability to deal with complexity, yet it did not change the fundamental requirements for ASIC design. Therefore, ALF needs to provide models with *functional* information and *performance* information, primarily including timing and power. Signal integrity characteristics, such as noise margin can also be included under performance category. Such information is typically found in any databook for analog cells. At deep sub-micron levels, digital cells behave similar to analog cells as electronic devices bound by physical laws and therefore are not infinitely robust against noise.

Table 1-1 shows a list of applications used in ASIC design flow and their relationship to ALF. The boundary between supported and not supported applications can be defined by the *physical* information provided by ALF. Information needed for area and performance estimation and optimization, notably by synthesis and design planning tools, is provided by ALF. On the other hand, layout information is considered to be available in complementary libraries such as LEF.

Note: ALF covers *library* data, whereas *design* data needs to be provided in other formats.

**Table 1-1 Target applications and models supported by ALF**

| Application | Functional model | Performance model | Physical model |
|---|---|---|---|
| *simulation* | derived from ALF | N/A | N/A |
| *synthesis* | supported by ALF | supported by ALF | supported by ALF |
| *design for test* | supported by ALF | N/A | N/A |
| *design planning* | supported by ALF | supported by ALF | supported by ALF |
| *timing analysis* | N/A | supported by ALF | N/A |
| *power analysis* | N/A | supported by ALF | N/A |
| *signal integrity* | N/A | supported by ALF | N/A |
| *layout* | N/A | N/A | supported by ALF |

Historically, a functional model was virtually identical to a simulation model. A functional gate-level model was used by the proprietary simulator of the ASIC company and it was easy to lump it together with a rudimentary timing model. Timing analysis was done through dynamic functional simulation. However, with the advanced level of sophistication of both functional simulation and timing analysis, this is no longer the case. The capabilities of the functional simulators have evolved far beyond the gate-level and timing analysis has been decoupled from simulation.

RTL design planning is an emerging application type aiming to produce "virtual prototypes" of complex for system-on-chip (SOC) designs. RTL design planning is thought of as a combination of some or all of RTL floorplanning and global routing, timing budgeting, power estimation, and functional verification, as well as analysis of signal integrity, EMI, and thermal effects. The library components for RTL design planning range from simple logic gates to parameterizeable macro-functions, such as memories, logic building blocks, and cores.

From the point of view of library requirements, applications involved in RTL design planning need functional, performance, and physical data. The functional aspect of design planning includes RTL simulation and formal verification. The performance aspect covers timing and power as primary issues, while signal integrity, EMI, and thermal effects are emerging issues. The physical aspect is floorplanning. As stated previously, the functional and performance models of components can be described in ALF.

ALF also covers the requirements for physical data, including layout. This is important for the new generation of tools, where logical design merges with physical design. Also, all design steps involve optimization for timing, power, signal integrity, i.e. electrical correctness and physical correctness. EDA tools must be knowledgable about an increasing number of design aspects. For example, a place and route tool must consider congestion as well as timing, crosstalk, electromigration, antenna rules etc. Therefore it is a logical tep to combine the functional, electrical and physical models needed by such a tool in a unified library.

Figure 1-1 shows how ALF provides information to various design tools.

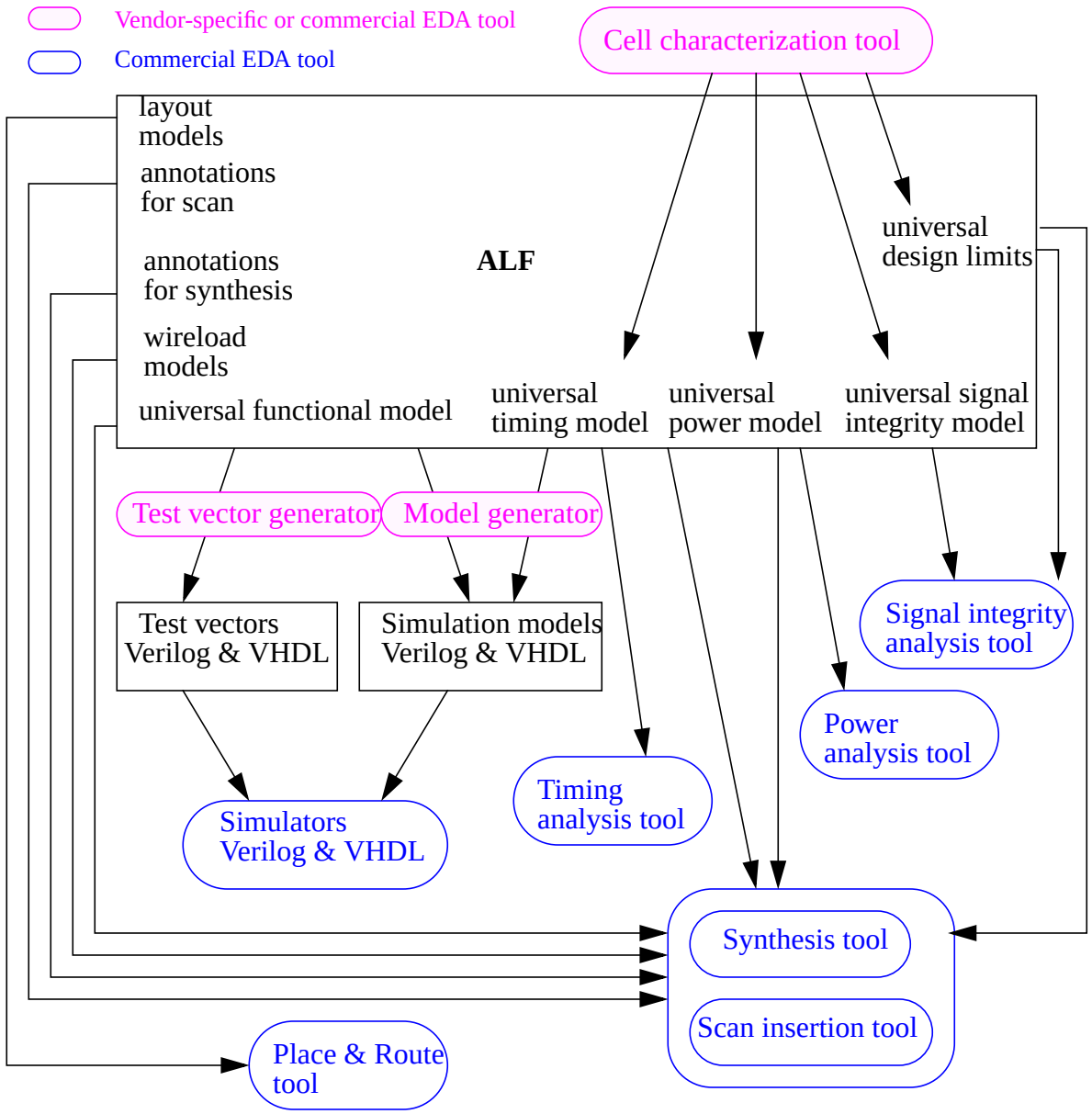



**Figure 1-1: ALF and its target applications**

The worldwide accepted standards for hardware description and simulation are VHDL and Verilog. Both languages have a wide scope of describing the design at various levels of abstraction: behavioral, functional, synthesizable RTL, and gate level. There are many ways to describe gate-level functions. The existing simulators are implemented in such a way that some constructs are more efficient for simulation run time than others. Also, how the simulation model handles timing constraints is a trade-off between efficiency and accuracy. Developing efficient simulation models which are functionally reliable (i.e., pessimistic for detecting timing constraint violation) is a major development effort for ASIC companies.

Hence, the use of a particular VHDL or Verilog simulation model as primary source of functional description of a cell is not very practical. Moreover, the existence of two simulation standards makes it difficult to pick one as a reference with respect to the other. The purpose of a generic functional model is to serve as an absolute reference for all applications that require functional information. Applications such as synthesis, which need functional information merely for recognizing and choosing cell types, can use the generic functional model directly. For other applications, such as simulation and test, the generic functional model enables automated simulation model and test vector generation and verification, which has a tremendous benefit for the ASIC industry.

With progress of technology, the set of physical constraints under which the design functions have increased dramatically, along with the cost constraints. Therefore, the requirements for detailed characterization and analysis of those constraints, especially timing and power in deep submicron design, are now much more sophisticated. Only a subset of the increasing amount of characterization data appears in today's databooks.

ALF provides a generic format for all type of characterization data, without restriction to state-of-the art timing models. Power models are the most immediate extension and they have been the starter and primary driver for ALF.

Detailed timing and power characterization needs to take into account the *mode of operation* of the ASIC cell, which is related to the functionality. ALF introduces the concept of *vector-based modeling*, which is a generalization and a superset of today's timing and power modeling approaches. All existing timing and power analysis applications can retrieve the necessary model information from ALF.

## 1.4 Conventions

The syntax for description of lexical and syntax rules uses following conventions.

| | |
|---|---|
| `::=` | definition of a syntax rule |
| `\|` | alternative definition |
| `[item]` | an optional item |
| `[item1 \| item2 \| ... ]` | optional item with alternatives |
| `{item}` | optional item that can be repeated |
| `{item1 \| item2 \| ... }` | optional items with alternatives which can be repeated |
| **`item`** | item in boldface font is taken verbatim |
| *`item`* | item in italic is for explanation purpose only |

The syntax for explanation of semantics of expressions uses the following conventions.

| | |
|---|---|
| `===` | left side and right side expressions are equivalent |
| `<item>` | a placeholder for an item in regular syntax |

## 1.5 Organization of this manual

This document presents the Advanced Library Format (ALF), a new standard library format for ASIC cells, blocks and cores, containing power, timing, functional, and physical information.

The first chapter defines the motivation and goals of ALF.

The second chapter describes the underlying modeling concepts: functional modeling, cell characterization for timing and power, and additional modeling features for synthesis and test.

The third chapter defines the object model.

The fourth chapter details the library organization within ALF.

The fifth chapter defines the formal constructs for functional modelling.

The sixth chapter defines supplementary constructs for synthesis and design for test.

The seventh chapter defines the general rules for arithmetic models.

The eight chapter defines electrical performance modeling, i.e., timing, power, signal integrity

The ninth chapter defines modeling for physical design.

The tenth chapter specifies the lexical rules.

The eleventh chapter specifies the syntactical rules.

The first appendix provides sample applications.

The second appendix provides an ALF/SDF cross-reference.

# Section 2
# Characterization and Modeling

This chapter elaborates on the basics of cell modeling and characterization, which is the primary source of library information.

## 2.1 Basic concepts

The functional models within an ASIC library describe functions and algorithms of hardware components, as opposed to synthesizeable functions or algorithms. The functional modeling language for the ASIC library is designed to make the description of existing hardware easy and efficient. The scope here is different from a hardware description language (HDL) or a programming language designed to specify functionality without other aspects of hardware implementation.

Functional description provides boolean functions or truth tables, including state variables for sequential logic. Boolean and arithmetic operators for scalars and vectors are also provided. Combinational and sequential logic cells, macrocells (e.g., adders, multipliers, and comparators), and atomic megacells (e.g., memories) can be modeled with these capabilities.

Vectors describe the stimuli for characterization. This encompasses both the concept of timing arcs and logical conditions. An exhaustive set of vectors can be generated from functional information, although the complexity of the exhaustive set precludes it from practical usage. The characterizer makes a choice of the relevant subset for characterization.

Power characterization is a superset of timing characterization using the same set and range of characterization variables: load, input slew rate, skew between multiple switching inputs, voltage, and temperature. Characterization measurements, such as delay, output slew rate, average current in time window, bounds of allowed skew for timing constraints, etc. can be described as functions of the characterization variables, by using equations or lookup tables. More complicated calculation algorithms cannot be described explicitly in the library, but can be referenced using templates.

A core is not an atomic megacell, since it can be split up into smaller components. Templates provide the capability of defining and reusing blocks consisting of atomic constructs or of other blocks. Thus a hierarchical description of the complete core can be created in a simple and efficient way.

Abstraction is required for the characterization of megacells: vectors describe events on buses rather than on scalar pins; number and range of switching pins within a bus become additional characterization variables. Characterization measurements are expandable and can be extrapolated from scalar pin to bus.

# 2.2 Performance modeling for characterization

This section highlights modeling for timing, power, and signal integrity.

## 2.2.1 Modeling for timing

The timing models of cells consist of two types: *delay models* for combinational and sequential cells, and *timing constraint models* for sequential cells. Both types can be described by timing arcs. A timing arc is a sequence of two events that can be described by a vector expression "event *e1* is followed by event *e2*".

For example, a particular input to output delay of an inverting logic cell is identified by the following timing arc:

```
01 A -> 10 Z
```

which reads "rising edge on input `A` is followed by falling edge on output `Z`".

A setup constraint between data and clock input of a positive edge triggered flip-flop is identified by the following timing arc:

```
01 D -> 01 CP
```

which reads "rising edge on input `D` is followed by rising edge on input `CP`".

A crucial part in ASIC cell development is to characterize a model that describes the behavior of each timing arc with sufficient accuracy in order to guarantee correct functional behavior under all required operational conditions.

A delay model usually needs two output variables:

- the i*ntrinsic delay*, measured between a well-defined threshold value of the input signal and a well-defined threshold value of the output signal, and
- the transition *delay*, measured between two well-defined threshold values of the output signal. Hence the transition delay is a fraction of the total output transition time, also called *slew rate* or *edge rate*.

A timing constraint model needs one output variable:

A timing constraint is the *minimum or maximum allowed elapsed time* between two signals, measured between well-defined threshold values between those two signals. This definition is similar to the *intrinsic delay,* except there is no input-output relationship between the two signals. Both signals are usually inputs to the cell.

The actual values of transition times and load capacitances seen by each pin of a cell instance are calculated by a delay predictor. Delay prediction can be separated into two tasks:

1. Acquisition of information on pin capacitance, then extracting or estimating layout parasitics for each net and fitting those into the load characterization model (lumped C, R, etc.).
2. Calculation of internal signal transition times based on the extracted internal load and on load and transition times at the boundaries of the system.

Lookup tables provide a general modeling capability without precluding any level of accuracy.

Equations can feature polynomial expressions, exponentials and logarithms, and arbitrary transcendent functions. For practical purpose, only the four basic arithmetic operations (`+`, `-`, `*`, `/`) and exponentiation and logarithm are supported for standard models.

Some models can require transcendent functions or complicated algorithms that cannot be expressed directly in equations. Other models and algorithms need protection from being visible. In order to address needs that go beyond standard modeling features, a template-reference scheme is used: any model which is not in table or equation format needs to be a pointer to a customer-defined model, which can reside outside the library. All these are defined further in Table 2-1.

**Table 2-1 Modeling choices for cell characterization library**

| Type of model | Features | Purpose |
|---|---|---|
| table | discrete points, multidimensional | direct storage of characterization data, direct accuracy control through mesh granularity |
| equation | expressions with +, -, *, /, exponent, logarithm | analytical model, well-suited for optimization purpose, more compact than table, also usable for arithmetic operations on tabulated data (scale, add, subtract ..) |
| reference | pointer to any type of model | reuse of predefined model (which can be table or equation), protection of user-defined model |

Regardless of which type of model is chosen, there is a need to explicitly specify the meaning of the variables and the units. The specification of variables and units can be made outside the model and independent of the chosen model.

Since the set of variables shall not be restrictive in order to allow any enhancements (e.g., move from a lumped capacitance to an RC model), *context-sensitive keywords* are proposed (e.g., `load` and `slewrate`). The application parser need not know the meaning of the context-sensitive keyword, except it is used as a variable in a model and has some unit attached to it, e.g., picofarad, nanosecond, etc.

### 2.2.2 Modeling for power

A power model is an extension of the delay model for each timing arc using a third variable:

> the *scaled average current*, measured by integrating and scaling the total transient current through the power supply of the cell for the specific timing arc or vector. The current measurement can start anytime before the first event of the vector starts and can end anytime after all transients of the vector have stabilized.

Variants of this model are scaled average power and energy, which are obtained by simple scaling of average current measurements:

```
power = current * Vdd
energy = current * Vdd * integration time
```

The set of vectors causing power consumption within a cell is a superset of those vectors causing the cell output to switch. While only the vectors with switching output are needed for delay characterization, more vectors are needed for accurate power characterization.

For example, consider a flip-flop, which consumes power at every edge of the clock, even if the output does not switch. The vectors for delay and power characterization can be described as follows:

```
01 CP -> 01 Q
01 CP -> 10 Q
```

The vectors for power characterization with only clock-switching can be described as follows:

```
01 CP && Q==D
10 CP && Q==D
```

The `D` input having the same value as the `Q` output is a necessary and sufficient condition that the output shall not switch at the rising edge of `CP` and that the value transferred to the master latch at the falling edge of `CP` is the same as already stored. Hence, those two vectors capture the actual power dissipation only within the clock buffers. Additional power vectors can be defined to capture the power dissipation within the data buffers and the master latch etc.

For a 2-input AND gate with input pins `A`, `B` and output pin `Z`, a *glitch* is observed if the event `01 A` is detected and then the event `10 B` is detected before the input-to-output delay elapses. It is possible to describe the glitch by a higher-order vector.

In dynamic simulation with *transport delay mode*, the glitch would appear as follows:

```
01 A -> 10 B -> 01 Z -> 10 Z
```

Simulation featuring *transport delay mode with invalid-value-detection* would exhibit the glitch as follows:[1]

```
01 A -> 10 B -> 'b0'bX Z -> 'bX'b0 Z
```

Simulation with *inertial delay mode* would suppress the output transitions:

```
(01 A -> 10 B) && !Z
```

The last expression can be used for each of the three simulation modes, since `!Z` is always *True* from beginning to end of the sequence `01 A -> 10 B`, in particular at the time when the sequence `01 A -> 10 B` is detected.

Each way of expressing vectors can be derived from the cell functionality. The different examples for delay vectors (i.e., timing arcs), power vectors, and glitch vectors emphasize the rich potential of modeling capabilities using vector expressions.

State-dependent *static power* is also within the scope of vector-based power models. Static power consumption is activated by a simulation model in the same way as level-sensitive logic in functional modeling by a boolean expression, whereas *transient power* consumption is activated similar to edge-sensitive logic by a vector expression.

The advantages of adding power models within each delay vector and providing extra power vectors are the following:

1. Use based-edge literals to avoid parser ambiguity.

- straightforward extension of delay characterization
- capable of yielding the most detailed and accurate model on gate-level
- each vector defines a comprehensive stimulus for power measurements

More abstract vector expressions are provided for power modeling of complex blocks, where simplification is needed in order to deal with the complexity of characterization vectors.

### 2.2.3 Modeling for signal integrity

The concept of vector-based cell characterization with multiple variables also accommodates the requirements for signal integrity modeling. Although signal integrity is closely related to interconnect parasitics, i.e., extracted *design* information, data in the cell *library* needs to exist in order to support signal integrity analysis.

- Crosstalk analysis needs characterization of *driver resistance* on output pins and *noise margin* on input pins.
- IR drop and electromigration analysis on power supply lines needs characterization of a*verage currents* for power analysis, *RMS currents*, and *current waveforms*.
- Electromigration (EM) analysis within cells needs characterization of *current limits*. In a direct evaluation approach, the current limits are checked against the actual currents. The latter data comes from the characterization for power and IR drop. In an indirect evaluation approach, the current limits can be expressed as *frequency-dependent load limits* and/or *slewrate limits*.
- Hot electron (HE) analysis within cells needs characterization of *flux* (charge density) or *fluence* (accumulated charge density over time) and its respective limits for performance degradation. In a direct evaluation approach, the flux or fluence limits are checked against the actual flux or fluence, respectively. In an indirect evaluation approach, the limits of performance degradation due to fluence can be expressed as *frequency-dependent load limits* and/or *slewrate limits*, in the same way as for electromigration.

The characterization vector set for driver resistance is a subset of delay characterization vectors. In buffered cells, the driving input does not matter, since the driver resistance seen at the output is the same. However, there is always a different driver resistance for rise and fall, which is also dependent on process, voltage, and temperature.

Noise margin characterization is especially important for control and data pins of sequential cells. The set of characterization vectors is complementary to the timing constraint characterization vectors. For instance, noise margin on a clock pin is complementary to the pulsewidth constraint. If pulsewidth corresponds to the smallest possible signal causing a *valid* functional reaction, noise margin corresponds to the largest possible signal causing *no* functional reaction.

The characterization vector set for IR drop and EM on power supply lines is essentially the same as for power analysis, only the set of data per vector is richer. IR drop analysis can use

average currents, peak currents, or current waveforms. EM analysis can use average, peak, RMS, or a combination of the above.

The characterization vector set for EM and HE effect occurring within cells is very similar to the characterization vector set for power analysis, depending whether a direct or indirect evaluation approach is used.

In summary, modeling for crosstalk is a natural extension of modeling for timing, whereas IR drop and EM and HE modeling are natural extensions of modeling for power.

## 2.3 Physical modeling for synthesis and test

This section highlights cell and wire modeling.

### 2.3.1 Cell modeling

Physical modeling of cells requires annotating cell properties (e.g., area, height, width, and aspect ratio). The set of annotated properties give an application, such as synthesis, a choice to pick one cell from a set of functionally equivalent cells, if one property is more desirable than another one under given synthesis goals and constraints.

Cell pins can also have annotated properties, such as pin capacitance, voltage swing, switching threshold, etc.

Most of the requirements for the modeling of test are already fulfilled by the functional model. Declaration of pins and their direction (input, output, or bidirectional) is already a generic requirement for cell modeling.

Scan insertion tools require specific annotations about cell and pin properties relevant for scan test. They also require reference to equivalent non-scan cells. An equivalent non-scan cell is a scan cell where all scan-specific hardware (e.g., a multiplexor or scan clock) is removed.

The variables used in the functional model shall have their counterpart in the pin declaration. Only primary input pins can be primary inputs of functions, while primary output pins, internal pins, or virtual pins can be primary or intermediate outputs of functions. Furthermore, test vectors for fault coverage can be derived from the functional model in a formal way.

The remainder of the modeling for test requirements can be covered by annotations of cell properties and cell pin properties. For instance, a cell can be labeled as a scan flip-flop and a pin can be labeled as scan input or mode select pin.

### 2.3.2 Wire modeling

The purpose of *wire modeling* is to get good estimates of *parasitic resistance* and *capacitance* as a function of *fanout*. These estimates are technology-specific and they depend on metal layer, sheet resistance, self-capacitance per unit wirelength, fringe capacitance per unit wirelength, and via resistance for wires routed through multiple layers.

The wires can be represented as a collection of models, in a similar way as cells. For example,

```
// wire with fanout ≤ 5 routed in metal 1, 2
WIRE small_wire {
   ATTRIBUTE { metal1 metal2 }
   LIMIT { FANOUT { MAX = 5; } }
   /* fill in data */
}
// wire with 10 ≤ fanout ≤ 20 routed in metal 1, 2, 3, 4, 5
WIRE big_wire {
   ATTRIBUTE { metal1 metal2 metal3 metal4 metal5 }
   LIMIT { FANOUT { MIN = 10; MAX = 20; } }
   /* fill in data */
}
```

From a modeling standpoint, no particular language is required for performance modeling of wires that is different from performance modeling of cells. The fanout shall be an input variable and the capacitance and resistance are output variables. The values can be expressed either in tables or in equations. Usually first order equations (with slope and intercept) are used for wire modeling.

# 2.4 Functional modeling

This section highlights the usage of combination and sequential logic.

## 2.4.1 Combinational logic

Combinational logic can be described by continuous assignments of boolean values (*True* or *False*) to output variables as a function of boolean values of input variables. Such functions can be expressed in either equation or table format.[2]

Consider an arbitrary continuous assignment:

$$z = f(a_1 .., .. a_n)$$

In a dynamic or simulation context, the left-hand side (LHS) variable *z* is evaluated whenever there is a change in one of the right-hand side (RHS) variables $a_i$. No storage of previous states is needed for dynamic simulation of combinational logic.

## 2.4.2 Level-sensitive sequential logic

In sequential logic, an output variable $z_j$ can also be a function of itself, i.e., of its previous state. The sequential assignment has the form

$$z_j = f(a_1 .., .. a_n , z_1 .., .. z_m)$$

The RHS cannot be evaluated continuously, since a change in the LHS as a result of a RHS evaluation shall trigger a new RHS evaluation repeatedly, unless the variables attain stable values. Modeling capabilities of sequential logic with continuous assignments are restricted to systems with oscillating or self-stabilizing behavior.

See Section 5.2.1 for more details.

2. This standard uses the existing boolean syntax notation described in the ANSI C standard.

### 2.4.3 Edge-sensitive sequential logic

In order to model *edge-sensitive sequential logic*, the concept of logical transitions and logical states are introduced here.

If the triggering function *g* is sensitive to logical transitions rather than to logical states, the function *g* evaluates to *True* only for an infinitely small time, exactly at the moment when the transition happens. The sole purpose of *g* is to trigger an assignment to the output variable through evaluation of the function *f* exactly at this time.

Edge-sensitive logic requires storage of the previous output state and the input state (to detect a transition). In fact, all implementations of edge-triggered flip-flops require at least two storage elements. For instance, the most popular flip-flop architecture features a master latch driving a slave latch.

See Section 5.2.2 for more details.

### 2.4.4 Vector-sensitive sequential logic

In order to model generalized higher order sequential logic, the concept of vector expressions is introduced, an extension of the boolean expressions.

A vector expression describes sequences of logical events or transitions in addition to static logical states. A vector expression represents a description of a logical stimulus without timescale. It describes the order of occurrence of events.

The -> operator *(followed by*) can be used to generally describe a sequence of events or a vector. For example, consider the following vector expression:

```
01 A -> 01 B
```

which reads “rising edge on `A` is followed by rising edge on `B`”.

A vector expression is evaluated by an event sequence detection function. Like a single event or a transition, this function evaluates *True* only at an infinitesimally short time when the event sequence is detected.

See Section 5.3.1 for more details.

# Section 3
# Object Model

This section discusses the object model used by ALF and provides the syntax rules for all objects. The syntax rules are provided in standard BNF form.

A *library* consists of one or more *objects*. Each object is defined by a keyword and an optional name for the object and an optional *value* of the object.

A *keyword* defines the type of the object. Section 3.2 and Section 3.3 define various types of objects used in ALF and related keywords.

An optional *identifier* (also called *name*) following the keyword defines the *name of the object*. This name shall be used while referencing an object inside other objects in the library. If an object is not referenced by name, then the object need not be named.

A *literal* defines an optional value associated with the object. An *expression* can be used when the value of the object cannot be expressed as a literal.

An object can contain one or more objects. The containing object is called a *hierarchical object*. The contained objects are called *children objects*. The children objects are defined and referenced inside curly braces ({}) in the description of the hierarchical object. An object without children is called an *atomic object*.

*Forward referencing* of objects is not allowed. Therefore, all objects shall be defined before they can be instantiated. This allows library parsers to be one-pass parsers.

## 3.1 Syntax conventions

In order to make ALF easy to parse, the syntax conventions follow the syntax rules defined in Section 1.4. These should also be followed for future extensions of the grammar.

The first token of the object is the object type identifier, followed by a name (mandatory or optional, depending on object type), followed by (mandatory or optional) `=` and value assignment, followed by (mandatory or optional) children objects enclosed by curly braces. Objects with more than one token (i.e., name and/or value) and without children are terminated with `a;`.

Examples:

1. Unnamed object without value assignment:

```
MY_OBJECT_TYPE
```

or

```
MY_OBJECT_TYPE {
   //fill in children objects
}
```

2. Unnamed object with value assignment:

```
MY_OBJECT_TYPE = my_object_value;
```

or

```
MY_OBJECT_TYPE = my_object_value {
   //fill in children objects
}
```

3. Named object without value assignment:

```
MY_OBJECT_TYPE my_object_name;
```

or

```
MY_OBJECT_TYPE my_object_name {
   //fill in children objects
}
```

4. Named object with value assignment:

```
MY_OBJECT_TYPE my_object_name = my_object_value;
```

or

```
MY_OBJECT_TYPE my_object_name = my_object_value {
   //fill in children objects
}
```

The objects in ALF can be divided into the following categories: *generic objects*, *library-specific objects*, *arithmetic models*, *geometric models*, and *library-specific singular objects*.

## 3.2 Generic objects

A generic object can appear at every level in the library within any scope. The semantics of a generic object need to be understood by any ALF compiler if the generic object is within the scope of application for that compiler.

The objects shown in Figure 3-1 shall be considered generic objects:

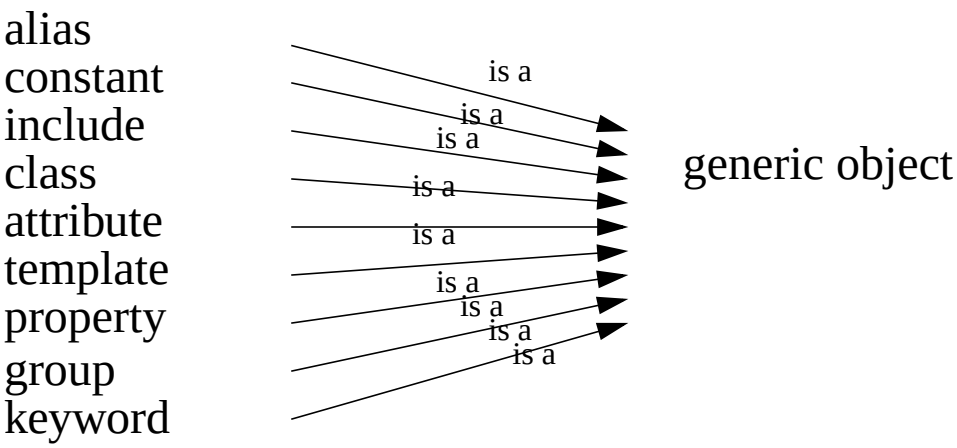



**Figure 3-1: Generic objects**

### 3.2.1 CONSTANT statement

A *CONSTANT* object is a named object with value assignment and without children objects. The value is a number.

Example:

```
CONSTANT vdd = 3.3;
```

### 3.2.2 ALIAS statement

An *ALIAS* object is a named object with value assignment and without children objects. The value is a string.

Example:

```
ALIAS RAMPTIME = SLEWRATE;
```

### 3.2.3 INCLUDE statement

An *INCLUDE* object is a named object without value assignment and without children. The name is a quoted string containing the name of a file to be included.

Example:

```
INCLUDE “primitives.alf”;
```

Since the file name is a quoted string, any special symbols (like ~ or *) are allowed within the filename. The interpretation of those (e.g., as a file search path) is up to the application.

### 3.2.4 CLASS statement

A *CLASS* object is a named object with optional value assignments and children objects. The name can be used by other objects to reference the class object.

Example:

```
CLASS my_class { ... }
...
MY_OBJECT_TYPE my_object {
   CLASS = my_class;
} // my_object belongs to my_class
```

### 3.2.5 ATTRIBUTE statement

An *ATTRIBUTE* object is an unnamed object without value, but containing children objects. The attribute object shall be the child object of another object. The children of the attribute object are unnamed objects that can have other unnamed objects as children objects. The purpose of an attribute object is to provide free association of objects with attributes when there is no special category available for the attributes.

Examples:

```
CELL rr_8x128 {
   ATTRIBUTE {ROM ASYNCHRONOUS STATIC}
}

PIN read_write_select {
   ATTRIBUTE {READ{POLARITY=low;} WRITE{POLARITY=high;}}
}
```

### 3.2.6 TEMPLATE statement

A *TEMPLATE* object is a named object with one or more children objects. Any valid ALF object can be a child object of a template object. Identifiers enclosed between < and > are recognized as *placeholders*. When a template object is used, each of its placeholders shall be referenced by order or by explicit name association.

Example:

```
TEMPLATE std_table {
   CAPACITANCE {PIN=<pin1>; UNIT=pF; TABLE {0.02 0.04 0.08 0.16}}
   SLEWRATE {PIN=<pin2>; UNIT=ns; TABLE {0.1 0.3 0.9}}
}
```

An instantiation of the above template object with explicit reference to placeholders by name:

```
std_table{pin1=out; pin2=in;}
```

An instantiation of the above template object with implicit reference to placeholders by order:

```
std_table{out in}
```

If a symbol within a placeholder appears more than once in the template definition, the order for implicit reference is defined by the first appearance of the symbol. Explicit referencing improves the readability and is the recommended usage.

A template instantiation can appear at any place within a hierarchical object, as long as the template object contains the structure of valid objects inside. Hierarchical templates contain other template objects.

### 3.2.7 PROPERTY statement

A *PROPERTY* object is a named or an unnamed *annotation container*. It can be used at any level in the library. It is used for arbitrary parameter-value assignment.

Example:

```
PROPERTY items {
   parameter1=value1;
   parameter2=value2;
}
```

A `PROPERTY` statement can also contain assignments with multiple values.

```
property ::=
          PROPERTY [ identifier ] { property_items }

property_items ::=
          property_item { property_item }

property_item ::=
          unnamed_assignment
     |    multi_value_assignment
```

Example:

```
PROPERTY {
   my_param1 = value1;
   my_param2 { val1 val2 val3 }
   my_param3 = value4;
}
```

### 3.2.8 GROUP statement

A *GROUP* object is a set of elements with commonality between them. Thus, the common characteristics can be defined once for the group instead of being repeated for each element.

Example:

```
GROUP time_measurements = {DELAY SLEWRATE SKEW JITTER}
```

The statement

```
time_measurements { UNIT = ns; }
```

replaces the following statements:

```
DELAY       { UNIT = ns; }
SLEWRATE    { UNIT = ns; }
SKEW        { UNIT = ns; }
JITTER      { UNIT = ns; }
```

### 3.2.9 KEYWORD statement

The ALF language allows the use of customized context-sensitive keywords for certain purposes. While the semantics of these custom keywords can only be known by the user of such

keywords, every ALF parser needs to have the capability to check the correct syntax of objects involving custom keywords.

Generic objects shall be augmented by using the KEYWORD statement. The KEYWORD statement shall be defined as:

```
generic_object ::=
          // set of current definitions in Section 11.7
   |      keyword_statement

keyword_statement ::=
   KEYWORD context_sensitive_keyword = syntax_item_identifier ;
```

The following *syntax_item*_identifiers , which are a subset of the objects defined in Section 11.8 are legal:

```
syntax_item_identifier ::=
          annotation
   |      annotation_container
   |      arithmetic_model
   |      arithmetic_submodel
   |      arithmetic_model_container
   |      vector_assignment
```

Example:

```
KEYWORD my_arithmetic_model = arithmetic_model;
KEYWORD my_annotation_for_capacitance = annotation;
KEYWORD my_annotation_for_resistance = annotation;
my_arithmetic_model {
   HEADER {
      CAPACITANCE { my_annotation_for_capacitance = foo; }
      RESITANCE { my_annotation_for_resistance = bar; }
   }
   EQUATION { 10*CAPACITANCE + 0.5*RESISTANCE }
}
```

It is illegal to redefine intrinsic ALF keywords.

Example:

```
KEYWORD vector = arithmetic_model; // THIS IS ILLEGAL!!!
```

## 3.3 Library-specific objects

The library-specific objects define their nature and their relationship to each other by containment rules. For example, a library can contain a cell, but a cell can not contain a library. However, both the library and the cell can contain any generic object. A generic object defined at the library level makes it visible inside the scope of that library, defining it on the cell level makes it visible inside the scope of that cell and its children objects.

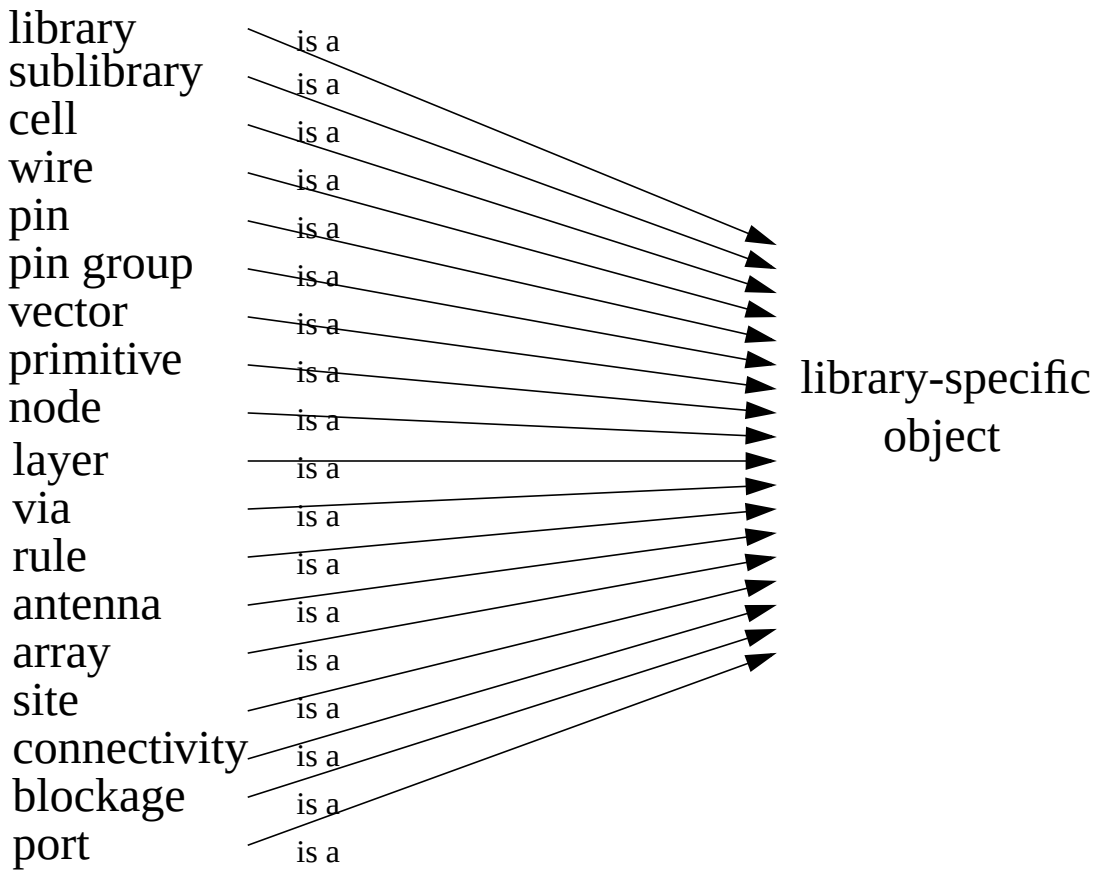



**Figure 3-2: Library-specific objects**

Multiple instances of named library-specific objects may appear in a given context. for example, a library may contain multiple cells, a cell may contain multiple pins etc.

## 3.4 Arithmetic models

An arithmetic model is an object that describes characterization data or a more abstract, measurable relationship between physical quantities. The modeling language allows tabulated data as well as linear and non-linear equations. The equations consist of arithmetic expressions based on the symbols defined in *IEEE 1364-1995*.

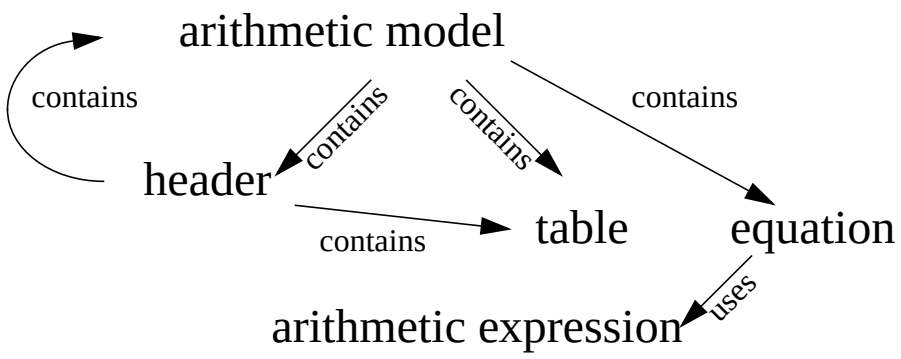



**Figure 3-3: Arithmetic model**

## 3.5 Geometric models

A geometric model describes the form of an object in a physical library. It is in the context of a pattern, which is associated with physical objects, such as via, blockage, port, rule. Patterns and other physical objects can also be subjected to geometric transformations.

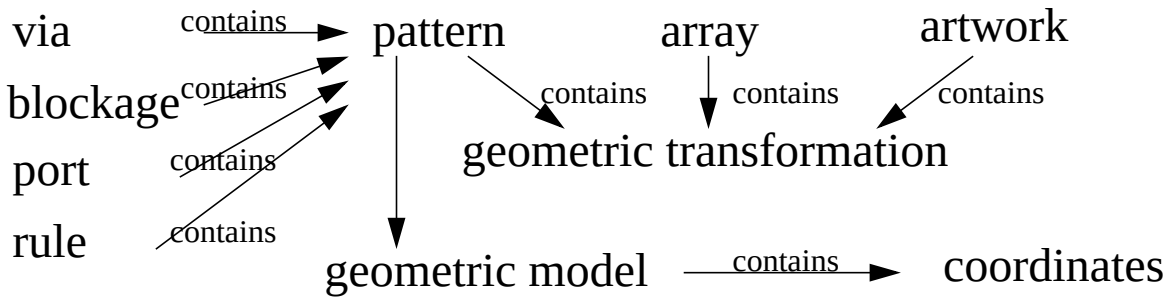



**Figure 3-4: Geometric model and its context**

## 3.6 Library-specific singular objects

Library-specific singular objects may only appear in one instance within a given context. For instance, a cell may contain at most one function and one test description.

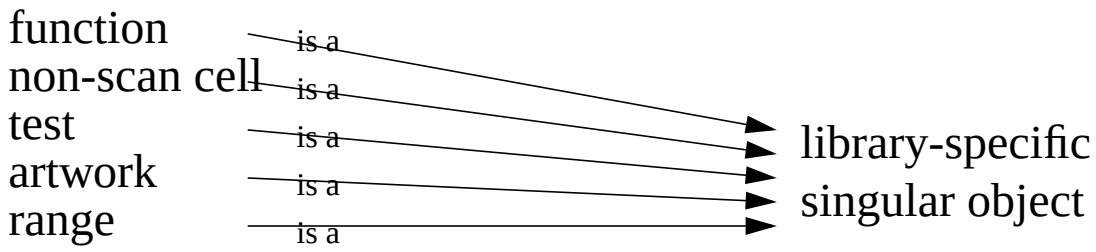



**Figure 3-5: Library-specific singular objects**

An object called FUNCTION describes the functional specification of a digital circuit (or a digital model of an analog or a mixed-signal circuit) in a canonical form. The modeling language allows behavioral models as well as statetables and structural models with primitives. The behavioral models contain boolean expressions, closely matching *IEEE 1364-1995*. Since boolean expressions are insufficient to describe sequential logic, ALF introduces new operators and symbols that can be used in conjunction with boolean operators and symbols. Expressions that use both the IEEE operators and the new operators are called *vector expressions*.

An object called TEST describes the specification for testing a digital circuit, using the same constructs as FUNCTION. However, TEST describes a stimulus generator for the circuit, whereas FUNCTION describes the circuit itself.

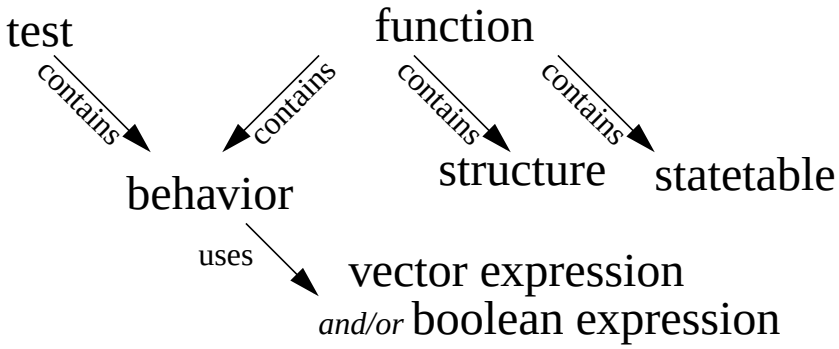



**Figure 3-6: FUNCTION and TEST**

## 3.7 Relationships between objects

The following figures describe the categories of objects and their relationships with each other.

Library-specifi objects, arithmetic models, geometric models and library-specific singular objects may contain auxiliary objects, called annotation and annotation container, respectively. The purpose of the latter is to qualify the semantics of the former.

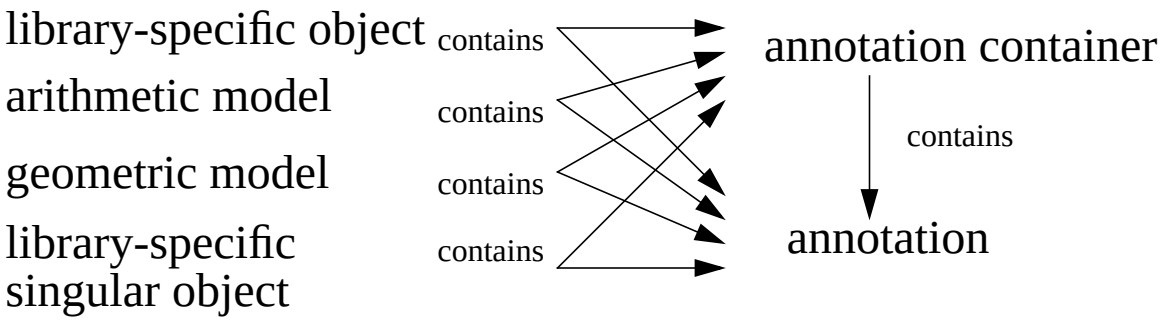



**Figure 3-7: Objects containing annotations or annotation containers**

All the above mentioned objects may contain generic objects.

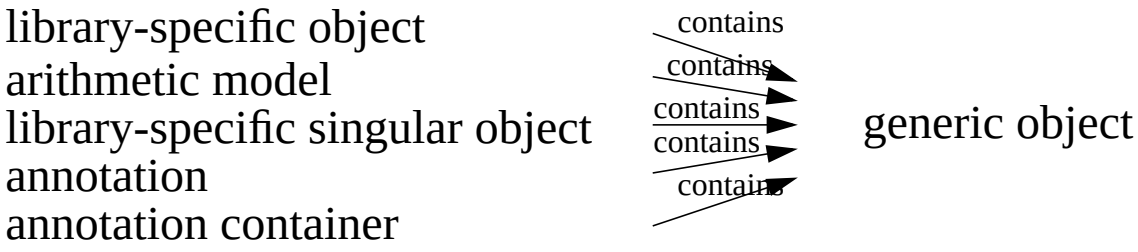



**Figure 3-8: Objects containing generic objects**

The following figures illustrate the relationship between objects in a library for functional and electrical design and for physical design, respectively.

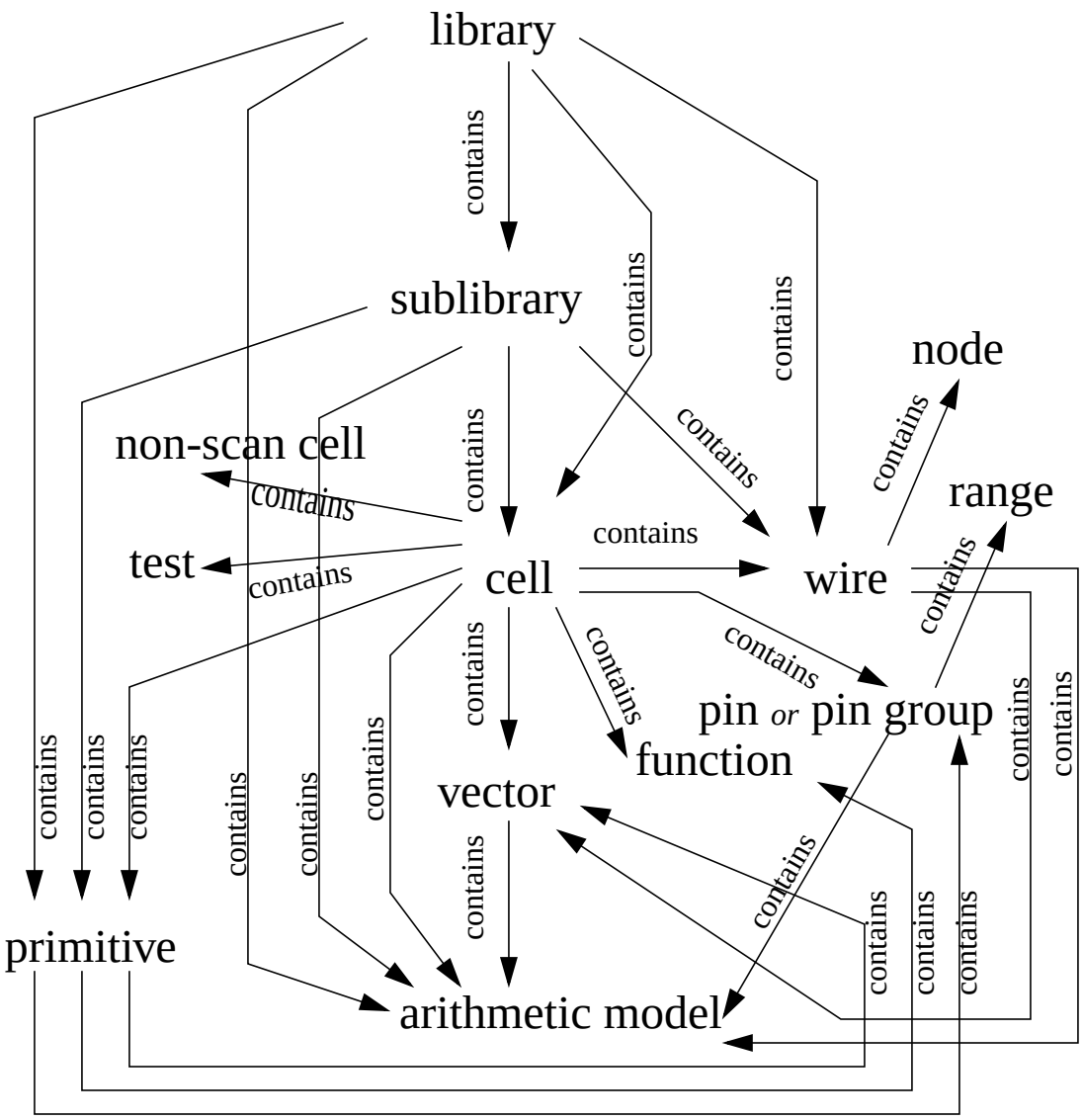



**Figure 3-9: Objects in a library for logical and electrical design and their relationships**

A library for functional and electrical design may contain sublibraries, cells, primitives, wires. Those cells which represent hierarchical blocks may also contain primitives and wires. Also, cells may contain pins, pin groups and vectors. Each object in the library may arithmetic models for electrical characteristics. In particular, electrical models which require a stimulus for characterization shall be in the context of a vector, which describes the stimulus.

Certain objects may also contain library-specific singular objects: A cell may contain function, test, and non-scan cell. A wire may contain node. A pin may contain range.

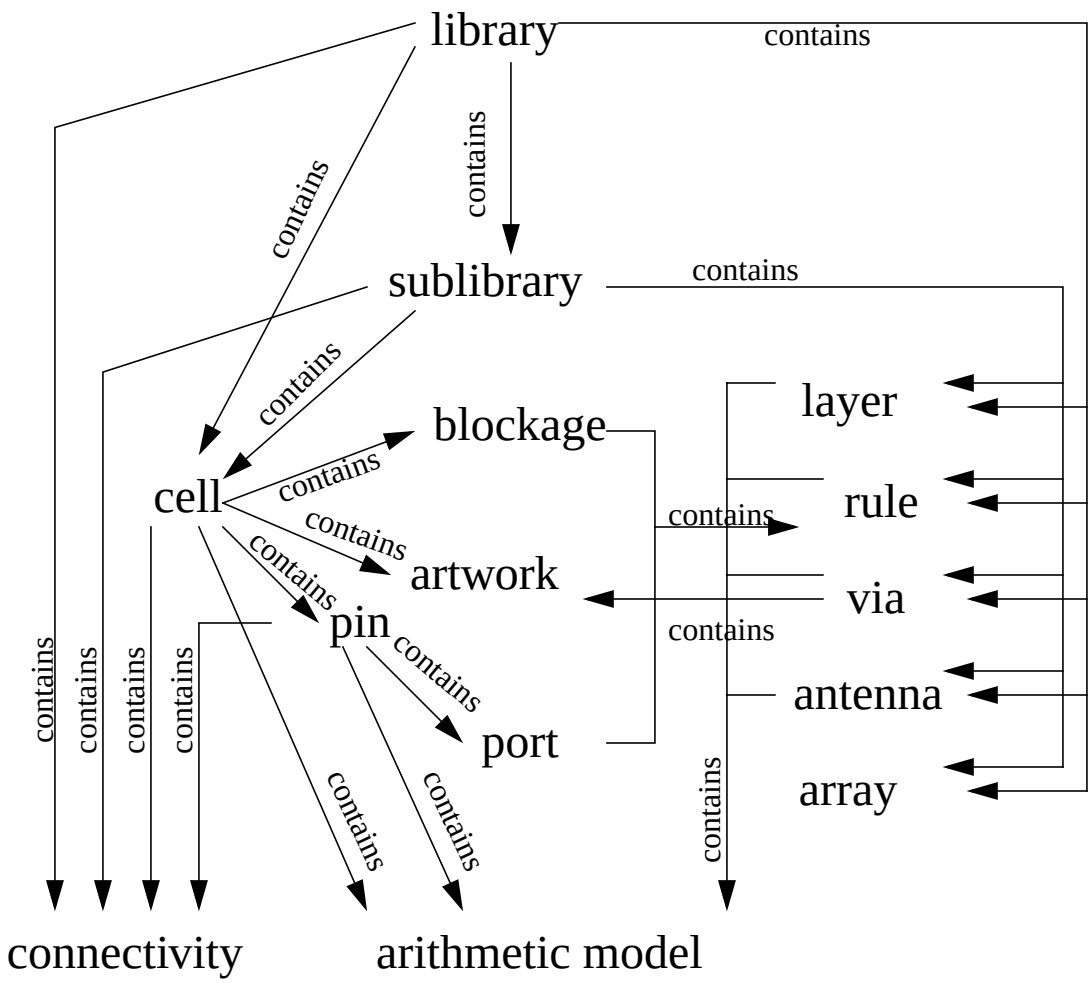



**Figure 3-10: Objects in a library for physical design and their relationships**

A library for physical design may contain sublibraries, cells, layers, vias, gneral rules, antenna rules, and arrays. Cells and vias may contain a reference to artwork. Cells may contain blockages. Pins may contain ports. almost each library may contain arithmetic models for physical characteristics. Library, sublibrary, cell and pin may also contein connectivity rules.

## 3.8 INFORMATION container

An `INFORMATION` container can be inside a `LIBRARY`, `SUBLIBRARY`, `CELL`, or `WIRE`. It can also be in `PRIMITIVE` objects inside a `LIBRARY` or `SUBLIBRARY`, but not in the locally defined primitives inside cells or functions. It can contain the annotations shown in Table 3-1.

**Table 3-1 : Information annotation container**

| Keyword | Value type | Description | Examples |
|---|---|---|---|
| VERSION | string | version of the object containing this INFORMATION block | “v1r3_2”<br>“1.3.2” |
| TITLE | string | title or comment related this object | “0.2u StdCell Library”<br>“2-input NAND, 4x drive”<br>“3-layer metal, best case, wireload model” |
| PRODUCT | string | product related to the object | “vsc1083”<br>“vsm10rs111”<br>“0.2u technology family” |
| AUTHOR | string | originator or modifier of the object | “user@system.com”<br>“Imn N. Gineer”<br>“An ASIC Vendor, Inc.” |
| DATETIME | string | date/time stamp related to the object | “Wed Aug 19 08:13:01 MST 1998”<br>“July 4, 1998” |

Example:

```
LIBRARY major_ASIC_vendor {
      INFORMATION {
            version = “v2.1.0”;
            title = “0.35 standard cell”;
            product = p35sc;
            author = “Major Asic Vendor, Inc.”;
            datetime = “Wed Jul 23 13:50:12 MST 1997”;
      }
}
```

## 3.9 Relations between objects

General referenceable objects within the scope of visibility are TEMPLATE and GROUP. Library-specific referenceable objects are PINs, PRIMITIVEs, and the arithmetic model. Figure 3-11 shows the relationships between these objects and where they can be referenced.

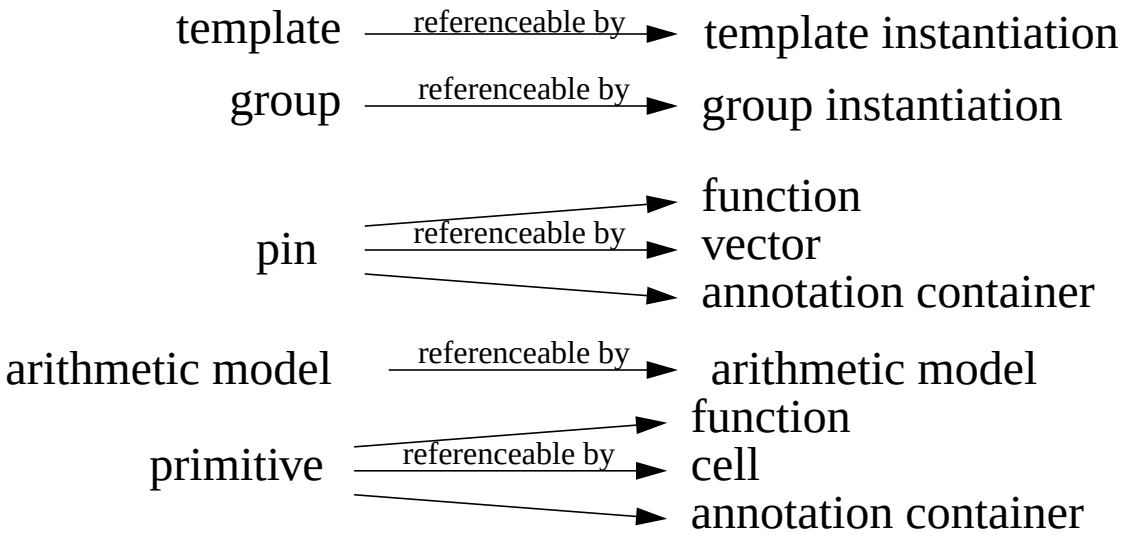



**Figure 3-11: Referencing rules for ALF objects**

The `TEMPLATE` and `GROUP` objects are referenceable only by their respective instantiation. The `TEMPLATE` definitions can contain instantiation of previously defined templates, which allows construction of reusable objects.

The arithmetic models can be referenced by other arithmetic models, if they are contained within each other. This allows hierarchical modeling and a mix of table- and equation-based models.

The `PIN` objects are referenced within `FUNCTION` and `VECTOR` objects and within any annotation container inside the same `CELL` object.

The `PRIMITIVE`s are referenceable by a `CELL`, to define pins and functionality, within a `FUNCTION`, to define functionality only, or within an annotation container, e.g., `SCAN`.

To use `PRIMITIVE`s and `PIN`s, see Section 5.6.1 and Section 5.3.7.

### 3.9.1 Keywords for referencing objects used as annotation

The object references hown in Table 3-2 can be used as annotations.

**Table 3-2 : Object references as annotation**

| Keyword | Value type | Description |
|---|---|---|
| `CELL` | `string` | reference to a declared CELL object |
| `PRIMITIVE` | `string` | reference to a declared PRIMITIVE object |
| `PIN` | `string` | reference to a declared PIN object |
| `CLASS` | `string` | reference to a declared CLASS object |

The syntax is:

```
object_keyword = string ;
```

### 3.9.2 Incremental definitions for VECTOR

In general, it is illegal to re-declare an ALF object (see Section 4.1, Rule 4). However, there are objects which merely define the context for other objects. When objects are incrementally added to the library, it is natural to re-declare the context as well. In this way, new objects can be added at the end of the library instead of being inserted somewhere in the middle within the already declared context. The classical example is the VECTOR object, which defines the context for timing and power models.

In a characterization process, timing models are always there in the first revision of a library, whereas power models are often added later. The new rule legalizes common practice within existing ALF libraries and tools. It makes it easier to add characterization data incrementally, e.g., power data. It also facilitates conversion from and to legacy library formats.

Multiple instances of the same VECTOR shall be legal for the purpose of incrementally adding children objects. The first instance of the VECTOR shall be interpreted as a declaration. All following instances shall be interpreted as supplemental definitions of the VECTOR. The rule of illegal re-declaration shall apply for the children objects within a VECTOR.

Example:

```
// the following is legal
VECTOR ( 01 A -> 01 Z ) {
   DELAY = 1 { FROM { PIN = A; } TO { PIN = Z; } }
}
VECTOR ( 01 A -> 01 Z ) {
   ENERGY = 25 ;
}

// the following is illegal
VECTOR ( 01 A -> 01 Z ) {
   DELAY = 1 { FROM { PIN = A; } TO { PIN = Z; } }
}
VECTOR ( 01 A -> 01 Z ) {
   DELAY = 2 { FROM { PIN = A; } TO { PIN = Z; } }
}
```

### 3.9.3 Other incremental definitions

Supplemental definitions of PROPERTY, ATTRIBUTE, LIBRARY, SUBLIBRARY shall also be legal.

# Section 4
# Library Organization

This section defines the scoping rules and use of multiple files within a library.

## 4.1 Scoping rules

The following scope rules shall apply to all library objects and its usage.

**Rule 1:** An object shall be defined before it is referenced.

**Rule 2:** An ALF object shall be known (referenceable) inside the parent object, inside all objects defined after that object within the same parent object, and inside all the children of those objects.

**Rule 3:** An object definition with only a keyword, but without an object identifier, implies the content of this definition shall be applied to all objects identified by this keyword at the current scope and the underlying levels of hierarchy.

Example:

```
LIBRARY my_library {
   CAPACITANCE {UNIT = pF;}     // default capacitance units for all
   ...                          // cells in my_library
   CELL cell1 {
      CAPACITANCE {UNIT = fF;}  // capacitance units specific to cell1
      PIN A {CAPACITANCE = 10.5;}
      ...
   }
   CELL cell2 {
      PIN A {CAPACITANCE = 0.010;} // default capacitance units
      ...
   }
}
```

Here, the capacitance of pin `A` of `cell1` is `10.5 fF`. The capacitance of pin `A` of `cell2` is `0.010 pF`.

**Rule 4:** An object shall not be defined again at the same level of scope. A definition of an object is considered duplicate, if both keyword and object identifier are identical.

Example:

It is illegal to write the following:

```
LIBRARY my_library {
   CAPACITANCE {UNIT = fF;}
   ...
   CELL cell1 {
      pin A {CAPACITANCE = 10.5;}
      ...
   }
   CAPACITANCE {UNIT = pF;}     // duplicate definition
   CELL cell2 {
      pin A {CAPACITANCE = 0.010;}
      ...
   }
}
```

There are three possible ways capacitance units can be set to `fF` for some of the cells in the library and `pF` for other cells in the same library:

1. Put each set of cells in a different sublibrary
2. Define templates for the different units and reference them appropriately
3. Define the units locally inside each cell

## 4.2 Use of multiple files

Sometimes it is inconvenient or impractical to include all of the data for a technology library in a single file. The *INCLUDE* keyword is used to compose a library from multiple files.

An `INCLUDE` statement can be used within any context, but any included file shall contain at least a valid object definition to be considered a legal ALF file. It needs to begin with a keyword, otherwise it can be ignored by a generic parser.

In general, the effect of using the `INCLUDE` statement is to be considered equivalent to inserting the contents of the included file at that point in the parent file.

For example, a top-level ALF library file can contain only the following statements, where each file contains appropriate data to make up the entire library.

```
LIBRARY mylib {
   INCLUDE “libdata.alf”;
   INCLUDE “templates.alf”;
   INCLUDE “cells.alf”;
   INCLUDE “wiremodels.alf”;
}
```

A complete ALF library definition shall begin with the `LIBRARY` keyword. A list of cell definitions shall not be considered a full, legal ALF library database.

# Section 5
# Functional Modeling

This chapter specifies the functional modeling for synthesis, formal verification, and simulation.

## 5.1 Combinational functions

This section defines the different types of combinational functions in ALF.

### 5.1.1 Combinational logic

Combinational logic can be described by continuous assignments of boolean values (*True* or *False*) to output variables as a function of boolean values of input variables. Such functions can be expressed in either boolean expression format or statetable format.

Let us consider an arbitrary continuous assignment

$z = f(a_1 .., .. a_n)$

In a dynamic or simulation context, the left-hand side (LHS) variable *z* is evaluated whenever there is a change in one of the right-hand side (RHS) variables $a_i$. No storage of previous states is needed for dynamic simulation of combinational logic.

### 5.1.2 Boolean operators on scalars

Table 5-1, Table 5-2, and Table 5-3 list unary, binary, and ternary boolean operators on scalars.

**Table 5-1 : Unary boolean operators**

| Operator | Description |
|---|---|
| !, ~ | logical inversion |

**Table 5-2 : Binary boolean operators**

| Operator | Description |
|---|---|
| &&, & | logical AND |
| \|\|, \| | logical OR |
| ~^ | logic equivalence (XNOR) |
| ^ | logic anti valence (XOR) |

**Table 5-3 : Ternary operator**

| Operator | Description |
|---|---|
| **?** | boolean condition operator for construction of combinational if-then-else clause |
| **:** | boolean else operator for construction of combinational if-then-else clause |

Combinational if-then-else clauses are constructed as follows:

```
<cond1>? <value1>: <cond2>? <value2>: <cond3>? <value3>: <default_value>
```

If `cond1` evaluates to boolean *True*, then `value1` is the result; else if `cond2` evaluates to boolean *True*, then `value2` is the result; else if `cond3` evaluates to boolean *True*, then `value3` is the result; else `default_value` is the result of this clause.

### 5.1.3 Boolean operators on words

Table 5-4 and Table 5-5 list unary and binary reduction operators on words (logic variables with one or more bits). The result of an expression using these operators shall be a logic value.

**Table 5-4 : Unary reduction operators**

| Operator | Description |
|---|---|
| **&** | AND all bits |
| **~&** | NAND all bits |
| **\|** | OR all bits |
| **~\|** | NOR all bits |
| **^** | XOR all bits |
| **~^** | XNOR all bits |

**Table 5-5 : Binary reduction operators**

| Operator | Description |
|---|---|
| **==** | equality for case comparison |
| **!=** | non-equality for case comparison |
| **>** | greater |
| **<** | smaller |
| **>=** | greater or equal |
| **<=** | smaller or equal |

Table 5-6 and Table 5-7 list unary and binary bitwise operators. The result of an expression using these operators shall be an array of bits.

**Table 5-6 : Unary bitwise operators**

| Operator | Description |
|---|---|
| ~ | bitwise inversion |

**Table 5-7 : Binary bitwise operators**

| Operator | Description |
|---|---|
| & | bitwise AND |
| \| | bitwise OR |
| ^ | bitwise XOR |
| ~^ | bitwise XNOR |

The following arithmetic operators, listed in Table 5-8, are also defined for boolean operations on words. The result of an expression using these operators shall be an extended array of bits.

**Table 5-8 : Binary operators**

| Operator | Description |
|---|---|
| << | shift left |
| >> | shift right |
| + | addition |
| - | subtraction |
| * | multiplication |
| / | division |
| % | modulo division |

The arithmetic operations addition, subtraction, multiplication, and division shall be *unsigned* if all the operands have the datatype *unsigned*. If any of the operands have the datatype signed, the operation shall be *signed*. See Table 6-25 for the `DATATYPE` definitions.

### 5.1.4 Operator priorities

The priority of binding operators to operands in boolean expressions shall be from strongest to weakest in the following order:

1. unary boolean operator (`!`, `~`, `&`, `~&`, `|`, `~|`, `^`, `~^`)
2. `XNOR` (`~^`), `XOR` (`^`), relational (`>`, `<`, `>=`, `<=`, `==`, `!=`), shift (`<<`, `>>`)
3. `AND` (`&`, `&&`), `NAND` (`~&`), multiply (`*`), divide (`/`), modulus (`%`)

4. OR (|, ||), NOR (~|), add (+), subtract (-)
5. ternary operators (?, :)

### 5.1.5 Datatype mapping

Logical operations can be applied to scalars and words. For that purpose, the values of the operands are reduced to a system of three logic values in the following way:

H has the logic value 1
L has the logic value 0
W, Z, U have the logic value X
A word has the logic value 1, if the unary OR reduction of all bits results in 1
A word has the logic value 0, if the unary OR reduction of all bits results in 0
A word has the logic value X, if the unary OR reduction of all bits results in X

Case comparison operations can also be applied to scalars and words. For scalars, they are defined in Table 5-9.

**Table 5-9 : Case comparison operators**

| A | B | A==B | A!=B | A>B | A<B |
|---|---|---|---|---|---|
| 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | H | 0 | 1 | X | X |
| 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | L | 0 | 1 | 1 | 0 |
| 1 | W, U, Z, X | 0 | 1 | X | 0 |
| H | 1 | 0 | 1 | X | X |
| H | H | 1 | 0 | 0 | 0 |
| H | 0 | 0 | 1 | 1 | 0 |
| H | L | 0 | 1 | 1 | 0 |
| H | W, U, Z, X | 0 | 1 | X | 0 |
| 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | H | 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | L | 0 | 1 | X | X |
| 0 | W, U, Z, X | 0 | 1 | 0 | X |
| L | 1 | 0 | 1 | 0 | 1 |
| L | H | 0 | 1 | 0 | 1 |
| L | 0 | 0 | 1 | X | X |
| L | L | 1 | 0 | 0 | 0 |
| L | W, U, Z, X | 0 | 1 | 0 | X |
| X | X | 1 | 0 | X | X |
| X | U | X | X | X | X |
| X | 0, 1, H, L, W, Z | 0 | 1 | X | X |

**Table 5-9 : Case comparison operators, *continued***

| A | B | A==B | A!=B | A>B | A<B |
|---|---|---|---|---|---|
| W | W | 1 | 0 | X | X |
| W | U | X | X | X | X |
| W | 0, 1, H, L, X, Z | 0 | 1 | X | X |
| Z | Z | 1 | 0 | X | X |
| Z | U | X | X | X | X |
| Z | 0, 1, H, L, X, W | 0 | 1 | X | X |
| U | 0, 1, H, L, X,W, Z, U | X | X | X | X |

For word operands, the operations > and < are performed after reducing all bits to the 3-value system first and then interpreting the resulting number according to the datatype of the operands. For example, if datatype is *signed*, `'b1111` is smaller than `'b0000`; if datatype is *unsigned*, `'b1111` is greater than `'b0000`. If two operands have the same value `'b1111` and a different datatype, the unsigned `'b1111` is greater than the signed `'b1111`.

The operations >= and <= are defined in the following way:

```
(a >= b) === (a > b) || (a == b)
(a <= b) === (a < b) || (a == b)
```

### 5.1.6 Rules for combinational functions

If a boolean expression evaluates *True*, the assigned output value is `1`. If a boolean expression evaluates *False*, the assigned output value is `0`. If the value of a boolean expression cannot be determined, the assigned output value is `X`. Assignment of values other than `1`, `0`, or `X` needs to be specified explicitly.

For evaluation of the boolean expression, input value `'bH` shall be treated as `'b1`. Input value `'bL` shall be treated as `'b0`. All other input values shall be treated as `'bX`.

Examples:

In equation form, these rules can be expressed as follows.

```
BEHAVIOR {
      Z = A;
}
```

is equivalent to

```
BEHAVIOR {
      Z = A ? 'b1 : 'b0;
}
```

More explicitly, this is also equivalent to

```
BEHAVIOR {
      Z = (A=='b1 || A=='bH)? 'b1 : (A=='b0 || A=='bL)? 'b0 : 'bX;
}
```

In table form, this can be expressed as follows:

```
STATETABLE {
      A     :     Z;
      ?     :     (A);
}
```

which is equivalent to

```
STATETABLE {
      A     :     Z;
      0     :     0;
      1     :     1;
}
```

More explicitly, this is also equivalent to

```
STATETABLE {
      A     :     Z;
      0     :     0;
      L     :     0;
      1     :     1;
      H     :     1;
      X     :     X;
      W     :     X;
      Z     :     X;
      U     :     X;
}
```

### 5.1.7 Concurrency in combinational functions

Multiple boolean assignments in combinational functions are understood to be concurrent. The order in the functional description does not matter, as each boolean assignment describes a piece of a logic circuit. This is illustrated in Figure 5-1.

```
BEHAVIOR {
      Q1 = <1st_boolean_expression(D1..Di)> ;
      ...
      Qn = <nth_boolean_expression(D1..Di)> ;
}
```

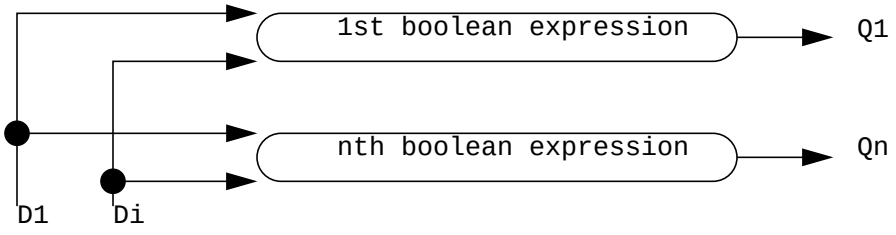



**Figure 5-1: Concurrency for combinational logic**

# 5.2 Sequential functions

This section defines the different types of sequential functions in ALF.

## 5.2.1 Level-sensitive sequential logic

In sequential logic, an output variable $z_j$ can also be a function of itself, i.e., of its previous state. The sequential assignment has the form

$$z_j = f(a_1 ..,.. a_n , z_1 ..,.. z_m)$$

The RHS cannot be evaluated continuously, since a change in the LHS as a result of a RHS evaluation shall trigger a new RHS evaluation repeatedly, unless the variables attain stable values. Modeling capabilities of sequential logic with continuous assignments are restricted to systems with oscillating or self-stabilizing behavior.

However, using the concept of *triggering conditions* for the LHS enables everything which is necessary for modeling *level-sensitive* sequential logic. The expression of a triggered assignment can look like this:

$$@\ g(b_1 ..,.. b_k)\ z_j = f(a_1 ..,.. a_n , z_1 ..,.. z_m)$$

The evaluation of *f* is activated whenever the *triggering function* $g$ is *True*. The evaluation of $g$ is self-triggered, i.e. at each time when an argument of $g$ changes its value. If *g* is a boolean expression like *f*, we can model all types of *level-sensitive sequential logic*.

During the time when *g* is *True*, the logic cell behaves exactly like combinational logic. During the time when *g* is *False*, the logic cell holds its value. Hence, one memory element per state bit is needed.

## 5.2.2 Edge-sensitive sequential logic

In order to model *edge-sensitive sequential logic*, notations for logical transitions and logical states are needed.

If the triggering function *g* is sensitive to logical transitions rather than to logical states, the function *g* evaluates to *True* only for an infinitely small time, exactly at the moment when the transition happens. The sole purpose of *g* is to trigger an assignment to the output variable through evaluation of the function *f* exactly at this time.

Edge-sensitive logic requires storage of the previous output state and the input state (to detect a transition). In fact, all implementations of edge-triggered flip-flops require at least two storage elements. For instance, the most popular flip-flop architecture features a master latch driving a slave latch.

Using transitions in the triggering function for value assignment, the functionality of a positive edge triggered flip-flop can be described as follows in ALF:

```
@ (01 CP) {Q = D;}
```

which reads “at rising edge of `CP`, assign `Q` the value of `D`”.

If the flip-flop also has an asynchronous direct clear pin (`CD`), the functional description consists of either two concurrent statements or two statements ordered by priority, as shown in Figure 5-2.

```
// concurrent style
@ (!CD) {Q = 0;}
@ (01 CP && CD) {Q = D;}
// priority (if-then-else) style
@ (!CD) {Q = 0;} : (01 CP) {Q = D;}
```

**Figure 5-2: Model of a flip-flop with asynchronous clear in ALF**

The following two examples show corresponding simulation models in Verilog and VHDL.

```
// full simulation model
always @(negedge CD or posedge CP) begin
   if ( ! CD ) Q <= 0;
   else if (CP && !CP_last_value) Q <= D;
   else Q <= 1'bx;
end
always @ (posedge CP or negedge CP) begin
   if (CP===0 | CP===1'bx) CP_last_value <= CP ;
end
// simplified simulation model for synthesis
always @(negedge CD or posedge CP) begin
   if ( ! CD ) Q <= 0;
   else Q <= D;
end
```

**Figure 5-3: Model of a flip-flop with asynchronous clear in Verilog**

```
// full simulation model

process (CP, CD) begin
   if (CD = '0') then
      Q <= '0';
   elsif (CP'last_value = '0' and CP = '1' and CP'event) then
      Q <= D;
   elsif (CP'last_value = '0' and CP = 'X' and CP'event) then
      Q <= 'X';
   elsif (CP'last_value = 'X' and CP = '1' and CP'event) then
      Q <= 'X';
   end if;
end process;

// simplified simulation model for synthesis

process (CP, CD) begin
   if (CD = '0') then
      Q <= '0';
   elsif (CP = '1' and CP'event) then
      Q <= D;
   end if;
end process;
```

**Figure 5-4: Model of a flip-flop with asynchronous clear in VHDL**

The following differences in modeling style can be noticed: VHDL and Verilog provide the list of sensitive signals at the beginning of the `process` or `always` block, respectively. The information of level-or edge-sensitivity shall be inferred by `if-then-else` statements inside the block. ALF shows the level-or-edge sensitivity as well as the priority directly in the triggering expression. Verilog has another particularity: The sensitivity list indicates whether at least one of the triggering signals is edge-sensitive by the use of `negedge` or `posedge`. However, it does not indicate which one, since either none or all signals shall have `negedge` or `posedge` qualifiers.

Furthermore, `posedge` is any transition with `0` as initial state *or* `1` as final state. A positive-edge triggered flip-flop shall be inferred for synthesis, yet this flip-flop shall only work correctly if both the initial state is `0` *and* the final state is `1`. Therefore, a simulation model for verification needs to be more complex than the model in the synthesizeable RTL code.

In Verilog, the extra non-synthesizeable code needs to also reproduce the relevant previous state of the clock signal, whereas VHDL has built-in support for `last_value` of a signal.

### 5.2.3 Unary operators for vector expressions

A transition operation is defined using unary operators on a scalar net. The scalar constants (see Figure 10-6) shall be used to indicate the start and end states of a transition on a scalar net.

```
bit bit          // apply transition from bit value to bit value
```

For example,

`01` is a transition from `0` to `1`.

No whitespace shall be allowed between the two scalar constants. The transition operators shown in Table 5-10 shall be considered legal.

**Table 5-10 : Unary vector operators on bits**

| Operator | Description |
|---|---|
| **01** | signal toggles from `0` to `1` |
| **10** | signal toggles from `1` to `0` |
| **00** | signal remains `0` |
| **11** | signal remains `1` |
| **0?** | signal remains `0` or toggles from `0` to arbitrary value |
| **1?** | signal remains `1` or toggles from `1` to arbitrary value |
| **?0** | signal remains `0` or toggles from arbitrary value to `0` |
| **?1** | signal remains `1` or toggles from arbitrary value to `1` |
| **??** | signal remains constant or toggles between arbitrary values |
| **0*** | a number of arbitrary signal transitions, including possibility of constant value, with the initial value `0` |
| **1*** | a number of arbitrary signal transitions, including possibility of constant value, with the initial value `1` |
| **?*** | a number of arbitrary signal transitions, including possibility of constant value, with arbitrary initial value |
| ***0** | a number of arbitrary signal transitions, including possibility of constant value, with the final value `0` |
| ***1** | a number of arbitrary signal transitions, including possibility of constant value, with the final value `1` |
| ***?** | a number of arbitrary signal transitions, including possibility of constant value, with arbitrary final value |

Unary operators for transitions can also appear in the `STATETABLE`.

Transition operators are also defined on words (and can appear the in `STATETABLE` as well):

`'`*`base word`* `'`*`base word`*

In this context, the transition operator shall apply transition from first word value to second word value.

For example,

`'hA'h5` is a transition of a 4-bit signal from `'b1010` to `'b0101`.

No whitespace shall be allowed between *base* and *word*.

The unary and binary operators for transition, listed in Table 5-11 and Table 5-12 respectively, are defined on bits and words.

**Table 5-11 : Unary vector operators on bits or words**

| Operator | Description |
|---|---|
| **?-** | no transition occurs |
| **??** | apply arbitrary transition, including possibility of constant value |
| **?!** | apply arbitrary transition, excluding possibility of constant value |
| **?~** | apply arbitrary transition with all bits toggling |

### 5.2.4 Basic rules for sequential functions

A sequential function is described in equation form by a boolean assignment with a condition specified by a boolean expression or a vector expression. If the condition evaluates to `1` (*True*), the boolean assignment is activated and the assigned output values follows the rules for combinational functions. If the vector expression evaluates to `0` (*False*), the output variables hold their assigned value from the previous evaluation.

For evaluation of a condition, the value `'bH` shall be treated as *True*, the value `'bL` shall be treated as *False*. All other values shall be treated as the unknown value `'bX`.

Example:

The following behavior statement

```
BEHAVIOR {
      @ (E) {Z = A;}
}
```

is equivalent to

```
BEHAVIOR {
      @ (E=='b1 || E=='bH) {Z = A;}
}
```

The following statetable statement, describing the same logic function

```
STATETABLE {
      E     A     :     Z;

      0     ?     :     (Z);
      1     ?     :     (A);
}
```

is equivalent to

```
STATETABLE {
      E     A     :     Z;

      0     ?     :     (Z);
      L     ?     :     (Z);
      1     ?     :     (A);
      H     ?     :     (A);
}
```

For edge-sensitive and higher-order event sensitive functions, transitions from or to `'bL` shall be treated like transitions from or to `'b0`, and transitions from or to `'bH` shall be treated like transitions from or to `'b1`.

Not every transition can trigger the evaluation of a function. The set of vectors triggering the evaluation of a function are called *active vectors*. From the set of active vectors, a set of *inactive vectors* can be derived, which shall clearly not trigger the evaluation of a function. There are is also a set of ambiguous vectors, which can trigger the evaluation of the function.

The set of active vectors is the set of vectors for which both observed states before and after the transition are known to be logically equivalent to the corresponding states defined in the vector expression.

The set of inactive vectors is the set of vectors for which at least one of the observed states before or after the transition is known to be not logically equivalent to the corresponding states defined in the vector expression.

Example:

For the following sequential function

```
@ (01 CP) { Z = A; }
```

the active vectors are

```
('b0'b1 CP)
('b0'bH CP)
('bL'b1 CP)
('bL'bH CP)
```

and the inactive vectors are

```
(’b1’b0 CP)
(’b1’bL CP)
(’b1’bX CP)
(’b1’bW CP)
(’b1’bZ CP)
(’bH’b0 CP)
(’bH’bL CP)
(’bH’bX CP)
(’bH’bW CP)
(’bH’bZ CP)
(’bX’b0 CP)
(’bX’bL CP)
(’bW’b0 CP)
(’bW’bL CP)
(’bZ’b0 CP)
(’bZ’bL CP)
(’bU’b0 CP)
(’bU’bL CP)
```

and the ambiguous vectors are

```
('b0'bX CP)
('b0'bW CP)
('b0'bZ CP)
('bL'bX CP)
('bL'bW CP)
('bL'bZ CP)
('bX'b1 CP)
('bW'b1 CP)
('bZ'b1 CP)
('bX'bH CP)
('bW'bH CP)
('bZ'bH CP)
('bX'bW CP)
('bX'bZ CP)
('bW'bX CP)
('bW'bZ CP)
('bZ'bX CP)
('bZ'bW CP)
('bU'bX CP)
('bU'bW CP)
('bU'bZ CP)
```

For vectors using exclusively based literals, the set of active vectors is the vector itself, the set of inactive vectors is any vector with at least one different literal, and the set of ambiguous vectors is empty.

Therefore, ALF does not provide a default behavior for ambiguous vectors, since the behavior for each vector can be explicitly defined in vectors using based literals.

### 5.2.5 Concurrency in sequential functions

The principle of concurrency applies also for edge-sensitive sequential functions, where the triggering condition is described by a vector expression rather than a boolean expression. In edge-sensitive logic, the target logic variable for the boolean assignment (LHS) can also be an operand of the boolean expression defining the assigned value (RHS). Concurrency implies that the RHS expressions are evaluated immediately *before* the triggering edge, and the values are assigned to the LHS variables immediately *after* the triggering edge. This is illustrated in Figure 5-5.

```
BEHAVIOR {
      @ ( <vector_expression(E1..Em)> ) {
            Q1 =
<1st_boolean_expression(D1..Di)> ;
            ...
            Qn =
<nth_boolean_expression(D1..Di)> ; } }
```

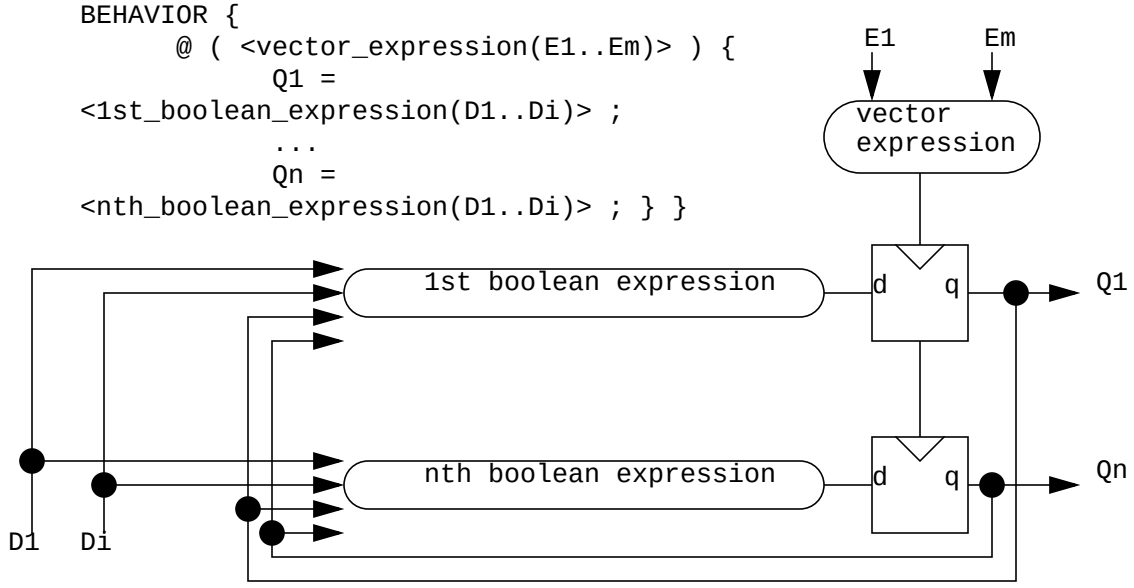



**Figure 5-5: Concurrency for edge-sensitive sequential logic**

Statements with multiple concurrent conditions for boolean assignments can also be used in sequential logic. In that case conflicting values can be assigned to the same logic variable. A default conflict resolution is not provided for the following reasons:

- Conflict resolution might not be necessary, since the conflicting situation is prohibited by specification.
- For different types of analysis (e.g., logic simulation), a different conflict resolution behavior might be desirable, while the physical behavior of the circuit shall not change. For instance, pessimistic conflict resolution always assigns X, more accurate conflict resolution first checks whether the values are conflicting. Different choices can be motivated by a trade-off in analysis accuracy and runtime.
- If complete library control over analysis is desired, conflict resolution can be specified explicitly.

Example:

```
BEHAVIOR {
      @ ( <condition_1> ) { Q = <value_1>; }
      @ ( <condition_2> ) { Q = <value_2>; }
}
```

Explicit pessimistic conflict resolution can be described as follows:

```
BEHAVIOR {
      @ ( <condition_1> && <condition_2>  ) { Q = 'bX; }
      @ ( <condition_1> && ! <condition_2>) { Q = <value_1>; }
      @ ( <condition_2> && ! <condition_1>) { Q = <value_2>; }
}
```

Explicit accurate conflict resolution can be described as follows:

```
BEHAVIOR {
      @ ( <condition_1> && <condition_2>  ) {
            Q = (<value_1>==<value_2>)? <value_1> : 'bX;
      }
      @ ( <condition_1> && ! <condition_2>) { Q = <value_1>; }
      @ ( <condition_2> && ! <condition_1>) { Q = <value_2>; }
}
```

Since the conditions are now rendered mutually exclusive, equivalent descriptions with priority statements can be used. They are more elegant than descriptions with concurrent statements.

```
BEHAVIOR {
      @ ( <condition_1> && <condition_2>  ) {
            Q = <conflict_resolution_value>;
      }
      : ( <condition_1> ) { Q = <value_1>; }
      : ( <condition_2> ) { Q = <value_2>; }
}
```

Given the various explicit description possibilities, the standard does not prescribe a default behavior. The model developer has the freedom of incomplete specification.

### 5.2.6 Initial values for logic variables

Per definition, all logic variables in a behavioral description have the initial value `U` which means “uninitialized”. This value cannot be assigned to a logic variable, yet it can be used in a behavioral description in order to assign other values than `U` after initialization.

Example:

```
BEHAVIOR {
      @ ( Q1 == 'bU ) { Q1 = 'b1 ; }
      @ ( Q2 == 'bU ) { Q2 = 'b0 ; }
      // followed by the rest of the behavioral description
}
```

A template can be used to make the intent more obvious, for example:

```
TEMPLATE VALUE_AFTER_INITIALIZATION {
      @ ( <logic_variable> == 'bU ) { <logic_variable> = <initial_value>
; }
}
BEHAVIOR {
      VALUE_AFTER_INITIALIZATION ( Q1 'b1' )
      VALUE_AFTER_INITIALIZATION ( Q2 'b0' )
      // followed by the rest of the behavioral description
}
```

Logic variables in a vector expression shall be declared as `PIN`s. It is possible to annotate initial values directly to a pin. Such variables shall never take the value `U`. Therefore vector expressions involving `U` for such variables (see the previous example) are meaningless.

Example:

```
PIN Q1 { INITIAL_VALUE = 'b1 ; }
PIN Q2 { INITIAL_VALUE = 'b0 ; }
```

# 5.3 Higher-order sequential functions

This section defines the different types of higher-order sequential functions in ALF.

## 5.3.1 Vector-sensitive sequential logic

Vector expressions can be used to model generalized higher order sequential logic; they are an extension of the boolean expressions. A *vector expression* describes sequences of logical events or transitions in addition to static logical states. A vector expression represents a description of a logical stimulus without timescale. It describes the order of occurrence of events.

The -> operator *(followed by*) gives a general capability of describing a sequence of events or a vector. For example, consider the following vector expression:

```
01 A -> 01 B
```

which reads “rising edge on `A` is followed by rising edge on `B`”.

A vector expression is evaluated by an event sequence detection function. Like a single event or a transition, this function evaluates *True* only at an infinitely short time when the event sequence is detected, as shown in Figure 5-6.

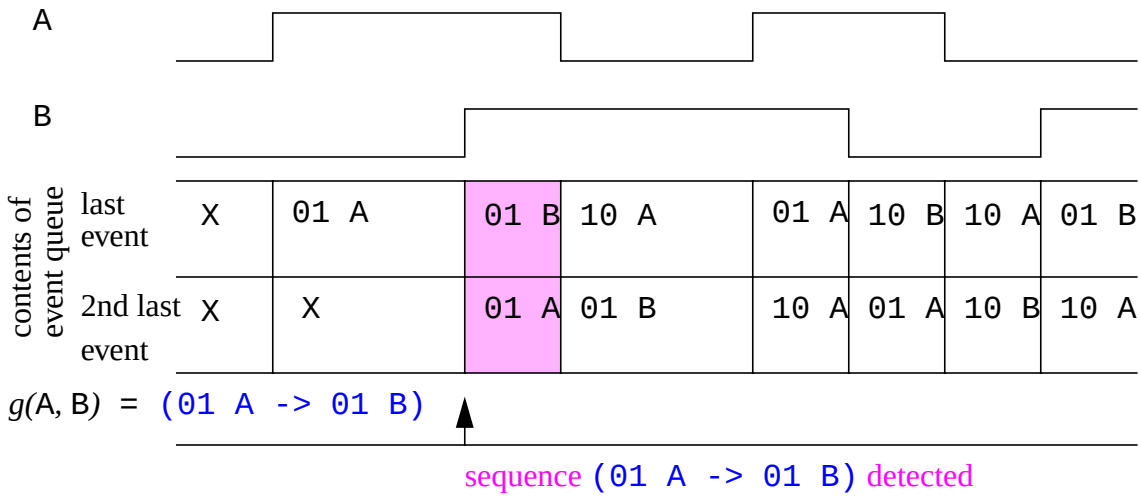



**Figure 5-6: Example of event sequence detection function**

The event sequence detection mechanism can be described as a queue that sorts events according to their order of arrival. The event sequence detection function evaluates *True* at exactly the time when a new event enters the queue and forms the required sequence, i.e., *the sequence specified by the vector expression* with its preceding events.

A vector-sensitive sequential logic can be called *(N+1) order sequential logic*, where *N* is the number of events to be stored in the queue. The implementation of *(N+1)* order sequential logic requires *N* memory elements for the event queue and one memory element for the output itself.

A sequence of events can also be gated with static logical conditions. In the example,

```
(01 CP -> 10 CP) && CD
```

the pin `CD` shall have `state 1` from some time before the rising edge at `CP` to some time after the falling edge of `CP`. The pin `CD` can not go low (`state 0`) after the rising edge of `CP` and go high again before the falling edge of `CP` because this would insert events into the queue and the sequence "rising edge on `CP` followed by falling edge on `CP`" would not be detected.

The formal calculation rules for general vector expressions featuring both states and transitions are detailed in Section 5.3.2 and Section 5.3.3.

The concept of vector expression supports functional modeling of devices featuring digital communication protocols with arbitrary complexity.

### 5.3.2 Canonical binary operators for vector expressions

The following canonical binary operators are necessary to define sequences of transitions:

- `vector_followed_by` for completely specified sequence of events
- `vector_and` for simultaneous events
- `vector_or` for alternative events
- `vector_followed_by` for incompletely specified sequence of events

The symbols for the boolean operators for `AND` and `OR` are overloaded for `vector_and` and `vector_or`, respectively. The new symbols for the `vector_followed_by` operators are shown in Table 5-12.

**Table 5-12 : Canonical binary vector operators**

| Operator | Operands | LHS, RHS commutative | Description |
|---|---|---|---|
| `->` | 2 vector expressions | no | Left-hand side (LHS) transition *is followed by* Right-hand side (RHS) transition, no transition can occur in-between |
| `&&, &` | 2 vector expressions | yes | LHS *and* RHS transition *occur simultaneously* |
| `\|\|, \|` | 2 vector expressions | yes | LHS *or* RHS transition *occur alternatively* |
| `~>` | 2 vector expressions | no | Left-hand side (LHS) transition *is followed by* Right-hand side (RHS) transition, other transitions can occur in-between |

Per definition, the `->` and `~>` operators shall not be commutative, whereas the `&&` and `||` operators on events shall be commutative.

```
01 a && 01 b === 01 b && 01 a
01 a || 01 b === 01 b || 01 a
```

The `->` and `~>` operators shall be freely associative.

```
01 a -> 01 b -> 01 c === (01 a -> 01 b) -> 01 c === 01 a -> (01 b -> 01 c)
01 a ~> 01 b ~> 01 c === (01 a ~> 01 b) ~> 01 c === 01 a ~> (01 b ~> 01 c)
```

The && operator is defined for single events and for event sequences with the same number of -> operators each.

```
(01 A1 .. -> ... 01 AN) & (01 B1 .. -> ... 01 BN)
===
01 A1 & 01 B1 ... -> ... 01 AN & 01 BN
```

The || operator reduces the set of edge operators (unary vector operators) to canonical and non-canonical operators.

```
(?? a)       === (?! a)||(?- a)  //a does or does not change its value
```

Hence ?? is non-canonical, since it can be defined by other operators.

If <value1><value2> is an edge operator consisting of two based literals value1 and value2 and word is an expression which can take the value value1 or value2, then the following vector expressions are considered equivalent:

```
<value1><value2> <word>
      ===         10 (<word> == <value1>) && 01 (<word> == <value2>)
      ===         01 (<word> != <value1>) && 01 (<word> == <value2>)
      ===         10 (<word> == <value1>) && 10 (<word> != <value2>)
      ===         01 (<word> != <value1>) && 10 (<word> != <value2>)
// all expressions describe the same event:
// <word> makes a transition from <value1> to <value2>
```

Hence vector expressions with edge operators using based literals can be reduced to vector expressions using only the edge operators 01 and 10.

### 5.3.3 Complex binary operators for vector expressions

Table 5-13 defines the complex binary operators for vector operators.

**Table 5-13 : Complex binary vector operators**

| Operator | Operands | LHS, RHS commutative | Description |
|---|---|---|---|
| **<->** | 2 vector expressions | yes | LHS transition follows or is followed by RHS transition |
| **&>** | 2 vector expressions | no | LHS transition *is followed by or occurs simultaneously* with RHS transition |
| **<&>** | 2 vector expressions | yes | LHS transition *follows or is followed by or occurs simultaneously* with RHS transition |

The following expressions shall be considered equivalent:

```
(01 a <-> 01 b) === (01 a -> 01 b)||(01 b -> 01 a)
(01 a  &> 01 b) === (01 a -> 01 b)||(01 a && 01 b)
(01 a <&> 01 b) === (01 a -> 01 b)||(01 b -> 01 a)||(01 a && 01 b)
```

By their symmetric definition, the <-> and <&> operators are commutative.

```
01 a <-> 01 b === 01 b <-> 01 a
01 a <&> 01 b === 01 b <&> 01 a
```

The commutative complex binary vector operators are defined in Table 5-12. The commutativity rules are only defined for two operands:

- commutative “followed by”:

```
vect_expr1 <-> vect_expr2 ===
   vect_expr1 -> vect_expr2 // vect_expr1 occurs first
|  vect_expr2 -> vect_expr1 // vect_expr2 occurs first
```

- commutative “followed by or simultaneously occurring”:

```
vect_expr1 <&> vect_expr2 ===
   vect_expr1 -> vect_expr2 // vect_expr1 occurs first
|  vect_expr2 -> vect_expr1 // vect_expr2 occurs first
|  vect_expr1 && vect_expr2 // both occur simultaneously
```

#### 5.3.3.1 Extension to N operands

This section defines how to use *N* operands.

A `complex_vector_expression` of the form

```
vector_expression { <-> vector_expression }
```

shall be commutative for all operands. The `complex_vector_expression` describes alternative event sequences in which the temporal order of each constituent `vector_expression` is completely permutable, excluding simultaneous occurrence of each constituent `vector_expression`.

A `complex_vector_expression` of the form

```
vector_expression { <&> vector_expression }
```

shall be commutative for all operands. The `complex_vector_expression` describes alternative event sequences in which the temporal order of each constituent `vector_expression` is completely permutable, including simultaneous occurrence of each constituent `vector_expression`.

Example:

```
01 A <-> 01 B <-> 01 C ===
   01 A -> 01 B -> 01 C
|  01 B -> 01 C -> 01 A
|  01 C -> 01 A -> 01 B
|  01 C -> 01 B -> 01 A
|  01 B -> 01 A -> 01 C
|  01 A -> 01 C -> 01 B
```

```
01 A <&> 01 B <&> 01 C ===
   01 A -> 01 B -> 01 C
|  01 B -> 01 C -> 01 A
|  01 C -> 01 A -> 01 B
|  01 C -> 01 B -> 01 A
|  01 B -> 01 A -> 01 C
|  01 A -> 01 C -> 01 B
|  01 A && 01 B -> 01 C
|  01 A -> 01 B && 01 C
|  01 B && 01 C -> 01 A
|  01 B -> 01 C && 01 A
|  01 C && 01 A -> 01 B
|  01 C -> 01 A && 01 B
|  01 A && 01 B && 01 C
```

#### 5.3.3.2 Boolean rules

The following rule applies for a boolean `AND` operation with three operands:

```
rule 1:
   A & B & C === (A & B) & C | A & (B & C)
```

A corresponding rule also applies to the commutative followed-by operation with three operands:

```
rule 2:
   01 A <-> 01 B <-> 01 C ===
      (01 A <-> 01 B) <-> 01 C
   |  01 A <-> (01 B <-> 01 C)
```

The alternative boolean expressions `(A & B) & C` and `A & (B & C)` in rule 1 are equivalent. Therefore, rule 1 can be reduced to the following:

```
rule 3:
   A & B & C === (A & B) & C === (B & C) & A
```

A corresponding rule does *not* apply to complex vector operands, since each expression with associated operands generates only a subset of permutations:

```
(01 A <-> 01 B) <-> 01 C ===
      ((01 A <-> 01 B) -> 01 C)
|     (01 C -> (01 A <-> 01 B)) ===
      01 A -> 01 B -> 01 C
|     01 B -> 01 A -> 01 C
|     01 C -> 01 A -> 01 B
|     01 C -> 01 B -> 01 A
```

The permutations

```
01 A -> 01 C -> 01 B
01 B -> 01 C -> 01 A
```

are missing.

```
01 A <-> (01 B <-> 01 C) ===
   (01 A -> (01 B <-> 01 C))
|  ((01 B <-> 01 C) -> 01 A) ===
   01 A -> 01 B -> 01 C
|  01 A -> 01 C -> 01 B
|  01 B -> 01 C -> 01 A
|  01 C -> 01 B -> 01 A
```

The permutations

```
|  01 B -> 01 A -> 01 C
|  01 C -> 01 A -> 01 B
```

are missing.

### 5.3.4 Operators for conditional vector expressions

The definitions of the `&&`, `?`, and `:` operators are also overloaded to describe a *conditional vector expression* (involving boolean expressions and vector expressions), as shown in Table 5-14. The clauses are boolean expressions; while vector expressions are subject to those clauses.

**Table 5-14 : Operators for conditional vector expressions**

| Operator | Operands | LHS, RHS commutative | Description |
|---|---|---|---|
| **&&, &** | 1 vector expression, 1 boolean expression | yes | boolean expression (LHS or RHS) is *True* while sequence of transitions, defined by vector expression (RHS or LHS) occurs |
| **?** | 1 vector expression, 1 boolean expression | no | boolean condition operator for construction of if-then-else clause involving vector expressions |
| **:** | 1 vector expression, 1 boolean expression | no | boolean else operator for construction of if-then-else clause involving vector expressions |

An example for conditional vector expression using `&&` is given below:

```
(01 a && !b)                    // a rises while b==0
```

The order of the operands in a conditional vector expression using `&&` shall not matter.

```
<vector_exp> && <boolean_exp> === <boolean_exp> && <vector_exp>
```

The `&&` operator is still commutative in this case, although one operand is a boolean expression defining a static state, the other operand is a vector expression defining an event or a sequence of events. However, since the operands are distinguishable per se, it is not necessary to impose a particular order of the operands.

An example for conditional vector expression using `?` and `:` is given below.

```
!b ? 01 a : c ? 10 b : 01 d
===
!b & 01 a | !(!b) & c & 10 b | !(!b) & !c & 01 d
```

This example shows how a conditional vector expression using ternary operators can be expressed with alternative conditional vector expressions.

A conditional vector expression can be reduced to a non-conditional vector expression in some cases (see Section 5.4.11).

Every binary vector operator can be applied to a conditional vector expression.

### 5.3.5 Operators for sequential logic

Table 5-15 defines the complex binary operators for vector operators.

**Table 5-15 : Operators for sequential logic**

| Operator | Description |
|---|---|
| @ | sequential `if` operator, followed by a boolean logic expression (for level-sensitive assignment) or by a vector expression (for edge-sensitive assignment) |
| : | sequential `else if` operator, followed by a boolean logic expression (for level-sensitive assignment) or by a vector expression (for edge-sensitive assignment) with lower priority |

Sequential assignments are constructed as follows:

```
@ ( <trigger1> ) { <action1> } : ( <trigger2> ) { <action2> } :
  ( <trigger3> ) { <action3> }
```

If `trigger1` event is detected, then `action1` is performed; else if `trigger2` event is detected, then `action2` is performed; else if `trigger3` event is detected, then `action3` is performed as a result of this clause.

### 5.3.6 Operator priorities

The priority of binding operators to operands in non-conditional vector expressions shall be from strongest to weakest in the following order:

1. unary vector operators (edge literals)
2. complex binary vector operators (`<->`, `&>`, `<&>`)
3. vector `AND` (`&`, `&&`)
4. vector_followed_by operators (`->`, `~>`)
5. vector `OR` (`|`, `||`)

### 5.3.7 Using PINs in VECTORs

A `VECTOR` defines state, transition, or sequence of transitions of pins that are controllable and observable for characterization.

Within a CELL, the set of PINs with SCOPE=behavior or SCOPE=measure or SCOPE=both is the default set of variables in the event queue for vector expressions relevant for BEHAVIOR or VECTOR statements or both, respectively.

For detection of a sequence of transitions it is necessary to observe the set of variables in the event queue. For instance, if the set of pins consists of A, B, C, D, the vector expression

```
(01 A -> 01 B)
```

implies no transition on A, B, C, D occurs between the transitions 01 A and 01 B.

The default set of pins applies only for vector expressions without conditions. The conditional event AND operator limits the set of variables in the event queue. In this case, only the state of the condition and the variables appearing in the vector expression are observed.

Example:

```
(01 A -> 01 B) && (C | D)
```

No transition on A, B occurs between 01 A and 01 B, and (C | D) needs to stay *True* in-between 01 A and 01 B as well. However, C and D can change their values as long as (C | D) is satisfied.

## 5.4 Modeling with vector expressions

Vector expressions provide a formal language to describe digital waveforms. This capability can be used for functional specification, for timing and power characterization, and for timing and power analysis.

In particular, vector expressions add value by addressing the following modeling issues:

- *Functional specification*: complex sequential functionality, e.g., bus protocols.
- *Timing analysis*: complex timing arcs and timing constraints involving more than two signals.
- *Power analysis*: temporal and spatial correlation between events relevant for power consumption.
- *Circuit characterization and test*: specification of characterization and/or test vectors for particular timing, power, fault, or other measurements within a circuit.

Like boolean expressions, vector expressions provide the means for describing the functionality of digital circuits in various contexts without being self-sufficient. Vector expressions enrich this functional description capability by adding a “dynamic” dimension to the otherwise “static” boolean expressions.

The following subsections explain the semantics of vector expressions step-by-step. The vector expression concept is explained using terminology from simulation event reports. However, the application of vector expressions is not restricted to post-processing event reports.

Some application tools (e.g., power analysis tools) can actually evaluate vector expressions during post-processing of event reports from simulation. Other application tools, especially simulation model generators, need to respect the causality between the triggering events and the actions to be triggered. While it is semantically impossible to describe cause and effect in the same vector expression for the purpose of functional modeling, both cause and effect can appear in a vector expression used for a timing arc description.

ALF does not make assumption about the physical nature of the event report. Vector expressions can be applied to an actual event report written in a file, to an internal event queue within a simulator, or to a hypothetical event report which is merely a mathematical concept.

### 5.4.1 Event reports

This section describes the terminology of event reports from simulation, which is used to explain the concept of ALF vector expressions. The intent of ALF vector expressions is not to *replace* existing event report formats. Non-pertinent details of event report formats are not described here.

Simulation events (e.g., from Verilog or VHDL) can be reported in a value change dump (VCD) file, which has the following general form:

```
<time1>
      <variableA> <stateU>
      <variableB> <stateV>
      ...
<time2>
      <variableC> <stateW>
      <variableD> <stateX>
      ...
<time3> ...
```

The set of variables for which simulation events are reported, i.e., the *scope* of the event report needs to be defined beforehand. Each variable also has a definition for the *set of states* it can take. For instance, there can be binary variables, 16-bit integer variables, 1-bit variables with drive-strength information, etc. Furthermore, the initial state of each variable shall be defined as well. In an ALF context, the terms *signal* and *variable* are used interchangeably. In VHDL, the corresponding term is *signal*. In Verilog, there is no single corresponding term. All `input`, `output`, `wire`, and `reg` variables in Verilog correspond to a `signal` in VHDL.

The time values `<time1>`, `<time2>`, `<time3>`, etc. shall be in increasing order. The order in which simultaneous events are reported does not matter. The number of time points and the number of simultaneous events at a certain time point are unlimited.

In the physical world, each event or change of state of a variable takes a certain amount of time. A variable cannot change its state more than once at a given point in time. However, in simulation, this time can be smaller than the resolution of the time scale or even zero (`0`). Therefore, a variable can change its state more than once at a given point in simulation time. Those events are, strictly speaking, not simultaneous. They occur in a certain order, separated by an infinitely small delta-time. Multiple simultaneous events of the same variable are not reported in the VCD. Only the final state of each variable is reported.

A VCD file is the most compact format that allows reconstruction of entire waveforms for a given set of variables. A more verbose form is the test pattern format.

```
<TIME>  <variableA> <variableB> <variableC> <variableD>
<time1> <stateU>    <stateV>    ...         ...
<time2> <stateU>    <stateV>    <stateW>    <stateX>
<time3> ...         ...         ...         ...
```

The test pattern format reports the state of each variable at every point in time, regardless of whether the state has changed or not. Previous and following states are immediately available in the previous and next row, respectively. This makes the test pattern format more readable than the VCD and well-suited for taking a snapshot of events in a time window.

An example of an event report in VCD format:

```
// initial values
A 0   B 1   C 1   D X   E 1
// event dump
109   A 1   D 0
258   B 0
573   C 0
586   A 0
643   A 1
788   A 0   B 1   C 1
915   A 1
1062  E 0
1395  B 0   C 0
1640  A 0   D 1
// end of event dump
```

An example of an event report in test pattern format:

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     0     1     0     1
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     1
1062  1     1     1     0     0
1395  1     0     0     0     0
1640  0     0     0     1     0
```

Both VCD and test pattern formats represent the same amount of information and can be translated into each other.

### 5.4.2 Event sequences

For specification of a functional waveform (e.g., the write cycle of a memory), it is not practical to use an event report format, such as a VCD or test pattern format. In such waveforms, there is no absolute time. And the relative time, for example, the setup time between address change and write enable change, can vary from one instance to the other.

The main purpose of `vector_expressions` is waveform specification capability. The following operators can be used:

- `vector_unary` (also called *edge operator* or *unary vector operator*)

  The edge operator is a prefix to a variable in a vector expression. It contains a pair of states, the first being the previous state, the second being the new state. Edge operators can describe a change of state or no change of state.

- vector_and (also called *simultaneous event operator*)
  This operator uses the overloaded symbol & or && interchangeably. The & operator is the separator between simultaneously occurring events
- vector_followed_by (also called *followed-by operator*)
  The "immediately followed-by operator" using the symbol -> is treated first. The -> operator is the separator between consecutively occurring events.

These operators are necessary and sufficient to describe the following subset of vector_expressions:

- vector_single_event
  A change of state in a single variable, for example:
  ```
  01 A
  ```
- vector_event
  A simultaneous change of state in one or more variables, for example:
  ```
  01 A & 10 B
  ```
- vector_event_sequence
  Subsequently occurring changes of state in one or more variables, for example:
  ```
  01 A & 10 B -> 10 A
  ```

The vector_and operator has a higher binding priority than the vector_followed_by operator.

We can now express the pattern of the sample event report in a vector_event_sequence expression:

```
01 A & X0 D -> 10 B -> 10 C -> 10 A -> 01 A
-> 10 A & 01 B & 01 C -> 01 A -> 10 E -> 10 B & 10 C -> 10 A & 01 D
```

We can define the *length* of a vector_event_sequence expression as the number of subsequent events described in the vector_event_sequence expression. The length is equal to the number of -> operators plus one (1).

Although the vector expression format contains an inherent redundancy, since the old state of each variable is always the same as the new state of the same variable in a previous event, it is more human-readable, especially for waveform description. On the other hand, it is more compact than the test pattern format. For short event sequences, it is even more compact than the VCD, since it eliminates the declaration of initial values. To be accurate, for variables with exactly one event the vector expression is more compact than the VCD. For variables with more than one event the VCD is more compact than the vector expression. In summary, the vector expression format offers readability similar to the test pattern format and compactness close to the VCD format.

### 5.4.3 Scope and content of event sequences

The *scope* applicable to a vector expression defines the set of variables in the event report. The *content* of a vector expression is the set of variables that appear in the vector expression itself. The content of a vector expression shall be a subset of variables within scope.

- PINs with the annotation SCOPE = BEHAVIOR are applicable variables for vector expressions within the context of BEHAVIOR.

- `PIN`s with the annotation `SCOPE = MEASURE` are applicable variables for vector expressions within the context of `VECTOR`.
- `PIN`s with the annotation `SCOPE = BOTH` are applicable variables for all vector expressions.

A `vector_event_sequence` expression is an event pattern without time, containing only the variables within its own content. This event pattern is evaluated against the event report containing all variables within scope. The vector expression is *True* when the event pattern matches the event report.

Example:

```
time  A     B     C     D     E     // scope is A, B, C, D, E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     0     1     0     1
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     1
1062  1     1     1     0     0
1395  1     0     0     0     0
1640  0     0     0     1     0
```

Consider the following vector expressions in the context of the sample event report:

```
01 A                                          //(1)  content is A

//event pattern expressed by (1):
//    A
//    0
//    1
```

`(1)` is *True* at time 109, time 643, and time 915.

```
10 B -> 10 C                                  //(2)  content is B, C

//event pattern expressed by (2):
//    B     C
//    1     1
//    0     1
//    0     0
```

`(2)` is *True* at time 573.

```
10 A -> 01 A                                  //(3)  content is A

//event pattern expressed by (3):
//    A
//    1
//    0
//    1
```

`(3)` is *True* at time 643 and time 915.

```
01 D                                          //(4)  content is D
```

```
//event pattern expressed by (4):
//    D
//    0
//    1
```

`(4)` is *True* at time 1640.

```
01 A -> 10 C                                    //(5)  content is A, C

//event pattern expressed by (5):
//    A     C
//    0     1
//    1     1
//    1     0
```

`(5)` is not be *True* at any time, since the event pattern expressed by `(5)` does not match the event report at any time.

### 5.4.4 Alternative event sequences

The following operator can be used to describe alternative events:

`vector_or`, also called *event-or operator* or *alternative-event operator*, using the overloaded symbol `|` or `||` interchangeably. The `|` operator is the separator between alternative events or alternative event sequences.

In analogy to boolean operators, `|` has a lower binding priority than `&` and `->`. Parentheses can be used to change the binding priority.

Example:

```
(01 A -> 01 B) | 10 C === 01 A -> 01 B | 10 C
01 A -> (01 B | 10 C) === 01 A -> 01 B | 01 A -> 10 C
```

Consider the following vector expressions in the context of the sample event report:

```
01 A | 10 C                                     //(6)

//event pattern expressed by (6):
//    A
//    0
//    1

//alternative event pattern expressed by (6):
//    C
//    1
//    0
```

`(6)` is *True* at time 109, time 573, time 643, time 915, and time 1395.

```
10 B -> 10 C | 10 A -> 01 A                     //(7)

//event pattern expressed by (7):
//    B     C
//    1     1
//    0     1
//    0     0
```

```
//alternative event pattern expressed by (7):
//    A
//    1
//    0
//    1
```

(7) is *True* at time 573, time 643, and time 915.

```
01 D | 10 B -> 10 C                                    //(8)

//event pattern expressed by (8):
//    D
//    0
//    1

//alternative event pattern expressed by (8):
//    B     C
//    1     1
//    0     1
//    0     0
```

(8) is *True* at time 573 and time 1640.

```
10 B -> 10 C | 10 A                                    //(9)

//event pattern expressed by (9):
//    B     C
//    1     1
//    0     1
//    0     0

//alternative event pattern expressed by (9):
//    A
//    1
//    0
```

(9) is *True* at time 573, time 586, time 788, and time 1640.

The following operators provide a more compact description of certain alternative event sequences:

- &> events occur simultaneously or follow each other in the order RHS after LHS
- <-> a LHS event followed by a RHS event or a RHS event followed by a LHS event
- <&> events occur simultaneously or follow each other in arbitrary order

Example:

```
01 A &> 01 C      ===   01 A & 01 C | 01 A -> 01 C
01 A <-> 01 C     ===   01 A -> 01 C | 01 C -> 01 A
01 A <&> 01 C     ===   01 A <-> 01 C | 01 A & 01 C
```

The binding priority of these operators is higher than of & and ->.

### 5.4.5 Symbolic edge operators

Alternative events of the same variable can be described in a even more compact way through the use of edge operators with symbolic states. The symbol ? stands for “any state”.

- edge operator with ? as the previous state:
  transition from any state to the defined new state

- edge operator with `?` as the next state:
  transition from the defined previous state to any state.

Both edge operators include the possibility no transition occurred at all, i.e., the previous and the next state are the same. This situation can be explicitly described with the following operator:

edge operator with next state = previous state, also called *non-event operator*
The operand stays in the state defined by the operator.

The following symbolic edge operators also can be used:

- `?-` no transition on the operand
- `?!` transition from any state to any state different from the previous state
- `??` transition from any state to any state or no transition on the operand
- `?~` transition from any state to its bitwise complementary state

Example: Let `A` be a logic variable with the possible states `1`, `0`, and `X`.

```
?0 A === 00 A | 10 A | X0 A
?1 A === 01 A | 11 A | X1 A
?X A === 0X A | 1X A | XX A
0? A === 00 A | 01 A | 0X A
1? A === 10 A | 11 A | 1X A
X? A === X0 A | X1 A | XX A
?! A === 01 A | 0X A | 10 A | 1X A | X0 A | X1 A
?~ A === 01 A | 10 A | XX A
?? A === 00 A | 01 A | 0X A | 10 A | 11 A | 1X A | X0 A | X1 A | XX A
?- A === 00 A | 11 A | XX A
```

For variables with more possible states (e.g., logic states with different drive strength and multiple bits) the explicit description of alternative events is quite verbose. Therefore the symbolic edge operators are useful for a more compact description.

This completes the set of `vector_binary` operators necessary for the description of a subset of `vector_expressions` called `vector_complex_event` expressions. All `vector_binary` operators have two `vector_complex_event` expressions as operands. The set of `vector_event_sequence` expressions is a subset of `vector_complex_event` expressions. Every `vector_complex_event` expression can be expressed in terms of alternative `vector_event_sequence` expressions. The latter could be called *minterms*, in analogy to boolean algebra.

### 5.4.6 Non-events

A `vector_single_event` expression involving a non-event operator is called a *non-event*. A rigorous definition is required for `vector_complex_event` expressions containing non-events. Consider the following example of a flip-flop with clock input `CLK` and data output `Q`.

```
01 CLK -> 01 Q    // (i)
01 CLK -> 00 Q    // (ii)
```

The vector expression `(i)` describes the situation where the output switches from `0` to `1` after the rising edge of the clock. The vector expression `(ii)` describes the situation where the output remains at `0` after the rising edge of the clock.

How is it possible to decide whether `(i)` or `(ii)` is *True*, without knowing the delay between `CLK` and `Q`? The only way is to wait until any event occurs after the rising edge of `CLK`. If the event is not on `Q` and the state of `Q` is `0` during that event, then `(ii)` is *True*.

Hence, a non-event is *True* every time when another event happens and the state of the variable involved in the non-event satisfies the edge operator of the non-event.

Example:

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     0     1     0     1
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     1
1062  1     1     1     0     0
1395  1     0     0     0     0
1640  0     0     0     1     0
```

The test pattern format represents an event, for example `01 A`, in no different way than a non-event, for example `11 E`. This non-event is *True* at times 109, 258, 573, 586, 643, 788, and 915; in short, every time when an event happens while `E` is constant `1`.

### 5.4.7 Compact and verbose event sequences

A `vector_event_sequence` expression in a compact form can be transformed into a verbose form by padding up every `vector_event` expression with non-events. The next state of each variable within a `vector_event` expression shall be equal to the previous state of the same variable in the subsequent `vector_event` expression.

Example:

```
01 A -> 10B === 01 A & 11 B -> 11 A & 10 B
```

A vector expression for a complete event report in compact form resembles the VCD, whereas the verbose form looks like the test pattern.

```
// compact form
01 A & X0 D -> 10 B -> 10 C -> 10 A -> 01 A
-> 10 A & 01 B & 01 C -> 01 A -> 10 E
-> 10 B & 10 C -> 10 A & 01 D
===
// verbose form
?0 A & ?1 B & ?1 C & ?X D & ?1 E->
01 A & 11 B & 11 C & X0 D & 11 E->
11 A & 10 B & 11 C & 00 D & 11 E->
11 A & 00 B & 10 C & 00 D & 11 E->
```

```
10 A & 00 B & 00 C & 00 D & 11 E->
01 A & 00 B & 00 C & 00 D & 11 E->
10 A & 01 B & 01 C & 00 D & 11 E->
01 A & 11 B & 11 C & 00 D & 11 E->
11 A & 11 B & 11 C & 00 D & 10 E->
11 A & 10 B & 10 C & 00 D & 00 E->
10 A & 00 B & 00 C & 01 D & 00 E
```

The transformation rule needs to be slightly modified in case the compact form contains a `vector_event` expression consisting only of non-events. By definition, the non-event is *True* only if a real event happens simultaneously with the non-event. Padding up a `vector_event` expression consisting of non-events with other non-events make this impossible. Rather, this `vector_event` expression needs to be padded up with unspecified events, using the `??` operator. Eventually, unspecified events can be further transformed into partly specified events, if a former or future state of the involved variable is known.

Example:

```
01 A -> 00 B
=== 01 A & 00 B -> ?? A & 00 B
```

In the first transformation step, the unspecified event `?? A` is introduced.

```
01 A & 00 B -> ?? A & 00 B
=== 01 A & 00 B -> 1? A & 00 B
```

In the second step, this event becomes partly specified. `?? A` is bound to be `1? A` due to the previous event on `A`.

### 5.4.8 Unspecified simultaneous events within scope

Variables which are within the scope of the vector expression yet do not appear in the vector expression, can be used to pad up the vector expression with unspecified events as well. This is equivalent to omitting them from the vector expression.

Example:

```
01 A -> 10 B // let us assume a scope containing A, B, C, D, E
===
01 A & 10 B & ?? C & ?? D & ?? E -> 11 A & 10 B & ?? C & ?? D & ?? E
```

This definition allows unspecified events to occur *simultaneously* with specified events or specified non-events. However, it disallows unspecified events to occur *in-between* specified events or specified non-events.

At first sight, this distinction seems to be arbitrary. Why not disallow unspecified events altogether? Yet there are several reasons why this definition is practical.

If a vector expression disallows simultaneously occurring unspecified events, the application tool has the burden not only to match the pattern of specified events with the event report but also to check whether the other variables remain constant. Therefore, it is better to specify this extra pattern matching constraint explicitly in the vector expression by using the `?-` operator.

There are many cases where it actually does not matter whether simultaneously occurring unspecified events are allowed or disallowed:

- *Case 1*: Simultaneous events are impossible by design of the flip-flop. For instance, in a flip-flop it is impossible for a triggering clock edge `01 CK` and a switch of the data output `? Q` to occur at the same time. Therefore, such events can not appear in the event report. It makes no difference whether `01 CK & ?- Q`, `01 CK & ?? Q`, or `01 CK` is specified. The only occurring event pattern is `01 CK & ?- Q` and this pattern can be reliably detected by specifying `01 CK`.
- *Case 2*: Simultaneous events are prohibited by design. For instance, in a flip-flop with a positive setup time and positive hold time, the triggering clock edge `01 CK` and a switch of the data input `?! D` is a timing violation. A timing checker tool needs the violating pattern specified explicitly, i.e., `01 CK & ?! D`. In this context, it makes sense to specify the non-violating pattern also explicitly, i.e., `01 CK & ?- D`. The pattern `01 CK` by itself is not applicable.
- *Case 3*: Simultaneous events do not occur in correct design. For instance, power analysis of the event `01 CK` needs no specification of `?! D` or `?- D`. In the analysis of an event report with timing violations, the power analysis is less accurate anyway. In the analysis of the event report for the design without timing violation, the only occurring event pattern is `01 CK & ?- D` and this pattern can be reliably detected by specifying `01 CK`.[1]
- *Case 4*: The effects of simultaneous events are not modeled accurately. This is the case in static timing analysis and also to some degree in dynamic timing simulation. For instance, a `NAND` gate can have the inputs `A` and `B` and the output `Z`. The event sequence exercising the timing arc `01 A -> 10 Z` can only happen if `B` is constant `1`. No event on `B` can happen in-between `01 A` and `10 Z`. Likewise, the timing arc `01 B -> 10 Z` can only happen if `A` is constant `1` and no event happens in-between `01 B` and `10 Z`. The timing arc with simultaneously switching inputs is commonly ignored. A tool encountering the scenario `01 A & 01 B -> 10 Z` has no choice other than treating it arbitrarily as `01 A -> 10 Z` or as `01 B -> 10 Z`.
- *Case 5*: The effects of simultaneous events are modeled accurately. Here it makes sense to specify all scenarios explicitly, e.g., `01 A & ?- B -> 10 Z`, `01 A &?! B -> 10 Z`, `?- A & 01 B -> 10 Z`, etc., whereas the patterns `01 A -> 10 Z` and `01 B -> 10 Z` by themselves apply only for less accurate analysis (see *Case 4*).

There is also a formal argument why unspecified events on a vector expression need to be allowed rather than disallowed. Consider the following vector expressions within the scope of two variables A and B.

```
01 A              // (i)
01 B              // (ii)
01 A & 01 B       // (iii)
```

The natural interpretation here is `(iii)` === `(i)` & `(ii)`. This interpretation is only possible by allowing simultaneously occurring unspecified events.

*Allowing* simultaneously occurring unspecified events, the vector expressions `(i)` and `(ii)`, respectively, are interpreted as follows:

1. The power analysis tool relates to a timing constraint checker in a similar way as a parasitic extraction tool relates to a DRC tool. If the layout has DRC violations, for instance shorts between nets, the parasitic extraction tool shall report inaccurate wire capacitance for those nets. After final layout, the DRC violations shall be gone and the wire capacitance shall be accurate.

```
01 A & ?? B       // (i')
?? A & 01 B       // (ii')
```

*Disallowing* simultaneously occurring unspecified events, the vector expressions `(i)` and `(ii)`, respectively, are interpreted as follows:

```
01 A & ?- B       // (i'')
?- A & 01 B       // (ii'')
```

The vector expressions `(i')` and `(ii')` are compatible with `(iii)`, whereas `(i'')` and `(ii'')` are not.

### 5.4.9 Simultaneous event sequences

The semantic meaning of the "simultaneous event operator" can be extended to describe simultaneously occurring *event sequences*, by using the following definition:

```
(01 A#1 .. -> ... 01 A#N) & (01 B#1 .. -> ... 01 B#N)
=== 01 A#1 & 01 B#1 ... -> ... 01 A#N & 01 B#N
```

This definition is analogous to scalar multiplication of vectors with the same number of indices. The number of indices corresponds to the number of `vector_event` expressions separated by `->` operators. If the number of `->` in both vector expressions is not the same, the shorter vector expression can be left-extended with unspecified events, using the `??` operator, in order to align both vector expressions.

Example:

```
(01 A -> 01 B -> 01 C) & (01 D -> 01 E)
=== (01 A -> 01 B -> 01 C) & (?? D -> 01 D -> 01 E)
=== 01 A & ?? D -> 01 B & 01 D -> 01 C & 01 E
=== 01 A -> 01 B & 01 D -> 01 C & 01 E
```

The easiest way to understand the meaning of "simultaneous event sequences" is to consider the event report in test pattern format. If each `vector_event_sequence` expression matches the event report in the same time window, then the event sequences happen simultaneously.

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     0     1     0     1
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     1
1062  1     1     1     0     0
1395  1     0     0     0     0
1640  0     0     0     1     0
```

Example:

```
01 A -> 10 B === 01 A & 11 B -> 11 A & 10 B     // (10a)
```

```
// event pattern expressed by (10a):
//    A     B
//    0     1
//    1     1
//    1     0

X0 D -> 00 D                                    // (10b)

// event pattern expressed by (10b):
//    D
//    X
//    0
//    0

(01 A -> 10 B) & (X0 D -> 00 D)                 // (10) === (10a)&(10b)
```

Both `(10a)` and `(10b)` are *True* at time 258. Therefore `(10)` is *True* at time 258.

```
10 C
=== ?? C -> ?? C -> 10 C
=== ?? C -> ?1 C -> 10 C                        // (11a)

// event pattern expressed by (11a):
//    C
//    ?
//    ?
//    1
//    0
```

`(11a)` is left-extended to match the length of the following `(11b)`.

```
01 A -> 00 D -> 11 E ===
   01 A & 00 D & ?? E
-> ?? A & 00 D & ?? E
-> ?? A & ?? D & 11 E
===
   01 A & 00 D & ?? E
-> 1? A & 00 D & ?1 E
-> ?? A & 0? D & 11 E                           // (11b)

// event pattern expressed by (11b):
//    A     D     E
//    0     0     ?
//    1     0     ?
//    ?     0     1
//    ?     ?     1
```

`(11b)` contains explicitly specified non-events. The non-event `00 D` calls for the unspecified events `?? A` and `?? E`. The non-event `00 E` calls for the unspecified events `?? A` and `?? D`. By propagating well-specified previous and next states to subsequent events, some unspecified events become partly specified.

```
10 C & (01 A -> 00 D -> 11 E)                   // (11) === (11a)&(11b)
```

`(11a)` is *True* at time 573 and time 1395. `(11b)` is *True* at time 573 and time 915. Therefore, `(11)` is *True* at time 573.

### 5.4.10 Implicit local variables

Until now, vector expressions are evaluated against an event report containing all variables within the scope of a cell. It is practical for the application to work with only one event report per cell or, at most, two event reports if the set of variables for `BEHAVIOR` (`scope=behavior`) and `VECTOR` (`scope=measure`) is different. However, for complex cells and megacells, it is sometimes necessary to change the scope of event observation, depending on operation modes. Different modes can require a different set of variables to be observed in different event reports.

The following definition allows to *extend* the scope of a vector expression locally:

> Edge operators apply not only to variables, but also to boolean expressions involving those variables. Those boolean expressions represent *implicit local variables* that are visible only within the vector expression where they appear.

Suppose the local variables `(A & B)`, `(A | B)` are inserted into the event report:

```
time  A     B     C     D     E     A&B   A|B
0     0     1     1     X     1     0     1
109   1     1     1     0     1     1     1
258   1     0     1     0     1     0     1
573   1     0     0     0     1     0     1
586   0     0     0     0     1     0     0
643   1     0     0     0     1     0     1
788   0     1     1     0     1     0     1
915   1     1     1     0     1     1     1
1062  1     1     1     0     0     1     1
1395  1     0     0     0     0     0     1
1640  0     0     0     1     0     0     0
```

Example:

```
01 (A & B)                                          // (12)

// event pattern expressed by (12):
//    A&B
//    0
//    1
```

`(12)` is *True* at time 109 and time 915.

```
10 (A | B)                                          // (13)

// event pattern expressed by (13):
//    A|B
//    1
//    0
```

`(13)` is *True* at time 586 and time 1640.

```
01 (A & B) -> 10 B                                  // (14)

// event pattern expressed by (14):
//    B     A&B
//    1     0
//    1     1
//    0     1
```

(14) is *True* at time 258.

```
10 ( A & B) & 10 B -> 10 C                          // (15)

// event pattern expressed by (15):
//    B     C     A&B
//    1     1     1
//    0     1     0
//    0     0     0
```

(15) is *True* at time 573.

```
10 (A & B) -> 10 (A | B)                            // (16)

// event pattern expressed by (16):
//    A&B   A|B
//    1     1
//    0     1
//    0     0
```

(16) is *True* at time 1640.

### 5.4.11 Conditional event sequences

The following definition *restricts* the scope of a vector expression locally:

> vector_boolean_and, also called *conditional event operator*
> This operator is defined between a vector expression and a boolean expression, using the overloaded symbol & or &&. The scope of the vector expression is restricted to the variables and eventual implicit local variables appearing within that vector expression. The boolean expression shall be *True* during the entire vector expression. The boolean expression is called the *Existence Condition* of the vector expression.[2]

Vector expressions using the vector_boolean_and operator are called vector_conditional_event expressions. Scope and contents of such expressions are identical, as opposed to non-conditional vector_complex_event expressions, where the content is a subset of the scope.

Example:

```
(10 (A & B) -> 10 (A | B)) & !D                          // (17)

// event pattern expressed by (17):
//    A&B   A|B
//    1     1
//    0     1
//    0     0
```

2. An Existence Condition can also appear as annotation to a VECTOR object instead of appearing in the vector expression. This enables recognition of existence conditions by application tools which can not evaluate vector expressions (e.g., static timing analysis tools). However, for tools that can evaluate vector expressions, there is no difference between existence condition as a co-factor in the vector expression or as an annotation.

```
// event report without C, E:
time  A     B     D     A&B   A|B
0     0     1     X     0     1
109   1     1     0     1     1
258   1     0     0     0     1
586   0     0     0     0     0
643   1     0     0     0     1
788   0     1     0     0     1
915   1     1     0     1     1
1062  1     1     0     1     1
1395  1     0     0     0     1
1640  0     0     1     0     0
```

`(17)` contains the same `vector_complex_event` expression as `(16)`. However, although `(16)` is not *True* at time 586, `(17)` is *True* at time 586, since the scope of observation is narrowed to `A`, `B`, `A&B`, and `A|B` by the existence condition `!D`, which is statically *True* while the specified event sequence is observed.

Within, and only within, the narrowed scope of the `vector_conditional_event` expression, `(17)` can be considered equivalent to the following:

```
(10 (A & B) -> 10 (A | B)) & !D
===
(10 (A & B) -> 10 (A | B)) & (11 (!D) -> 11 (!D))
===
10 (A & B) & 11 (!D) -> 10 (A | B) & 11 (!D)
```

The transformation consists of the following steps:

1. Transform the boolean condition into a non-event.
   For example, `!D` becomes `11 (!D)`.

2. Left-extend the `vector_single_event` expression containing the non-event in order to match the length of the `vector_complex_event` expression.
   For example, `11 (!D)` becomes `11 (!D) -> 11 (!D)` to match the length of `10 (A & B) -> 10 (A | B)`.

3. Apply scalar multiplication rule for simultaneously occurring event sequences.

Thus, a `vector_conditional_event` expression can be transformed into an equivalent `vector_complex_event` expression, but the change of scope needs to be kept in mind. An operator which can express the change of scope in the vector expression language is defined in Section 5.4.13. This can make the transformation more rigorous.

Regardless of scope, the transformation from `vector_conditional_event` expression to `vector_complex_event` expression also provides the means of detecting ill-specified `vector_conditional_event` expressions.

Example:

```
(10 A -> 01 B -> 01 A) & A
===
10 A & 11 A -> 01 B & 11 A -> 01 A & 11 A
```

The first expression `10 A & 11 A` and the third expression `01 A & 11 A` within the `vector_complex_event` expression are contradictory. Hence, the `vector_conditional_event` expression can never be *True*.

### 5.4.12 Alternative conditional event sequences

All `vector_binary` operators, in particular the `vector_or` operator, can be applied to `vector_conditional_event` expressions as well as to `vector_complex_event` expressions.

Consider again the event report:

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     0     1     0     1
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     1
1062  1     1     1     0     0
1395  1     0     0     0     0
1640  0     0     0     1     0
```

Concurrent alternative `vector_conditional_event` expressions can be paraphrased in the following way:

```
IF <boolean_expression1> THEN <vector_expression1>
OR IF <boolean_expression2> THEN <vector_expression2>
... OR IF <boolean_expressionN> THEN <vector_expressionN>
```

The conditions can be *True* within overlapping time windows and thus the vector expressions are evaluated concurrently. The `vector_boolean_and` operator and `vector_or` operator describe such vector expressions.

Example:

```
C&(01 A -> 10 B) | !D&(10 B -> 10 A) | E&(10 B -> 10 C)     // (18)

// Event pattern expressed by (18):
//    A     B     C
//    0     1     1
//    1     1     1
//    1     0     1
```

`(18)` is *True* at time 258 because of `C & (01 A -> 10 B)`.

```
// Alternative event pattern expressed by (18):
//    A     B     D
//    1     1     0
//    1     0     0
//    0     0     0
```

`(18)` is also *True* at time 586 because of `!D & (10 B -> 10 A)`.

```
// Alternative event pattern expressed by (18):
//    B     C     E
//    1     1     1
```

```
//    0     1     1
//    0     0     1
```

`(18)` is also *True* at time 573 because of `E & (10 B -> 10 C)`.

Prioritized alternative `vector_conditional_event` expressions can be paraphrased in the following way:

```
IF <boolean_expression1> THEN <vector_expression1>
ELSE IF <boolean_expression2> THEN <vector_expression2>
... ELSE IF <boolean_expressionN> THEN <vector_expressionN>
(optional) ELSE <vector_expressiondefault>
```

Only the vector expression with the highest priority *True* condition is evaluated. The `vector_boolean_cond` operator and `vector_boolean_else` operator are used in ALF to describe such vector expressions.

Example:

```
C? (01 A -> 10 B): !D? (10 B -> 10 A): E? (10 B -> 10 C)    // (19)
```

The prioritized alternative `vector_conditional_event` expression can be transformed into concurrent alternative `vector_conditional_event` expression as shown:

```
C ? (01 A -> 10 B) : !D ? (10 B -> 10 A) : E ? (10 B -> 10 C)
===
C & (01 A -> 10 B)
| !C & !D & (10 B -> 10 A)
| !C & !(!D) & E & (10 B -> 10 C)
```

`(19)` is *True* at time 258 because of `C & (01 A -> 10 B)`, but not at time 586 because of higher priority `C` while `!D & (10 B -> 10 A)`, nor at time 573 because of higher priority `!D` while `E & (10 B -> 10 C)`.

### 5.4.13 Change of scope within a vector expression

Conditions on vector expressions redefine the scope of vector expressions locally. The following definition can be used to change the scope even within a part of a vector expression. For this purpose, the symbolic state `*` can be used, which means “don’t care about events”. This is different from the symbolic state `?` which means “don’t care about state”. When the state of a variable is `*`, arbitrary events occurring on that variable are disregarded.

- Edge operator with `*` as next state:
  The variable to which the operator applies is no longer within the scope of the vector expression.
- Edge operator with `*` as previous state:
  The variable to which the edge operator applies is now within the scope of the vector expression.

As opposed to `?`, `*` stands for an infinite variety of possibilities.

Example:

Let `A` be a logic variable with the possible states `1`, `0`, and `X`.

```
*0 A ===
00 A | 10 A | X0 A
| 00 A -> 00 A | 10 A -> 00 A | X0 A -> 00 A
| 01 A -> 10 A | 11 A -> 10 A | X1 A -> 10 A
| 0X A -> X0 A | 1X A -> X0 A | XX A -> X0 A
| 00 A -> 00 A -> 00 A | ...
```

```
0* A ===
00 A | 01 A | 0X A
| 00 A -> 00 A | 00 A -> 01 A | 00 A -> 0X A
| 01 A -> 10 A | 01 A -> 11 A | 01 A -> 1X A
| 0X A -> X0 A | 0X A -> X1 A | 0X A -> XX A
| 00 A -> 00 A -> 00 A | ...
```

A vector expression with an infinite variety of possible event sequences cannot be directly matched with an event report. However, there are feasible ways to implement event sequence detection involving `*`. In principle, there is a “static” and “dynamic” way. The following parts of the vector expression are separated by `*` *sub-sequences* of events.

- “Static” event sequence detection with `*`:
  The event report with all variables can be maintained, but certain variables are masked for the purpose of detection of certain sub-sequences.
- “Dynamic” event sequence detection with `*`:
  The event report shall contain the set of variables necessary for detection of a relevant sub-sequence. When such a sub-sequence is detected, the set of variables in the event report shall change until the next sub-sequence is detected, etc.

Examples:

```
01 A -> 1* B -> 10 C                                    // (20)
```

```
// Event pattern expressed by (20):
//    A     B     C
//    0     1     1
//    1     1     1
//    1     *     1
//    1     *     0
```

```
// pattern for 1st sub-sequence:
//    A     B     C
//    0     1     1
//    1     1     1
//    1     *     1
```

```
// pattern for 2nd sub-sequence:
//    A     B     C
//    1     *     1
//    1     *     0
```

The event report with masking relevant for `(20)`:

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     *     1     0     1     // detection of 1st sub-sequence
573   1     *     0     0     1     // detection of 2nd sub-sequence
```

```
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     1
1062  1     *     1     0     0     // detection of 1st sub-sequence
1395  1     *     0     0     0     // detection of 2nd sub-sequence
1640  0     0     0     1     0
```

(20) is *True* at time 573 and time 1395. The first sub-sequence `01 A -> 1* B` is detected at time 258, since * maps to any state. From time 258 onwards, `B` is masked. The second sub-sequence `10 C` is detected at time 573. From time 573 onwards, `B` is unmasked. The first sub-sequence is detected again at time 1062. The second sub-sequence is detected again at time 1395.

```
01 A & 1* E -> 10 C                                    // (21)

// Event pattern expressed by (21):
//    A     C     E
//    0     1     1
//    1     1     *
//    1     0     *

// pattern for 1st sub-sequence:
//    A     C     E
//    0     1     1
//    1     1     *

// pattern for 2nd sub-sequence:
//    A     C     E
//    1     1     *
//    1     0     *
```

The event report with masking relevant for (21):

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     *     // detection of 1st sub-sequence
258   1     0     1     0     *     // abortion of detection process
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     1     1     0     1
915   1     1     1     0     *     // detection of 1st sub-sequence
1062  1     1     1     0     *     // disregard event out of scope
1395  1     0     0     0     0     // detection of 2nd sub-sequence
1640  0     0     0     1     0
```

(21) is *True* at time 1395. The first sub-sequence `01 A & 1* E` is detected at time 109. From time 109 onwards, `E` is masked. The event on `B` at time 258 aborts continuation of the detection process and triggers restart from the beginning. The first sub-sequence is detected again at time 915. From time 915 onwards, `E` is masked. The event at time 1062 is therefore out of scope. The second sub-sequence `10 C` is detected at time 1395.

```
01 A -> *1 B -> 10 B & 10 C                            // (22)
```

```
// Event pattern expressed by (22):
//    A     B     C
//    0     *     1
//    1     *     1
//    1     1     1
//    1     0     0

// pattern for 1st sub-sequence:
//    A     B     C
//    0     *     1
//    1     *     1

// pattern for 2nd sub-sequence:
//    A     B     C
//    1     *     1
//    1     1     1
//    1     0     0
```

The event report with masking relevant for `(22)`:

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1     // detection of 1st sub-sequence
258   1     0     1     0     1     // abort
573   1     *     0     0     1
586   0     *     0     0     1
643   1     *     0     0     1
788   0     *     1     0     1
915   1     *     1     0     1     // detection of 1st sub-sequence
1062  1     1     1     0     0     // continue
1395  1     0     0     0     0     // detection of 2nd sub-sequence
1640  0     0     0     1     0
```

`(22)` is *True* at time 1395. The first sub-sequence `01 A` is detected at time 109. Therefore, `B` is unmasked. Since `B=0` at time 258, the attempt to detect the second sub-sequence is aborted and the detection process restarts from the beginning. The first sub-sequence `01 A` is detected again at time 109. The second sub-sequence `*1 B -> 10 B & 10 C` is detected at time 1395.

```
01 A -> 1? A & 0* B & 1* E -> 10 C                    // (23)

// Event pattern expressed by (23):
//    A     B     C     E
//    0     0     1     1
//    1     0     1     1
//    1     *     1     *
//    1     *     0     *

// pattern for 1st sub-sequence:
//    A     B     C     E
//    0     0     1     1
//    1     0     1     1
//    ?     *     1     *

// pattern for 2nd sub-sequence:
//    A     B     C     E
//    ?     *     1     *
//    ?     *     0     *
```

The event report with masking relevant for (23):

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1
258   1     0     1     0     1
573   1     0     0     0     1
586   0     0     0     0     1
643   1     0     0     0     1
788   0     *     1     0     *     // detection of 1st sub-sequence
915   1     *     1     0     *     // abort
1062  1     1     1     0     0
1395  1     0     0     0     0
1640  0     0     0     1     0
```

`(23)` is not *True* at any time. The first sub-sequence is detected at time 788. The event at time 915 does not match the expected second sub-sequence.

### 5.4.14 Sequences of conditional event sequences

The symbol `*` can be used to describe the scope of a vector expression directly in the vector expression language. This is particularly useful for sequences of `vector_conditional_event` expressions.

In reusing `(17)` as example:

```
(10 (A & B) -> 10 (A | B)) & !D
```

the scope of the sample event report contains contain the variables `A`, `B`, `C`, `D`, and `E`. The `vector_conditional_event` expression `(17)` contains only the variables `A`, `B`, and `D` and the implicit local variables `A&B` and `A|B`. Therefore, the global variables `C` and `E` are out of scope within `(17)`. The implicit local variables `A&B` and `A|B` are in scope within, and only within, `(17)`.

Now consider a *sequence* of `vector_conditional_event` expressions, where variables move in and out of scope. With the following formalism, it is possible to transform such a sequence into an equivalent `vector_complex_event` expression, allowing for a change of scope within each `vector_conditional_event` expression.

```
<vector_conditional_event#1> .. -> .. <vector_conditional_event#N>
```

where

```
<vector_conditional_event#i>
=== <vector_complex_event#i> & <boolean_expression#i>// 1 ≤ i ≤ N
```

The principle is to decompose each `vector_conditional_event` expression into a sequence of three vector expressions *prefix*, *kernel*, and *postfix* and then to reassemble the decomposed expressions.

```
<vector_conditional_event#i>
=== <prefix#i> -> <kernel#i> -> <postfix#i>// 1 ≤ i ≤ N
```

1. Define the prefix for each `vector_conditional_event` expression.
   The *prefix* is a `vector_event` expression defining all implicit local variables.

Example:

```
*? (A&B) & *? (A|B)
```

2. Define the kernel for each `vector_conditional_event` expression.
   The *kernel* is the `vector_complex_event` expression equivalent to the `vector_conditional_event` expression.

```
<vector_complex_event#i> & <boolean_expression#i>
===   <vector_complex_event#i>
&     (11 <boolean_expression#i> ..->.. 11 <boolean_expression#i>)
```

The kernel can consist of one or several alternative `vector_event_sequence` expressions. Within each `vector_event_sequence` expression, the same set of global variables are pulled out of scope at the first `vector_event` expression and pushed back in scope at the last `vector_event` expression.
Example:

```
?* C & ?* E             // global variables out of scope
& 10 (A & B) & 11 (!D) -> 10 (A | B) & 11 (!D)
& *? C & *? E           // global variables back in scope
```

3. Define the postfix for each `vector_conditional_event` expression.
   The *postfix* is a `vector_event` expression removing all implicit local variables.

Example:

```
?* (A&B) & ?* (A|B)
```

4. Join the subsequent `vector_complex_event` expressions with the `vector_and` operator between prefix#i+1and kernel#i and also between postfix#i and kernel#i+1.

```
   .. <vector_conditional_event#i> -> <vector_conditional_event#i+1>
..
   ===   .. <prefix#i>
         -> <postfix#i-1> & <kernel#i> & <prefix#i+1>
         -> <postfix#i> & <kernel#i+1> & <prefix#i+2>
         -> <postfix#i+1> ..
```

The complete example:

```
(10 (A & B) -> 10 (A | B)) & !D
===
*? (A&B) & *? (A|B)
-> ?* C & ?* E
& 10 (A & B) & 11 (!D) -> 10 (A | B) & 11 (!D)
& *? C & *? E
-> ?* (A&B) & ?* (A|B)
```

Note: The in-and-out-of-scope definitions for global variables are within the kernel, whereas the in-and-out-of-scope definitions for global variables are within the prefix and postfix. In this way, the resulting `vector_complex_event` expression contains the same uninterrupted sequence of events as the original sequence of `vector_conditional_event` expressions.

### 5.4.15 Incompletely specified event sequences

So far the vector expression language has provided support for *completely specified event sequences* and also the capability to put variables temporarily in and out of scope for event observation. As opposed to changing the scope of event observation, *incompletely specified event sequences* require continuous observation of all variables while allowing the occurrence of intermediate events between the specified events. The following operator can be used for that purpose:

`vector_followed_by`, also called *followed-by operator*, using the symbol `~>`.
The `~>` operator is the separator between consecutively occurring events, with possible unspecified events in-between.

Detection of event sequences involving `~>` requires detection of the sub-sequence before `~>`, setting a flag, detection of the sub-sequence after `~>`, and clearing the flag.

This can be illustrated with a sample event report:

```
time  A     B     C     D     E
0     0     1     1     X     1
109   1     1     1     0     1     // 01 A detected, set flag
258   1     0     1     0     1
573   1     0     0     0     1     // 10 C detected, clear flag
586   0     0     0     0     1
643   1     0     0     0     1     // 01 A detected, set flag
788   0     1     1     0     1
915   1     1     1     0     1     // 01 A detected again
1062  1     1     1     0     0
1395  1     0     0     0     0     // 10 C detected, clear flag
1640  0     0     0     1     0
```

Example:

```
01 A ~> 10 C                                          // (24)
// as opposed to previous example (5):01 A -> 10 C
```

`(24)` is *True* at time 573 because of `01 A` at time 109 and `10 C` at time 573. It is *True* again at time 1395 because of `01 A` at time 643 and `10 C` at 1395. On the other hand, `(5)` is never *True* because there are always events in-between `01 A` and `10 C`.

Vector expressions consisting of `vector_event` expressions separated by `->` or by `~>` are called `vector_event_sequence` expressions, using the same syntax rules for the two different `vector_followed_by` operators. Consequently, all vector expressions involving `vector_event_sequence` expressions and `vector_binary` operators are called `vector_complex_event` expressions.

However, only a subset of the semantic transformation rules can be applied to vector expressions containing `~>`.

Associative rule applies for both `->` and `~>`.

```
(01 A ~> 01 B) ~> 01 C === 01 A ~> (01 C ~> 01 B ~> 01 C)

(01 A -> 01 B) -> 01 C === 01 A -> (01 C -> 01 B -> 01 C)

(01 A ~> 01 B) -> 01 C === 01 A ~> (01 C ~> 01 B -> 01 C)

(01 A -> 01 B) ~> 01 C === 01 A -> (01 C -> 01 B ~> 01 C)
```

Distributive rule applies for both -> and ~>.

```
(01 A | 01 B) -> 01 C === 01 A ~> 01 C | 01 B -> 01 C

(01 A | 01 B) ~> 01 C === 01 A ~> 01 C | 01 B ~> 01 C

(01 A | 01 B) -> 01 C === 01 A ~> 01 C | 01 B -> 01 C
```

Scalar multiplication rule applies only for ->. The transformation involving ~> is more complicated.

```
(01 A -> 01 B) & (01 C -> 01 D)
=== (01 A & 01 C) -> (01 B & 01 D)

(01 A ~> 01 B) & (01 C -> 01 D)
=== (01 A & 01 C) -> (01 B & 01 D)
|    01 A ~> 01 C -> (01 B & 01 D)

(01 A ~> 01 B) & (01 C ~> 01 D)
=== (01 A & 01 C) ~> (01 B & 01 D)
|    01 A ~> 01 C ~> (01 B & 01 D)
|    01 C ~> 01 A ~> (01 B & 01 D)
```

Transformation of vector_conditional_event expressions into vector_complex_event expressions applies only for ->.

```
(01 A -> 01 B) & C
=== 01 A & 11 C -> 01 B & 11 C

(01 A ~> 01 B) & C
=== 01 A & 11 C ~> 01 B & 11 C
```

Since the ~> operator allows intermediate events, there is no way to express the continuously *True* condition C.

### 5.4.16 How to determine well-specified vector expressions

By defining semantics for

alternative vector_event_sequence expressions

and establishing calculation rules for

transforming vector_complex_event expressions into alternative vector_event_sequence expressions

and for

transforming alternative vector_conditional_event expressions into alternative vector_complex_event expressions,

semantics are now defined for all vector expressions.

The calculation rules also provide means to determine whether a vector expression is well-specified or ill-specified. An ill-specified vector expression is contradictory in itself and can therefore never be *True*.

Once a vector expression is reduced to a set of alternative vector_event_sequence expressions, two criteria define whether a vector expression is well-defined or not.

- Compatibility between subsequent events on the same variable:
  The next state of earlier event shall be compatible with previous state of later event. This check applies only if no ~> operator is found between the events.
- Compatibility between simultaneous events on the same variable:
  Both the previous and next state of both events shall be compatible. Such events commonly occur as intermediate calculation results within vector expression transformation.

The following compatibility rules apply:

- `?` is compatible with any other state. If the other state is `*`, the resulting state is `?`. Otherwise, the resulting state is the other state.
- `*` is compatible with any other state. The resulting state is the other state.
- Any other state is only compatible with itself.

Examples:

```
01 A -> 01 B -> 10 A
```

The next state of `01 A` is compatible with the previous state of `10 A`.

```
0X A -> 01 B -> 10 A
```

The next state of `0X A` is not compatible with the previous state of `10 A`.

```
0X A ~> 01 B -> 10 A
```

Compatibility check does not apply, since intermediate events are allowed.

```
01 A & 10 A
```

Both the previous and next state of **`A`** are contradictory; this results in an impossible event.

```
?1 A & 1? A
```

Both previous and next state of `A` are compatible; this results in the non-event `11 A`.

## 5.5 Variable declarations

Inside a `CELL` object, the `PIN` objects with the `PINTYPE digital` define variables for `FUNCTION` objects inside the same `CELL`. A *primary input variable* inside a `FUNCTION` shall be declared as a `PIN` with `DIRECTION=input` or `both` (since `DIRECTION=both` is a bidirectional pin). However, it is not required that all declared pins are used in the function. Output variables inside a `FUNCTION` need not be declared pins, since they are implicitly declared when they appear at the left-hand side (LHS) of an assignment.

Example:

```
CELL my_cell {
      PIN A {DIRECTION = input;}
      PIN B {DIRECTION = input;}
      PIN C {DIRECTION = output;}
```

```
        FUNCTION {
              BEHAVIOR {
                    D = A && B;
                    C = !D;
              }
        }
  }
```

C and D are output variables that need not be declared prior to use. After implicit declaration, D is reused as an input variable. A and B are primary input variables.

### 5.5.1 BEHAVIOR

Inside BEHAVIOR, variables that appear at the LHS of an assignment conditionally controlled by a vector expression, as opposed to an unconditional continuous assignment, hold their values, when the vector expression evaluates *False*. Those variables are considered to have latch-type behavior.

Examples:

```
BEHAVIOR {
      @(G){
            Q = D;  // both Q and QN have latch-type behavior
            QN = !D;
      }
}
BEHAVIOR {
      @(G){
            Q = D;  // only Q has latch-type behavior
      }
      QN = !Q;
}
```

### 5.5.2 STATETABLE

The functional description can be supplemented by a STATETABLE, the first row of which contains the arguments that are object IDs of the declared PINs. The arguments appear in two fields, the first is input and the second is output. The fields are separated by a :. The rows are separated by a ;. The arguments can appear in both fields if the PINs have attribute direction=output or direction=both. If direction=output, then the argument has latch-type behavior. The argument on the input field is considered previous state and the argument on the output field is considered the next state. If direction=both, then the argument on the input field applies for input direction and the argument on the output field applies for output direction of the bidirectional PIN.

Example:

```
CELL ff_sd {
      PIN  q {DIRECTION=output;}
      PIN  d {DIRECTION=input;}
      PIN cp {DIRECTION=input;
                  SIGNALTYPE=clock;
                  POLARITY=rising_edge;}
```

```
        PIN cd {DIRECTION=input; SIGNALTYPE=clear; POLARITY=low;}
        PIN sd {DIRECTION=input; SIGNALTYPE=set; POLARITY=low;}
        FUNCTION {
              BEHAVIOR {
                    @(!cd) {q = 0;} :(!sd) {q = 1;} :(01 cp) {q = d;}
              }
              STATETABLE {
                    cd sd  cp  d   q  : q ;
                    0  ?   ??  ?   ?  : 0 ;
                    1  0   ??  ?   ?  : 1 ;
                    1  1   1?  ?   0  : 0 ;
                    1  1   ?0  ?   1  : 1 ;
                    1  1   1?  ?   0  : 0 ;
                    1  1   ?0  ?   1  : 1 ;
                    1  1   01  ?   ?  :(d);
              }
        }
  }
```

If the output variable with latch-type behavior depends only on the previous state of itself, as opposed to the previous state of other output variables with latch-type behavior, it is not necessary to use that output variable in the input field. This allows a more compact form of the `STATETABLE`.

Example:

```
  STATETABLE {
        cd sd  cp  d  : q ;
        0  ?   ??  ?  : 0 ;
        1  0   ??  ?  : 1 ;
        1  1   1?  ?  :(q);
        1  1   ?0  ?  :(q);
        1  1   01  ?  :(d);
  }
```

A generic ALF parser shall make the following semantic checks:

- Are all variables of a `FUNCTION` declared either by declaration as `PIN` names or through assignment?
- Does the `STATETABLE` exclusively contain declared `PIN`s?
- Is the format of the `STATETABLE`, i.e., the number of elements in each field of each row, consistent?
- Are the values consistently either state or transition digits?
- Is the number of digits in each `TABLE` entry compatible with the signal bus width?

A more sophisticated checker for complete verification of logical consistency of a `FUNCTION` given in both equation and tabular representation is out of scope for a generic ALF parser, which checks only syntax and compliance to semantic rules. However, formal verification algorithms can be implemented in special-purpose ALF analyzers or model generators/ compilers.

### 5.5.3 Multi-dimensional variables

A group of pins of a cell can be logically considered together by declaring a PIN with a range. A pin can be declared with one dimension or two dimensions. For example,

```
PIN                    A ;   // declares a scalar pin A
PIN [1:8]              A1 ;  // declares pin A1 with bits numbered 1
                             // through 8
PIN [1:8]              A2[1:4] ;// declares pin A2 with two dimensions
```

When a pin is declared with one dimension, the left number in the range shall specify the most significant bit number and the right number shall specify the least significant bit number. If the pin is declared with two dimensions, the second dimension shall specify the index of the first and the last rows of the two-dimension pin object.

A PIN object can be referenced in one of the four forms:

- Individual bit - the pin name shall be followed by an index of the bit.
- Contiguous group of bits - the pin name shall be followed by the contiguous range of bits. The most significant and least significant bit numbers shall follow the same relationship as given in the declaration.
- Entire PIN object - only the pin name shall be used. It shall be illegal to reference the entire two-dimension pin object in any operation.
- One row of a PIN object - for a two-dimension pin object, the name of the pin shall be followed by the row index of that pin. It shall be illegal to reference the individual bit or a group of bits of a two-dimension pin object directly in an operation.

When a PIN object is referenced on the left-hand side of an assignment, the result of the right-hand side expression is copied from the least significant bit towards the most significant bit. If the right-hand side value has lesser number of bits than the referenced PIN object in an assignment, the right-hand side value shall be zero-extended to fill the remaining bits of the referenced PIN object. If the right-hand side value has more bits than the referenced PIN object in an assignment, the right-hand side value shall be truncated to the size of the referenced PIN object.

Example:

```
pin [1:8] A1;
pin [1:8] A2[1:32] ;

A1[8]   = 'b0 ;
A1[1:6] = 'o75 ;             // is equivalent to A1[1:6] = 'b111_101
A1[1:5] = 'o75 ;             // is equivalent to A1[1:5] = 'b11_101,
                             // left most bit is truncated
A2[18]  = 'h5 ;              // is equivalent to A2[18] = 'b0000_0101
                             // entire row 18 of A2 is assigned a value.
```

Two-dimension PIN objects shall be referenced with the row index. It shall be illegal to directly reference an individual bit or a contiguous group of bits of a two-dimension PIN object. It shall be illegal to reference the entire PIN object as a two-dimension PIN object.

Example:

```
pin [1:8] A2[1:32] ;
pin [1:8] B1 ;
pin C ;

                        // legal references and assignments

A2[10]  = 'h45 ;        // assign 'h45 to row 10 of A2  ('b0100_0101)
B1      = A2[10] ;      // copies whole row A2[10] to B1
C       = B1[3] ;       // c = 'b0

// Illegal references and assignments
// B1[3]   = A2[10][3] ;illegal reference to bit 3 of A2[10]
// A2      = B1 ;       illegal reference to entire A2
```

It shall be legal to use identifiers as an index, but expressions shall not be permitted.

Example:

```
pin [4:1] ADDR;

ADDR       = 'd 10;
A2[ADDR]   = 'h45 ;     // assign 'h45 to row 10 of A2  ('b0100_0101)

// A2[ADDR+1] = 'h45 ;  illegal
```

### 5.5.4 ROM initialization

The STATETABLE statement can be used to describe the contents of a ROM, as far as this content is fixed in the library.

Example:

```
CELL my_rom {
   CELLTYPE = memory;
   ATTRIBUTE { rom asynchronous }
   PIN[1:2] addr { DIRECTION = input; SIGNALTYPE = address; }
   PIN[3:0] dout { DIRECTION = output; SIGNALTYPE = data; }
   PIN[3:0] mem[1:4] { DIRECTION=none; VIEW=none; SIGNALTYPE=data; }
   FUNCTION {
      BEHAVIOR { dout = mem[addr]; }
      STATETABLE {
         addr: mem ;
         'h0:  'h5 ;
         'h1:  'hA ;
         'h2:  'h5 ;
         'h3:  'hA ;
      }
   }
}
```

For flexibility, a separate included file can be used:

```
CELL my_rom {
   CELLTYPE = memory;
   ATTRIBUTE { rom asynchronous }
   PIN[1:2] addr { DIRECTION = input; SIGNALTYPE = address; }
   PIN[3:0] dout { DIRECTION = output; SIGNALTYPE = data; }
   PIN[3:0] mem[1:4] { DIRECTION=none; VIEW=none; SIGNALTYPE=data; }
   FUNCTION {
      BEHAVIOR { dout = mem[addr]; }
      INCLUDE “rom_initialization_file.alf” ;
      }
   }
}
```

The contents of the included file `rom_initialization_file.alf` are:

```
STATETABLE {
   addr: mem ;
   ‘h0:  ‘h5 ;
   ‘h1:  ‘hA ;
   ‘h2:  ‘h5 ;
   ‘h3:  ‘hA ;
}
```

## 5.6 Predefined models

This section defines the use of predefined models in ALF.

### 5.6.1 Usage of PRIMITIVEs

A `PRIMITIVE` referenced in a `CELL` can replace the complete set of `PIN` and `FUNCTION` definition. `PIN`s can be declared before the reference to the `PRIMITIVE`, in order to provide supplementary annotations that cannot be inherited from the `PRIMITIVE`. However, the `CELL` shall be pin-compatible with the `PRIMITIVE`.

If the `PRIMITIVE` or a `CELL` is referenced in an annotation container such as `SCAN`, only the subset of `PIN`s used in the non-scan cell shall be compatible with the `PIN`s of the cell.

The pin names can be referenced by order or by name. In the latter case, the LHS is the pin name of the referenced `PRIMITIVE` or `CELL` (e.g., the non-scan cell), the RHS is the pin name of the actual cell. A constant logic value can also appear at the LHS or RHS, indicating a pin needs to be tied to a constant value. If this information is already specified in an annotation inside the `PIN` object itself, referencing between a pin name and a constant value is not necessary.

`PRIMITIVE`s can also be instantiated inside `BEHAVIOR`.

### 5.6.2 Concept of user-defined and predefined primitives

Primitives are described in ALF syntax. Primitives are generic cells containing `PIN` and FUNCTION objects only, i.e., no characterization data. The primitives are used for structural functional modeling.

Example:

```
PRIMITIVE MY_PRIMITIVE {
      PIN x { ... }
      PIN y { ... }
      PIN z { ... }
      FUNCTION { ... }
}

CELL MY_CELL {
      PIN a { ... }
      PIN b { ... }
      PIN c { ... }
      FUNCTION {
            BEHAVIOR { MY_PRIMITIVE { x=a; y=b; z=c; } }
      }
      ...
}
```

Extensible primitives, i.e., primitives with variable number of pins can be modeled using a `TEMPLATE`.

Example:

```
TEMPLATE EXTENSIBLE_PRIMITIVE{
      PRIMITIVE <primitive_name> {
            PIN [0:<max_index>] pin_name {  ... }
            ...
      }
}

// instantiation of the template creates a primitive
EXTENSIBLE_PRIMITIVE {
      primitive_name = MY_EXTENSIBLE_PRIMITIVE;
      max_index = 2;
}
```

The set of statements above is equivalent to the following statement:

```
PRIMITIVE MY_EXTENSIBLE_PRIMITIVE {
      PIN [0:2] pin_name {  ... }
            ...
}
```

The primitive can be used as shown in the following example:

```
CELL MY_MEGACELL {
      PIN a { ... }
      PIN b { ... }
      PIN c { ... }
      FUNCTION {
            BEHAVIOR {
                  // reference to the primitive
                  MY_EXTENSIBLE_PRIMITIVE {
                        pin_name[0] = a;
```

```
                        pin_name[1] = b;
                        pin_name[2] = c;
                  }
            }
      }
      ...
}
```

Primitives can be freely defined by the user. For convenience, ALF provides a set of predefined primitives with the reserved prefix `ALF_` in their name, which cannot be used by user-defined primitives.

For all `PIN`s of predefined primitives, the following annotations are defined by default:

```
VIEW = functional;
SCOPE = behavioral;
```

For predefined extensible primitives, a placeholder can be directly in the `PRIMITIVE` definition:

```
PRIMITIVE ALF_EXTENSIBLE_PRIMITIVE {
            PIN [0:<max_index>] pin_name {  ... }
      ...
}
```

This is equivalent to the following more verbose set of statements:

```
TEMPLATE EXTENSIBLE_PRIMITIVE{
      PRIMITIVE <primitive_name> {
            PIN [0:<max_index>] pin_name {  ... }
            ...
      }
}

EXTENSIBLE_PRIMITIVE {
      primitive_name = ALF_EXTENSIBLE_PRIMITIVE;
      max_index = <max_index>;
}
```

## 5.6.3 Predefined combinational primitives

This section defines the use of predefined combinational primitives.

### 5.6.3.1 One input, multiple output primitives

There are two combinational primitives with one input pin and multiple output pins:

`ALF_BUF` and `ALF_NOT`

A GROUP statement is used to define the behavior of all output pins in one statement.

The output pins are indexed starting with `0`. If `0` is the only index used, the index can be omitted when referencing the output pin, e.g., `out` refers to `out[0]`.

```
PRIMITIVE ALF_BUF {
      GROUP index {0:<max_index>}
      PIN[0:<max_index>] out {
            DIRECTION = output ;
      }
      PIN in {
            DIRECTION = input ;
      }
      FUNCTION {
            BEHAVIOR {
                  out[index] = in;
            }
      }
}
```

**Figure 5-7: Primitive model of ALF_BUF**

```
PRIMITIVE ALF_NOT {
      GROUP index {0:<max_index>}
      PIN[0:<max_index>] out {
            DIRECTION = output ;
      }
      PIN in {
            DIRECTION = input ;
      }
      FUNCTION {
            BEHAVIOR {
                  out[index] = !in;
            }
      }
}
```

**Figure 5-8: Primitive model of ALF_NOT**

#### 5.6.3.2 One output, multiple input primitives

There are six combinational primitives with one output pin and multiple input pins:

`ALF_AND`, `ALF_NAND`, `ALF_OR`, `ALF_NOR`, `ALF_XOR`, and `ALF_XNOR`

The input pins are indexed starting with `0`. If `0` is the only index used, the index can be omitted when referencing the input pin, e.g., `in` refers to `in[0]`.

```
PRIMITIVE ALF_AND {
      PIN out {
            DIRECTION = output;
      }
      PIN[0:<max_index>] in {
            DIRECTION = input;
      }
      FUNCTION {
            BEHAVIOR {
                  out = & in;
            }
      }
}
```

**Figure 5-9: Primitive model of ALF_AND**

```
PRIMITIVE ALF_NAND {
      PIN out {
            DIRECTION = output;
      }
      PIN[0:<max_index>] in {
            DIRECTION = input;
      }
      FUNCTION {
            BEHAVIOR {
                  out = ~& in;
            }
      }
}
```

**Figure 5-10: Primitive model of ALF_NAND**

```
PRIMITIVE ALF_OR {
      PIN out {
            DIRECTION = output;
      }
      PIN[0:<max_index>] in {
            DIRECTION = input;
      }
      FUNCTION {
            BEHAVIOR {
                  out = | in;
            }
      }
}
```

**Figure 5-11: Primitive model of ALF_OR**

```
PRIMITIVE ALF_NOR {
      PIN out {
            DIRECTION = output;
      }
      PIN[0:<max_index>] in {
            DIRECTION = input;
      }
      FUNCTION {
            BEHAVIOR {
                  out = ~| in;
            }
      }
}
```

**Figure 5-12: Primitive model of ALF_NOR**

```
PRIMITIVE ALF_XOR {
      PIN out {
            DIRECTION = output;
      }
      PIN[0:<max_index>] in {
            DIRECTION = input;
      }
      FUNCTION {
            BEHAVIOR {
                  out = ^in;
            }
      }
}
```

**Figure 5-13: Primitive model of ALF_XOR**

```
PRIMITIVE ALF_XNOR {
      PIN out {
            DIRECTION = output;
      }
      PIN[0:<max_index>] in {
            DIRECTION = input;
      }
      FUNCTION {
            BEHAVIOR {
                  out = ~^in;
            }
      }
}
```

**Figure 5-14: Primitive model of ALF_XNOR**

### 5.6.4 Predefined tristate primitives

There are four tristate primitives:

ALF_BUFIF1, ALF_BUFIF0, ALF_NOTIF1, and ALF_NOTIF0

```
PRIMITIVE ALF_BUFIF1 {
      PIN out {
            DIRECTION  = output;
            ENABLE_PIN = enable;
            ATTRIBUTE {TRISTATE}
      }
      PIN in  {
            DIRECTION  = input;
      }
      PIN enable {
            DIRECTION  = input;
            SIGNALTYPE = out_enable;
      }

      FUNCTION {
            BEHAVIOR {
                  out = (enable)? in : 'bZ;
            }
            STATETABLE {
                  enable in : out;
                   0     ?  : Z;
                   1     ?  : (in);
            }
      }
}
```

**Figure 5-15: Primitive model of ALF_BUFIF1**

```
PRIMITIVE ALF_BUFIF0 {
      PIN out {
            DIRECTION  = output;
            ENABLE_PIN = enable;
            ATTRIBUTE {TRISTATE}
      }
      PIN in  {
            DIRECTION  = input;
      }
      PIN enable {
            DIRECTION  = input;
            SIGNALTYPE = out_enable;
      }
      FUNCTION {
            BEHAVIOR {
                  out = (!enable)? in : 'bZ;
            }
```

```
            STATETABLE {
                  enable in : out;
                   1     ?  : Z;
                   0     ?  : (in);
            }
      }
}
```

**Figure 5-16: Primitive model of ALF_BUFIF0**

```
PRIMITIVE ALF_NOTIF1 {
      PIN out {
            DIRECTION  = output;
            ENABLE_PIN = enable;
            ATTRIBUTE {TRISTATE}
      }
      PIN in  {
            DIRECTION  = input;
      }

      PIN enable {
            DIRECTION  = input;
            SIGNALTYPE = out_enable;
      }
      FUNCTION {
            BEHAVIOR {
                  out = (enable)? !in : 'bZ;
            }
            STATETABLE {
                  enable in : out;
                   0     ?  : Z;
                   1     ?  : (!in);
            }
      }
}
```

**Figure 5-17: Primitive model of ALF_NOTIF1**

```
PRIMITIVE ALF_NOTIF0 {
      PIN out {
            DIRECTION  = output;
            ENABLE_PIN = enable;
            ATTRIBUTE {TRISTATE}
      }
      PIN in  {
            DIRECTION  = input;
      }
      PIN enable {
            DIRECTION  = input;
            SIGNALTYPE = out_enable;
      }
```

```
        FUNCTION {
              BEHAVIOR {
                    out = (!enable)? !in : 'bZ;
              }
              STATETABLE {
                    enable in : out;
                     1     ?  : Z;
                     0     ?  : (!in);
              }
        }
}
```

**Figure 5-18: Primitive model of ALF_NOTIF0**

### 5.6.5 Predefined multiplexor

The predefined multiplexor has a known output value if either the select signal and the selected data inputs are known or both data inputs have the same known value while the select signal is unknown.

```
PRIMITIVE ALF_MUX {
      PIN Q  {
            DIRECTION  = output;
            SIGNALTYPE = data;
      }
      PIN[1:0] D  {
            DIRECTION  = input;
            SIGNALTYPE = data;
      }
      PIN S  {
            DIRECTION  = input;
            SIGNALTYPE = select;
      }
      FUNCTION {
            BEHAVIOR {
                  Q = (S || (d[0] ~^ d[1]) )? d[1] : d[0];
            }
            STATETABLE {
                  D[0] D[1] S  : Q ;
                  ?    ?    0  : (D[0]);
                  ?    ?    1  : (D[1]);
                  0    0    ?  : 0;
                  1    1    ?  : 1;
            }
      }
}
```

**Figure 5-19: Primitive model of ALF_MUX**

### 5.6.6 Predefined flip-flop

A dual-rail output D-flip-flop with asynchronous set and clear pins is a generic edge-sensitive sequential device. Simpler flip-flops can be modeled using this primitive by setting input pins to appropriate constant values. More complex flip-flops can be modeled by adding combinational logic around the primitive.

A particularity of this model is the use of the last two pins `Q_CONFLICT` and `QN_CONFLICT`, which are virtual pins. They specify the state of `Q` and `QN` in the event `CLEAR` and `SET` become active simultaneously.

```
PRIMITIVE ALF_FLIPFLOP {
      PIN Q     {
            DIRECTION  = output;
            SIGNALTYPE = data;
            POLARITY   = non_inverted;
      }

      PIN QN    {
            DIRECTION  = output;
            SIGNALTYPE = data;
            POLARITY   = inverted;
      }
      PIN D     {
            DIRECTION  = input;
            SIGNALTYPE = data;
      }
      PIN CLOCK {
            DIRECTION  = input;
            SIGNALTYPE = clock;
            POLARITY   = rising_edge;
      }
      PIN CLEAR {
            DIRECTION  = input;
            SIGNALTYPE = clear;
            POLARITY   = high;
            ACTION     = asynchronous;
      }
      PIN SET   {
            DIRECTION  = input;
            SIGNALTYPE = set;
            POLARITY   = high;
            ACTION     = asynchronous;
      }
      PIN Q_CONFLICT   {
            DIRECTION  = input;
            VIEW       = none;
      }
      PIN QN_CONFLICT  {
            DIRECTION  = input;
            VIEW       = none;
      }
      FUNCTION {
            ALIAS QX  = Q_CONFLICT;
```

```
            ALIAS QNX = QN_CONFLICT;
            BEHAVIOR {
                  @ (CLEAR && SET) {
                        Q  = QX;
                        QN = QNX;
                  }
                  : (CLEAR) {
                        Q  = 0;
                        QN = 1;
                  }
                  : (SET) {
                        Q  = 1;
                        QN = 0;
                  }
                  : (01 CLOCK) {          // edge-sensitive behavior
                        Q  = D;
                        QN = !D;
                  }
            }
            STATETABLE {
                  D CLOCK CLEAR SET QX  QNX :  Q    QN ;
                  ?  ??    1     1  ?   ?   : (QX) (QNX);
                  ?  ??    0     1  ?   ?   :  1    0 ;
                  ?  ??    1     0  ?   ?   :  0    1 ;
                  ?  1?    0     0  ?   ?   : (Q)  (QN) ;
                  ?  ?0    0     0  ?   ?   : (Q)  (QN) ;
                  ?  01    0     0  ?   ?   : (D)  (!D) ;
            }
      }
}
```

**Figure 5-20: Primitive model of ALF_FLIPFLOP**

### 5.6.7 Predefined latch

The dual-rail D-latch with set and clear pins has the same functionality as the flip-flop, except the level-sensitive clock (ENABLE pin) is used instead of the edge-sensitive clock.

```
PRIMITIVE ALF_LATCH {
      PIN Q     {
            DIRECTION  = output;
            SIGNALTYPE = data;
            POLARITY   = non_inverted;
      }
      PIN QN    {
            DIRECTION  = output;
            SIGNALTYPE = data;
            POLARITY   = inverted;
      }
      PIN D     {
            DIRECTION  = input;
            SIGNALTYPE = data;
      }
```

```
PIN ENABLE {
      DIRECTION  = input;
      SIGNALTYPE = clock;
      POLARITY   = high;
}
PIN CLEAR {
      DIRECTION  = input;
      SIGNALTYPE = clear;
      POLARITY   = high;
      ACTION     = asynchronous;
}
PIN SET   {
      DIRECTION  = input;
      SIGNALTYPE = set;
      POLARITY   = high;
      ACTION     = asynchronous;
}

PIN Q_CONFLICT   {
      DIRECTION = input;
      VIEW      = none;
}
PIN QN_CONFLICT  {
      DIRECTION = input;
      VIEW      = none;
}
FUNCTION {
      ALIAS QX  = Q_CONFLICT;
      ALIAS QNX = QN_CONFLICT;
      BEHAVIOR {
            @ (CLEAR && SET) {
                  Q  = QX;
                  QN = QNX;
            }
            : (CLEAR) {
                  Q  = 0;
                  QN = 1;
            }
            : (SET) {
                  Q  = 1;
                  QN = 0;
            }
            : (ENABLE) {            // level-sensitive behavior
                  Q  = D;
                  QN = !D;
            }
      }
      STATETABLE {
            D  ENABLE CLEAR SET QX  QNX :  Q    QN ;
            ?  ?      1     1   ?   ?   : (QX) (QNX);
```

```
                ?  ?      0     1   ?   ?   :  1    0 ;
                ?  ?      1     0   ?   ?   :  0    1 ;
                ?  0      0     0   ?   ?   : (Q)  (QN) ;
                ?  1      0     0   ?   ?   : (D)  (!D) ;
            }
        }
}
```

**Figure 5-21: Primitive model of ALF_LATCH**

### 5.6.8 Parameterizeable cells

The concept of describing primitives with variable bus size shall be extended to parameterizeable cells. Dynamic template instantiations can be used for that purpose.

Template definitions can incorporate any type of object. Placeholders in the template definition are the equivalent of parameters. Hence, the definition of parameterizeable cells is already supported within the support of general template definitions.

In a *static template instantiation*, which is identified by the name of the template and by the optional value assignment `static`, placeholders are replaced by fixed values or by complex objects containing fixed values. Non-referenced placeholders stay in place and eventually result in semantically unrecognizable objects, which cannot be processed by downstream applications. Such unrecognizable objects shall be disregarded.

In a *dynamic template instantiation*, which is identified by the name of the template and by the mandatory value assignment `dynamic`, some placeholders can not be replaced. Those placeholders are application parameters. The template definition can already contain certain relationships between parameters (e.g., arithmetic model and its arguments in the header). Therefore the template instantiation determines which parameters need application values in order to calculate values for other parameters.

Going one step further, even the relationship between parameters can be defined in the dynamic template instantiation rather than in the template definition. In this case, the identifiers inside the placeholders become variables for arithmetic assignments. This definition of variables shall only be recognized within the context of the dynamic template instantiation.

Arithmetic assignments provide a shorter syntax for equation-based arithmetic models where only placeholder-parameters are involved.

```
param1 = 1.5 + 0.4 * param2 ** 3 - 2.7 / param3
```

is equivalent to

```
param1 {
      HEADER { param2 param3 }
      EQUATION { 1.5 + 0.4 * param2 ** 3 - 2.7 / param3 }
}
```

For table-based models or for models where the arguments have children objects attached to them, the verbose syntax with HEADER needs to be used.

Example:

```
TEMPLATE adder {
   CELL <cellname> {
      PIN [ <bitwidth> : 1 ] A { DIRECTION = input; }
      PIN [ <bitwidth> : 1 ] B { DIRECTION = input; }
      PIN Cin { DIRECTION = input; }
      PIN [ <bitwidth> : 1 ] S { DIRECTION = output; }
      PIN Cout { DIRECTION = output; }
      FUNCTION {
         BEHAVIOR {
            S = A + B + Cin;
            Cout = (A + B + Cin >= (‘b1 << (<bitwidth> - 1)));
         }
      }
      AREA = <areavalue>;
      VECTOR (?! Cin -> ?! Cout) {
         DELAY {
            HEADER {
               CAPACITANCE {PIN = Cout; }
               SLEWRATE {PIN = Cin; }
            }
            EQUATION { <D0> + <D1>*CAPACITANCE + <D2>*SLEWRATE }
         }
      }
   }
}
```

The template is used for instantiation of a hard macro:

```
adder { /* a hard macro */
   cellname = ripple_carry_adder_16_bit;
   bitwidth = 16;
   areavalue = 500;
   // D0, D1, D2 are undefined. DELAY cannot be calculated.
}
```

The static instantiation of the hard macro is equivalent to the following static object:

```
CELL ripple_carry_adder_16_bit {
   PIN [ 16 : 1 ] A { DIRECTION = input; }
   PIN [ 16 : 1 ] B { DIRECTION = input; }
   PIN Cin { DIRECTION = input; }
   PIN [ 16 : 1 ] S { DIRECTION = output; }
   PIN Cout { DIRECTION = output; }

   FUNCTION {
      BEHAVIOR {
         S = A + B + Cin;
         Cout = (A + B + Cin >= ’b1000000000000000);
      }
   }
   AREA = 500 ;
```

```
   VECTOR (?! Cin -> ?! Cout) {
//    DELAY {
//       HEADER {
//          CAPACITANCE {PIN = Cout; }
//          SLEWRATE {PIN = Cin; }
//       }
//       EQUATION { <D0> + <D1>*CAPACITANCE + <D2>*SLEWRATE }
//    }
   }
}
```

Now the template is used for instantiation of a soft macro:

```
adder = dynamic { /* a soft macro */
   cellname = ripple_carry_adder_N_bit;
   areavalue = 20 + 30 * bitwidth;
   }
   D0 {
      HEADER { AREA { TABLE { 10 20 30 } } }
      TABLE { 15.6 34.3 50.7 }
   }
   D1 = 0.29;
   D2 = 0.08;
}
```

The dynamic instantiation of the soft macro results in an object for which certain data depend on the runtime-values of the placeholder-parameters, as indicated in *italic* below. The calculation method for such data, however, can be compiled statically (e.g., the equation for `AREA` is a function of bitwidth and the lookup table for `D0` is a function of `AREA`).

```
CELL ripple_carry_adder_N_bit {
   PIN [ bitwidth : 1 ] A { DIRECTION = input; }
   PIN [ bitwidth : 1 ] B { DIRECTION = input; }
   PIN Cin { DIRECTION = input; }
   PIN [ bitwidth : 1 ] S { DIRECTION = output; }
   PIN Cout { DIRECTION = output; }

   FUNCTION {
      BEHAVIOR {
         S = A + B + Cin;
         Cout = (A + B + Cin >= (‘b1 << (bitwidth - 1)));
      }
   }

   AREA = 20 + 30 * bitwidth ;

   VECTOR (?! Cin -> ?! Cout) {
      DELAY {
         HEADER {
            CAPACITANCE {PIN = Cout; }
            SLEWRATE {PIN = Cin; }
            D0 {
```

```
                HEADER { AREA { TABLE { 10 20 30 } } }
                TABLE { 15.6 34.3 50.7 }
             }
          }
          EQUATION { D0 + 0.29*CAPACITANCE + 0.08*SLEWRATE }
       }
    }
}
```

# Section 6

# Modeling for Synthesis and Test

## 6.1 Annotations and attributes for a CELL

This section defines various CELL annotations and attributes.

### 6.1.1 CELLTYPE annotation

CELLTYPE classifies the functionality of cells into broad categories. This is useful for information purpose, for tools which do not need the exact specification of functionality, and for tools which can interpret the exact specification of functionality only for certain categories of cells. The exact specification of the functionality is described in the FUNCTION statement.

```
CELLTYPE = string ;
```

which can take the values shown in Table 6-1.

**Table 6-1 : CELLTYPE annotations for a CELL object**

| Annotation string | Description |
|---|---|
| buffer | cell is a buffer, inverting or non-inverting |
| combinational | cell is a combinational logic element |
| multiplexor | cell is a multiplexor |
| flipflop | cell is a flip-flop |
| latch | cell is a latch |
| memory | cell is a memory or a register file |
| block | cell is a hierarchical block, i.e., a complex element which can be represented as a netlist. All instances of the netlist are library elements, i.e., there is a CELL model for each of them in the library. |
| core | cell is a core, i.e., a complex element which can be represented as a netlist. At least one instance of the netlist is not a library element, i.e., there is no CELL model, but a PRIMITIVE model for that instance. |
| special | cell is a special element, which can only be used in certain application contexts not describable by the FUNCTION statement. Examples: busholders, protection diodes, and fillcells. |

### 6.1.2 ATTRIBUTE within a CELL object

An ATTRIBUTE within a CELL classifies the functionality given by CELLTYPE in more detail.

The attributes shown in Table 6-2 can be used within a CELL with CELLTYPE=memory.

**Table 6-2 : Attributes within a CELL with CELLTYPE=memory**

| Attribute item | Description |
|---|---|
| RAM | Random Access Memory |
| ROM | Read Only Memory |
| CAM | Content Addressable Memory |
| static | static memory (e.g., static RAM) |
| dynamic | dynamic memory (e.g., dynamic RAM) |
| asynchronous | asynchronous memory |
| synchronous | synchronous memory |

The attributes shown in Table 6-3 can be used within a CELL with CELLTYPE=block.

**Table 6-3 : Attributes within a CELL with CELLTYPE=block**

| Attribute item | Description |
|---|---|
| counter | cell is a complex sequential cell going through a predefined sequence of states in its normal operation mode where each state represents an encoded control value. |
| shift_register | cell is a complex sequential cell going through a predefined sequence of states in its normal operation mode, where each subsequent state can be obtained from the previous one by a shift operation. Each bit represents a data value. |
| adder | cell is an adder, i.e., a combinational element performing an addition of two operands. |
| subtractor | cell is a subtractor, i.e., a combinational element performing a subtraction of two operands. |
| multiplier | cell is a multiplier, i.e., a combinational element performing a multiplication of two operands. |
| comparator | cell is a comparator, i.e., a combinational element comparing the magnitude of two operands. |
| ALU | cell is an arithmetic logic unit, i.e., a combinational element combining the functionality of adder, subtractor, comparator in a selectable way. |

The attributes shown in Table 6-4 can be used within a CELL with CELLTYPE=core.

**Table 6-4 : Attributes within a CELL with CELLTYPE=core**

| Attribute item | Description |
|---|---|
| PLL | CELL is a phase-locked loop |
| DSP | CELL is a digital signal processor |
| CPU | CELL is a central processing unit |
| GPU | CELL is a graphical processing unit |

The attributes shown in Table 6-5 can be used within a CELL with CELLTYPE=special.

**Table 6-5 : Attributes within a CELL with CELLTYPE=special**

| Attribute item | Description |
|---|---|
| busholder | CELL enables a tristate bus to hold its last value before all drivers went into high-impedance state (detail see FUNCTION statement) |
| clamp | CELL connects a net to a constant value (logic value and drive strength see FUNCTION statement) |
| diode | CELL is a diode (no FUNCTION statement) |
| capacitor | CELL is a capacitor (no FUNCTION statement) |
| resistor | CELL is a resistor (no FUNCTION statement) |
| inductor | CELL is an inductor (no FUNCTION statement) |
| fillcell | CELL is merely used to fill unused space in layout (no FUNCTION statement) |

### 6.1.3 SWAP_CLASS annotation

```
SWAP_CLASS = string ;
```

The value is the name of a declared CLASS. Multi-value annotation can be used. Cells referring to the same CLASS can be swapped for certain applications.

Cell-swapping is only allowed under the following conditions:

- the RESTRICT_CLASS annotation (see Section 6.1.4) authorizes usage of the cell
- the cells to be swapped are compatible from an application standpoint (functional compatibility for synthesis and physical compatibility for layout)

### 6.1.4 RESTRICT_CLASS annotation

```
RESTRICT_CLASS = string ;
```

The value is the name of a declared CLASS. Multi-value annotation can be used. Cells referring to a particular class can be used in design tools identified by the value. The restricted annotations are shown in Table 6-6.

**Table 6-6 : Predefined values for RESTRICT_CLASS**

| Annotation string | Description |
|---|---|
| synthesis | use restricted to logic synthesis |
| scan | use restricted to scan synthesis |
| datapath | use restricted to datapath synthesis |
| clock | use restricted to clock tree synthesis |
| layout | use restricted to layout, i.e., place & route |

User-defined values are also possible. If a cell has no or only unknown values for RESTRICT_CLASS, the application tool shall not modify any instantiation of that cell in the design. However, the cell shall still be considered for analysis.

### 6.1.5 Independent SWAP_CLASS and RESTRICT CLASS

SWAP_CLASS and RESTRICT_CLASS may be defined for cells, independent of each other. In this case, the set of cells that can be swapped with each other is the set of cells with a non-empty intersection of both SWAP_CLASS and RESTRICT_CLASS.

Example:

```
CLASS foo;
CLASS bar;
CLASS whatever;
CLASS my_tool;
CELL cell1 {
   SWAP_CLASS { foo bar }
   RESTRICT_CLASS { synthesis datapath }
}
CELL cell2 {
   SWAP_CLASS { foo whatever }
   RESTRICT_CLASS { synthesis scan my_tool }
}
```

The cells cell1 and cell2 can be used for synthesis, where they can be swapped which each other. Cell cell1 can be also used for datapath. Cell cell2 can be also used for scan insertion and for the user-defined application my_tool. Figure 6-1 depicts this scenario.

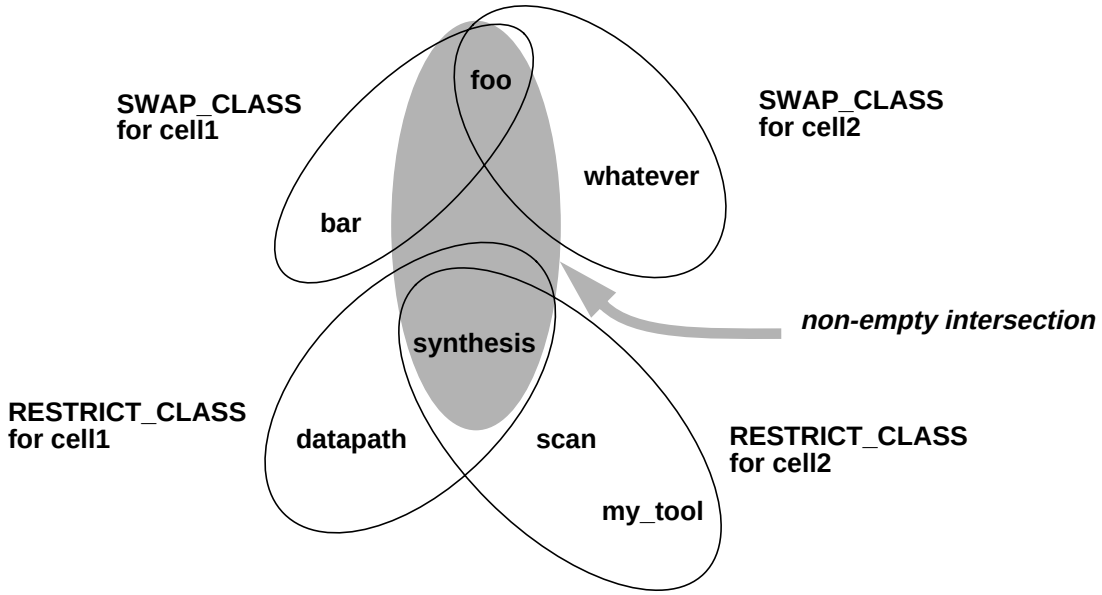



**Figure 6-1: Illustration of independent SWAP_CLASS and RESTRICT_CLASS**

### 6.1.6 SWAP_CLASS with inherited RESTRICT_CLASS

The definition of a CLASS may contain a RESTRICT_CLASS annotation. In this case, the RESTRICT_CLASS is inherited by the SWAP_CLASS. Cells can only be swapped if the intersection of their SWAP_CLASS and the inherited RESTRICT_CLASS is non-empty.

Example :

A combination of `SWAP_CLASS` and `RESTRICT_CLASS` can be used to emulate the concept of "logically equivalent cells" and "electrically equivalent cells". A synthesis tool needs to know about "logically equivalent cells" for swapping. A layout tool needs to know about "electrically equivalent cells" for swapping.

```
CLASS all_nand2 { RESTRICT_CLASS { synthesis } }
CLASS all_high_power_nand2 { RESTRICT_CLASS { layout } }
CLASS all_low_power_nand2 { RESTRICT_CLASS { layout } }

CELL cell1 {
   SWAP_CLASS { all_nand2 all_low_power_nand2 }
}
CELL cell2 {
   SWAP_CLASS { all_nand2 all_high_power_nand2 }
}
CELL cell3 {
   SWAP_CLASS { all_low_power_nand2 }
}
CELL cell4 {
   SWAP_CLASS { all_high_power_nand2 }
}
```

`all_nand2` encompasses a set of logically equivalent cells.
`all_high_power_nand2` encompasses a set of electrically equivalent cells.
`all_low_power_nand2` encompasses another set of electrically equivalent cells.

The synthesis tool can swap `cell1` with `cell2`. The layout tool can swap `cell1` with `cell3` and `cell2` with `cell4`. Figure 6-2 depicts this scenario.

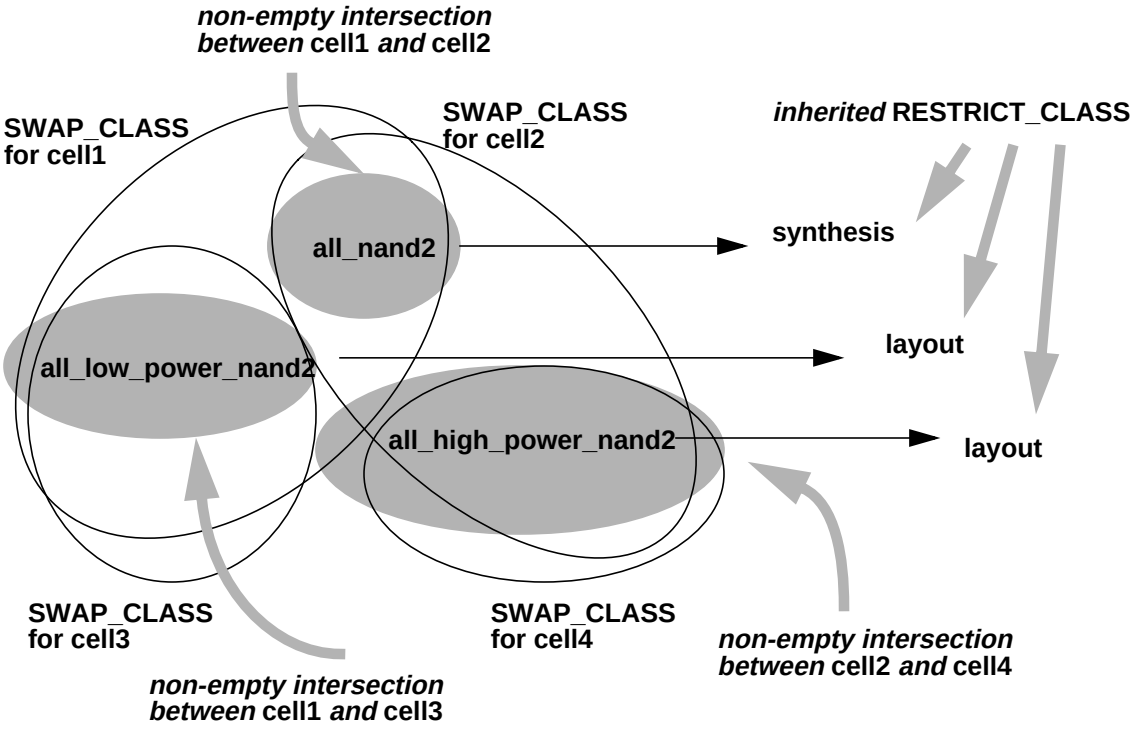



**Figure 6-2: Illustration of SWAP_CLASS with inherited RESTRICT_CLASS**

### 6.1.7 SCAN_TYPE annotation

**SCAN_TYPE =** `string` **;**

can take the values shown in Table 6-7.

**Table 6-7 : SCAN_TYPE annotations for a CELL object**

| Annotation string | Description |
|---|---|
| `muxscan` | a multiplexor for normal data and scan data |
| `clocked` | a special scan clock |
| `lssd` | combination between flip-flop and latch with special clocking (level sensitive scan design) |
| `control_0` | combinational scan cell, controlling pin shall be `0` in scan mode |
| `control_1` | combinational scan cell, controlling pin shall be `1` in scan mode |

### 6.1.8 SCAN_USAGE annotation

**SCAN_USAGE =** `string` **;**

can take the values shown in Table 6-8.

**Table 6-8 : SCAN_USAGE annotations for a CELL object**

| Annotation string | Description |
|---|---|
| `input` | primary input in a chain of cells |
| `output` | primary output in a chain of cells |
| `hold` | holds intermediate value in the scan chain |

### 6.1.9 BUFFERTYPE annotation

**`BUFFERTYPE =`** `string` **`;`**

can take the values shown in Table 6-9.

**Table 6-9 : BUFFERTYPE annotations for a CELL object**

| Annotation string | Description |
|---|---|
| `input` | cell has at least one external (off-chip) input pin |
| `output` | cell has at least one external (off-chip) output pin |
| `inout` | cell has at least one external (off-chip) bidirectional pin |
| `internal` | cell has only internal (on-chip) pins |

### 6.1.10 DRIVERTYPE annotation

**`DRIVERTYPE =`** `string` **`;`**

can take the values shown in Table 6-10.

**Table 6-10 DRIVERTYPE annotations for a CELL object**

| Annotation string | Description |
|---|---|
| `predriver` | cell is a predriver |
| `slotdriver` | cell is a slotdriver |
| `both` | cell is both a predriver and a slot driver |

Note: DRIVERTYPE applies only for cells with `BUFFERTYPE = input | output | inout`.

### 6.1.11 PARALLEL_DRIVE annotation

**`PARALLEL_DRIVE =`** `unsigned` **`;`**

specifies the number of parallel drivers. Must be greater than zero.

## 6.2 NON_SCAN_CELL statement

```
non_scan_cell ::=
   NON_SCAN_CELL { non_scan_cell_instantiations }
```

```
non_scan_cell_instantiations ::=
   non_scan_cell_instantiation { non_scan_cell_instantiation }

non_scan_cell_instantiation ::=
          cell_identifier { pin_assignments }
   |      primitive_identifier { pin_assignments }
```

In case of a single non-scan cell, the following syntax shall also be valid:

**NON_SCAN_CELL =** *non_scan*_cell_instantiation

This statement shall define non-scan cell equivalency to the scan cell in which this annotation is contained. A cell instantiation form is used to reference the library cell that defines the non-scan functionality of the current cell. If no such cell is available or defined, or if an explicit reference to such a cell is not desired, then a primitive instantiation form can reference a primitive, either ALF- or user- defined, for such use. In either case, constant values can appear on either the left-hand side or right-hand side of the pin connectivity relationships. A constant on the left-hand side defines the value the scan cell pins (appearing on the right-hand side) shall have in order for the primitive to perform with the same functionality as does the instantiated reference. A statement containing multiple non-scan cells shall indicate a choice between alternative non-scan cells.

Example:

```
CELL my_flip_flop {
      PIN q     { DIRECTION=output; }
      PIN d     { DIRECTION=input;  }
      PIN clk   { DIRECTION=input;  }
      PIN clear { DIRECTION=input; polarity=low; }
      // followed by function, vectors etc.
}

CELL my_other_flip_flop {
      // declare the pins
      // followed by function, vectors etc.
}

CELL my_scan_flip_flop {
      PIN data_out { DIRECTION=output; }
      PIN data_in  { DIRECTION=input;  }
      PIN clock    { DIRECTION=input;  }
      PIN scan_in  { DIRECTION=input;  }
      PIN scan_sel { DIRECTION=input;  }
      NON_SCAN_CELL {
            my_flip_flop {
                  q = data_out;
                  d = data_in;
                  clk = clock;
                  clear = 'b1;      // scan cell has no clear
```

```
                    'b0 = scan_in;    // non-scan cell has no scan_in
                    'b0 = scan_sel;   // non-scan cell has no scan_sel
                }
                my_other_flip_flop {
                // put in the pin assignments
                }
        }
        // followed by function, vectors etc.
}
```

Note: Both scan cells and the referenced non-scan cells must have at least the RESTRICT_CLASS value “scan”.

## 6.3 STRUCTURE statement

An optional STRUCTURE statement shall be legal in the context of a FUNCTION. The purpose of the STRUCTURE statement is to describe the structure of a complex cell composed of atomic cells, for example I/O buffers, LSSD flip-flops, or clock trees.

The syntax for the FUNCTION statement shall be augmented as follows:

```
function ::=
          FUNCTION [ identifier ] { [ all_purpose_items ] [primitives]
                [ behavior ] [ structure ] [ statetables ] }
   |      function_template_instantiation

structure ::=
          STRUCTURE { named_cell_instantiations }

named_cell_instantiations ::=
          named_cell_instantiation { named_cell_instantiation }

named_cell_instantiation ::=
          cell_identifier instance_identifier { logic_values }
   |      cell_identifier instance_identifier { pin_instantiations }
```

The STRUCTURE statement shall describe a netlist of components inside the CELL. The STRUCTURE statement shall not be a substitute for the BEHAVIOR statement. If a FUNCTION contains only a STRUCTURE statement and no BEHAVIOR statement, a behavior description for that particular cell shall be meaningless (e.g., fillcells, diodes, vias, or analog cells).

Timing and power models shall be provided for the CELL, if such models are meaningful. Application tools are not expected to use function, timing, or power models from the instantiated components as a substitute of a missing function, timing, or power model at the top-level. However, tools performing characterization, construction, or verification of a top-level model shall use the models of the instantiated components for this purpose.

Test synthesis applications can use the structural information in order to define a one-to-many mapping for scan cell replacement, such as where a single flip-flop is replaced by a pair of

master/slave latches. A macro cell can be defined whose structure is a netlist containing the master and slave latch and this shall contain the NON_SCAN_CELL annotation to define which sequential cells it is replacing. No timing model is required for this macro cell, since it should be treated as a transparent hierarchy level in the design netlist after test synthesis.

Notes:

1. Every *instance*_identifier within a STRUCTURE statement shall be different from each other.
2. The STRUCTURE statement provides a directive to the application (e.g., synthesis and DFT) as to how the CELL is implemented. A CELL referenced in named_cell_instantiation can be replaced by another CELL within the same SWAP_CLASS and RESTRICT_CLASS (recognized by the application).
3. The *cell*_identifier within a STRUCTURE statement can refer to actual cells as well as to primitives. The usage of primitives is recommended in fault modeling for DFT.
4. BEHAVIOR statements also provide the possibility of instantiating primitives. However, those instantiations are for modeling purposes only; they do not necessarily match a physical structure. The STRUCTURE statement always matches a physical structure.

Example 1:

iobuffer = pre buffer + main buffer

```
CELL my_main_driver {
   DRIVERTYPE = slotdriver ;
   BUFFERTYPE = output ;
   PIN i { DIRECTION = input; }
   PIN o { DIRECTION = output; }
   FUNCTION { BEHAVIOR { o = i ; } }
}
CELL my_pre_driver {
   DRIVERTYPE = predriver ;
   BUFFERTYPE = output ;
   PIN i { DIRECTION = input; }
   PIN o { DIRECTION = output; }
   FUNCTION { BEHAVIOR { o = i ; } }
}
CELL my_buffer {
   DRIVERTYPE = both ;
   BUFFERTYPE = output ;
   PIN A { DIRECTION = input; }
   PIN Z { DIRECTION = output; }
   PIN Y { VIEW = physical; }
   FUNCTION {
```

```
      BEHAVIOR { Z = A ; }
      STRUCTURE {
         my_pre_driver pre { A Y }// pin by order
         my_main_driver main { i=Y; o=Z; }// pin by name
      }
   }
}
```

Example 2:

lssd flip-flop = latch + flip-flop + mux

```
CELL my_latch {
   RESTRICT_CLASS { synthesis scan }
   PIN enable { DIRECTION = input;  }
   PIN d      { DIRECTION = input;  }
   PIN d      { DIRECTION = output; }
   FUNCTION { BEHAVIOR {
      @ ( enable ) { q = d ; }
   }  }
}
CELL my_flip-flop {
   RESTRICT_CLASS { synthesis scan }
   PIN clock  { DIRECTION = input;  }
   PIN d      { DIRECTION = input;  }
   PIN q      { DIRECTION = output; }
   FUNCTION { BEHAVIOR {
      @ ( 01 clock ) { q = d ; }
   }  }
}
CELL my_mux {
   RESTRICT_CLASS { synthesis scan }
   PIN dout   { DIRECTION = output; }
   PIN din0   { DIRECTION = input;  }
   PIN din1   { DIRECTION = input;  }
   PIN select { DIRECTION = input;  }
   FUNCTION { BEHAVIOR {
      dout = select ? din1 : din0 ;
   }  }
}
CELL my_lssd_flip-flop {
   RESTRICT_CLASS { scan }
   CELLTYPE = block;
   SCAN_TYPE = lssd;
   PIN clock         { DIRECTION = input;  }
   PIN master_clock  { DIRECTION = input;  }
   PIN slave_clock   { DIRECTION = input;  }
   PIN scan_data     { DIRECTION = input;  }
   PIN din           { DIRECTION = input;  }
   PIN dout          { DIRECTION = output; }
   PIN scan_master   { VIEW = physical; }
   PIN scan_slave    { VIEW = physical; }
   PIN d_internal    { VIEW = physical; }
   FUNCTION { BEHAVIOR {
```

```
            @ ( master_clock ) {
               scan_data_master = scan_data ;
            }
            @ ( slave_clock & ! clock ) {
               dout = scan_data_master ;
            } : ( 01 clock ) {
               dout = din ;
         }  }
         STRUCTURE {
            my_latch U0 {
               enable = master_clock;
               din    = scan_data;
               dout   = scan_data_master;
            }
            my_flip-flop U1 {
               clock  = clock;
               d      = din;
               q      = d_internal;
            }
            my_mux U2 {
               select = slave_clock;
               din1   = scan_data_master;
               din0   = dout;
               dout   = scan_data_slave;
            }
            my_mux U3 {
               select = clock;
               din1   = d_internal;
               din0   = scan_data_slave;
               dout   = dout;
         }  }
      }
      NON_SCAN_CELL {
         my_flip_flop {
            clock = clock;
            d     = din;
            q     = dout;
            'b0   = slave_clock;
         }
      }
   }
```

Example 3:

clock tree = chains of clock buffers

```
   CELL my_root_buffer {
      RESTRICT_CLASS { clock }
      PIN i0 { DIRECTION = input; }
      PIN o0 { DIRECTION = output; }
      FUNCTION { BEHAVIOR { o0 = i0 ; } }
   }
```

```
CELL my_level1_buffer {
   RESTRICT_CLASS { clock }
   PIN i1 { DIRECTION = input; }
   PIN o1 { DIRECTION = output; }
   FUNCTION { BEHAVIOR { o1 = i1 ; } }
}
CELL my_level2_buffer {
   RESTRICT_CLASS { clock }
   PIN i2 { DIRECTION = input; }
   PIN o2 { DIRECTION = output; }
   FUNCTION { BEHAVIOR { o2 = i2 ; } }
}
CELL my_level3_buffer {
   RESTRICT_CLASS { clock }
   PIN i3 { DIRECTION = input; }
   PIN o3 { DIRECTION = output; }
   FUNCTION { BEHAVIOR { o3 = i3 ; } }
}
CELL my_tree_from_level2 {
   RESTRICT_CLASS { clock }
   PIN in  { DIRECTION = input; }
   PIN out { DIRECTION = output; }
   PIN[1:2] level3  { DIRECTION = output; }
   FUNCTION {
      BEHAVIOR { out = in ; }
      STRUCTURE {
         my_level2_buffer U1 { i2=in; o2=out; }
         my_level3_buffer U2 { i3=out; o3=level3[1]; }
         my_level3_buffer U3 { i3=out; o3=level3[2]; }
      }
   }
}
CELL my_tree_from_level1 {
   RESTRICT_CLASS { clock }
   PIN in  { DIRECTION = input; }
   PIN out { DIRECTION = output; }
   PIN[1:4] level2  { DIRECTION = output; }
   FUNCTION {
      BEHAVIOR { out = in ; }
      STRUCTURE {
         my_level1_buffer U1 { i1=in; o1=out; }
         my_tree_from_level2 U2 { i2=out; o2=level2[1]; }
         my_tree_from_level2 U3 { i2=out; o2=level2[2]; }
         my_tree_from_level2 U4 { i2=out; o2=level2[3]; }
         my_tree_from_level2 U5 { i2=out; o2=level2[4]; }
      }
   }
}
```

```
CELL my_tree_from_root {
   RESTRICT_CLASS { clock }
   PIN in  { DIRECTION = input; }
   PIN out { DIRECTION = output; }
   PIN[1:4] level1  { DIRECTION = output; }
   FUNCTION {
      BEHAVIOR { out = in ; }
      STRUCTURE {
         my_root_buffer U1 { i0=in; o0=out; }
         my_tree_from_level1 U2 { i1=o; o1=level1[1]; }
         my_tree_from_level1 U3 { i1=o; o1=level1[2]; }
         my_tree_from_level1 U4 { i1=o; o1=level1[3]; }
         my_tree_from_level1 U5 { i1=o; o1=level1[4]; }
      }
   }
}
```

Example 4:

Multiplexor, showing the conceptional difference between BEHAVIOR and STRUCTURE.

```
CELL my_multiplexor {
   PIN a { DIRECTION = input; }
   PIN b { DIRECTION = input; }
   PIN s { DIRECTION = input; }
   PIN y { DIRECTION = output; }
   FUNCTION {
      BEHAVIOR {
// s_a and s_b are virtual internal nodes
         ALF_AND { out = s_a; in[0] = !s; in[1] = a; }
         ALF_AND { out = s_b; in[0] =  s; in[1] = b; }
         ALF_OR  { out = y; in[0] = s_a; in[1] = s_b; }
      }
      STRUCTURE {
// sbar, sel_a, sel_b are physical internal nodes
         ALF_NOT  { out = sbar; in = s; }
         ALF_NAND { out = sel_a; in[0] = sbar; in[1] = a; }
         ALF_NAND { out = sel_b; in[0] =  s;   in[1] = b; }
         ALF_NAND { out = y; in[0] = sel_a; in[1] = sel_b; }
      }
   }
}
```

## 6.4 Annotations and attributes for a PIN

This section defines various `PIN` annotations and attributes.

### 6.4.1 VIEW annotation

**`VIEW =`** `string` **`;`**

annotates the view where the pin appears, which can take the values shown in Table 6-11.

**Table 6-11 : VIEW annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `functional` | pin appears in functional netlist |
| `physical` | pin appears in physical netlist |
| `both` (default) | pin appears in both functional and physical netlist |
| `none` | pin does not appear in netlist |

### 6.4.2 PINTYPE annotation

**`PINTYPE =`** `string` **`;`**

annotates the type of the pin, which can take the values shown in Table 6-12.

**Table 6-12 : PINTYPE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `digital` (default) | digital signal pin |
| `analog` | analog signal pin |
| `supply` | power supply or ground pin |

### 6.4.3 DIRECTION annotation

**`DIRECTION =`** `string` **`;`**

annotates the direction of the pin, which can take the values shown in Table 6-13.

**Table 6-13 : DIRECTION annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `input` | input pin |
| `output` | output pin |
| `both` | bidirectional pin |
| `none` | no direction can be assigned to the pin |

Table 6-14 gives a more detailed semantic interpretation for using `DIRECTION` in combination with `PINTYPE`.

**Table 6-14 : DIRECTION in combination with PINTYPE**

| DIRECTION | PINTYPE=digital | PINTYPE=analog | PINTYPE=supply |
|---|---|---|---|
| `input` | pin receives a digital signal | pin receives an analog signal | pin is a power sink |
| `output` | pin drives a digital signal | pin drives an analog signal | pin is a power source |

**Table 6-14 : DIRECTION in combination with PINTYPE, *continued***

| DIRECTION | PINTYPE=digital | PINTYPE=analog | PINTYPE=supply |
|---|---|---|---|
| `both` | pin drives or receives a digital signal, depending on the operation mode | pin drives or receives an analog signal, depending on the operation mode | pin is both power sink and source |
| `none` | pin represents either an internal digital signal with no external connection or a feed through | pin represents either an internal analog signal with no external connection or a feed through | pin represents either an internal power pin with no external connection or a feed through |

Examples:

- The power and ground pins of regular cells shall have `DIRECTION=input`.
- A level converter cell shall have a power supply pin with `DIRECTION=input` and another power supply pin with `DIRECTION=output`.
- A level converter can have separate ground pins on the input and output side or a common ground pin with `DIRECTION=both`.
- The power and ground pins of a feed through cell shall have `DIRECTION=none`.

### 6.4.4 SIGNALTYPE annotation

`SIGNALTYPE` classifies the functionality of a pin. The currently defined values apply for pins with `PINTYPE=DIGITAL`.

Conceptually, a pin with `PINTYPE = ANALOG` can also have a `SIGNALTYPE` annotation. However, no values are currently defined.

**`SIGNALTYPE =`** `string` **`;`**

annotates the type of the signal connected to the pin.

The fundamental `SIGNALTYPE` values are defined in Table 6-15.

**Table 6-15 : Fundamental SIGNALTYPE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `data` (default) | general data signal, i.e., a signal that carries information to be transmitted, received, or subjected to logic operations within the `CELL`. |
| `address` | address signal of a memory, i.e., an encoded signal, usually a bus or part of a bus, driving an address decoder within the `CELL`. |
| `control` | general control signal, i.e., an encoded signal that controls at least two modes of operation of the `CELL`, eventually in conjunction with other signals. The signal value is allowed to change during real-time circuit operation. |

**Table 6-15 : Fundamental SIGNALTYPE annotations for a PIN object, *continued***

| Annotation string | Description |
|---|---|
| `select` | select signal of a multiplexor, i.e., a decoded or encoded signal that selects the data path of a multiplexor or de-multiplexor within the `CELL`. Each selected signal has the same `SIGNALTYPE`. |
| `enable` | general enable signal, i.e., a decoded signal which enables and disables a set of operational modes of the `CELL`, eventually in conjunction with other signals. The signal value is expected to change during real-time circuit operation. |
| `tie` | the signal needs to be tied to a fixed value statically in order to define a fixed or programmable mode of operation of the `CELL`, eventually in conjunction with other signals. The signal value is not allowed to change during real-time circuit operation. |
| `clear` | clear signal of a flip-flop or latch, i.e., a signal that controls the storage of the value `0` within the `CELL`. |
| `set` | set signal of a flip-flop or latch, i.e., a signal that controls the storage of the value `1` within the `CELL`. |
| `clock` | clock signal of a flip-flop or latch, i.e., a timing-critical signal that triggers data storage within the `CELL`. |

"Flipflop", "latch", "multiplexor", and "memory" can be standalone cells or embedded in larger cells. In the former case, the celltype is `flipflop`, `latch`, `multiplexor`, and `memory`, respectively. In the latter case, the celltype is `block` or `core`.

Composite values for `SIGNALTYPE` shall be constructed using one or more prefixes in combination with certain fundamental values, separated by the underscore (_) character, as shown in Table 6-16 through Table 6-20.

The scheme for this is shown in Figure 6-3.

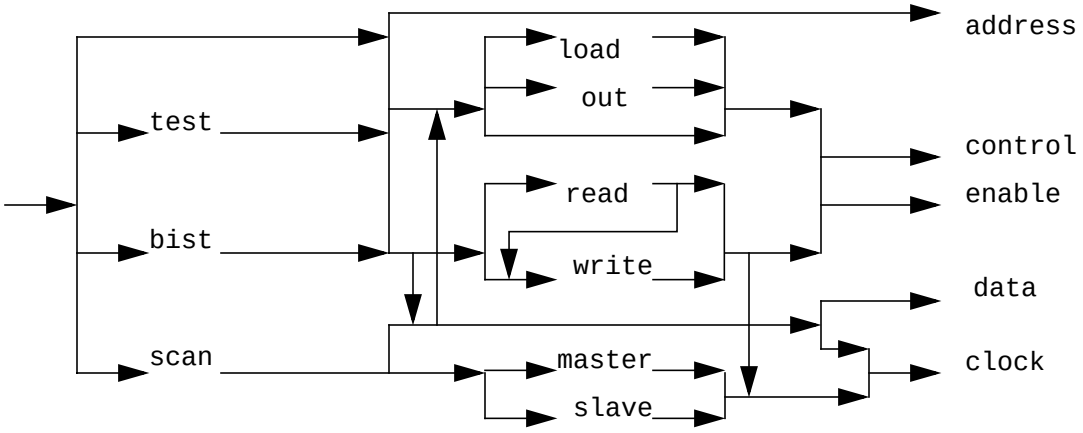



**Figure 6-3: Construction scheme for composite SIGNALTYPE values**

**Table 6-16 : Composite SIGNALTYPE annotations based on DATA**

| Annotation string | Description |
|---|---|
| scan_data | data signal for scan mode |
| test_data | data signal for test mode |
| bist_data | data signal in BIST mode |

**Table 6-17 : Composite SIGNALTYPE annotations based on ADDRESS**

| Annotation string | Description |
|---|---|
| test_address | address signal for test mode |
| bist_address | address signal for BIST mode |

**Table 6-18 : Composite SIGNALTYPE annotations based on CONTROL**

| Annotation string | Description |
|---|---|
| load_control | control signal for switching between load mode and normal mode |
| scan_control | control signal for switching between scan mode and normal mode |
| test_control | control signal for switching between test mode and normal mode |
| bist_control | control signal for switching between BIST mode and normal mode |
| read_write_control | control signal for switching between read and write operation |
| test_read_write_control | control signal for switching between read and write operation in test mode |
| bist_read_write_control | control signal for switching between read and write operation in BIST mode |

**Table 6-19 : Composite SIGNALTYPE annotations based on ENABLE**

| Annotation string | Description |
|---|---|
| load_enable | signal enables load operation in a counter or a shift register |
| out_enable | signal enables the output stage of an arbitrary cell |
| scan_enable | signal enables scan mode of a flip-flop or latch only |
| scan_out_enable | signal enables the output of a flip-flop or latch in scan mode only |
| test_enable | signal enables test mode only |

**Table 6-19 : Composite SIGNALTYPE annotations based on ENABLE, *continued***

| Annotation string | Description |
|---|---|
| bist_enable | signal enables BIST mode only |
| test_out_enable | signal enables the output stage in test mode only |
| bist_out_enable | signal enables the output stage in BIST mode only |
| read_enable | signal enables the read operation of a memory |
| write_enable | signal enables the write operation of a memory |
| test_read_enable | signal enables the read operation in test mode only |
| test_write_enable | signal enables the write operation in test mode only |
| bist_read_enable | signal enables the read operation in BIST mode only |
| bist_write_enable | signal enables the write operation in BIST mode only |

:

**Table 6-20 : Composite SIGNALTYPE annotations based on CLOCK**

| Annotation string | Description |
|---|---|
| scan_clock | signal is clock of a flip-flop or latch in scan mode |
| master_clock | signal is master clock of a flip-flop or latch |
| slave_clock | signal is slave clock of a flip-flop or latch |
| scan_master_clock | signal is master clock of a flip-flop or latch in scan mode |
| scan_slave_clock | signal is slave clock of a flip-flop or latch in scan mode |
| read_clock | clock signal triggers the read operation in a synchronous memory |
| write_clock | clock signal triggers the write operation in a synchronous memory |
| read_write_clock | clock signal triggers both read and write operation in a synchronous memory |
| test_clock | signal is clock in test mode |
| test_read_clock | clock signal triggers the read operation in a synchronous memory in test mode |
| test_write_clock | clock signal triggers the write operation in a synchronous memory in test mode |
| test_read_write_clock | clock signal triggers both read and write operation in a synchronous memory in test mode |
| bist_clock | signal is clock in BIST mode |
| bist_read_clock | clock signal triggers the read operation in a synchronous memory in BIST mode |
| bist_write_clock | clock signal triggers the write operation in a synchronous memory in BIST mode |
| bist_read_write_clock | clock signal triggers both read and write operation in a synchronous memory in BIST mode |

### 6.4.5 ACTION annotation

**ACTION =** string **;**

annotates the action of the signal, which can take the values shown in Table 6-21.

**Table 6-21 : ACTION annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `synchronous` | signal acts in synchronous way, i.e., self-triggered |
| `asynchronous` | signal acts in asynchronous way, i.e., triggered by a signal with `SIGNALTYPE CLOCK` or a composite `SIGNALTYPE` with postfix `_CLOCK`. |

The `ACTION` annotation applies only to pins with certain `SIGNALTYPE` values, as shown in Table 6-22. The rule applies also to any composite `SIGNALTYPE` values based on the fundamental values.

**Table 6-22 : ACTION applicable in conjunction with fundamental SIGNALTYPE values**

| fundamental SIGNALTYPE | applicable ACTION |
|---|---|
| `data` | N/A |
| `address` | N/A |
| `control` | synchronous or asynchronous |
| `select` | N/A |
| `enable` | synchronous or asynchronous |
| `tie` | N/A |
| `clear` | synchronous or asynchronous |
| `set` | synchronous or asynchronous |
| `clock` | N/A, but the presence of `SIGNALTYPE=clock` conditions the validity of `ACTION=synchronous` for other signals |

### 6.4.6 POLARITY annotation

**`POLARITY =`** `string` **`;`**

annotates the polarity of the pin signal.

The polarity of an input pin (i.e., `DIRECTION = input;`) takes the values shown in Table 6-23.

**Table 6-23 : POLARITY annotations for a PIN**

| Annotation string | Description |
|---|---|
| `high` | signal active high or to be driven high |
| `low` | signal active low or to be driven low |
| `rising_edge` | signal sensitive to rising edge |
| `falling_edge` | signal sensitive to falling edge |
| `double_edge` | signal sensitive to any edge |

The POLARITY annotation applies only to pins with certain SIGNALTYPE values, as shown iTable 6-24. The rule applies also to any composite SIGNALTYPE values based on the fundamental values.

**Table 6-24 : POLARITY applicable in conjunction with fundamental SIGNALTYPE values**

| fundamental SIGNALTYPE | applicable POLARITY value |
|---|---|
| data | N/A |
| address | N/A |
| control | mode-specific high or low for composite signaltype |
| select | N/A |
| enable | Mandatory high or low |
| tie | Optional high or low |
| clear | Mandatory high or low |
| set | Mandatory high or low |
| clock | Mandatory high, low, rising_edge, falling_edge, or double_edge, can be mode-specific for composite signaltype. |

Signals with composite signaltypes *mode*_CLOCK can have a single polarity or mode-specific polarities.

Example:

```
PIN rw {
   SIGNALTYPE = READ_WRITE_CONTROL;
   POLARITY { READ=high; WRITE=low; }
}
PIN rwc {
   SIGNALTYPE = READ_WRITE_CLOCK;
   POLARITY { READ=rising_edge; WRITE=falling_edge; }
}
```

### 6.4.7 DATATYPE annotation

**DATATYPE =** string **;**

annotates the datatype of the pin, which can take the values shown in Table 6-25.

**Table 6-25 : DATATYPE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| signed | result of arithmetic operation is signed 2's complement |
| unsigned | result of arithmetic operation is unsigned |

DATATYPE is only relevant for bus pins.

### 6.4.8 INITIAL_VALUE annotation

**INITIAL_VALUE =** logic_constant **;**

shall be compatible with the buswidth and DATATYPE of the signal.

INITIAL_VALUE is used for a downstream behavioral simulation model, as far as the simulator (e.g., a VITAL-compliant simulator) supports the notion of initial value.

### 6.4.9 SCAN_POSITION annotation

**SCAN_POSITION =** unsigned **;**

annotates the position of the pin in scan chain, starting with 1. Value 0 (default) indicates that the PIN is not on the scan chain.

### 6.4.10 STUCK annotation

**STUCK =** string **;**

annotates the stuck-at fault model as shown in Table 6-26.

**Table 6-26 : STUCK annotations for a PIN object**

| Annotation string | Description |
|---|---|
| stuck_at_0 | pin can have stuck-at-0 fault |
| stuck_at_1 | pin can have stuck-at-1 fault |
| both (default) | pin can have both stuck-at-0 and stuck-at-1 faults |
| none | pin can not have stuck-at faults |

### 6.4.11 SUPPLYTYPE

A PIN with PINTYPE = SUPPLY shall have a SUPPLYTYPE annotation.

```
supplytype_assignment ::=
   SUPPLYTYPE = supplytype_identifier ;

supplytype_identifier ::=
   power
|  ground
|  reference
```

### 6.4.12 SIGNAL_CLASS

The following new keyword for class reference shall be defined:

**SIGNAL_CLASS**
a PIN referring to the same SIGNAL_CLASS belong to the same logic port.
For example, the ADDRESS, WRITE_ENABLE, and DATA pin of a logic port of a memory have the same SIGNAL_CLASS.
SIGNAL_CLASS applies to a PIN with PINTYPE=DIGITAL | ANALOG.
SIGNAL_CLASS is orthogonal to SIGNALTYPE.

Example:

```
CLASS portA;
CLASS portB;
CELL my_memory {
   PIN[1:4] addrA { DIRECTION = input;
      SIGNALTYPE = address;
      SIGNAL_CLASS = portA;
   }
   PIN[7:0] dataA { DIRECTION = output;
      SIGNALTYPE = data;
      SIGNAL_CLASS = portA;
   }
   PIN[1:4] addrB { DIRECTION = input;
      SIGNALTYPE = address;
      SIGNAL_CLASS = portB;
   }
   PIN[7:0] dataB { DIRECTION = input;
      SIGNALTYPE = data;
      SIGNAL_CLASS = portB;
   }
   PIN weB { DIRECTION = input;
      SIGNALTYPE = write_enable;
      SIGNAL_CLASS = portB;
   }
}
```

Note: The combination of `SIGNAL_CLASS` and `SIGNALTYPE` identifies the port type. `CLASS portA` represents a read port, since it consists of a `PIN` with `SIGNALTYPE = address` and a `PIN` with `SIGNALTYPE = data` and `DIRECTION = output`. `CLASS portB` represents a write port, since it consists of a `PIN` with `SIGNALTYPE = address`, a `PIN` with `SIGNALTYPE = data` and `DIRECTION = input`, and a `PIN` with `SIGNALTYPE = write_enable`.

### 6.4.13 SUPPLY_CLASS

The following new keyword for class reference shall be defined:

**SUPPLY_CLASS**
a `PIN` referring to the same `SUPPLY_CLASS` belongs to the same power terminal.
For example, digital `VDD` and digital `VSS` have the same `SUPPLY_CLASS`.
`SIGNAL_CLASS` applies to a `PIN` with `PINTYPE=SUPPLY`.
`SUPPLY_CLASS` is orthogonal to `SUPPLYTYPE`.

Example:

```
CELL my_core {
   PIN vdd_dig { SUPPLYTYPE = power; SUPPLY_CLASS = digital; }
   PIN vss_dig { SUPPLYTYPE = ground; SUPPLY_CLASS = digital; }
   PIN vdd_ana { SUPPLYTYPE = power; SUPPLY_CLASS = analog; }
   PIN vss_ana { SUPPLYTYPE = ground; SUPPLY_CLASS = analog; }
}
```

### 6.4.14 Driver CELL and PIN specification

The keywords `CELL` and `PIN` can be used as references to existing objects to define a driver cell and pin in a macro, i.e., a cell with `CELLTYPE=block`.

Example:

```
// this is a standard ASIC cell
CELL my_inv {
   CELLTYPE = buffer;
   PIN in { DIRECTION = input; }
   PIN out { DIRECTION = output; }
}

// this is a macro, synthesized from standard ASIC cells
CELL my_macro {
   CELLTYPE = block;
   PIN my_output {
      DIRECTION = output;
      CELL = my_inv { PIN = out; }
   }
   /* fill in other pins and stuff */
}
```

### 6.4.15 DRIVETYPE annotation

```
DRIVETYPE = string ;
```

annotates the drive type for the pin, which can take the values shown in Table 6-27.

**Table 6-27 : DRIVETYPE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `cmos` (default) | standard cmos signal |
| `nmos` | nmos or pseudo nmos signal |
| `pmos` | pmos or pseudo pmos signal |
| `nmos_pass` | nmos passgate signal |
| `pmos_pass` | pmos passgate signal |
| `cmos_pass` | cmos passgate signal, i.e., the full transmission gate |
| `ttl` | TTL signal |
| `open_drain` | open drain signal |
| `open_source` | open source signal |

### 6.4.16 SCOPE annotation

```
SCOPE = string ;
```

annotates the modeling scope of a pin, which can take the values shown in Table 6-28.

**Table 6-28 : SCOPE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `behavior` | the in is used for modeling functional behavior and events on the pin are monitored for vector expressions in `BEHAVIOR` statements |
| `measure` | measurements related to the pin can be described, e.g., timing or power characterization, and events on the pin are monitored for vector expressions in `VECTOR` statements |
| `both` (default) | the pin is used for functional behavior as well as for characterization measurements |
| `none` | no model; only the pin exists |

### 6.4.17 PULL annotation

```
PULL = string ;
```

annotates the pull type for the pin, which can take the values shown in Table 6-29.

**Table 6-29 : PULL annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `up` | pullup device connected to pin |
| `down` | pulldown device connected to pin |
| `both` | pullup and pulldown device connected to pin |
| `none` (default) | no pull device |

### 6.4.18 ATTRIBUTE for PIN objects

The attributes shown in Table 6-30 can be used within a `PIN` object.

**Table 6-30 : Attributes within a PIN object**

| Attribute item | Description |
|---|---|
| `SCHMITT` | Schmitt trigger signal |
| `TRISTATE` | tristate signal |
| `XTAL` | crystal/oscillator signal |
| `PAD` | pad going off-chip |

The attributes shown in Table 6-31 are only applicable for pins within cells with `CELLTYPE=memory` and certain values of `SIGNALTYPE`.

**Table 6-31 : Attributes for pins of a memory**

| Attribute item | SIGNALTYPE | Description |
|---|---|---|
| `ROW_ADDRESS_STROBE` | `clock` | samples the row address of the memory |
| `COLUMN_ADDRESS_STROBE` | `clock` | samples the column address of the memory |

**Table 6-31 : Attributes for pins of a memory, *continued***

| Attribute item | SIGNALTYPE | Description |
|---|---|---|
| ROW | address | selects an addressable row of the memory |
| COLUMN | address | selects an addressable column of the memory |
| BANK | address | selects an addressable bank of the memory |

The attributes shown in Table 6-32 are only applicable for pins representing double-rail signals.

**Table 6-32 : Attributes for pins representing double-rail signals**

| Attribute item | Description |
|---|---|
| INVERTED | represents the inverted value within a pair of signals carrying complementary values |
| NON_INVERTED | represents the non-inverted value within a pair of signals carrying complementary values |
| DIFFERENTIAL | signal is part of a differential pair, i.e., both the inverted and non-inverted values are always required for physical implementation |

The following restrictions apply for double-rail signals:

- The PINTYPE, SIGNALTYPE, and DIRECTION of both pins shall be the same.
- One PIN shall have the attribute INVERTED, the other NON_INVERTED.
- Either both pins or no pins shall have the attribute DIFFERENTIAL.
- POLARITY, if applicable, shall be complementary as follows:
  HIGH is paired with LOW
  RISING_EDGE is paired with FALLING_EDGE
  DOUBLE_EDGE is paired with DOUBLE_EDGE

# 6.5 Definitions for bus pins

This section defines how to specify bus pins and group pins.

## 6.5.1 RANGE for bus pins

A one-dimensional bus pin can contain a RANGE statement, defined as follows:

```
range ::=
        RANGE { unsigned : unsigned }
```

The RANGE statement applies only if the range of valid indices is contiguous. The range is limited by the width of the bus. The possible range for a N-bit wide bus is between 0 and $2^N$. The possible range of values shall also be the default range.

Example:

A 4-bit wide bus has the following possible range of indices: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15.

`RANGE { 3 : 13 }` specifies the indices `0`, `1`, `2`, `14`, and `15` are invalid.

In the case where non-contiguous indices are valid, for example 1, 2, 3, 5, 6, 7, 9, 10, 11, 13, 14, 15, the `RANGE` statement does not apply.

### 6.5.2 Scalar pins inside a bus

A `PIN` declared as a bus shall contain the optional `pin_instantiation` statement, defined as follows:

```
pin_instantiation ::=
          pin_identifier [ index ] {
                pin_items
          }
```

where `index` and `pin_items` are defined in Section 11.5 and Section 11.11, respectively.

A `pin_instantiation` statement can also refer to a part of the bus.

Annotations within the scope of the `PIN` or a higher-level `pin_instantiation` shall be inherited by a lower-level `pin_instantiation` (see Section 6.4), as long as their values are applicable for both the bus and each scalar pin within the bus. Values of `VIEW`, `INITIAL_VALUE`, and arithmetic models such as `CAPACITANCE` shall not be inherited, since a particular value cannot apply at the same time to the bus and to its scalar pins.

Example:

```
PIN [1:4] my_address {
   DIRECTION = input;
   SIGNALTYPE = address;
   VIEW = functional;
   CAPACITANCE = 0.07;
   my_address [1:2] { ATTRIBUTE { ROW } CAPACITANCE = 0.03; }
      my_address[1] { VIEW = physical; CAPACITANCE = 0.01; }
      my_address[2] { VIEW = physical; CAPACITANCE = 0.02; }
      my_address [3:4] { ATTRIBUTE { COLUMN } CAPACITANCE = 0.04; }
      my_address[3] { VIEW = physical; CAPACITANCE = 0.02; }
      my_address[4] { VIEW = physical; CAPACITANCE = 0.02; }
   }
}
```

### 6.5.3 PIN_GROUP statement

A pin group shall be defined as follows:

```
pin_group ::=
          PIN_GROUP [ index ] pin_group_identifier {
                pin_items
                MEMBERS { pins }
          }
```

where pin_items is defined in Section 11.11.

The pins in the MEMBERS field shall refer to previously defined pins. The range of the index, if defined, shall match the number and range of pins in the MEMBERS field.

Annotations within the scope of the PIN contained in the MEMBERS field shall be inherited by the PIN_GROUP, as long as their values are applicable for both the pin and the pin group. Values of VIEW, INITIAL_VALUE, and arithmetic models such as CAPACITANCE shall not be inherited, since a particular value cannot apply at the same time to the pin and the pin group.

A pin group with VIEW=functional shall be treated like a bus pin in the functional netlist. It shall appear in the netlist in place of the first defined pin within the MEMBERS field.

Example 1:

```
PIN my_address_1 {DIRECTION = input; VIEW = physical; CAPACITANCE = 0.01;}
PIN my_address_2 {DIRECTION = input; VIEW = physical; CAPACITANCE = 0.02;}
PIN my_address_3 {DIRECTION = input; VIEW = physical; CAPACITANCE = 0.02;}
PIN my_address_4 {DIRECTION = input; VIEW = physical; CAPACITANCE = 0.02;}
PIN_GROUP [1:2] my_address_1_2 {
   ATTRIBUTE { ROW }
   CAPACITANCE = 0.03;
   MEMBERS { my_address_1 my_address_2 }
}
PIN_GROUP [1:2] my_address_3_4 {
   ATTRIBUTE { COLUMN }
   CAPACITANCE = 0.03;
   MEMBERS { my_address_3 my_address_4 }
}
PIN_GROUP [1:4] my_address {
   VIEW = functional;
   CAPACITANCE = 0.07;
   MEMBERS { my_address_1 my_address_2 my_address_3 my_address_4 }
}
```

Pairs of complementary pins, differential pins in particular, are special cases of pin groups.

Example 2:

```
CELL my_flip-flop {
   PIN CLK { DIRECTION=input; SIGNALTYPE=clock; POLARITY=rising_edge; }
   PIN D { DIRECTION=input; SIGNALTYPE=data; }
   PIN Q { DIRECTION=output; SIGNALTYPE=data; ATTRIBUTE { NON_INVERTED } }
   PIN Qbar { DIRECTION=output; SIGNALTYPE=data; ATTRIBUTE { INVERTED } }
   PIN_GROUP [0:1] Q_double_rail { RANGE { 1 : 2 } MEMBERS { Q Qbar } }
}
```

The pins Q and Qbar are complementary. Their valid set of data comprises 'b01==='d1 and 'b10==='d2. The values 'b00==='d0 and 'b11==='d3 are invalid.

```
CELL my_differential_buffer {
   PIN DIN { DIRECTION=input; ATTRIBUTE { DIFFERENTIAL NON_INVERTED } }
   PIN DINN { DIRECTION=input; ATTRIBUTE { DIFFERENTIAL INVERTED } }
   PIN DOUT { DIRECTION=output; ATTRIBUTE { DIFFERENTIAL NON_INVERTED } }
   PIN DOUTN { DIRECTION=output; ATTRIBUTE { DIFFERENTIAL INVERTED } }
   PIN_GROUP [0:1] DI { RANGE { 1 : 2 } MEMBERS { DIN DINN } }
   PIN_GROUP [0:1] DO { RANGE { 1 : 2 } MEMBERS { DOUT DOUTN } }
}
```

The pins DIN and DINN represent a pair of differential input pins. The pins DOUT and DOUTN represent a pair of differential output pins.

## 6.6 Annotations for CLASS and VECTOR

This section defines the annotations for CLASS and VECTOR.

### 6.6.1 PURPOSE annotation

A CLASS is a generic object which can be referenced inside another object. An object referencing a class inherits all children object of that class. In addition to this general reference, the usage of the keyword CLASS in conjunction with a predefined prefix (e.g., CONNECT_CLASS, SWAP_CLASS, RESTRICT_CLASS, EXISTENCE_CLASS, or CHARACTERIZATION_CLASS) also carries a specific semantic meaning in the context of its usage. Note the keyword <prefix>_CLASS is used for referencing a class, whereas the definition of the class always uses the keyword CLASS. Thus a class can have multiple purposes. With the growing number of usage models of the class concept, it is useful to include the purpose definition in the class itself in order to make it easier for specific tools to identify the classes of relevance for that tool.

A CLASS object can contain the PURPOSE annotation, which can take one or multiple values. A VECTOR entitled to inherit the PURPOSE annotation from the CLASS can also contain the PURPOSE annotation as follows.

```
vector_purpose_assignment ::=
        PURPOSE { purpose_identifier { purpose_identifier } }

vector_purpose_identifier :: =
        bist
   |    test
   |    timing
   |    power
   |    integrity
```

### 6.6.2 OPERATION annotation

The OPERATION statement inside a VECTOR shall be used to indicate the combined definition of signal values or signal changes for certain operations which are not entirely controlled by a single signal.

```
operation_assignment ::=
        OPERATION = operation_identifier ;
```

An OPERATION within the context of a VECTOR indicates certain a function of a cell, such as a memory write, or change to some state, such as test mode. Many functions are not controlled by a single pin and are therefore not able to be defined by the use of SIGNALTYPE alone. The VECTOR shall describe the complete operation, including the sequence of events on input and expected output signals, such that one operation can be followed seamlessly by the next.

The following values shall be predefined:

```
operation_identifier ::=
   read
|  write
|  read_modify_write
|  write_through
|  start
|  end
|  refresh
|  load
|  iddq
```

Their definitions are:

- *read*: read operation at one address
- *write*: write operation at one address
- *read_modify_write*: read followed by write of different value at same address
- *start*: first operation required in a particular mode
- *end*: last operation required in a particular mode
- *refresh*: operation required to maintain the contents of the memory without modifying it
- *load*: operation for loading control registers
- *iddq*: operation for supply current measurements in quiescent state

With exception of “iddq”, all values apply for only cells with CELLTYPE=memory.

The EXISTENCE_CLASS (see Section 6.6.5) within the context of a VECTOR shall be used to identify which operations can be combined in the same mode. OPERATION is orthogonal to

EXISTENCE_CLASS. The EXISTENCE_CLASS statement is only necessary, if there is more than one mode of operation.

Example 1:

```
CLASS normal_mode { PURPOSE = test; }
CLASS fast_page_mode { PURPOSE = test; }
VECTOR ( ! WE && (
      ?! addr -> 01 RAS -> 10 RAS ->
      ?! addr -> 01 CAS -> X? dout -> 10 CAS -> ?X dout
   ) ) {
   OPERATION = read; EXISTENCE_CLASS = normal_mode;
}
VECTOR ( WE && (
      ?! addr -> 01 RAS -> 10 RAS ->
      ?! addr -> ?? din -> 01 CAS -> 10 CAS
   ) ) {
   OPERATION = write; EXISTENCE_CLASS = normal_mode;
}
VECTOR ( ! WE && (?! addr -> 01 CAS -> X? dout -> 10 CAS -> ?X dout ) ) {
   OPERATION = read; EXISTENCE_CLASS = fast_page_mode;
}
VECTOR ( WE && ( ?! addr -> ?? din -> 01 CAS -> 10 CAS ) ) {
   OPERATION = write; EXISTENCE_CLASS = fast_page_mode;
}
VECTOR ( ?! addr -> 01 RAS -> 10 RAS ) {
   OPERATION = start; EXISTENCE_CLASS = fast_page_mode;
}
```

Note: The complete description of a “read” operation also contains the behavior after “read” is disabled.

Example 2:

```
VECTOR ( 01 read_enb -> X? dout -> 10 read_enb -> ?X dout) {
   OPERATION = read; // output goes to X in read-off
}
VECTOR ( 01 read_enb -> ?? dout -> 10 read_enb -> ?- dout) {
   OPERATION = read; // output holds is value in read-off
}
```

### 6.6.3 LABEL annotation

```
LABEL = string ;
```

ensures SDF matching with conditional delays across Verilog, VITAL, etc.

### 6.6.4 EXISTENCE_CONDITION annotation

```
EXISTENCE_CONDITION = boolean_expression ;
```

For false-path analysis tools, the existence condition shall be used to eliminate the vector from further analysis if, and only if, the existence condition evaluates to *False*. For applications other than false-path analysis, the existence condition shall be treated as if the boolean expression was a co-factor to the vector itself. The default existence condition is *True*.

Example:

```
VECTOR (01 a -> 01 z & (c | !d) ) {
   EXISTENCE_CONDITION = !scan_select;
   DELAY { FROM { PIN=a; } TO { PIN=z; } /* data */ }
}
VECTOR (01 a -> 01 z & (!c | d) ) {
   EXISTENCE_CONDITION = !scan_select;
   DELAY { FROM { PIN=a; } TO { PIN=z; } /* data */ }
}
```

Each vector contains state-dependent delay for the same timing arc. If `!scan_select` evaluates *True*, both vectors are eliminated from timing analysis.

### 6.6.5 EXISTENCE_CLASS annotation

```
EXISTENCE_CLASS = string ;
```

Reference to the same existence class by multiple vectors has the following effects:

- A common mode of operation is established between those vectors, which can be used for selective analysis, for instance mode-dependent timing analysis. The name of the mode is the name of the class.
- A common existence condition is inherited from that existence class, if there is one.

Example:

```
CLASS non_scan_mode {
   EXISTENCE_CONDITION = !scan_select;
}
VECTOR (01 a -> 01 z & (c | !d) ) {
   EXISTENCE_CLASS = non_scan_mode;
   DELAY { FROM { PIN=a; } TO { PIN=z; } /* data */ }
}
VECTOR (01 a -> 01 z & (!c | d) ) {
   EXISTENCE_CLASS = non_scan_mode;
   DELAY { FROM { PIN=a; } TO { PIN=z; } /* data */ }
}
```

Each vector contains state-dependent delay for the same timing arc. If the mode `non_scan_mode` is turned off or if `!scan_select` evaluates *True*, both vectors are eliminated from timing analysis.

### 6.6.6 CHARACTERIZATION_CONDITION annotation

```
CHARACTERIZATION_CONDITION = boolean_expression ;
```

For characterization tools, the characterization condition shall be treated as if the boolean expression was a co-factor to the vector itself. For all other applications, the characterization condition shall be disregarded. The default characterization condition is *True*.

Example:

```
VECTOR (01 a -> 01 z & (c | !d) ) {
   CHARACTERIZATION_CONDITION = c & !d;
   DELAY { FROM { PIN=a; } TO { PIN=z; } /* data */ }
}
```

The delay value for the timing arc applies for any of the following conditions: (`c & !d`), (`c & d`), or (`!c & !d`), since they all satisfy (`c | !d`) . However, the only condition chosen for delay characterization is (`c & !d`).

### 6.6.7 CHARACTERIZATION_VECTOR annotation

```
CHARACTERIZATION_VECTOR = ( vector_expression ) ;
```

The characterization vector is provided for the case where the vector expression cannot be constructed using the vector and a boolean co-factor. The use of the characterization vector is restricted to characterization tools in the same way as the use of the characterization condition. Either a characterization condition or a characterization vector can be provided, but not both. If none is provided, the vector itself shall be used by the characterization tool.

Example:

```
VECTOR (01 A -> 01 Z) {
   CHARACTERIZATION_VECTOR = ((01 A & 10 inv_A) -> (01 Z & 10 inv_Z));
}
```

Analysis tools see the signals `A` and `Z`. The signals `inv_A` and `inv_Z` are visible to the characterization tool only.

### 6.6.8 CHARACTERIZATION_CLASS annotation

```
CHARACTERIZATION_CLASS = string ;
```

Reference to the same characterization class by multiple vectors has the following effects:

- A commonality is established between those vectors, which can be used for selective characterization in a way defined by the library characterizer, for instance, to share the characterization task between different teams or jobs or tools.
- A common characterization condition or characterization vector is inherited from that characterization class, if there is one.

## 6.7 ILLEGAL statement for VECTOR

For complex cells, especially multi-port memories, it is useful to define the behavior as a consequence of illegal operations, for example when several ports try to access the same address.

A `VECTOR` statement shall contain the optional `ILLEGAL` statement, defined as follows:

```
illegal ::=
   ILLEGAL [ identifier ] { illegal_items }
```

```
illegal_items ::=
   illegal_item { illegal_item }

illegal_item ::=
   all_purpose_item
|  violation
```

where `all_purpose_item` and `violation` are defined in Section 11.7 and Section 11.16, respectively.

The `vector_expression` within the `VECTOR` statement describes a state or a sequence of events which define an illegal operation. The `VIOLATION` statement describes the consequence of such an illegal operation.

Example 1:

```
VECTOR ( (addr_A == addr_B) && write_enable_A && write_enable_B ) {
   ILLEGAL write_A_write_B {
      VIOLATION {
         MESSAGE = “write conflict between port A and B“;
         MESSAGE_TYPE = error;
         BEHAVIOR { data[addrA] = ‘bxxxxxxxx; }
      }
   }
}
```

Note: An illegal operation can be legalized by using `MESSAGE_TYPE=INFORMATION` or `MESSAGE_TYPE=WARNING`.

This statement can also be used to define the behavior when an address is out of range. Sometimes the address space is not continuous, i.e., it can contain holes in the middle. In this case, a `MIN` or `MAX` value for legal addresses would not be sufficient. On the other hand, a `boolean_expression` can always exactly describe the legal and illegal address space.

Example 2:

```
VECTOR ( (addr > ‘h3) && write_enb ) {
   ILLEGAL {
      VIOLATION {
         MESSAGE = “write address out of range“;
         MESSAGE_TYPE = error;
         BEHAVIOR { data[addr] = ‘bxxxxxxxx; }
      }
   }
}
```

## 6.8 TEST statement

A `CELL` can contain a `TEST` statement, which is defined as follows:

```
test ::=
          TEST { behavior }
```

The purpose is to describe the interface between an externally applied test algorithm and the CELL. The behavior statement within the TEST statement uses the same syntax as the behavior statement within the FUNCTION statement. However, the set of used variables is different. Both the TEST and the FUNCTION statement shall be self-contained, complete and complementary to each other.

## 6.9 Physical bitmap for memory BIST

This section defines the physical bitmap for memory BIST. This is a particular case of the usage of the TEST statement.

### 6.9.1 Definition of concepts

The physical architecture of a memory can be described by the following parameters:

*BANK index*: A memory can be arranged in one or several banks, each of which constitutes a two-dimensional array of rows and columns

*ROW index*: A row of memory cells within one bank shares the same row decoder line.

*COLUMN index*: A column of memory cells within one bank shares the same data bit line and, if applicable, the same sense amplifier.

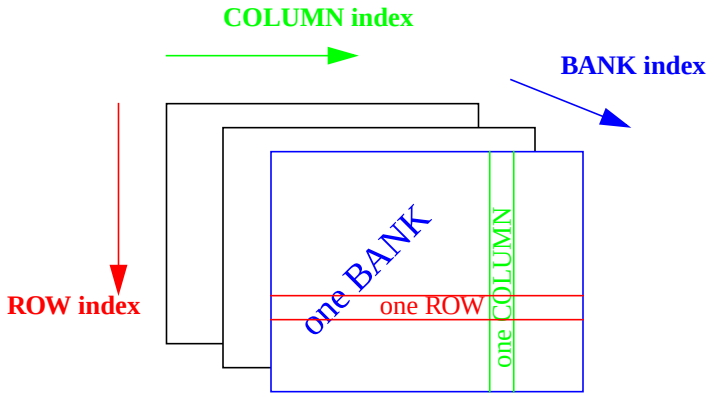



**Figure 6-4: Illustration of a physical memory architecture, arranged in banks, rows, columns**

The physical memory architecture is not evident from the functional description and the pins involved in the functional description of the memory. Those pins are called logical pins, e.g., logical address and logical data.

A memory BIST tool needs to know which logical address and data corresponds to a physical row, column, or bank in order to write certain bit patterns into the memory and read expected bit patterns from the memory. Also, the tool needs to know whether the physical data in a specific location is inverted or not with respect to the corresponding logical data.

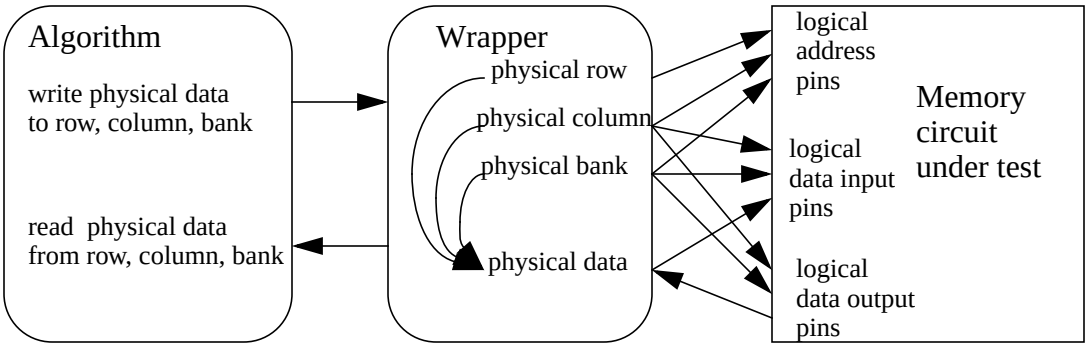



**Figure 6-5: Illustration of the memory BIST concept**

A mapper between physical rows, columns, banks, data and logical addresses, and data pins shall be part of the library description of a memory cell.

The physical row, column, and bank indices can be modeled as virtual inputs to the memory circuit. The data to be written to a physical memory location can also be modeled as a virtual input. The data to be read from a physical memory location can be modeled as a virtual output. Since every data that is written for the purpose of test also needs to be read, the data can be modeled as a virtual bidirectional pin. A virtual pin is a pin with `VIEW=none`, i.e., the pin is not visible in any netlist.

### 6.9.2 Definitions of pin ATTRIBUTE values for memory BIST

The special pin `ATTRIBUTE` values shown in Table 6-33 shall be defined for memory BIST.

**Table 6-33 : PIN attributes for memory BIST**

| Attribute item | Description |
|---|---|
| `ROW_INDEX` | pin is a bus with a contiguous range of values, indicating a physical row of a memory |
| `COLUMN_INDEX` | pin is a bus with a contiguous range of values, indicating a physical column of a memory |
| `BANK_INDEX` | pin is a bus with a contiguous range of values, indicating a physical bank of a memory |
| `DATA_INDEX` | pin is a bus with a contiguous range of values, indicating the bit position within a data bus of a memory |
| `DATA_VALUE` | pin represents a value stored in a physical memory location |

These attributes apply to the pins of the BIST wrapper around the memory rather than to the pins of the memory itself.

The BEHAVIOR statement within TEST shall involve the variables declared as PINs with ATTRIBUTE ROW_INDEX, COLUMN_INDEX, BANK_INDEX, DATA_INDEX, or DATA_VALUE.

### 6.9.3 Explanatory example

One-dimensional arrays with SIGNALTYPE=address (here: PIN[3:0] addr) shall be recognized as address pins to be mapped, involving other one-dimensional arrays with ATTRIBUTE { ROW_INDEX } (here: PIN[1:0] row ) and ATTRIBUTE { COLUMN_INDEX } (here: PIN[3:0] col). This memory has only one bank. Therefore, no one-dimensional array with ATTRIBUTE { BANK_INDEX } exists here.

One-dimensional arrays with SIGNALTYPE=data (here: PIN[3:0] Din and PIN[3:0] Dout) shall be recognized as data pins to be mapped, involving other one-dimensional arrays with ATTRIBUTE { DATA_INDEX } (here: PIN[1:0] dat ) and scalar pins with ATTRIBUTE { DATA_VALUE } (here: PIN bit ).

Note: Since the data buses are 4-bits wide, the data index is 2-bits wide, since 2=log2(4).

Base Example:

```
CELL my_memory {
   PIN[3:0] addr { DIRECTION=input;  SIGNALTYPE=address; }
   PIN[3:0] Din  { DIRECTION=input;  SIGNALTYPE=data; }
   PIN[3:0] Dout { DIRECTION=output; SIGNALTYPE=data; }
   PIN[3:0] bits[0:15] { DIRECTION=none; VIEW=none; SCOPE=behavior; }
   PIN write_enb { DIRECTION=input; SIGNALTYPE=write_enable;
      POLARITY=high; ACTION=asynchronous;
   }
   PIN[1:0] dat { ATTRIBUTE { DATA_INDEX } DIRECTION=none; VIEW=none; }
   PIN bit { ATTRIBUTE { DATA_VALUE } DIRECTION=both; VIEW=none; }
   PIN[1:0] row {
      ATTRIBUTE { ROW_INDEX }  RANGE { 0: 3 }
      DIRECTION=input; VIEW=none;
   }
   PIN[3:0] col {
      ATTRIBUTE { COLUMN_INDEX } RANGE { 0 : 15 }
      DIRECTION=input; VIEW=none;
   }
   FUNCTION {
      BEHAVIOR {
         Dout = bits[addr];
         @ (write_enb) { bits[addr] = Din; }
   }  }
/*different physical architectures are shown in the following examples*/
}
```

Example 1

| addr[3:2] | 00 | 00 | 00 | 00 | 01 | 01 | 01 | 01 | 10 | 10 | 10 | 10 | 11 | 11 | 11 | 11 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| physical column | 'h0 | 'h1 | 'h2 | 'h3 | 'h4 | 'h5 | 'h6 | 'h7 | 'h8 | 'h9 | 'hA | 'hB | 'hC | 'hD | 'hE | 'hF |
| addr[1:0] physical row | | | | | | | | | | | | | | | | |
| 00 'h0 | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] |
| 01 'h1 | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] |
| 10 'h2 | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] |
| 11 'h3 | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] | D[0] | D[1] | D[2] | D[3] |

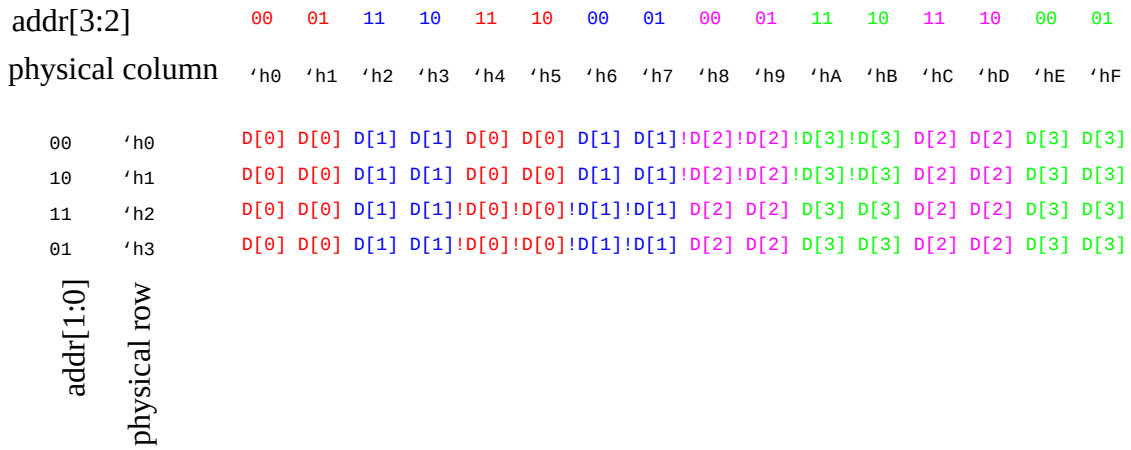
```
   TEST {
      BEHAVIOR {
// map row and column index to logical address
         addr[1:0] = row[1:0];
         addr[3:2] = col[3:2];
// map column index to logical data index
         dat[1:0] = col[1:0];
// map physical data to input and output data
         Din[dat] = bit;
         bit = Dout[dat];
      }
   }
```

Example 2

| addr[3:2] | 00 | 01 | 10 | 11 | 00 | 01 | 10 | 11 | 00 | 01 | 10 | 11 | 00 | 01 | 10 | 11 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| physical column | 'h0 | 'h1 | 'h2 | 'h3 | 'h4 | 'h5 | 'h6 | 'h7 | 'h8 | 'h9 | 'hA | 'hB | 'hC | 'hD | 'hE | 'hF |
| addr[1:0] physical row | | | | | | | | | | | | | | | | |
| 00 'h0 | D[0] | D[0] | D[0] | D[0] | D[1] | D[1] | D[1] | D[1] | D[2] | D[2] | D[2] | D[2] | D[3] | D[3] | D[3] | D[3] |
| 01 'h1 | D[0] | D[0] | D[0] | D[0] | D[1] | D[1] | D[1] | D[1] | D[2] | D[2] | D[2] | D[2] | D[3] | D[3] | D[3] | D[3] |
| 10 'h2 | D[0] | D[0] | D[0] | D[0] | D[1] | D[1] | D[1] | D[1] | D[2] | D[2] | D[2] | D[2] | D[3] | D[3] | D[3] | D[3] |
| 11 'h3 | D[0] | D[0] | D[0] | D[0] | D[1] | D[1] | D[1] | D[1] | D[2] | D[2] | D[2] | D[2] | D[3] | D[3] | D[3] | D[3] |

```
   TEST {
      BEHAVIOR {
// map row and column index to logical address
         addr[1:0] = row[1:0];
         addr[3:2] = col[1:0];
// map column index to logical data index
         dat[1:0] = col[3:2];
// map physical data to input and output data
         Din[dat] = bit;
         bit = Dout[dat];
      }
   }
```

Example 3




```
   TEST {
      BEHAVIOR {
// map row and column index to logical address
         addr[0] = row[1];
         addr[1] = row[0] ^ row[1]
         addr[2] = col[0] ^ col[1] ^ col[2];
         addr[3] = col[2] ^ col[3];
// map column index to logical data index
         dat[0] = col[1];
         dat[1] = col[3];
// map physical data to input and output data
         Din[dat]=bit^(row[1]&col[2]&!col[3] | !row[1]&!col[2]&col[3]);
         bit=Dout[dat]^(row[1]&col[2]&!col[3] | !row[1]&!col[2]&col[3]);
      }
   }
```

Notes:

1. This enables the description of a complete bitmap of a memory in a compact way.

2. The RANGE feature is not restricted to BIST. It can be used to describe a valid contiguous range on any bus. This alleviates the need for interpreting a VECTOR with ILLEGAL statement to get the valid range. However, the VECTOR with ILLEGAL statement is still necessary to describe the behavior of a device when illegal values are driven on a bus.

3. The TEST statement with BEHAVIOR allows for generalization from memory BIST to any test vector generation requirement, e.g., logic BIST. The only necessary additions would be other PIN ATTRIBUTEs describing particular features to be recognized by the test vector generation algorithm for the target test algorithm.

# Section 7
# General Rules for Arithmetic Models

This chapter defines the general rules for arithmetic models.

## 7.1 Principles of arithmetic models

The purpose of arithmetic models is to specify calculable mathematical relationships between objects representing physical quantities in the library. Arithmetic models are identified by context-sensitive keywords, because how these quantities are measured, extracted, or interpreted depends on the context in which the objects are placed.

The quantity identified by the keyword `CAPACITANCE` can serve as example. In the context of a `PIN`, it represents pin capacitance. In the context of a `WIRE`, it represents wire capacitance. In the context of a `RULE`, it represents the calculation method for a capacitance formed by a layout pattern described within the rule. The context-specific semantics of each arithmetic model are specified in Section 8 for electrical models and Section 9 for physical models.

In certain cases, the context alone does not completely specify the semantics of an arithmetic model. Auxiliary definitions within the arithmetic model are needed; these are represented by using annotations or annotation containers.

A simple example is the `UNIT` annotation, which is applicable for most arithmetic models. It specifies the unit in terms of which the arithmetic model data is represented. The applicable auxiliary objects for each arithmetic model are specified in Section 8 for electrical models and Section 9 for physical models.

### 7.1.1 Global definitions for arithmetic models

In many cases, auxiliary definitions apply globally to all arithmetic models within a certain context, for instance, the `UNIT` can apply for all `CAPACITANCE` objects within a library. In order to specify such global definitions, the arithmetic model construct can be used without data.

```
model_definition ::=
          model_keyword [ identifier ] { all_purpose_items }
```

This construct has the syntactical form of an `annotation_container` (see Section 11.7).

### 7.1.2 Trivial arithmetic model

The simplest form of an arithmetic model contains just constant data.

```
trivial_model ::=
          model_keyword [ identifier ] = number ;
   |      model_keyword [ identifier ] = number { all_purpose_items }
```

This construct has the syntactical form of an annotation (see Section 11.7).

### 7.1.3 Arithmetic model using EQUATION

The arithmetic model data can be represented as an EQUATION. In this case, a HEADER defines the arguments of the equation. It is also possible to use other arithmetic models, which are visible within the context of this arithmetic model, as arguments. Those arguments need not appear in the HEADER.

```
equation_based_model ::=
          model_keyword [ identifier ] {
                [ all_purpose_items ]
                [ equation_based_header ]
                equation
          }

equation_based_header ::=
          HEADER { model_keyword { model_keyword } }
   |      HEADER { model_definition { model_definition } }

equation ::=
          EQUATION { arithmetic_expression }
```

The syntax of arithmetic_expression is explained in Section 7.2.

### 7.1.4 Arithmetic model using TABLE

The arithmetic model data can be represented as a lookup table. In this case, a TABLE is necessary for the data itself and for each argument.

```
table_based_model ::=
          model_keyword [ identifier ] {
                [ all_purpose_items ]
                table_based_header
                table
                [ equation ]
          }

table_based_header ::=
          HEADER { table_model_definition { table_model_definition }
   }

table_model_definition ::=
          model_keyword [ identifier ] { all_purpose_items table }

table ::=
          TABLE { symbol { symbol } }
   |      TABLE { number { number } }
```

Tables containing symbols are only meant for lookup of discrete datapoints. Tables containing numbers are for calculation and, eventually, interpolation of datapoints. The model_keyword

(see Section 8 and Section 9) defines whether symbols or numbers are legal for a particular table.

The size of the table inside the `table_based_model` shall be the product of the size of the tables inside the `table_header`. In order to support interpolation, the numbers in each table inside the `table_header` shall be in strictly monotonic ascending order. See Section 7.3 for more details.

The `table_model_definition` can also be used outside the context of a `table_header`, very much like a `model_definition`. In this case, the `model_definition` supplies the same information as the `table_model_definition`, plus the additional information of a discrete set of valid numbers applicable for the model.

For example, the `WIDTH` of a physical layout object can contain only a discrete set of legal values. Those can be specified using a `table_model_definition`.

However, the table in a `table_model_definition` *outside* a `table_header` shall not substitute the table *inside* the `table_header`. The former defines a legal set of values, the latter defines the table-lookup indices.

If all table data are numbers, the `table_based_model` can also have an optional equation. This equation is to be used when the argument data are out of interpolation range. Without the equation, extrapolation shall be applied for data which are out of range.

### 7.1.5 Complex arithmetic model

A complex arithmetic model can be constructed by defining a nested arithmetic model within another arithmetic model. The data of the inner arithmetic model is calculated first. Then the result is applied for calculation of the data of the outer arithmetic model.

```
complex_model ::=
          model_keyword [ identifier ] {
                [ all_purpose_items ]
                HEADER { model { model } }
                equation
          }
   |      model_keyword {
                all_purpose_items
                HEADER { header_model { header_model } }
                table
                [ equation ]
          }

header_model ::=
          model_definition
   |      table_model_definition
   |      equation_based_model
   |      table_based_model
   |      header_table_model

header_table_model ::=
          model_keyword [ identifier ] {
```

```
                all_purpose_items
                HEADER { symbol { symbol } }
                TABLE { number { number } }
          }
```

If any header_model is either model_definition or table_model_definition, then the complex_model reduces to the previously defined equation_based_model and table_based_model, respectively. In order to support a table in the general_model, any header_model shall be either a table_model_definition or table_based_model, and the numbers in each table inside each header_model shall be strictly monotonically increasing.

The header_table_model construct can be used to associate symbols with numbers. For example, process corners can be defined as discrete symbols and associated with process derating factors. The numbers can be used in equations and for interpolation, whereas the symbols cannot.

### 7.1.6 Containers for arithmetic models and submodels

Containers for arithmetic models can supplement the context-specific semantics of the arithmetic model. Therefore, arithmetic models can be placed in the context of arithmetic model containers, using the following construct.

```
model_container ::=
          model_container_keyword  {
                [ all_purpose_items ]
                model_container_contents { model_container_contents }
          }

model_container_contents ::=
          model_container
   |      trivial_model
   |      complex_model
```

There is a dedicated set of *model_container*_keywords. In addition, *model*_keywords can also be used as *model_container*_keywords and dedicated *submodel*_keywords can be used as *model*_keywords. The number of levels in nested arithmetic model containers is restricted by the set of allowed combinations between *model_container*_keywords, *model*_keywords and *submodel*_keywords (see Section 7.6).

## 7.2 Arithmetic expressions

Arithmetic expressions define the contents of an EQUATION. Variables used in the EQUATION are the identifiers of the header_model, if present, or else the *model*_keywords of the header_model.

### 7.2.1 Syntax of arithmetic expressions

The syntax of arithmetic expressions is:

```
arithmetic_expression ::=
     ( arithmetic_expression )
   | number
   | [ arithmetic_unary ] identifier
   | arithmetic_expression arithmetic_binary arithmetic_expression
   | arithmetic_function_operator
          ( arithmetic_expression { , arithmetic_expression } )
   | boolean_expression ? arithmetic_expression :
          { boolean_expression ? arithmetic_expression : }
          arithmetic_expression
```

Examples:

```
1.24
- Vdd
C1 + C2
MAX ( 3.5*C , -Vdd/2 , 0.0 )
(C > 10) ? Vdd**2 : 1/2*Vdd - 0.5*C
```

### 7.2.2 Arithmetic operators

Table 7-1, Table 7-2, and Table 7-3 list unary, binary, and function arithmetic operators.

**Table 7-1 : Unary arithmetic operators**

| Operator | Description |
|---|---|
| + | positive sign (for integer or number) |
| - | negative sign (for integer or number) |

**Table 7-2 : Binary arithmetic operators**

| Operator | Description |
|---|---|
| **+** | addition (integer or number) |
| **-** | subtraction (integer or number) |
| ***** | multiplication (integer or number) |
| **/** | division (integer or number) |
| ****** | exponentiation (integer or number) |
| **%** | modulo division (integer or number) |

**Table 7-3 : Function arithmetic operators**

| Operator | Description |
|---|---|
| **LOG** | natural logarithm (argument is + integer or number) |
| **EXP** | natural exponential (argument is integer or number) |
| **ABS** | absolute value (argument is integer or number) |

**Table 7-3 : Function arithmetic operators, *continued***

| Operator | Description |
|---|---|
| **MIN** | minimum (all arguments are integer or number) |
| **MAX** | maximum (all arguments are integer or number) |

Function operators with one argument (such as `log`, `exp`, and `abs`) or multiple arguments (such as `min` and `max`) shall have their arguments within parenthesis, e.g., `min(1.2,-4.3,0.8)`.

### 7.2.3 Operator priorities

The priority of binding operators to operands in arithmetic expressions shall be from strongest to weakest in the following order:

1. unary arithmetic operator (+, -)
2. exponentiation (**)
3. multiplication (*), division (/), modulo division (%)
4. addition (+), subtraction (-)

## 7.3 Construction of arithmetic models

Input variables, also called *arguments of arithmetic models*, appear in the `HEADER` of the model. In the simplest case, the `HEADER` is just a list of arguments, each being a context-sensitive keyword. The model itself is also defined with a context-sensitive keyword.

The model can be in equation form. All arguments of the equation shall be in the `HEADER`. The ALF parser shall issue an error if the `EQUATION` uses an argument not defined in the `HEADER`. A warning shall be issued if the `HEADER` contains arguments not used in the `EQUATION`.

Example:

```
DELAY {
      ...
      HEADER {
            CAPACITANCE {...}
            SLEWRATE {...}
      }
      EQUATION {
            0.01 + 0.3*SLEWRATE + (0.6 + 0.1*SLEWRATE)*CAPACITANCE
      }
}
```

If the model uses a `TABLE`, then each argument in the `HEADER` also needs a table defining the format. The order of arguments decides how the index to each entry is calculated. The first argument is the innermost index, the following arguments are outer indices.

```
DELAY {
      HEADER {
            CAPACITANCE {
                  TABLE {0.03 0.06 0.12 0.24}
            }
            SLEWRATE {
                  TABLE {0.1 0.3 0.9}
            }
      }
      TABLE {
            0.07 0.10 0.14 0.22
            0.09 0.13 0.19 0.30
            0.10 0.15 0.25 0.41
      }
}
```

The first argument `CAPACITANCE` has four entries. The second argument `SLEWRATE` has three entries. Thus, `DELAY` has 4*3=12 entries. For readability, comments can be inserted in the table.

```
TABLE {
//capacitance:0.03 0.06 0.12 0.24
//            ------------------   slewrate:
              0.07 0.10 0.14 0.22 // 0.1
              0.09 0.13 0.19 0.30 // 0.3
              0.10 0.15 0.25 0.41 // 0.9
}
```

Comments have no significance for the ALF parser nor does the arrangement of rows and columns. Only the order of values is important for index calculation. The table can be made more compact by removing newlines.

```
TABLE { 0.07 0.10 0.14 0.22 0.09 0.13 0.19 0.30 0.10 0.15 0.25 0.41 }
```

For readability, the models and arguments can also have names, i.e., object IDs. For named objects, the name is used for referencing, rather than the keyword.

```
DELAY rise_out{
      ...
      HEADER {
            CAPACITANCE c_out {...}
            SLEWRATE fall_in {...}
      }
      EQUATION {
            0.01 + 0.3 * fall_in + (0.6 + 0.1* fall_in) * c_out
      }
}
```

The arguments of an arithmetic model can be arithmetic models themselves. In this way, combinations of `TABLE`- and `EQUATION`-based models can be used, for instance, in derating.

Analogous with `FUNCTION`, both `EQUATION` and `TABLE` representation of an arithmetic model are allowed. The `EQUATION` is intended to be used when the values of the arguments fall out of range, i.e., to avoid extrapolation.

# 7.4 Annotations for arithmetic models

Annotations and annotation containers described in this chapter are relevant for the semantic interpretation of arithmetic models and their arguments.

Example:

```
DELAY=f(CAPACITANCE)
```

`DELAY` is the arithmetic model, `CAPACITANCE` is the argument.

Arguments of arithmetic models have the form of annotation containers. They can also have the form of arithmetic models themselves, in which case they represent nested arithmetic models.

## 7.4.1 DEFAULT annotation

*Default annotation* promotes use of the default value instead of the arithmetic model if the arithmetic model is beyond the scope of the application tool.

**`DEFAULT =`** `number` **`;`**

Restrictions can apply for the allowed type of `number`. For instance, if the arithmetic model allows only `non_negative_number`, then the default is restricted to `non_negative_number`.

## 7.4.2 UNIT annotation

*Unit annotation* associates units with the value computed by the arithmetic model.

**`UNIT =`** `string | non_negative_number` **`;`**

A unit specified by a `string` can take the values (`*` indicates a wild card) shown in Table 7-4.

**Table 7-4 : UNIT annotation**

| Annotation string | Description |
|---|---|
| `f*` or `F*` | equivalent to `1E-15` |
| `p*` or `P*` | equivalent to `1E-12` |
| `n*` or `N*` | equivalent to `1E-9` |
| `u*` or `U*` | equivalent to `1E-6` |
| `m*` or `M*` | equivalent to `1E-3` |
| `1*` | equivalent to `1E+0` |
| `k*` or `K*` | equivalent to `1E+3` |
| `meg*` or `MEG*`[a] | equivalent to `1E+6` |
| `g*` or `G*` | equivalent to `1E+9` |

a. or any uppercase/lowercase combination of these three characters

Arithmetic models are context-sensitive, i.e., the units for their values can be determined from the context. If the `UNIT` annotation for such a context does not exist, default units are applied to the value (see Section 7.6).

Example:

```
TIME { UNIT = ns; }
FREQUENCY { UNIT = gigahz; }
```

If the unit is a string, then only the first character (the first three characters in case of `MEG`) is interpreted. The reminder of the string can be used to define base units. Metric base units are assumed, but not verified, in ALF.

There is no semantic difference between

```
unit = 1sec;
```

and

```
unit = 1volt;
```

Therefore, if the unit is specified as

```
unit = meg;
```

the interpretation is `1E+6`. However, for

```
unit = 1meg;
```

the interpretation is `1` and not `1E+6`.

Units in a non-metric system can only be specified with numbers, not with strings. For instance, if the intent is to specify an inch instead of a meter as the base unit, the following specification does not meet the intent:

```
unit = 1inch;
```

since the interpretation is `1` and meters are assumed.

The correct way of specifying inch instead of meter is

```
unit = 25.4E-3;
```

since 1 inch is (approximately) 25.4 millimeters.

### 7.4.3 CALCULATION annotation

An arithmetic model in the context of a `VECTOR` can have the `CALCULATION` annotation defined as follows:

```
calculation_annotation ::=
        CALCULATION = calculation_identifier ;

calculation_identifier ::=
        absolute
   |    incremental
```

It shall specify whether the data of the model are to be used by themselves or in combination with other data. The default is **`absolute`**.

The **incremental** data from one VECTOR shall be added to **absolute** data from another VECTOR under the following conditions:

- The model definitions are compatible, i.e., measurement specifications shall be the same. Units are allowed to be different.
  Example: slewrate measurements at the same pin, same switching direction, and same threshold values.
- The model definitions for common arguments are compatible, i.e., the same range of values for table-based models and measurement specifications are the same. Units can be different.
  Example: same values for derate_case and same threshold definitions for input slewrate.
- The vector definitions are compatible, i.e, the vector_or_boolean_expression of the VECTOR containing **incremental** data matches the vector_or_boolean_expression of the VECTOR containing **absolute** data by removing all variables appearing exclusively in the former expression.

Example:

```
VECTOR ( 01 A -> 01 Z ) {
   DELAY {
      CALCULATION = absolute;
      FROM { PIN = A; } TO { PIN = Z; }
      HEADER {
         CAPACITANCE load { PIN = Z; }
         SLEWRATE slew { PIN = A; }
      }
      EQUATION { 0.5 + 0.3*slew + 1.2*load }
   }
}
VECTOR ( 01 A &> 01 B &> 01 Z ) {
   DELAY {
      CALCULATION = incremental;
      FROM { PIN = A; } TO { PIN = Z; }
      HEADER {
         SLEWRATE slew_A { PIN = A; }
         SLEWRATE slew_B { PIN = B; }
         TIME time_A_B { FROM { PIN = A; } TO { PIN = B; } }
      }
      EQUATION {- 0.1 + (0.05+0.002*slew_A*slew_B)*time_A_B) }
   }
}
```

Both models describe the rise-to-rise delay from A to Z. The second delay model describes the incremental delay (here negative), when input B switches in a time window between A and Z.

### 7.4.4 INTERPOLATION annotation

An argument of a table-based arithmetic model, i.e., a model in the HEADER containing a TABLE statement, can have the INTERPOLATION annotation defined as follows:

```
interpolation_annotation ::=
        INTERPOLATION = interpolation_identifier ;

interpolation_identifier ::=
        fit
    |   linear
    |   floor
    |   ceiling
```

This also needs to specify the interpolation scheme for the values in-between the values of the TABLE.

- **fit**
  the data points in the table are supposed to be part of a smooth curve. Linear interpolation or other algorithms, e.g., cubic spline or polynomial regression can be used to fit the data points into the curve.
- **linear**
  the data points in the table are supposed to be part of a piece wise linear curve. Linear interpolation shall be used.
- **floor**
  the value to the left in the table, i.e., the smaller value is used.
- **ceiling**
  the value to the right in the table, i.e., the larger value is used.

The default is **fit**. For multi-dimensional tables, different interpolation schemes can be used for each dimension.

Example:

```
my_model {
   HEADER {
      dimension1 { INTERPOLATION = fit; TABLE { 1 2 4 8 }
      dimension2 { INTERPOLATION = floor; TABLE { 10 100 }
      dimension3 { INTERPOLATION = ceiling; TABLE { 10 100 }
   }
   TABLE {
      1 7 3 5
      10 20 60 40
      50 30 20 100
      0.8 0.4 0.2 0.9
   }
}
```

Consider the following values:

```
   dimension1 = 6
=> following subtable is chosen:
      3  5     // interpolation between 3 and 5
      60 40    // or between 60 and 40
      20 100   // or between 20 and 100
      0.20.9   // or between 0.2 and 0.9
   dimension2 = 50
=> following subtable is picked:
      3  5     // interpolation between 3 and 5
      20 100   // or between 20 and 100
   dimension3 = 50
=> following subtable is picked:
      20 100   // interpolation between 20 and 100
```

The following rules shall apply for each dimension of a table-based model:

For values outside the range of the table, extrapolation shall apply, using the table data points at the leftmost or rightmost side, respectively, as reference.

If the value is smaller than the smallest, i.e. leftmost, data point in the table, the extrapolation shall be calculated as if the value would fall in-between the leftmost and second leftmost value.

If the value is greater than the greatest, i.e. rightmost, data point in the table, the extrapolation shall be calculated as if the value would fall in-between the rightmost and second rightmost value.

Example:

```
my_model Y {
   HEADER {
      my_argument X {
         TABLE { 0    2    4    8   }
         //      X[0] X[1] X[2] X[3]
      }
   }
   TABLE { 0.5  0.6  1.0  1.5 }
   //      Y[0] Y[1] Y[2] Y[3]
}
```

For linear interpolation, the following equation is used:

$$Y = Y[N] + \frac{Y[N+1] - Y[N]}{X[N+1] - X[N]} \cdot X \qquad X[N] \le X \le X[N+1]$$

If X < X[0], the values X[0], X[1], Y[0], Y[1] are plugged into the equation.

If X > X[3], the values X[2], X[3], Y[2], Y[3] are plugged into the equation.

The following figure illustrates a non-linear interpolation scheme with the goal of fitting 3 neighboring points into a smooth curve.

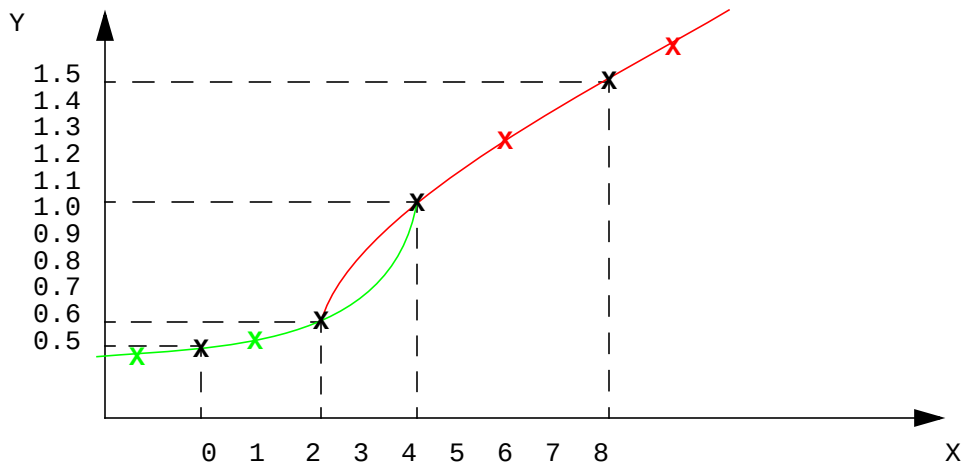



**Figure 7-1: Illustration of extrapolation rules**

The curve based on the 3 rightmost or the 3 leftmost points, respectively, is used for extrapolation to the right side or the left side, respectively.

## 7.5 Containers for arithmetic models

The keywords shown in Table 7-5 are defined for objects that can contain arithmetic models.

**Table 7-5 : Unnamed containers for arithmetic models**

| Objects | Description |
|---|---|
| FROM | contains start point of timing measurement or timing constraint |
| TO | contains end point of measurement or timing constraint |
| LIMIT | contains arithmetic models for limit values |
| EARLY | contains arithmetic models for timing measurements relevant for early signal arrival time |
| LATE | contains arithmetic models for timing measurements relevant for late signal arrival time |

The LIMIT container is for general use. The FROM, TO, EARLY, and LATE containers are only for use within the context of timing models.

### LIMIT container

A LIMIT container shall contain arithmetic models. The arithmetic models shall contain submodels identified by MIN and/or MAX.

Example:

```
PIN data_in {
      LIMIT {
            SLEWRATE { UNIT = ns; MIN = 0.05; MAX = 5.0;}
      }
}
```

The minimum slewrate allowed at pin `data_in` is `0.05 ns`, the maximum is `5.0 ns`.

```
PIN data_in {
      LIMIT {
            SLEWRATE {
                  UNIT = ns;
                  MAX {
                        HEADER { FREQUENCY { UNIT=megahz;} }
                        EQUATION { 250 / FREQUENCY  }
                  }
            }
      }
}
```

The maximum allowed slewrate is frequency-dependent, e.g., the value is `0.25ns` for `1GHz`.

## 7.6 Arithmetic submodels

Arithmetic submodels can be used to distinguish different measurement conditions for the same model. The root of an arithmetic model can contain nested arithmetic submodels. The header of an arithmetic model can contain nested arithmetic models, but not arithmetic submodels.

The arithmetic submodels shown in Table 7-6 are generally applicable.

**Table 7-6 : Generally applicable arithmetic submodels**

| Objects | Description |
|---|---|
| `MIN` | for measured or calculated data:<br>the data represents the minimal value / set of values within a statistical distribution<br>for data within LIMIT container:<br>the data represents the lower limit value / set of values |
| `TYP` | for measured or calculated data:<br>the data represents the typical value / set of values within a statistical distribution |
| `MAX` | for measured or calculated data:<br>the data represents the maximal value / set of values within a statistical distribution<br>for data within LIMIT container:<br>the data represents the lower limit value / set of values |
| `DEFAULT` | for measured or calculated data:<br>the data represents the default value / set of values to be used per default |

The arithmetic submodels shown in Table 7-7 are only applicable in the context of electrical modeling.

**Table 7-7 : Submodels restricted to electrical modeling**

| Objects | Description |
|---|---|
| HIGH | applicable for electrical data measured at a logic high state of a pin |
| LOW | applicable for electrical data measured at a logic low state of a pin |
| RISE | applicable for electrical data measured during a logic low to high transition of a pin |
| FALL | applicable for electrical data measured during a logic high to low transition of a pin |

The arithmetic submodels shown in Table 7-8 are only applicable in the context of physical modeling.

**Table 7-8 : Submodels restricted to physical modeling**

| Objects | Description |
|---|---|
| HORIZONTAL | applicable for layout measurements in horizontal direction |
| VERTICAL | applicable for layout measurements in vertical direction |

The semantics of the restricted submodels are explained in Section 8 and Section 9.

### 7.6.1 Semantics of MIN / TYP / MAX

MIN, TYP, and MAX indicate the data of the arithmetic model represent minimal, typical, or maximal values within a statistical distribution. No correlation is assumed or implied between MIN data, TYP data, or MAX data across different arithmetic models.

Example:

```
DELAY {
   FROM { PIN=A; } TO { PIN=Z; }
   MIN = 0.34; TYP = 0.38; MAX = 0.45;
}
POWER {
   MEASUREMENT = average; FREQUENCY = 1e6;
   MIN = 1.2; TYP = 1.4; MAX = 1.5;
}
```

The MIN value for DELAY could simultaneously apply with the MIN value for POWER. Typically, the case with smaller delay is also the case with larger power consumption.

Within the scope of a LIMIT container, MIN and MAX contain the data for a lower or upper limit, respectively. There shall be at least one limit, lower or upper, in each model, but not necessarily both.

Example:

```
LIMIT {
   SLEWRATE { PIN=A; MAX=5.0; }
   VOLTAGE { PIN=VDD; MIN=1.6; MAX=2.0; }
}
```

MIN, MAX as an annotation inside a model or inside a model argument within the HEADER define the validity range of the data. If MIN, MAX is not defined and the data is in a TABLE, the boundaries of the data in the TABLE shall be considered as validity limits.

Example:

```
POWER {
   HEADER {
      SLEWRATE { PIN=A; MIN=0.01; MAX=5.0; TABLE { 0.1 0.5 1.0 } }
      CAPACITANCE { PIN=Z; TABLE { 0.0 0.4 0.8 1.6 } }
   }
   TABLE { 0.2 0.3 0.6 0.4 0.5 0.7 0.8 0.8 1.0 1.5 1.5 1.6 }
}
```

The data for POWER is valid for SLEWRATE in the range between 0.01 and 5.0 (via extrapolation) and for CAPACITANCE in the range between 0.0 and 1.6.

### 7.6.2 Semantics of DEFAULT

Arithmetic submodels can be identified by `MIN`, `TYP`, and `MAX` or context-restricted keywords. For cases where the application tool cannot decide which qualifier applies, a supplementary arithmetic submodel with the qualifier `DEFAULT` can be used.

Example:

```
PIN my_pin {
   CAPACITANCE {
      MIN { HEADER { ... } TABLE { ... } }
      TYP { HEADER { ... } TABLE { ... } }
      MAX { HEADER { ... } TABLE { ... } }
      DEFAULT { HEADER { ... } TABLE { ... } }
   }
}
```

Note: The `DEFAULT` model can also degenerate to a single value; it represents a trivial arithmetic model.

In certain cases, there is no supplementary submodel. Instead, one of the already defined submodels is used by default. For this case, the `DEFAULT` annotation can be used to point to the applicable keyword.

Example:

```
PIN my_pin {
   CAPACITANCE {
      MIN { HEADER { ... } TABLE { ... } }
      TYP { HEADER { ... } TABLE { ... } }
      MAX { HEADER { ... } TABLE { ... } }
      DEFAULT = TYP;
   }
}
```

The trivial arithmetic model construct with `DEFAULT` can also be used for an argument in the context of the `HEADER` of an arithmetic model. This enables evaluation of the arithmetic model in case the data of the argument can not be supplied by the application tool.

Example:

```
PIN my_pin {
   CAPACITANCE {
      HEADER { TEMPERATURE { DEFAULT=50; TABLE { 0 50 100 } } }
      TABLE { 0.05 0.07 0.10 } }
   }
}
```

The `DEFAULT` value of the `CAPACITANCE` here is `0.07`.

# Section 8
# Electrical Performance Modeling

## 8.1 Overview of modeling keywords

This section details the keywords used for performance modeling.

### 8.1.1 Timing models

The following tables show the set of keywords used for timing measurements and constraints. All keywords have implied semantics that restrict their capability to describe general temporal relations between arbitrary signals. For unrestricted purposes, the keyword `TIME` shall be used.

**Table 8-1 Timing measurements**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| `DELAY` | number | Second | `n` (nano) | time between two threshold crossings within two consecutive events on two pins. A causal relationship between the two events is implied. |
| `RETAIN` | number | Second | `n` (nano) | time when an output pin shall retain its value after an event on the related input pin. `RETAIN` appears always in conjunction with `DELAY` for the same two pins. |
| `SLEWRATE` | non-negative number | Second | `n` (nano) | time between two threshold crossings within one event on one pin. |

**Table 8-2 Timing constraints**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| `HOLD` | number | Second | `n` (nano) | minimum time limit for hold between two threshold crossings within two consecutive events on two pins |
| `NOCHANGE` | optional[a] non-negative number | Second | `n` (nano) | minimum time limit between two threshold crossings within two arbitrary consecutive events on one pin, in conjunction with `SETUP` and `HOLD` |
| `PERIOD` | non-negative number | Second | `n` (nano) | minimum time limit between two identical events within a sequence of periodical events |

**Table 8-2 Timing constraints, *continued***

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| PULSEWIDTH | number | Second | n (nano) | minimum time limit between two threshold crossings within two consecutive and complementary events on one pin |
| RECOVERY | number | Second | n (nano) | minimum time limit for recovery between two threshold crossings within two consecutive events on two pins |
| REMOVAL | number | Second | n (nano) | minimum time limit for removal between two threshold crossings within two consecutive events on two pins |
| SETUP | number | Second | n (nano) | minimum time limit for setup between two threshold crossings within two consecutive events on two pins |
| SKEW | number | Second | n (nano) | absolute value is maximum time limit between two threshold crossings within two consecutive events on two pins; the sign indicates positive or negative direction |

a. The associated SETUP and HOLD measurements provide data. NOCHANGE itself need not provide data.

**Table 8-3 : Generalized timing measurements**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| TIME | number | Second | 1 (unit) | time point for waveform modeling, time span for average, RMS, and peak modeling |
| FREQUENCY | non-negative number | Hz | meg (mega) | frequency |
| JITTER | non-negative number | Second | n (nano) | uncertainty of arrival time |

**Table 8-4 : Normalized measurements**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| THRESHOLD | non-negative number between 0 and 1 | Normalized signal voltage swing | 1 (unit) | fraction of signal voltage swing, specifying a reference point for timing measurement data. The threshold is the voltage for which the timing measurement is taken. |
| NOISE_MARGIN | non-negative number between 0 and 1 | Normalized signal voltage swing | 1 (unit) | fraction of signal voltage swing, specifying the noise margin. The noise margin is a deviation of the actual voltage from the expected voltage for a specified signal level |

### 8.1.2 Analog models

This subsection defines the keywords for analog modeling.

**Table 8-5 : Analog measurements**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| CURRENT | number | Ampere | m (milli) | electrical current drawn by the cell. A pin can be specified as annotation.[a] |
| ENERGY | number | Joule | p (pico) | electrical energy drawn by the cell, including charge and discharge energy, if applicable. |
| POWER | number | Watt | u (micro) | electrical power drawn by the cell, including charge and discharge power, if applicable. |
| TEMPERATURE | number | o Celsius | 1 (unit) | temperature |
| VOLTAGE | number | Volt | 1 (unit) | voltage |
| FLUX | non-negative number | Coulomb per Square Meter | 1 (unit) | amount of hot electrons in units of electrical charge per gate oxide area |
| FLUENCE | non-negative number | Second times Coulomb per Square Meter | 1 (unit) | integral of FLUX over time |

a. If the annotated PIN has PINTYPE=supply, the CURRENT measurement qualifies for power analysis. In this case, the current includes charge/discharge current, if applicable.

**Table 8-6 : Electrical components**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| CAPACITANCE | non-negative number | Farad | p (pico) | pin, wire, load, or net capacitance |
| INDUCTANCE | non-negative number | Henry | n (nano) | pin, wire, load, or net inductance |
| RESISTANCE | non-negative number | Ohm | K (kilo) | pin, wire, load, or net resistance |

### 8.1.3 Supplementary models

This subsection defines the keywords for supplementary models.

**Table 8-7 : Abstract measurements**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| DRIVE_STRENGTH | non-negative number | None | 1 (unit) | drive strength of a pin, abstract measure for (drive resistance)$^{-1}$ |
| SIZE | non-negative number | None | 1 (unit) | abstract cost function for actual or estimated area of a cell or a block |

**Table 8-8 : Discrete measurements**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| SWITCHING_BITS | non-negative number | None | 1 | number of switching bits on a bus |
| FANOUT | non-negative number | None | 1 | number of receivers connected to a net |
| FANIN | non-negative number | None | 1 | number of drivers connected to a net |
| CONNECTIONS | non-negative number | None | 1 | number of pins connected to a net, where CONNECTIONS = FANIN+FANOUT |

The actual values for discrete measurements are always integer numbers, however, estimated values can be non-integer numbers (e.g., the average fanout of a net is 2.4).

Table 8-9 describes the arguments for arithmetic models to describe environmental dependency.

**Table 8-9 : Environmental data**

| Annotation string | Value type | Description |
|---|---|---|
| DERATE_CASE | string | derating case, i.e., the combination of process, supply voltage, and temperature |
| PROCESS | string | process corner |
| TEMPERATURE | number | environmental temperature |

# 8.2 Auxiliary statements for timing models

This section details the auxiliary statements used for timing modeling.

## 8.2.1 THRESHOLD definition

The THRESHOLD represents a reference voltage level for timing measurements, normalized to the signal voltage swing and measured with respect to the logic 0 voltage level, as shown in Figure 8-1.

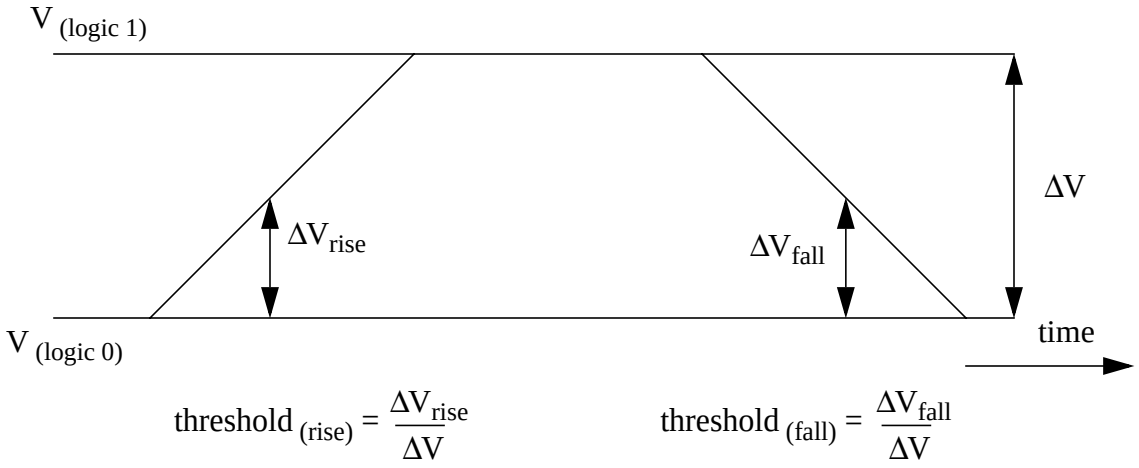



**Figure 8-1: THRESHOLD measurement definition**

The voltage levels for logic 1 and 0 represent a full voltage swing.

Different threshold data for RISE and FALL can be specified or else the data shall apply for both rising and falling transitions.

The THRESHOLD statement has the form of an arithmetic model. If the submodel keywords RISE and FALL are used, it has the form of an arithmetic model container.

Examples:

```
THRESHOLD = 0.4;
THRESHOLD { RISE = 0.3;  FALL = 0.5; }
THRESHOLD { HEADER { TEMPERATURE {TABLE{ 0 50 100 }}} TABLE { 0.5 0.4 0.3}}
```

### 8.2.2 FROM and TO container

A FROM container and a TO container shall be used inside timing measurements and timing constraints. Depending on the semantics of the timing model (see Section 8.3), they can contain a THRESHOLD statement, PIN annotation, and/or EDGE_NUMBER annotation. The data in the FROM and TO containers define the measurement start and end point, respectively.

Example:

```
DELAY {
   FROM {PIN = data_in;  THRESHOLD { RISE = 0.4; FALL = 0.6;} }
   TO   {PIN = data_out; THRESHOLD = 0.5;}
}
```

The delay is measured from pin data_in to pin data_out. The threshold for data_in is 0.4 for the rising signal and 0.6 for the falling signal. The threshold for data_out is 0.5, which applies for both the rising and falling signals.

### 8.2.3 PIN annotation

If the timing measurements or timing constraints, respectively, apply semantically for two pins (see Section 8.3.1), the FROM and TO containers shall each contain the PIN annotation.

Example:

```
DELAY {
   FROM { PIN = A ; }
   TO { PIN = Z ; }
}
```

Otherwise, if the timing measurements or timing constraints apply semantically only to one pin (see Section 8.3.3), the PIN annotation shall be outside the FROM or TO container.

Example:

```
SLEWRATE {
   PIN = A ;
}
```

### 8.2.4 EDGE_NUMBER annotation

The EDGE_NUMBER annotation within the context of a timing model shall specify the edge where the timing measurement applies. The timing model shall be in the context of a VECTOR. The EDGE_NUMBER shall have an unsigned value pointing to exactly one of subsequent vector_single_event expressions applicable to the referenced pin. The EDGE_NUMBER shall be counted individually for each pin which appears in the VECTOR, starting with zero (0).

If the timing measurements or timing constraints, apply semantically to two pins (see Section 8.3.1), the EDGE_NUMBER annotation shall be legal inside the FROM or TO container in conjunction with the PIN annotation.

Example:

```
DELAY {
   FROM { PIN = A ; EDGE_NUMBER = 0; }
   TO { PIN = Z ; EDGE_NUMBER = 0; }
}
```

Otherwise, if the timing measurements or timing constraints apply semantically only to one pin (see Section 8.3.3), the EDGE_NUMBER annotation shall be legal outside the FROM or TO container in conjunction with the PIN annotation.

Example:

```
SLEWRATE {
   PIN = A ;  EDGE_NUMBER = 0;
}
```

The default values for EDGE_NUMBER are specific for each timing model keyword (see Section 8.3).

The EDGE_NUMBER annotation is necessary for complex timing models involving multiple transitions on the same pin, as illustrated by the following figures and examples.

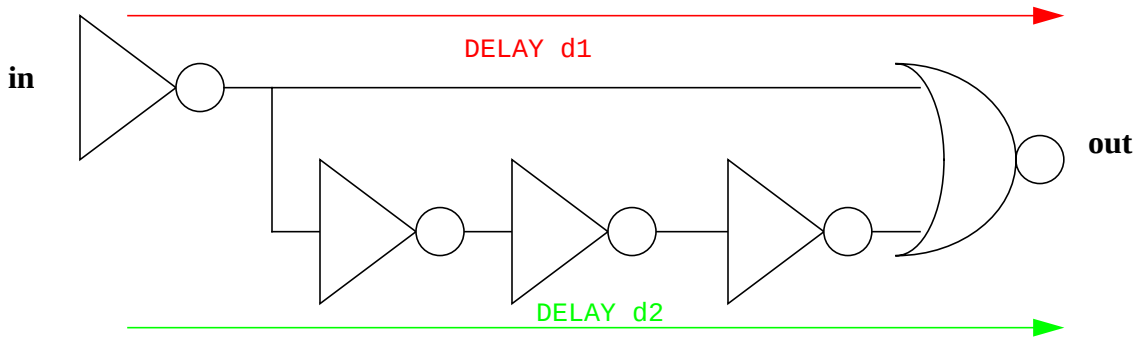



**Figure 8-2: Schematic of a pulse generator**

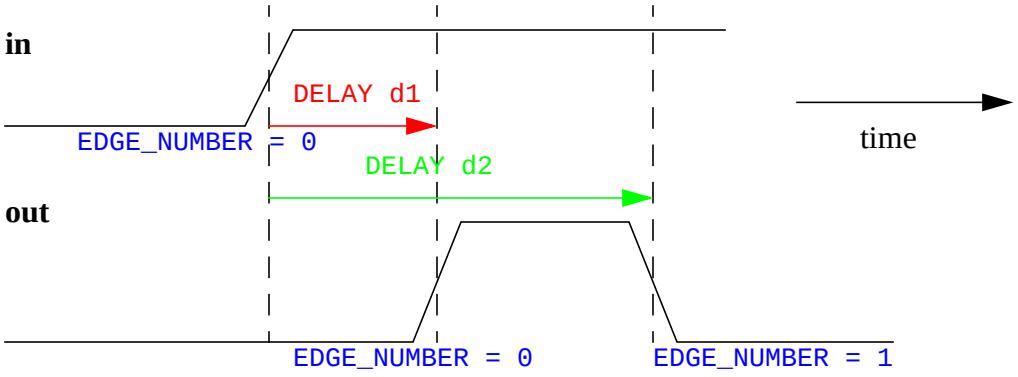



**Figure 8-3: Timing diagram of a pulse generator**

```
VECTOR ( 01 in -> 01 out -> 10 out ) {
   DELAY d1 {
      FROM { PIN = in; }
      TO { PIN = out; EDGE_NUMBER = 0; }
   }
   DELAY d2 {
      FROM { PIN = in; }
      TO { PIN = out; EDGE_NUMBER = 1; }
   }
}
```

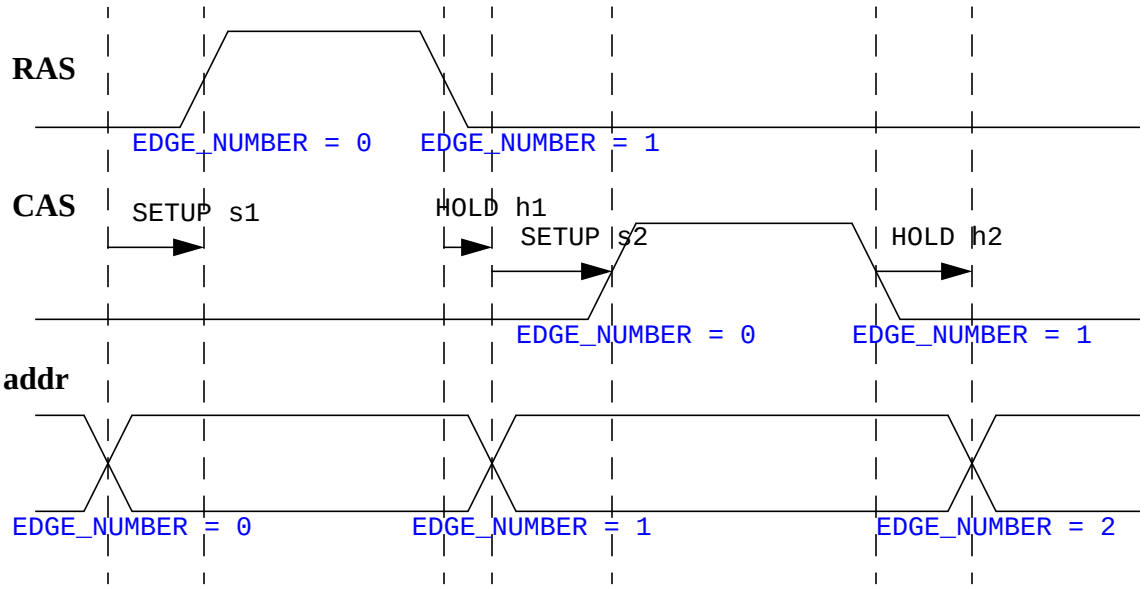



**Figure 8-4: Timing diagram of a DRAM cycle**

```
VECTOR(?! addr ->01 RAS ->10 RAS ->?! addr ->01 CAS ->10 CAS ->?! addr){
   SETUP s1 {
      FROM { PIN = addr; EDGE_NUMBER = 0; }
      TO { PIN = RAS; EDGE_NUMBER = 0; }
   }
   HOLD h1 {
      FROM { PIN = RAS; EDGE_NUMBER = 1; }
      TO { PIN = addr; EDGE_NUMBER = 1; }
   }
   SETUP s2 {
      FROM { PIN = addr; EDGE_NUMBER = 1; }
      TO { PIN = CAS; EDGE_NUMBER = 0; }
   }
   HOLD h2 {
      FROM { PIN = CAS; EDGE_NUMBER = 1; }
      TO { PIN = addr; EDGE_NUMBER = 2; }
   }
}
```

### 8.2.5 Context of THRESHOLD definitions

The THRESHOLD statement can appear in the context of a FROM or TO container. In this case, it specifies the applicable reference for the start and end point of the timing measurement, respectively.

Example:

```
SLEWRATE {
   FROM { THRESHOLD = 0.2; }
   TO { THRESHOLD = 0.8; }
}
```

The THRESHOLD statement can also appear in the context of a PIN. In this case, it specifies the applicable reference for the start or end point of timing measurements indicated by the PIN annotation inside a FROM or TO container, unless a THRESHOLD is specified explicitly inside the FROM or TO container.

If both the RISE and FALL thresholds are specified and the switching direction of the applicable pin is clearly indicated in the context of a VECTOR, the RISE or FALL data shall be applied accordingly.

Example:

```
PIN A { THRESHOLD { RISE = 0.3;  FALL = 0.5; } }
PIN Z { THRESHOLD = 0.4; }
// other statements ...
VECTOR ( 01 A -> 10 Z ) {
   DELAY { FROM { PIN=A; } TO { PIN=Z; } }
// the applicable threshold for A is 0.3
// the applicable threshold for Z is 0.4
```

If thresholds are needed for exact definition of the model data, the FROM and TO containers shall each contain an arithmetic model for THRESHOLD.

A THRESHOLD statement can also appear as argument of an arithmetic model for timing measurements. In this case, it shall contain a PIN annotation matching another PIN annotation in the FROM or TO container.

Example:

```
DELAY {
   FROM { PIN = A; THRESHOLD = 0.5; }
   TO { PIN = Z; }
   HEADER { THRESHOLD { PIN = Z; TABLE { 0.3 0.4 0.5 } }
   TABLE { 1.23 1.45 1.78 }
}
/* The measurement reference for pin A is always 0.5. The delay from A to
Z is expressed as a function of the measurement reference for pin Z. */
```

FROM and TO containers with THRESHOLD definitions, yet without PIN annotations, can appear within unnamed timing model definitions in the context of a VECTOR, CELL, WIRE, SUBLIBRARY, or LIBRARY object for the purpose of specifying global threshold definitions for all timing models within scope of the definition. The following priorities apply:

1. THRESHOLD in the HEADER of the timing model
2. THRESHOLD in the FROM or TO statement within the timing model
3. THRESHOLD for timing model definition in the context of the same VECTOR
4. THRESHOLD within the PIN definition
5. THRESHOLD for timing model definition in the context of the same CELL or WIRE
6. THRESHOLD for timing model definition in the context of the same SUBLIBRARY
7. THRESHOLD for timing model definition in the context of the same LIBRARY
8. THRESHOLD for timing model definition outside LIBRARY

Example:

```
LIBRARY my_library {
   DELAY {
      FROM { THRESHOLD = 0.4;  }
      TO   { THRESHOLD = 0.4;  }
   }
   SLEWRATE {
      FROM { THRESHOLD { RISE = 0.2; FALL = 0.8; } }
      TO   { THRESHOLD { RISE = 0.8; FALL = 0.2; } }
   }
   CELL my_cell {
      PIN A { DIRECTION=input; THRESHOLD { RISE = 0.3; FALL = 0.5; } }
      PIN Z { DIRECTION=output; }
      VECTOR (01 A -> 10 Z) {
```

```
            DELAY { FROM { PIN=A; } TO { PIN=Z; } }
            SLEWRATE { PIN = Z; }
         }
      }
   }
   // delay is measured from A (threshold=0.3) to Z (threshold=0.4)
   // slewrate on Z is measured from threshold=0.8 to threshold=0.2.
```

# 8.3 Specification of timing models

Timing models shall be specified in the context of a `VECTOR` statement.

## 8.3.1 TEMPLATE for timing measurements and timing constraints

The following templates show a general timing measurement and a general timing constraint description, respectively, applicable for two pins.

```
TEMPLATE TIMING_MEASUREMENT {
   <timeKeyword> = <timeValue> {
      FROM {
         PIN=<fromPin>;
         THRESHOLD=<fromThreshold>;
         EDGE_NUMBER=<fromEdge>;
      }
      TO {
         PIN=<toPin>;
         THRESHOLD=<toThreshold>;
         EDGE_NUMBER=<toEdge>;
      }
   }
}
TEMPLATE TIMING_CONSTRAINT {
   LIMIT {
      <timeKeyword> {
         FROM {
            PIN=<fromPin>;
            THRESHOLD=<fromThreshold>;
            EDGE_NUMBER=<fromEdge>;
         }
         TO {
            PIN=<toPin>;
            THRESHOLD=<toThreshold>;
            EDGE_NUMBER=<toEdge>;
         }
         MIN = <timeValueMin>;
         MAX = <timeValueMax>;
      }
   }
}
```

For simplicity, trivial arithmetic models shown here. In general, a `HEADER`, `TABLE`, or `EQUATION` construct can be used for calculation of `<timeValue>`, `<timeValueMin>`, or `<timeValueMax>`.

A particular timing constraint does not necessarily contain both <timeValueMin> and <timeValueMax>.

The <fromThreshold> and <toThreshold> can be globally predefined as explained in Section 8.2.4.

The vector_expression in the context where the <timeKeyword> appears shall contain at least two expressions of the type vector_single_event with the <fromPin> and <toPin>, respectively, as operands. The <fromEdge> and <toEdge> point to their respective vector_single_event, as shown in Figure 8-5.

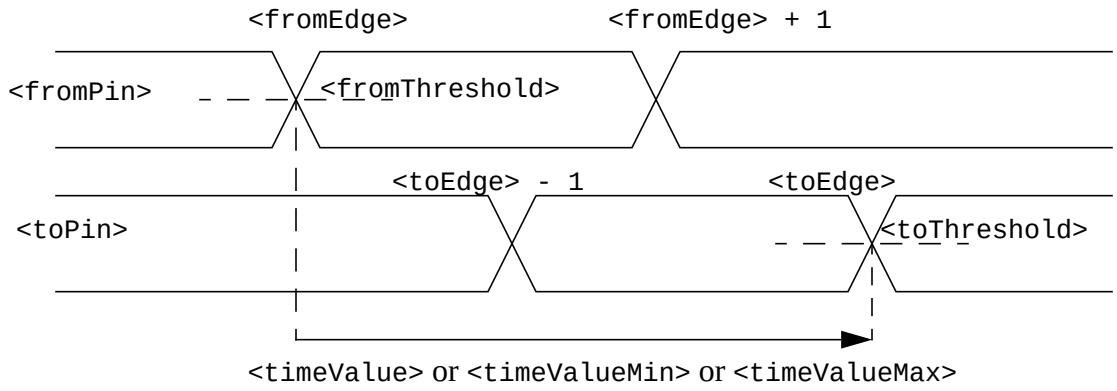



**Figure 8-5: General timing measurement or timing constraint**

The direction of the respective transition shall be identified by the respective edge_literal, i.e., the operator of the respective vector_single_event.

The temporal order of the LHS and RHS vector_single_event expressions within the vector_expression is indicated by a vector_binary operator.

The implications on the range of <timeValue> or <refPin> or <timeValueMax> are shown in Table 8-10.

**Table 8-10 Range of time value depending on VECTOR**

| LHS | operand | RHS | range of <timeValue> or <timeValueMin> or <timeValueMax> |
|---|---|---|---|
| <fromPin> | -> or ~> | <toPin> | positive |
| <toPin> | -> or ~> | <fromPin> | negative |
| <fromPin> | &> | <toPin> | positive or zero |
| <toPin> | &> | <fromPin> | negative or zero |
| <fromPin> | <-> | <toPin> | positive or negative |
| <toPin> | <-> | <fromPin> | positive or negative |
| <fromPin> | <&> | <toPin> | positive or negative or zero |
| <toPin> | <&> | <fromPin> | positive or negative or zero |

Note: This table does not apply for models with CALCULATION=incremental. Incremental values can always be positive, negative, or zero.

### 8.3.2 Partially defined timing measurements and constraints

A partially defined timing measurement or timing constraint contains only a FROM statement or a TO statement, but not both. This construct can be used to specify measurements from any point to a specific point (only TO is specified) or from a specific point to any point (only FROM is specified).

This is summarized in Table 8-11.

**Table 8-11 Partially specified timing measurements and constraints**

| DIRECTION of PIN | FROM or TO specified | Specified model applicable for |
|---|---|---|
| input | FROM only | cell timing arcs starting at this pin |
| input | TO only | interconnect timing arcs ending at this pin |
| output | FROM only | interconnect timing arcs starting at this pin |
| output | TO only | cell timing arcs ending at this pin |

It is recommended to use the constructs for interconnect timing arcs only in conjunction with CALCULATION=incremental. The <timeValue>, <timeValueMin>, or <timeValueMax> from this model is added to the <timeValue>, <timeValueMin>, or <timeValueMax> from timing arcs starting or ending at this pin, respectively. If the construct is used with CALCULATION=absolute, the timing model can only be used if completely specified interconnect timing models are not available and the result is not be accurate in general.

### 8.3.3 TEMPLATE for same-pin timing measurements and constraints

The following templates show a timing measurement and a timing constraint description, respectively, applicable for the same pin.

```
TEMPLATE SAME_PIN_TIMING_MEASUREMENT {
   <timeKeyword> = <timeValue> {
      PIN=<refPin>;
      EDGE_NUMBER=<refEdge>;
      FROM { THRESHOLD=<fromThreshold>; }
      TO { THRESHOLD=<toThreshold>; }
   }
}
```

```
TEMPLATE SAME_PIN_TIMING_CONSTRAINT {
   LIMIT {
      <timeKeyword> {
         PIN=<refPin>;
         EDGE_NUMBER=<refEdge>;
         FROM { THRESHOLD=<fromThreshold>; }
         TO { THRESHOLD=<toThreshold>; }
         MIN = <timeValueMin>;
         MAX = <timeValueMax>;
      }
   }
}
```

Depending on the <timeKeyword>, the <timeValue>, <timeValueMin>, or <timeValueMax> is measured on the same <refEdge> or between <refEdge> and <refEdge> *plus 1*.
Only the -> or ~> operators are applicable between subsequent edges. Therefore, the <timeValue>, <timeValueMin>, or <timeValueMax> are positive by definition.

Note: The <fromThreshold> and <toThreshold> can be globally predefined as explained in Section 8.2.4. However, the THRESHOLD in the context of a PIN does not apply for *SAME_PIN_TIMING_MEASUREMENT* or *SAME_PIN_TIMING_CONSTRAINT*, since the <refPin> is not within a FROM or TO statement.

### 8.3.4 Absolute and incremental evaluation of timing models

As mentioned in the previous sections, the calculation models for *TIMING_MEASUREMENT*, *TIMING_CONSTRAINT*, *SAME_PIN_TIMING_MEASUREMENT*, and *SAME_PIN_TIMING_CONSTRAINT* can have the annotation CALCULATION=absolute (the default) or CALCULATION=incremental. These annotations are only relevant more than one calculation model for the same timing arc exists.

Calculation models for the same timing arc with CALCULATION=absolute shall be within the context of mutually exclusive VECTORs. The vector_expression specifies which model to use under which condition.

Example:

```
VECTOR ( (01 A -> 01 Z) && B & !C ) {
   DELAY { CALCULATION=absolute; FROM { PIN=A; } TO { PIN=Z; }
   /* fill in HEADER, TABLE */ }
}
VECTOR ( (01 A -> 01 Z) && !B & C ) {
   DELAY { CALCULATION=absolute; FROM { PIN=A; } TO { PIN=Z; }
   /* fill in HEADER, TABLE */ }
}
```

The vectors ( (01 A -> 01 Z) && B & !C ) and ( (01 A -> 01 Z) && !B & C ) are mutually exclusive. They describe the same timing arc with two mutually exclusive conditions.

In the case of a VECTOR containing a calculation model for a timing arc with CALCULATION=incremental, there shall be another VECTOR with a calculation model for the same timing arc with CALCULATION=absolute and both vectors shall be compatible. The vector_expression of the latter shall necessarily be true when the vector_expression of the former is true.

Example:

```
VECTOR (01 A -> 01 Z) {
   DELAY { CALCULATION=absolute; FROM { PIN=A; } TO { PIN=Z; }
   /* fill in HEADER, TABLE */ }
}
VECTOR ( (01 A -> 01 Z) && B & !C ) {
   DELAY { CALCULATION=incremental; FROM { PIN=A; } TO { PIN=Z; }
   /* fill in HEADER, TABLE */ }
}
VECTOR ( (01 A -> 01 Z) && !B & C ) {
   DELAY { CALCULATION=incremental; FROM { PIN=A; } TO { PIN=Z; }
   /* fill in HEADER, TABLE */ }
}
```

The vectors `( (01 A -> 01 Z) && B & !C )` and `( (01 A -> 01 Z) && !B & C )` are both compatible with the vector `(01 A -> 01 Z)` and mutually exclusive with each other. The latter describe the same timing arc with two mutually exclusive conditions. The former describes the same timing arc without conditions. This modeling style is useful for timing analysis tools with or without support for conditions. The vectors with conditions, if supported, add accuracy to the calculation. However, the vector without conditions is always available for basic calculation.

### 8.3.5 RISE and FALL submodels

For timing models in the context of a `VECTOR`, submodels for `RISE` and `FALL` are only applicable if the `vector_expression` does not specify the switching direction of the referenced `PIN` and `EDGE_NUMBER`. This is the case, when symbolic `vector_unary` operators are used, i.e., `?!`, `??`, `?*`, or `*?` instead of `01`, `10`, etc.

For *`SAME_PIN_TIMING_MEASUREMENT`* or *`SAME_PIN_TIMING_CONSTRAINT`*, the `RISE` and `FALL` submodels apply for the `<refEdge>`.

For a partially specified *`TIMING_MEASUREMENT`* or *`TIMING_CONSTRAINT`*, the `RISE` and `FALL` submodels apply for the `<fromEdge>` or `<toEdge>`, whichever is specified.

For a completely specified *`TIMING_MEASUREMENT`* or *`TIMING_CONSTRAINT`*, it is not possible to apply a `RISE` and `FALL` submodel for both `<fromEdge>` and `<toEdge>`. The `vector_unary` operator shall specify the switching direction for at least one edge. If the switching direction for both edges is unspecified, the `RISE` and `FALL` submodel shall apply for the `<toEdge>`.

Example:

```
VECTOR ( 01 CLK -> ?! Q ) {
   DELAY { FROM { PIN = CLK; } TO { PIN = Q; }
      RISE = 0.76; FALL = 0.58;
   }
}

// If Q is a scalar pin, the following construct is equivalent:
VECTOR ( 01 CLK -> 01 Q ) {
   DELAY = 0.76 { FROM { PIN = CLK; } TO { PIN = Q; } }
}
```

```
VECTOR ( 01 CLK -> 10 Q ) {
   DELAY = 0.58 { FROM { PIN = CLK; } TO { PIN = Q; } }
}
```

### 8.3.6 TIME

The <timeKeyword> TIME describes a general *TIMING_MEASUREMENT* or *TIMING_CONSTRAINT* without implying any particular relationship between <fromEdge> and <toEdge>.

In general, <fromPin> and <toPin> refer to two different pins. However, it is legal for <fromPin> and <toPin> to refer to the same pin.

The default value for <fromEdge> and <toEdge> shall be 0.

### 8.3.7 DELAY

The <timeKeyword> DELAY describes a *TIMING_MEASUREMENT* implying a causal relationship between <fromEdge> and <toEdge>.

Usually, <fromPin> refers to an input pin and <toPin> refers to an output pin. However, it is legal for <fromPin> and <toPin> to refer to an output pin.

The default value for <fromEdge> and <toEdge> shall be 0, unless the DELAY statement appears in conjunction with a RETAIN statement within the context of the same VECTOR.

### 8.3.8 RETAIN

The <timeKeyword> RETAIN describes a *TIMING_MEASUREMENT* implying a causal relationship between <fromEdge> and <toEdge> in the same way as DELAY.

RETAIN is used to describe the elapsed time until the output changes its old value, whereas DELAY is used to describe the elapsed time until the output settles to a stable new value, as shown in Figure 8-6.

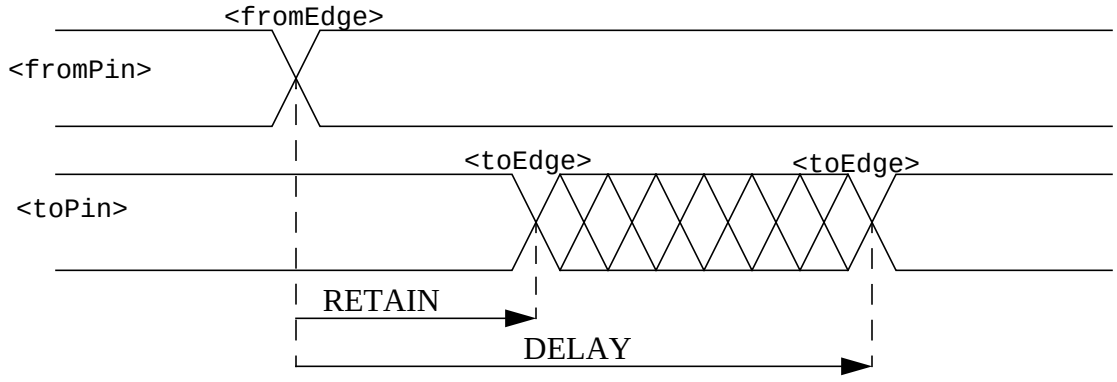



**Figure 8-6: RETAIN and DELAY**

When DELAY appears in conjunction with RETAIN, the <fromEdge> for both measurements shall be the same. The <toEdge> for DELAY shall be the <toEdge> for RETAIN *plus 1*.

The default value for <fromEdge> and <toEdge> for RETAIN shall be 0. The default value for <toEdge> for DELAY shall be 1.

### 8.3.9 SLEWRATE

The <timeKeyword> SLEWRATE describes a *SAME_PIN_TIMING_MEASUREMENT* for <timeValue> defining the duration of a signal transition or a fraction thereof.

The SLEWRATE applies for the <refEdge> on the <refPin>. The default value for <refEdge> shall be 0.

### 8.3.10 SETUP

The <timeKeyword> SETUP describes a *TIMING_CONSTRAINT* for <timeValueMin> defining the minimum stable time required for the data signal on the <fromPin> before it is sampled by the strobe signal on the <toPin>.

The <fromPin> usually is an input pin with SIGNALTYPE=data. The <toPin> is an input pin with SIGNALTYPE=clock.

The default value for <fromEdge> and <toEdge> for SETUP shall be 0.

### 8.3.11 HOLD

The <timeKeyword> HOLD describes a *TIMING_CONSTRAINT* for <timeValueMin> defining the minimum stable time required for the data signal on the <toPin> after it is sampled by the strobe signal on the <fromPin>.

The <toPin> usually is an input pin with SIGNALTYPE=data. The <fromPin> is an input pin with SIGNALTYPE=clock.

The default value for <fromEdge> shall be 0. The default value for <toEdge> shall be 0, unless HOLD appears in conjunction with SETUP in the context of the same VECTOR. In that case, the default value for <toEdge> shall be 1. All of this is depicted in Figure 8-7.

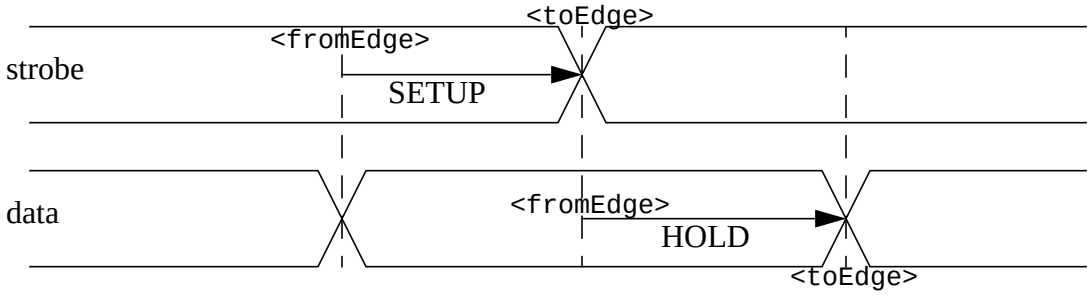



**Figure 8-7: SETUP and HOLD**

The <timeValueMin> for SETUP or the <timeValueMin> for HOLD with respect to the same strobe can be negative. However, the sum of both values shall be positive. The sum represents the minimum duration of a valid data signal around a strobe signal.

### 8.3.12 NOCHANGE

The <timeKeyword> NOCHANGE describes a *SAME_PIN_TIMING_CONSTRAINT* defining the requirement for a stable signal on a pin subjected to SETUP and HOLD on subsequent edges of a strobe signal., as shown in Figure 8-8.

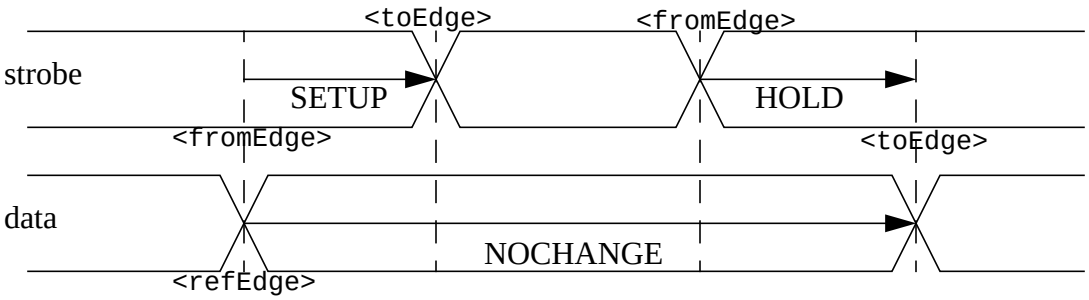



**Figure 8-8: NOCHANGE, SETUP, and HOLD**

The NOCHANGE applies between the <refEdge> and the subsequent edge, i.e., <refEdge> *plus 1* on the <refPin>. The default value for <refEdge> shall be 0.

When NOCHANGE appears in conjunction with SETUP and HOLD within the context of the same VECTOR, the default value for <fromEdge> and <toEdge> of SETUP shall be 0 and the default value for <fromEdge> and <toEdge> of HOLD shall be 1.

### 8.3.13 RECOVERY

The <timeKeyword> RECOVERY describes a *TIMING_CONSTRAINT* for <timeValueMin> defining the minimum stable time required for an asynchronous control signal on the <fromPin> to be inactive before a strobe signal on the <toPin> can be active.

The <fromPin> usually is an input pin with SIGNALTYPE=set|clear. The <toPin> is an input pin with SIGNALTYPE=clock.

The default value for <fromEdge> and <toEdge> for RECOVERY shall be 0.

### 8.3.14 REMOVAL

The <timeKeyword> REMOVAL describes a *TIMING_CONSTRAINT* for <timeValueMin> defining the minimum stable time required for an asynchronous control signal on the <toPin> to remain active after overriding a strobe signal on the <fromPin>.

The <toPin> usually is an input pin with SIGNALTYPE=set|clear. The <fromPin> is an input pin with SIGNALTYPE=clock.

The default value for <fromEdge> and <toEdge> for REMOVAL shall be 0.

REMOVAL can appear in conjunction with RECOVERY within the context of the same VECTOR, as shown in Figure 8-9.

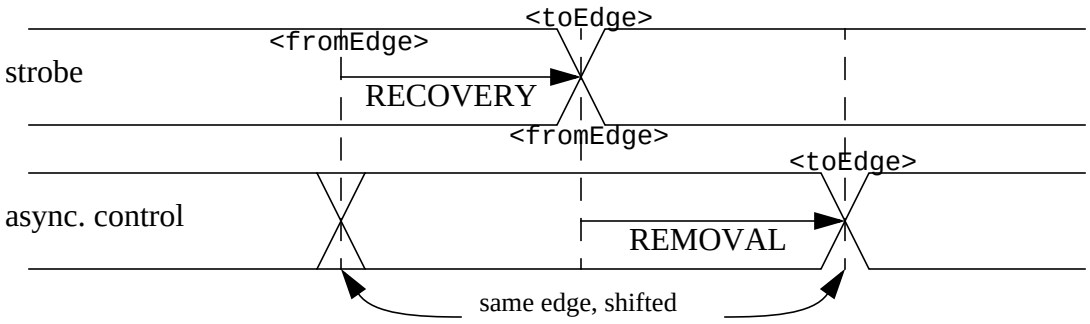



**Figure 8-9: RECOVERY and REMOVAL**

The <timeValueMin> for RECOVERY or the <timeValueMin> for REMOVAL with respect to the same strobe can be negative. However, the sum of both values shall be positive. The sum represents the time window around the clock signal when the asynchronous control signal shall not switch.

### 8.3.15 SKEW between two signals

The <timeKeyword> SKEW describes a *TIMING_CONSTRAINT* for <timeValueMax> defining the maximum allowed time separation between <fromEdge> on <fromPin> and <toEdge> on <toPin>.

The default value for <fromEdge> and <toEdge> for SKEW shall be 0.

### 8.3.16 SKEW between multiple signals

SKEW can also describe the maximum time distortion between signals on multiple pins. In this case, a list of pins appears in form of a multi-value annotation. No FROM or TO containers can be used here.

Example:

```
SKEW {
   PIN { <pinList> }
   EDGE_NUMBER { <edgeList> }
   <skewData>
}
```

The default for EDGE_NUMBER in SKEW for multiple signals shall be a list of 0s.

A special case of multiple pins is a single bus. In this case, the `unnamed_assignment` syntax is also valid as alternative to the `multi_value_assignment` syntax (see Section 8.15.3).

Example:

```
SKEW { PIN = my_bus_pin[8:1]; }
```

or

```
SKEW { PIN { my_bus_pin[8:1] } }
```

### 8.3.17 PULSEWIDTH

The `<timeKeyword> PULSEWIDTH` describes a *`SAME_PIN_TIMING_CONSTRAINT`* for `<timeValueMin>` defining the minimum duration of the signal before changing state.

The `PULSEWIDTH` statement is applicable for both input and output pins. In the case of an input pin, it represents a timing check against the minimum duration. In case of an output pin, it represents the minimum possible duration of the signal.

The `PULSEWIDTH` applies between the `<refEdge>` and the subsequent edge, i.e., `<refEdge>` *plus 1* on the `<refPin>`. The default value for `<refEdge>` shall be `0`.

### 8.3.18 PERIOD

The `<timeKeyword> PERIOD` describes a *`SAME_PIN_TIMING_CONSTRAINT`* for `<timeValueMin>` defining the minimum time between subsequent repetitions of a signal. Because of periodicity, `<fromThreshold>` and `<toThreshold>` are not required. Therefore, `FROM` and `TO` statements do not appear.

If the `VECTOR` describes a completely specified event sequence, `<refPin>` and `<refEdge>` are not required. `PERIOD` applies for the complete event sequence. If the `VECTOR` describes a partially specified event sequence, involving the `~>` operator, `<refPin>` and `<refEdge>` are required.

### 8.3.19 JITTER

The `<timeKeyword> JITTER` describes a *`SAME_PIN_TIMING_MEASUREMENT`* for `<timeValue>` defining the actual uncertainty of arrival time for a periodical signal at a pin.

The `JITTER` applies for the `<refEdge>` on the `<refPin>`. The default value for `<refEdge>` shall be `0`. Threshold definitions, i.e., `<fromThreshold>` or `<toThreshold>` do not apply.

A limit for tolerable jitter at a pin can be expressed using the `LIMIT` construct, as shown in the template for *`SAME_PIN_TIMING_CONSTRAINT`*.

## 8.4 VIOLATION container

A `VIOLATION` statement can appear within an `ILLEGAL` statement (see Section 6.7) and also within a *`TIMING_CONSTRAINT`* or a *`SAME_PIN_TIMING_CONSTRAINT`*. The `VIOLATION` statement can contain the `BEHAVIOR` object (see Section 11.17), since the behavior in case of timing

constraint violation cannot be described in the FUNCTION. The VIOLATION statement can also contain the annotations shown in Table 8-12.

**Table 8-12 : Annotations within VIOLATION**

| Keyword | Value type | Description |
|---|---|---|
| MESSAGE_TYPE | string | specifies the type of the message. It can be one of information, warning, or error. |
| MESSAGE | string | specifies the message itself. |

Example:

```
VECTOR (01 d <&> 01 cp) {
   SETUP {
      VIOLATION {
            MESSAGE_TYPE = error;
            MESSAGE = “setup violation 01 d <&> 01 cp“;
            BEHAVIOR {q = 'bx;}
      }
   }
}
```

## 8.5 EARLY and LATE container

The EARLY and LATE containers define the boundaries of timing measurements in one single analysis. They only apply to DELAY and SLEWRATE. Both of them need to appear in both containers.

The quadruple

```
EARLY {
   DELAY { FROM {...} TO { ...} /* data */ }
   SLEWRATE { /* data */ }
LATE {
   DELAY { FROM {...} TO { ...} /* data */ }
   SLEWRATE { /* data */ }
```

is used to calculate the envelope of the timing waveform at the TO point of a delay arc with respect to the timing waveform at the FROM point of a delay arc.

The EARLY DELAY is a smaller number (or a set of smaller numbers) than the LATE DELAY. However, the EARLY SLEWRATE is not necessarily smaller than the LATE SLEWRATE, since the SLEWRATE of the EARLY signal can be larger than the SLEWRATE of the LATE signal.

## 8.6 Environmental dependency for electrical data

This section defines the environmental dependencies for electrical data.

### 8.6.1 PROCESS

The following identifiers can be used as predefined process corners:

| | |
|---|---|
| `?n?p` | process definition with transistor strength |

where `?` can be

| | |
|---|---|
| `s` | strong |
| `w` | weak |

The possible process name combinations are shown in Table 8-13.

**Table 8-13 : Predefined process names**

| Process name | Description |
|---|---|
| `snsp` | strong NMOS, strong PMOS |
| `snwp` | strong NMOS, weak PMOS |
| `wnsp` | weak NMOS, strong PMOS |
| `wnwp` | weak NMOS, weak PMOS |

### 8.6.2 DERATE_CASE

The following identifiers can be used as predefined derating cases:

| | |
|---|---|
| `nom` | nominal case |
| `bc?` | prefix for best case |
| `wc?` | prefix for worst case |

where `?` can be

| | |
|---|---|
| `com` | suffix for commercial case |
| `ind` | suffix for industrial case |
| `mil` | suffix for military case |

The possible derating case combinations are defined in Table 8-14.

**Table 8-14 : Predefined derating cases**

| Derating case | Description |
|---|---|
| `bccom` | best case commercial |
| `bcind` | best case industrial |
| `bcmil` | best case military |
| `wccom` | worst case commercial |
| `wcind` | worst case military |
| `wcmil` | worst case military |

### 8.6.3 Lookup table without interpolation

The `PROCESS` or `DERATE_CASE` can be used in a `TABLE` within the `HEADER` of an arithmetic model for electrical data, e.g., `DELAY`. Data can not be interpolated in the dimension of this table.

Example:

```
DELAY {
   UNIT = ns;
   HEADER {
      PROCESS { TABLE { nom snsp wnwp } }
   }
   TABLE { 0.4  0.3  0.6 }
}
```

Here , the DELAY is 0.4 ns for nominal process, 0.3 ns for snsp, and 0.6 ns for wnwp. A delay "in-between" snsp and wnwp can not be interpolated.

### 8.6.4 Lookup table for process- or derating-case coefficients

A nested arithmetic model construct can be used to describe lookup tables for coefficients, based on PROCESS or DERATE_CASE. These coefficients can be used in an EQUATION to calculate electrical data, e.g., DELAY.

Example:

```
DELAY {
   UNIT = ns;
   HEADER {
      PROCESS { HEADER { nom snsp wnwp } TABLE {0.0 -0.25 0.5} }
   }
   EQUATION { (1 + PROCESS)*0.4 }
}
```

The equation uses the PROCESS coefficient 0.0 for nominal, -0.25 for snsp, and 0.5 for wnwp. Therefore the DELAY is 0.4 ns for the nominal process, 0.3 ns for snsp, and 0.6 ns for wnwp. Conceivably, the DELAY can be calculated for any value of the coefficient.

### 8.6.5 TEMPERATURE

TEMPERATURE can be used as argument in the HEADER of an arithmetic model for timing or electrical data. It can also be used as an arithmetic model with DERATE_CASE as argument, in order to describe what temperature applies for the specified derating case.

## 8.7 PIN-related arithmetic models for electrical data

This section details the PIN arithmetic models for electrical data.

### 8.7.1 Principles

Arithmetic models for electrical data can be associated with a pin of a cell. Their meaning is illustrated in Figure 8-10.

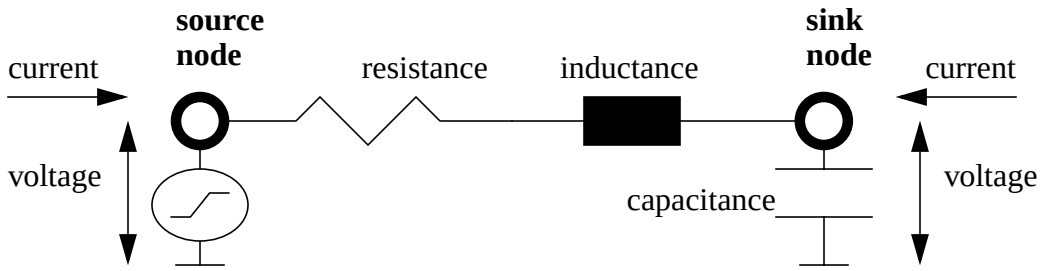



**Figure 8-10: General representation of electrical models around a pin**

A pin is represented as a source node and a sink node. For pins with `DIRECTION=input`, the source node is externally accessible. For pins with `DIRECTION=output`, the sink node is externally accessible.

### 8.7.2 CAPACITANCE, RESISTANCE, and INDUCTANCE

`RESISTANCE` and `INDUCTANCE` apply between the source and sink node. `CAPACITANCE` applies between the sink node and ground. By default, the values for resistance, inductance and capacitance shall be zero (`0`).

### 8.7.3 VOLTAGE and CURRENT

`VOLTAGE` and `CURRENT` can be measured at either source or sink node, depending on which node is externally accessible. However, a voltage source can only be connected to a source node. The sense of measurement for voltage shall be from the node to ground. The sense of measurement for current shall be *into* the node.

### 8.7.4 PIN-related timing models

*`SAME_PIN_TIMING_MEASUREMENT`* and *`SAME_PIN_TIMING_CONSTRAINT`* (see Section 8.3 and Section 8.7.6) are pin-related timing models. They are defined with reference to the externally accessible node.

### 8.7.5 Submodels for RISE, FALL, HIGH, and LOW

`RISE` and `FALL` contain data characterized in transient measurements. `HIGH` and `LOW` contain data characterized in static measurements.

```
<modelKeyword> { RISE=<modelValueRise>; FALL=<modelValueFall>; }
<modelKeyword> { HIGH=<modelValueHigh>;  LOW=<modelValueLow>; }
```

It is generally not required that both `RISE` and `FALL` or both `HIGH` and `LOW`, respectively, appear as an arithmetic submodel.

`HIGH` and `LOW` qualify states with the logic value `1` and `0`, respectively. `RISE` and `FALL` qualify transitions between states with initial logic value `0` and `1`, respectively and final values `1` and

0, respectively. For other states and their mapping to logic values, see Section 5.1.5. If the arithmetic model is within the scope of a vector which describes the logic values without ambiguity, the use of RISE and FALL or HIGH and LOW does not apply.

HIGH, LOW, RISE, and FALL apply for all pin-related arithmetic models with the following exceptions:

RISE and FALL do not apply for VOLTAGE.
HIGH and LOW do not apply for *SAME_PIN_TIMING_MEASUREMENT* and *SAME_PIN_TIMING_CONSTRAINT*.

Note: For states that cannot be mapped to logic 1 or 0, RISE and FALL, or HIGH and LOW cannot be used. The use of VECTOR with unambiguous description of the relevant states is mandatory in such cases.

### 8.7.6 Context-specific semantics

An arithmetic model for VOLTAGE, CURRENT, SLEWRATE, RESISTANCE, INDUCTANCE, and CAPACITANCE can be associated with a PIN in one of the following ways.

1. A model in the context of a PIN

Example:

```
PIN my_pin {
   CAPACITANCE = 0.025;
```

2. A model in the context of a CELL, WIRE, or VECTOR with PIN annotation

Example:

```
VOLTAGE = 1.8 { PIN = my_pin; }
```

The model in the context of a PIN shall be used if the data is completely confined to the pin. That means, no argument of the model shall make reference to any pin, since such reference implies an external dependency. A model with dependency only on environmental data not associated with a pin (e.g., TEMPERATURE, PROCESS, and DERATE_CASE) can be described within the context of the PIN.

A model with dependency on external data applied to a pin (e.g., load capacitance) shall be described outside the context of the PIN, using a PIN annotation. In particular, if the model involves a dependency on logic state or logic transition of other PINs, the model shall be described within the context of a VECTOR.

Figure 8-11 illustrates electrical models associated with input and output pins.

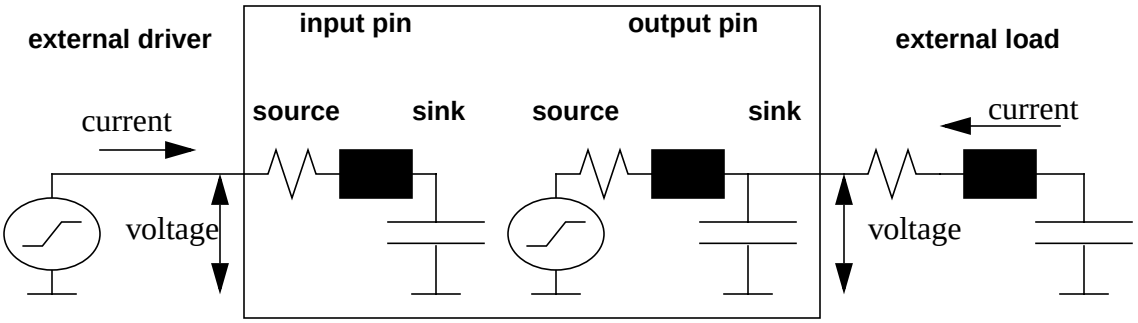



**Figure 8-11: Electrical models associated with input and output pins**

Table 8-15 and Table 8-16 define how models are associated with the pin, depending on the context.

**Table 8-15 Direct association of models with a PIN**

| Model | Model in context of PIN | Model in context of CELL, WIRE, and VECTOR with PIN annotation |
|---|---|---|
| CAPACITANCE | pin self-capacitance | externally controlled capacitance at the pin, e.g., voltage-dependent |
| INDUCTANCE | pin self-inductance | externally controlled inductance at the pin, e.g., voltage-dependent |
| RESISTANCE | pin self-resistance | externally controlled resistance at the pin, e.g., voltage-dependent, in the context of a VECTOR for timing-arc specific driver resistance |
| VOLTAGE | operational voltage measured at pin | externally controlled voltage at the pin |
| CURRENT | operational current measured into pin | externally controlled current into pin |
| *SAME_PIN_TIMING_MEASUREMENT* | for model definition, default, etc.; not for the timing arc | in context of VECTOR for timing arc, other context for definition, default, etc. |
| *SAME_PIN_TIMING_CONSTRAINT* | for model definition, default, etc.; not for the timing arc | in context of VECTOR for timing arc, other context for definition, default, etc. |

**Table 8-16 External association of models with a PIN**

| Model / Context | LIMIT within PIN or with PIN annotation | Model argument with PIN annotation |
|---|---|---|
| CAPACITANCE | min or max limit for applicable load | load for model characterization |
| INDUCTANCE | min or max limit for applicable load | load for model characterization |
| RESISTANCE | min or max limit for applicable load | load for model characterization |
| VOLTAGE | min or max limit for applicable voltage | voltage for model characterization |

**Table 8-16 External association of models with a PIN, *continued***

| Model / Context | LIMIT within PIN or with PIN annotation | Model argument with PIN annotation |
|---|---|---|
| CURRENT | min or max limit for applicable current | current for model characterization |
| *SAME_PIN_TIMING_MEASUREMENT* | currently applicable for min or max limit for SLEWRATE | stimulus with SLEWRATE for model characterization |
| *SAME_PIN_TIMING_CONSTRAINT* | N/A, since the keyword means a min or max limit by itself | N/A |

Example:

```
CELL my_cell {
   PIN pin1 { DIRECTION=input; CAPACITANCE = 0.05; }
   PIN pin2 { DIRECTION=output; LIMIT { CAPACITANCE { MAX=1.2; } } }
   PIN pin3 { DIRECTION=input; }
   PIN pin4 { DIRECTION=input; }
   CAPACITANCE {
      PIN=pin3;
      HEADER { VOLTAGE { PIN=pin4; } }
      EQUATION { 0.25 + 0.34*VOLTAGE }
   }
}
```

The capacitance on `pin1` is `0.05`. The maximum allowed load capacitance on `pin2` is `1.2`. The capacitance on `pin3` depends on the voltage on pin4.

# 8.8 Other PIN-related arithmetic models

This section details some other `PIN`-related arithmetic models.

## 8.8.1 DRIVE_STRENGTH

`DRIVE_STRENGTH` is a unit-less, abstract measure for the drivability of a `PIN`. It can be used as a substitute of driver `RESISTANCE`. The higher the `DRIVE_STRENGTH`, the lower the driver `RESISTANCE`. However, `DRIVE_STRENGTH` can only be used within a coherent system of calculation models, since it does not represent an absolute quantity, as opposed to `RESISTANCE`. For example, the weakest driver of a library can have drive strength 1, the next stronger driver can have drive strength 2 and so forth. This does not necessarily mean the resistance of the stronger driver is exactly half of the resistance of the weaker driver.

An arithmetic model for conversion from `DRIVE_STRENGTH` to `RESISTANCE` can be given to relate the quantity `DRIVE_STRENGTH` across technology libraries.

Example:

```
SUBLIBRARY high_speed_library {
   RESISTANCE {
      HEADER { DRIVE_STRENGTH } EQUATION { 800 / DRIVE_STRENGTH }
   }
```

```
      CELL high_speed_std_driver {
         PIN Z { DIRECTION = output; DRIVE_STRENGTH = 1; }
      }
   }
   SUBLIBRARY low_power_library {
      RESISTANCE {
         HEADER { DRIVE_STRENGTH } EQUATION { 1600 / DRIVE_STRENGTH }
      }
      CELL low_power_std_driver {
         PIN Z { DIRECTION = output; DRIVE_STRENGTH = 1; }
      }
   }
```

Drive strength 1 in the high speed library corresponds to 800 ohm. Drive strength 1 in the low power library corresponds to 1600 ohm.

Note: Any particular arithmetic model for `RESISTANCE` in either library shall locally override the conversion formula from drive strength to resistance.

### 8.8.2 SWITCHING_BITS

The quantity `SWITCHING_BITS` applies only for bus pins. The range is from `0` to the width of the bus. Usually, the quantity `SWITCHING_BITS` is not calculated by an arithmetic model, since the number of switching bits on a bus depends on the functional specification rather than the electrical specification. However, `SWITCHING_BITS` can be used as argument in the `HEADER` of an arithmetic model to calculate electrical quantities, for instance, energy consumption.

Example:

```
   CELL my_rom {
      PIN [3:0] addr { DIRECTION=input;  SIGNALTYPE=address; }
      PIN [7:0] dout { DIRECTION=output; SIGNALTYPE=data; }
      VECTOR ( ?! addr -> ?! dout ) {
         ENERGY {
            HEADER {
               SWITCHING_BITS addr_bits { PIN = addr; }
               SWITCHING_BITS dout_bits { PIN = dout; }
            }
            EQUATION { 0.45*LOG(addr_bits) + 2.6*dout_bits }
         }
      }
   }
```

The energy consumption of `my_rom` depends on the number of switching data bits and on the logarithm of the number of switching address bits.

## 8.9 Annotations for arithmetic models

This section defines the annotations for arithmetic models.

### 8.9.1 MEASUREMENT annotation

Arithmetic models describing analog measurements (see Table 8-5) can have a `MEASUREMENT` annotation. This annotation indicates the type of measurement used for the computation in arithmetic model.

```
MEASUREMENT = string ;
```

The string can take the values shown in Table 8-17.

**Table 8-17 : MEASUREMENT annotation**

| Annotation string | Description |
|---|---|
| `transient` | measurement is a transient value |
| `static` | measurement is a static value |
| `average` | measurement is an average value |
| `rms` | measurement is an root mean square value |
| `peak` | measurement is a peak value |

Their mathematical definitions are shown in Figure 8-12.

`transient` $\int_{(t=0)}^{(t=T)} dE(t)$

`average` $\frac{\int_{(t=0)}^{(t=T)} E(t)dt}{T}$

`static` $E = constant$

`rms` $\sqrt{\frac{\int_{(t=0)}^{(t=T)} E(t)^2 dt}{T}}$

`peak` $max(|E(t)|) \cdot \mathrm{sgn}E(t) \qquad t = T$

**Figure 8-12: Mathematical definitions for MEASUREMENT annotations**

Examples:

transient measurement of `ENERGY`
static measurement of `VOLTAGE`, `CURRENT`, and `POWER`
average measurement of `VOLTAGE`, `CURRENT`, and `POWER`
rms measurement of `VOLTAGE`, `CURRENT`, and `POWER`
peak measurement of `VOLTAGE`, `CURRENT`, and `POWER`

### 8.9.2 TIME and FREQUENCY annotation

Arithmetic models with certain values of `MEASUREMENT` annotation can also have *either* `TIME` *or* `FREQUENCY` as annotations.

The semantics are defined in Table 8-18.

**Table 8-18 : Semantic interpretation of MEASUREMENT, TIME, or FREQUENCY annotation**

| MEASUREMENT annotation | Semantic meaning of TIME annotation | Semantic meaning of FREQUENCY annotation |
|---|---|---|
| `transient` | integration of analog measurement is done during that time window | integration of analog measurement is repeated with that frequency |
| `static` | N/A | N/A |
| `average` | average value is measured over that time window | average value measurement is repeated with that frequency |
| `rms` | root-mean-square value is measured over that time window | root-mean-square measurement is repeated with that frequency |
| `peak` | peak value occurs at that time (only within context of `VECTOR`) | observation of peak value is repeated with that frequency |

In the case of `average` and `rms`, the interpretation `FREQUENCY = 1 / TIME` is valid. Either one of these annotations shall be mandatory. The values for `average` measurements and for `rms` measurements scale linearly with `FREQUENCY` and `1 / TIME`, respectively.

In the case of `transient` and `peak`, the interpretation `FREQUENCY = 1 / TIME` is not valid. Either one of these annotations shall be optional. The values do not necessarily scale with `TIME` or `FREQUENCY`. The `TIME` or `FREQUENCY` annotations for `transient` measurements are purely informational.

### 8.9.3 TIME to peak measurement

For a model in the context of a `VECTOR`, with a `peak` measurement, the `TIME` annotation shall define the time between a reference event within the `vector_expression` and the instant when the peak value occurs.

For that purpose, either the `FROM` or the `TO` statement shall be used in the context of the `TIME` annotation, containing a `PIN` annotation and, if necessary, a `THRESHOLD` and/or an `EDGE_NUMBER` annotation.

If the `FROM` statement is used, the start point shall be the reference event and the end point shall be the occurrence time of the peak, as shown in Figure 8-13.

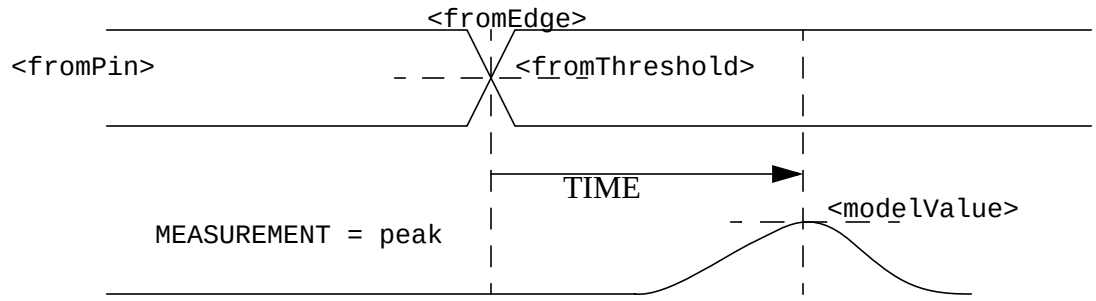



**Figure 8-13: Illustration of time to peak using FROM statement**

If the `TO` statement is used, the start point shall be the occurrence time of the peak and the end point shall be the reference event, as shown in Figure 8-14.

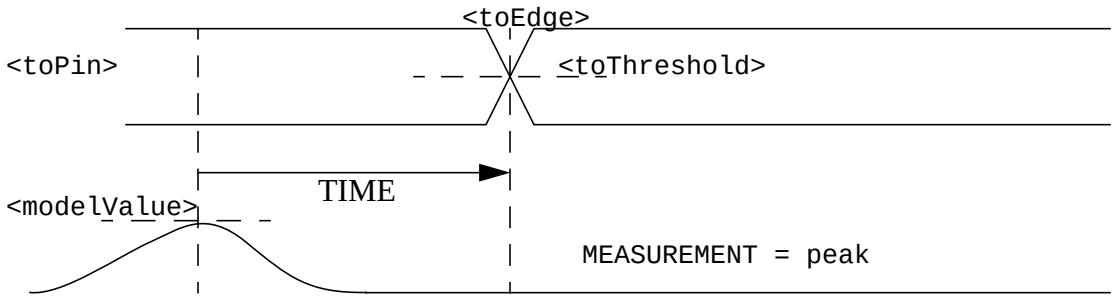



**Figure 8-14: Illustration of time to peak using TO statement**

Example:

```
VECTOR (01 A -> 01 B -> 10 B) {
   CURRENT peak1 = 10.8 {
      PIN = Vdd;
      MEASUREMENT = peak;
      TIME = 3.0 { UNIT=ns; FROM { PIN=A; EDGE_NUMBER=0; } }
   }
   CURRENT peak2 = 12.3 {
      PIN = Vdd;
      MEASUREMENT = peak;
      TIME = 2.0 { UNIT=ns; TO { PIN=B; EDGE_NUMBER=1; } }
   }
}
```

Here, the peak with magnitude `10.8` occurs `3` nanoseconds after the event `01 A`.

The peak with magnitude `12.3` occurs `2` nanoseconds before the event `10 B`.

### 8.9.4 Rules for combinations of annotations

Cumulative values of arithmetic models can be calculated for models which are cumulative in nature (e.g., ENERGY or POWER) or by the usage of CALCULATION=incremental (e.g., CURRENT or VOLTAGE). The MEASUREMENT annotation can be used in conjunction with the calculation of cumulative values under the following restrictions:

- Data with MEASUREMENT=average for each model can be combined, provided the TIME annotation value is the same.
- Data with MEASUREMENT=peak for each model can be combined, provided the TIME annotation or a complementary TIME model within the same context specify that the peak values can occur at the same time.
- Data with MEASUREMENT=rms for each model can not be combined.
- Data with different MEASUREMENT annotations can not be combined.
- Data with MEASUREMENT=transient|static can be combined with each other.

All data that can be combined under the abovementioned restrictions, must be in a compatible context, e.g., mutually non-exclusive VECTORs within a CELL.

## 8.10 Waveform description

This section specifies waveform descriptions.

### 8.10.1 Principles

In order to describe an arithmetic model representing a waveform, TIME shall be an argument in the HEADER. Other arguments can appear in the HEADER as well. The model can be described as a TABLE or EQUATION.

Example for TABLE:

```
VOLTAGE {
   HEADER {
      TIME {
         UNIT = ns;
         INTERPOLATION=linear;
         TABLE { 0.0  1.0  1.5  2.0  3.0 }
      }
   }
   TABLE { 0.0  0.0  5.0  0.0  0.0 }
}
```

Example for EQUATION:

```
VOLTAGE {
   HEADER {
      TIME { UNIT = ns; }
   }
```

```
      EQUATION {
         (TIME < 1.0) ? 0 :
         (TIME < 1.5) ? 5.0*(TIME - 1.0) :
         (TIME < 2.0) ? 5.0*(2.0 - TIME) :
         0.0
      }
   }
```

Both models describe the same piece-wise linear waveform, as shown in Figure 8-15.

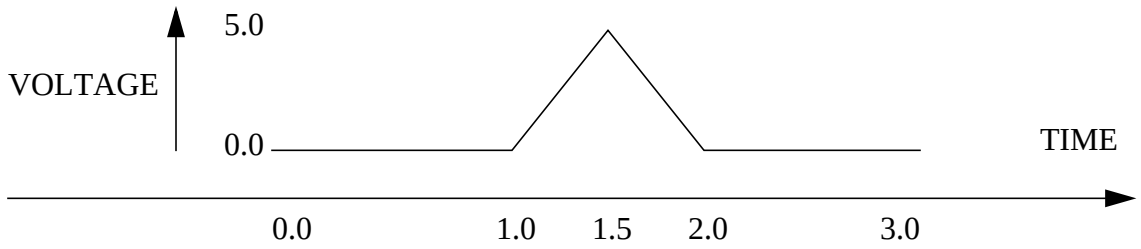



**Figure 8-15: Illustration of a piece-wise linear waveform**

If the model is within the context of a `VECTOR`, either the `FROM` or the `TO` statement can be used in the context of `TIME`, pointing to a reference event which occurs at `TIME = 0` relative to the waveform description. See Section 8.9 for the definition of start and end points of measurements.

Example:

```
VECTOR (01 A -> 01 B -> 10 B) {
   VOLTAGE {
      HEADER {
         TIME {
            FROM { PIN = B; EDGE_NUMBER = 1; }
            TABLE { 0.0  1.0  1.5  2.0  3.0 }
// alternative description:
//          TO { PIN = B; EDGE_NUMBER = 1; }
//          TABLE { -3.0  -2.0  -1.5  -1.0  0.0 }
         }
      }
      TABLE { 0.0  0.0  5.0  0.0  0.0 }
   }
}
```

Note: Use the `FROM` statement. If the `TO` statement is used, `TIME` is measured backwards, which is counter-intuitive. For dynamic analysis, use the last event in the `vector_expression` as the reference. Otherwise, the analysis tool remembers the occurrence time of previous events in order to place the waveform into the context of absolute time.

### 8.10.2 Annotations within a waveform

The `MEASUREMENT` annotation `transient` shall apply as a default for waveforms.

The `FREQUENCY` annotation can be used to specify a repetition frequency of the waveform. The following boundary restrictions are imposed in order to make the waveform repeatable:

- The initial value and the final value of waveform shall be the same.
- The extrapolation beyond the initial and the final value of the waveform shall yield the same result. Thus, the first, second, last, and second-to-last point of the waveform shall be the same.
- The time window between the first and the last measurement shall be smaller or equal to `1 / FREQUENCY`.

This is illustrated in Figure 8-16.

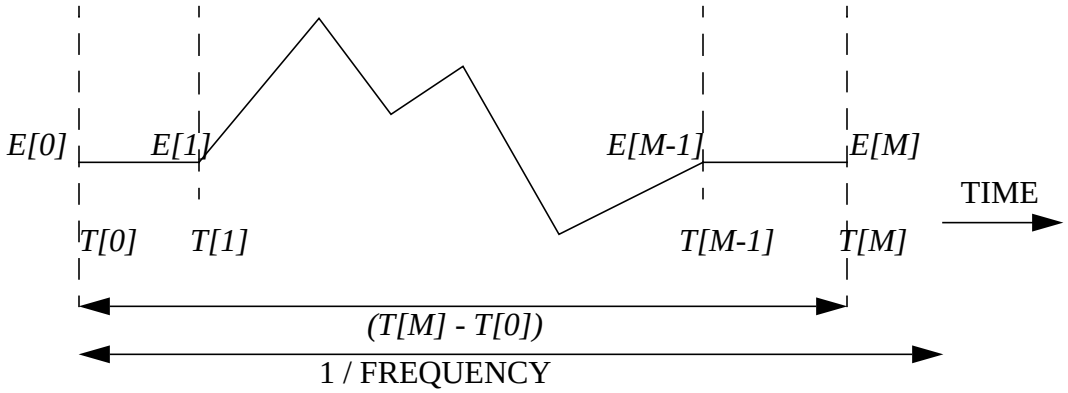



**Figure 8-16: TIME and FREQUENCY in a waveform**

## 8.11 Arithmetic models for power calculation

This section defines the arithmetic models used for power calculation.

### 8.11.1 Principles

The purpose of power calculation is to evaluate the electrical power supply demand and electrical power dissipation of an electronic circuit. In general, both power supply demand and power dissipation are the same, due to the energy conservation law. However, there are scenarios where power is supplied and dissipated locally in different places. The power models in ALF shall be specified in such a way that the total power supply and dissipation of a circuit adds up correctly to the same number.

Example: A capacitor `C` is charged from `0` volt to `V` volt by a switched DC source. The energy supplied by the source is $C*V^2$. The energy stored in the capacitor is $1/2*C*V^2$. Hence the

dissipated energy is also `1/2*C*V`$^2$. Later the capacitor is discharged from `V` volt to `0` volt. The supplied energy is `0`. The dissipated energy is `1/2*C*V`$^2$. A supply-oriented power model can associate the energy $E_1$=`C*V`$^2$ with the charging event and $E_2$=`0` with the discharging event. The total energy is $E=E_1+E_2$=`C*V`$^2$. A dissipation-oriented power model can associate the energy $E_3$=`1/2*C*V`$^2$ with both the charging and discharging event. The total energy is also $E=2*E_3$=`C*V`$^2$.

In many cases, it is not so easy to decide when and where the power is supplied and where it is dissipated. The choice between a supply-oriented and dissipation-oriented model or a mixture of both is subjective. Hence the ALF language provides no means to specify, which modeling approach is used. The choice is up to the model developer, as long as the energy conservation law is respected.

### 8.11.2 POWER and ENERGY

`POWER` and/or `ENERGY` models shall be in the context of a `CELL` or within a `VECTOR`. The total energy and/or power of a cell shall be calculated by combining the data of all models within the scope of the `CELL` or the `VECTOR`s within the cell.

The data for `POWER` and/or `ENERGY` shall be positive when energy is actually supplied to the `CELL` and/or dissipated within the `CELL`. The data shall be negative when energy is actually supplied or restored by the `CELL`.

Table 8-19 shows the mathematical relationship between ENERGY and POWER and the applicable MEASUREMENT annotations.

**Table 8-19 Relations between ENERGY and POWER**

| MEASUREMENT for ENERGY | MEASUREMENT for POWER | Formula to calculate POWER from ENERGY | Formula to calculate ENERGY from POWER |
|---|---|---|---|
| `transient` | `transient` | $\frac{d}{dt}\text{ENERGY}$ | $\int \text{POWER}\,dt$ |
| `transient` | `average` | $\frac{\text{ENERGY}}{\text{TIME}}$ | $\text{POWER} \cdot \text{TIME}$ |
| `transient` | `peak` | $\max\left(\left\lvert\frac{d}{dt}\text{ENERGY}\right\rvert\right)$ | N/A |
| `transient` | `rms` | $\frac{1}{\text{TIME}} \cdot \int\left(\frac{d}{dt}\text{ENERGY}\right)^2 dt$ | N/A |

**Table 8-19 Relations between ENERGY and POWER, *continued***

| MEASUREMENT for ENERGY | MEASUREMENT for POWER | Formula to calculate POWER from ENERGY | Formula to calculate ENERGY from POWER |
|---|---|---|---|
| N/A | `static` | N/A | POWER · TIME |
| `static` | N/A | 0 | N/A |

To establish a meaningful relationship between energy and power, the measurement for energy shall be `transient`. A `static` measurement for energy is conceivable, modeling a state with constant energy, but no power is dissipated during such a state. A `static` measurement for power models a state during which constant power dissipation occurs. Although it is not meaningful to describe an energy model for such a state, it is conceivable to calculate the energy by multiplying the power with the duration of the state. A 1-to-1 correspondence between power and energy can be established for `transient` and `average` power measurements, modeling instantaneous and average power, respectively. Therefore, it is redundant to specify both energy and power in such case. Also, `peak` and `rms` power can be conceivably calculated from a transient energy or power waveform, but transient energy can not be calculated from a `peak` or `rms` power measurement.

# 8.12 Arithmetic models for hot electron calculation

This section defines arithmetic models for hot electron calculation.

## 8.12.1 Principles

The purpose of hot electron calculation is to evaluate the damage done to the performance of an electronic device due to the hot electron effect. The hot electron effect consists in accumulation of electrons trapped in the gate oxide of a transistor. The more electrons are trapped, the more the device slows down. At a certain point, the performance specification no longer is met and the device is considered to be damaged.

## 8.12.2 FLUX and FLUENCE

`FLUX` and/or `FLUENCE` models shall be in the context of a `CELL` or within a `VECTOR`. Total fluence and/or flux of a cell shall be calculated by combining the data of all models within the scope of the `CELL` or the `VECTOR`s within the cell.

Both `FLUX` and `FLUENCE` are measures for hot electron damage. `FLUX` relates to `FLUENCE` in the same way as `POWER` relates to `ENERGY`.

Table 8-20 shows the mathematical relationship between FLUENCE and FLUX and the applicable MEASUREMENT annotations.

**Table 8-20 Relations between FLUENCE and FLUX**

| MEASUREMENT for FLUENCE | MEASUREMENT for FLUX | Formula to calculate FLUX from FLUENCE | Formula to calculate FLUENCE from FLUX |
|---|---|---|---|
| transient | transient | $\frac{d}{dt}\text{FLUENCE}$ | $\int \text{FLUX}\, dt$ |
| transient | average | $\frac{\text{FLUENCE}}{\text{TIME}}$ | FLUX · TIME |
| N/A | static | N/A | FLUX · TIME |
| static | N/A | 0 | N/A |

Since hot electron damage is purely cumulative, the only meaningful MEASUREMENT annotations are transient, average, and static.

## 8.13 Reliability calculation

In general, reliability is modeled by arithmetic models using the LIMIT construct.

### 8.13.1 TIME within the LIMIT construct

Within a LIMIT construct, TIME can be used in the following ways:

1. TIME itself is subjected to a LIMIT (see Section 8.14.2)
2. TIME is the argument of a model subjected to a LIMIT

When TIME is used as argument of a model within the LIMIT construct, it shall mean the amount of time during which the device is exposed to the quantity modeled within the LIMIT construct. This amount of time is also called a *lifetime*.

Example:

```
LIMIT {
   CURRENT {
      PIN = my_pin;
      MEASUREMENT = static;
```

```
      MAX {
         HEADER { TIME TEMPERATURE }
         EQUATION { 6.5*EXP(-10/(TEMPERATURE+273))*TIME**(-0.3) }
      }
   }
}
```

The limit for maximum current depends on the temperature and the expected lifetime of the device.

### 8.13.2 FREQUENCY within a LIMIT construct

Within a `LIMIT` construct, `FREQUENCY` can be used in the following ways:

1. `FREQUENCY` itself is subjected to a `LIMIT`
2. `FREQUENCY` is the argument of a model subjected to a `LIMIT`

`FREQUENCY` can be subjected to a `LIMIT` within the context of a `VECTOR`. The `LIMIT` construct specifies an upper and/or lower limit for the repetition frequency of the event sequence described by the `vector_expression`.

Example:

```
VECTOR ( 01 A -> 01 Z ) {
   LIMIT {
      FREQUENCY {
         MAX {
            HEADER {
               SLEWRATE { PIN = A; TABLE { 0.1 0.5 1.0 5.0 } }
               CAPACITANCE { PIN = Z; TABLE { 0.1 0.4 1.6 } }
            }
            TABLE {
               200 190 180 120
               150 150 145 130
                80  80  80  70
            }
         }
      }
   }
}
```

The maximum allowed switching frequency for a rising edge on `A`, followed by a rising edge on `Z`, depends on the slewrate on `A` and the load capacitance on `Z`.

A `LIMIT` for a quantity with `MEASUREMENT` annotation `average`, `rms`, or `peak` can be frequency-dependent. The `FREQUENCY` specifies the repetition frequency for the measurement.

Example:

```
LIMIT {
   CURRENT {
      PIN = Vdd;
      MEASUREMENT = average;
```

```
         MAX {
            HEADER { FREQUENCY TIME TEMPERATURE }
            EQUATION {
               (FREQUENCY<1)? 6.5*EXP(-10/(TEMPERATURE+273))*TIME**(-0.3) :
               7.8*EXP(-9/(TEMPERATURE+273))*TIME**(-0.2) :
            }
         }
      }
   }
```

The limit for average current is specified for low frequencies ( < 1MHz) and for higher frequencies. In both cases, the limit depends on temperature and lifetime.

### 8.13.3 Global LIMIT specifications

Global limits can be specified for electrical quantities, even if they are related to CELLs, PINs, or VECTORs. Such global limits apply, unless local limits are specified within the context of CELLs, PINs, or VECTORs. The priorities are given below.

1. LIMIT within the context of the VECTOR
2. LIMIT within the context of a PIN (if the LIMIT in the VECTOR has PIN annotation)
3. LIMIT within the context of the CELL
4. LIMIT within the context of the SUBLIBRARY
5. LIMIT within the context of the LIBRARY
6. LIMIT outside LIBRARY

The arguments in the HEADER of the LIMIT model can only be items that are visible within the scope of the LIMIT model. In particular, arguments with PIN annotations are only legal for LIMIT models in the context of a CELL or a VECTOR within the CELL.

### 8.13.4 LIMIT specification and model specification in the same context

An arithmetic model for a physical quantity and a limit specification for the same physical quantity can appear within the same context, for example, an arithmetic model for FLUENCE calculation and a LIMIT for FLUENCE within the context of a VECTOR. In such a case, the calculated quantity shall be checked against the limit of the quantity within that context.

On the other hand, if multiple arithmetic models are given within the context for which the limit applies, the limit shall be checked against the combination of all arithmetic models in the case of cumulative quantities, or against the minimum or maximum calculated value in the case of non-cumulative or mutually exclusive quantities.

For example, a LIMIT for FLUENCE can be given in the context of a CELL. Calculation models for FLUENCE can be given for multiple VECTORs within the context of the CELL. The LIMIT for FLUENCE shall be checked against the accumulated FLUENCE calculated for all VECTORs.

Example:

```
CELL my_cell {
   PIN A { DIRECTION = input; }
   PIN B { DIRECTION = input; }
   PIN C { DIRECTION = input; }
   PIN Z { DIRECTION = output; }
   LIMIT { FLUENCE { MAX = 1e20; } } }
   VECTOR ( 01 A -> 10 Z ) {
      FLUENCE = 1e-5;
   }
   VECTOR ( 01 B -> 10 Z ) {
      FLUENCE = 1e-5;
   }
   VECTOR ( 01 C -> 10 Z ) {
      FLUENCE = 1e-6;
      LIMIT { FLUENCE { MAX = 1e18; } }
   }
}
```

The fluence limit for the cell is reached after $10^{25}$ occurrences of `VECTOR ( 01 A -> 10 Z )` or `VECTOR ( 01 B -> 10 Z )` counted together.
The fluence limit for the `VECTOR ( 01 C -> 10 Z )` is reached after $10^{24}$ occurrences of that vector.

An example for a non-cumulative quantity is `SLEWRATE`. The `VECTOR`s in the context of which `SLEWRATE` is modeled describe timing arcs with mutually exclusive conditions. Therefore, if a minimum or maximum `LIMIT` for `SLEWRATE` is given for a `PIN` in the context of a `CELL`, this `SLEWRATE` shall be checked against the minimum or maximum value of any calculated `SLEWRATE` applicable to that `PIN`.

Example:

```
CELL my_cell {
   PIN A { DIRECTION = input; }
   PIN B { DIRECTION = input; }
   PIN C { DIRECTION = input; }
   PIN Z { DIRECTION = output; LIMIT { SLEWRATE { MAX = 5; } } }
   VECTOR ( 01 A -> 10 Z ) {
      SLEWRATE { PIN = Z; /* fill in HEADER, TABLE */ }
   }
   VECTOR ( 01 B -> 10 Z ) {
      SLEWRATE { PIN = Z; /* fill in HEADER, TABLE */ }
   }
   VECTOR ( 01 C -> 10 Z ) {
      SLEWRATE { PIN = Z; /* fill in HEADER, TABLE */ }
   }
}
```

Here the slewrate on pin `Z` calculated in the context of any vector is checked against the same maximum limit.

### 8.13.5 Model specification and argument specification in the same context

An cumulative quantity can also be an argument in the HEADER of an arithmetic model. If the model for calculation of that quantity is within the same context as the argument of the other model, then the value of the calculated quantity shall be used. Otherwise, the value of the accumulated quantity shall be used.

For example, SLEWRATE can be modeled as a function of FLUENCE in the context of a VECTOR. If a calculation model for FLUENCE appears in the context of the same VECTOR, the value for FLUENCE shall be used for the SLEWRATE calculation. On the other hand, if there is no calculation model for FLUENCE in the context of the same VECTOR, but there is one in the context of other VECTORs, then the accumulated value of FLUENCE from the other calculation models shall be used for SLEWRATE calculation.

Example:

```
CELL my_cell {
   PIN A { DIRECTION = input; }
   PIN B { DIRECTION = input; }
   PIN C { DIRECTION = input; }
   PIN Z { DIRECTION = output; }
   VECTOR ( (01 A | 01 B) -> 10 Z ) { FLUENCE = 1e-5; }
   VECTOR ( 01 A -> 10 Z ) {
      SLEWRATE { CALCULATION=incremental; PIN = Z;
         HEADER { FLUENCE } EQUATION { 1e-8 * FLUENCE }
      }
   }
   VECTOR ( 01 B -> 10 Z ) {
      SLEWRATE { CALCULATION=incremental; PIN = Z;
         HEADER { FLUENCE } EQUATION { 1e-8 * FLUENCE }
      }
   }
   VECTOR ( 01 C -> 10 Z ) {
      FLUENCE = 1e-6;
      SLEWRATE { CALCULATION=incremental; PIN = Z;
         HEADER { FLUENCE } EQUATION { 1e-9 * FLUENCE }
      }
   }
}
```

After $10^{13} = 10^5 * 10^8$ occurrences of VECTOR ( (01 A | 01 B) -> 10 Z ), the slewrate at pin Z for VECTOR ( 01 A -> 10 Z ) and VECTOR ( 01 B -> 10 Z ) is increased by 1 unit. After $10^{15} = 10^6 * 10^9$ occurrences of VECTOR ( 01 C -> 10 Z ), the slewrate at pin Z for VECTOR ( 01 C -> 10 Z ) is increased by 1 unit.

## 8.14 Noise calculation

This section details the noise calculation definitions.

### 8.14.1 NOISE_MARGIN definition

*Noise margin* is defined as the maximal allowed difference between the ideal signal voltage under a well-specified operation condition and the actual signal voltage normalized to the ideal voltage swing. This is illustrated in Figure 8-17.

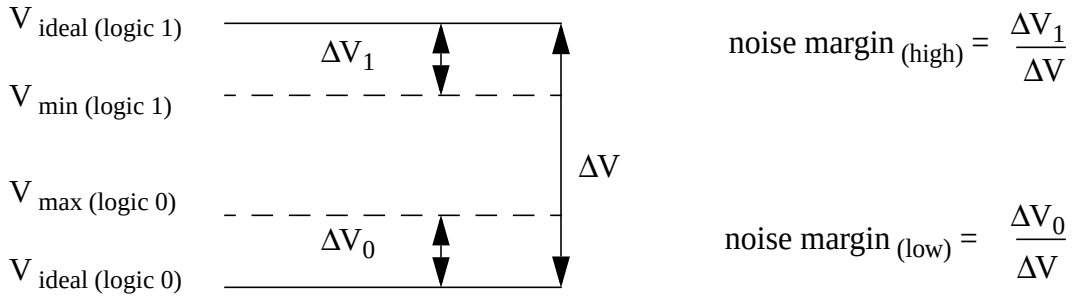



**Figure 8-17: Definition of noise margin**

Noise margin is measured at a signal input pin of a digital cell. The terms *ideal signal voltage* and *actual signal voltage* apply from the standpoint of that particular pin. In CMOS technology, the ideal signal voltage at a pin is the actual supply voltage of the cell, which is not necessarily identical to the nominal supply voltage of the chip.

The `NOISE_MARGIN` statement has the form of an arithmetic model. If the submodel keywords `HIGH` and `LOW` are used, it has the form of an arithmetic model container.

Examples:

```
NOISE_MARGIN = 0.3;

NOISE_MARGIN { HIGH = 0.2;  LOW = 0.4; }

NOISE_MARGIN {
   HEADER { TEMPERATURE { TABLE { 0 50 100 } } }
   TABLE { 0.4 0.3 0.2 }
}
```

`NOISE_MARGIN` can be related to signal `VOLTAGE` by using the following statement:

```
   VOLTAGE {
      LOW = 0;
      HIGH = 2.5;
   }
   NOISE_MARGIN {
      LOW = 0.4;
      HIGH = 0.3;
   }
}
```

In this example, the valid signal voltage levels are bound by 1 volt = `2.5 volt * 0.4` for logic `0` and 1.75 volt = `2.5 volt * (1 - 0.3)` for logic `1`.

### 8.14.2 Representation of noise in a VECTOR

In order to describe timing diagrams involving noisy signals, the symbolic state "*" (see Section 5.4.13) shall be used. This state represents arbitrary transitions between arbitrary states, which corresponds to the nature of noise, as shown in Figure 8-18.

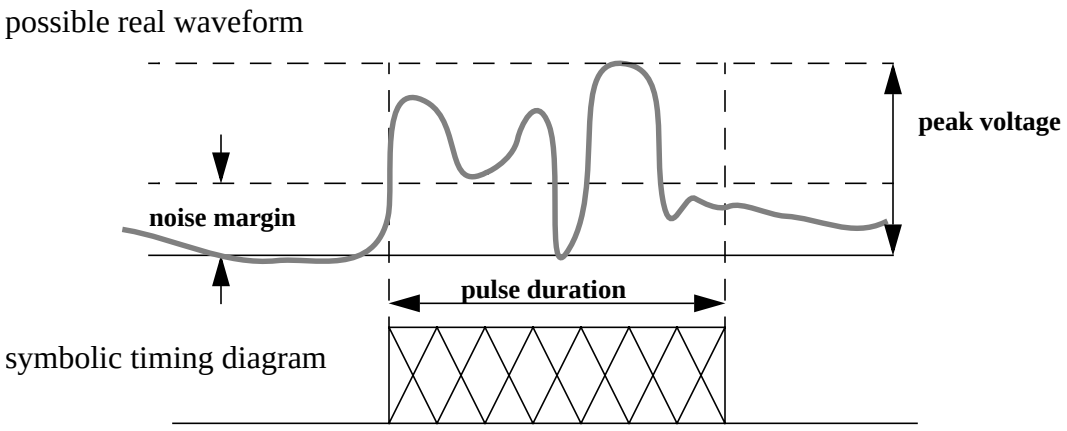



**Figure 8-18: Timing diagram of a noisy signal**

The signal can be above or below noise margin during the state "*", but it shall be within noise margin during the state `0` or `1`. During the state "*", the signal is bound by an envelope defined by the pulse duration and the peak voltage.

A description of the noisy signal is given in the following template:

```
VECTOR ( 0* my_pin -> *0 my_pin ) {
   TIME = <pulse_duration> {
      FROM { PIN=my_pin; EDGE_NUMBER=0; }
      TO   { PIN=my_pin; EDGE_NUMBER=1; }
   }
   VOLTAGE = <peak_voltage> {
      CALCULATION = incremental;
      MEASUREMENT = peak;
      PIN = my_pin;
   }
}
```

The `VECTOR` describes the symbolic timing diagram. The `TIME` statement specifies the duration of the pulse. The `VOLTAGE` statement specifies the peak voltage. The annotation `CALCULATION=incremental` specifies that the voltage is measured from the nominal signal voltage level rather than from an absolute reference level and that noise voltage can add up.

It is also necessary to specify whether a noisy signal (which can oscillate above and below the noise margin) is considered as one symbolic noise pulse or separated into multiple symbolic noise pulses.

The LIMIT statement for TIME shall be used for that purpose, as shown in the following example and illustrated by the timing diagram shown in Figure 8-19.

```
VECTOR ( *0 my_pin -> 0* my_pin ) {
   LIMIT {
      TIME {
         FROM { PIN = my_pin; EDGE_NUMBER = 0; }
         TO   { PIN = my_pin; EDGE_NUMBER = 1; }
         MIN = <minimum_pulse_separation> ;
      }
   }
}
```

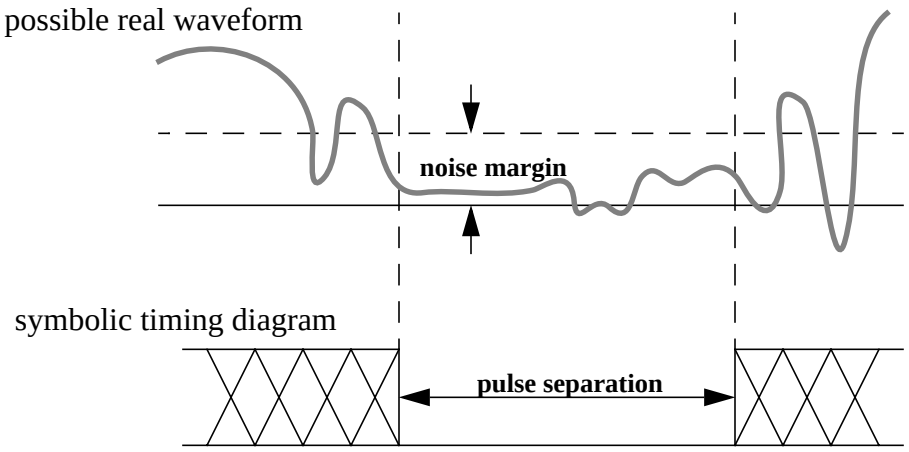



**Figure 8-19: Separation between two noise pulses**

When the minimum pulse separation is not met, consecutive noise pulses shall be symbolically merged into one pulse.

### 8.14.3 Context of NOISE_MARGIN

NOISE_MARGIN is a pin-related quantity. It can appear either in the context of a PIN statement or in the context of a VECTOR statement with PIN annotation. It can also appear in the global context of a CELL, SUBLIBRARY, or LIBRARY statement.

If a NOISE_MARGIN statement appears in multiple contexts, the following priorities apply:

1. NOISE_MARGIN with PIN annotation in the context of the VECTOR, NOISE_MARGIN with PIN annotation in the context of the CELL, or NOISE_MARGIN in the context of the PIN
2. NOISE_MARGIN without PIN annotation in the context of the CELL
3. NOISE_MARGIN in the context of the SUBLIBRARY
4. NOISE_MARGIN in the context of the LIBRARY

5. NOISE_MARGIN outside the LIBRARY

If the noise margin is constant or depends only on environmental quantities, the NOISE_MARGIN statement shall appear within the context of the PIN. The noise margin shall relate to the signal VOLTAGE levels applicable for that pin.

Example:

```
PIN my_signal_pin {
   PINTYPE = digital;
   DIRECTION = input;
   VOLTAGE { LOW = 0; HIGH = 2.5; }
   NOISE_MARGIN { LOW = 0.4; HIGH = 0.3; }
}
```

If the noise margin depends on electrical quantities related to other pins, e.g., the supply voltage, the NOISE_MARGIN statement shall have a PIN annotation and appear in the context of the CELL.

Example:

```
CELL my_cell {
   PIN my_signal_pin { PINTYPE = digital; DIRECTION = input; }
   PIN my_power_pin  { PINTYPE = supply; SUPPLYTYPE = power; }
   PIN my_ground_pin { PINTYPE = supply; SUPPLYTYPE = ground; }
   NOISE_MARGIN {
      PIN = my_signal_pin;
      HEADER {
         VOLTAGE vdd { PIN = my_power_pin;  }
         VOLTAGE vss { PIN = my_ground_pin; }
      }
      EQUATION { 0.16 * (vdd - vss ) }
   }
}
```

If the noise margin depends on the logical states and/or the timing of other pins, the NOISE_MARGIN statement shall have a PIN annotation and appear in the context of a VECTOR, describing the state-and/or timing dependency.

Example for state-dependent noise margin:

```
CELL my_latch {
   PIN Q { DIRECTION = output; SIGNALTYPE = data; }
   PIN D { DIRECTION = input; SIGNALTYPE = data; }
   PIN CLK { DIRECTION = input; SIGNALTYPE = clock; POLARITY = high; }
   VECTOR ( CLK && ! D ) { NOISE_MARGIN = 0.4 { PIN = D; } }
   VECTOR ( CLK && D ) { NOISE_MARGIN = 0.3 { PIN = D; } }
}
```

Here, the pin D is only noise-sensitive when CLK is high. No noise margin is given for the case when CLK is low.

In the case of timing-dependency, the vector_expression shall indicate the time window where noise is allowed and not allowed for the applicable pin. The symbolic state * (see Section 5.4.13) shall be used to indicate a noisy signal.

Example for timing-dependent noise margin:

```
VECTOR ( *? D -> 10 CLK -> ?* D ) {
   TIME T1 = 0.35 {
      FROM { PIN = D;   EDGE_NUMBER = 0; }
      TO   { PIN = CLK; EDGE_NUMBER = 0; }
   }
   TIME T2 = 0.28 {
      FROM { PIN = CLK; EDGE_NUMBER = 0; }
      TO   { PIN = D;   EDGE_NUMBER = 1; }
   }
   NOISE_MARGIN = 0.44 { PIN = D; }
}
```

This example corresponds to the timing diagram shown in Figure 8-20.

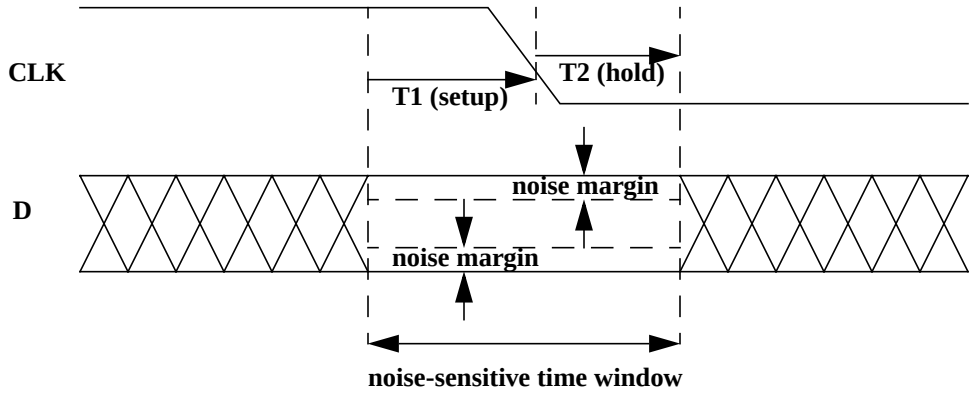



**Figure 8-20: Example for timing-dependent noise margin**

Noise on pin `D` is allowed `0.35` time-units before and `0.28` time-units after the falling edge of `CLK`. During the time window in-between, the noise margin is `0.44`.

### 8.14.4 Noise propagation

*Noise propagation* from input to output can be modeled in a similar way as signal propagation, using the concept of timing arcs. This is illustrated in Figure 8-21.

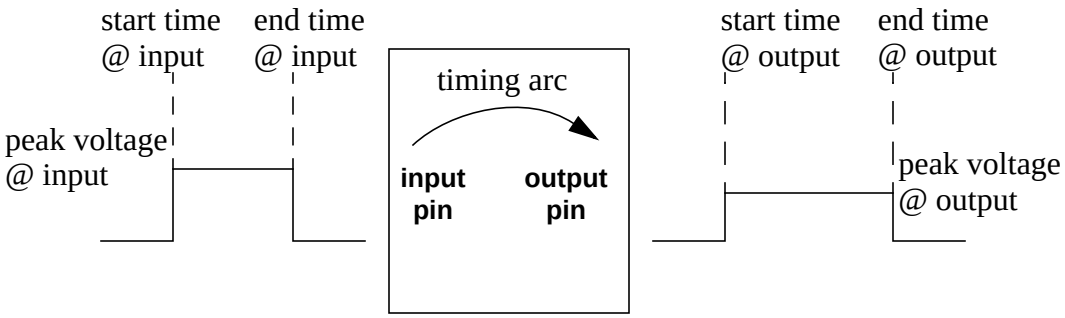



**Figure 8-21: Principle of noise propagation**

The principle of *signal propagation* is to calculate the output arrival time and slewrate from the input arrival time and slewrate. In a more abstract way, two points in time propagate from input to output. The same principle applies for noise propagation. Two points in time, start and end time of the noise waveform, propagate from input to output. In addition, the noise peak voltage also propagates from input to output. This is illustrated in Figure 8-22.

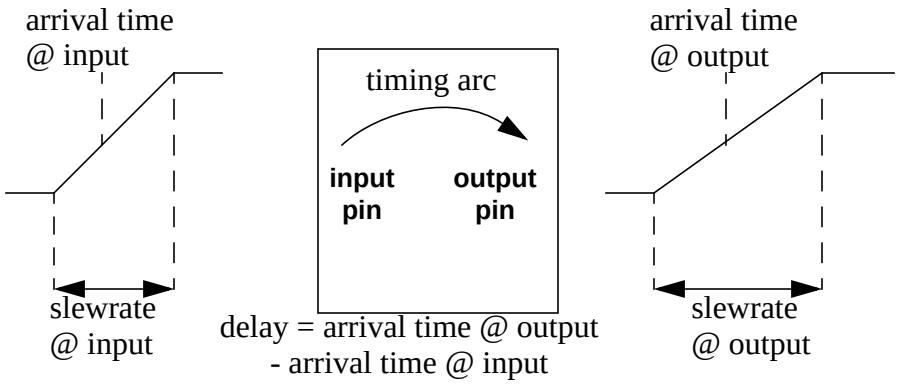



**Figure 8-22: Principle of signal propagation**

A `VECTOR` shall be used to describe the timing of the noise waveform. Again, the symbolic state `*` (see Section 5.4.13) shall be used to indicate a noisy signal.

Example:

```
CELL my_cell {
   PIN A { DIRECTION = input; }
   PIN Z { DIRECTION = output; }
   VECTOR ( 0* A -> *0 A <&> 0* Z -> *0 Z ) {
      DELAY T1 {
         FROM { PIN = A; EDGE_NUMBER = 0; }
         TO   { PIN = Z; EDGE_NUMBER = 0; }
         /* fill in HEADER, TABLE or EQUATION */
      }
```

```
      DELAY T2 {
         FROM { PIN = A; EDGE_NUMBER = 1; }
         TO   { PIN = Z; EDGE_NUMBER = 1; }
         /* fill in HEADER, TABLE or EQUATION */
      }
      VOLTAGE { PIN = Z; MEASUREMENT = peak;
         /* fill in HEADER, TABLE or EQUATION */
      }
   }
```

This example corresponds to the timing diagram shown in Figure 8-23.

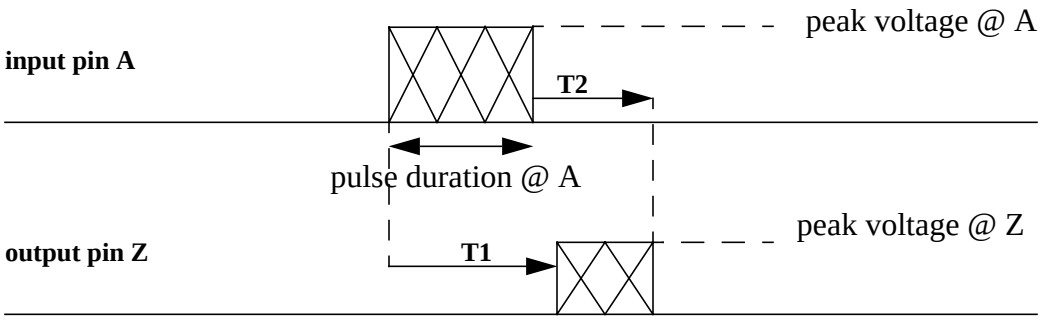



**Figure 8-23: Example of noise propagation**

The input to output delay of the leading edge of the noise pulse can depend on the peak voltage at pin A, the load capacitance at pin Z and other electrical quantities. In addition, the input to output delay of the trailing edge of the noise pulse as well as the peak voltage at pin Z can also depend on the duration of the pulse at pin A.

Note: The time measurement from start to end of the noise pulse shall be represented by the keyword TIME (no causality between start and end time), whereas the time measurement from input to output shall be represented by the keyword DELAY (causality between input and output arrival time).

### 8.14.5 Noise rejection

Noise rejection is a limit case for noise propagation, when the output peak voltage is so low the noise is considered rejected. In this case, the input peak voltage can still be above noise margin, whereas the output peak voltage is way below noise margin.

Example:

```
CELL my_cell {
   PIN A { DIRECTION = input; }
   PIN Z { DIRECTION = output; }
   VECTOR ( 0* A -> *0 A -> 00 Z ) {
```

```
        LIMIT {
           VOLTAGE {
              PIN = A; MEASUREMENT = peak;
              MAX { /* fill in HEADER, TABLE or EQUATION */ }
           }
        }
     }
  }
```

Note: The `vector_expression 00 Z` says explicitly a transition at pin `Z` does *not* happen.

This example corresponds to the timing diagram shown in Figure 8-24.

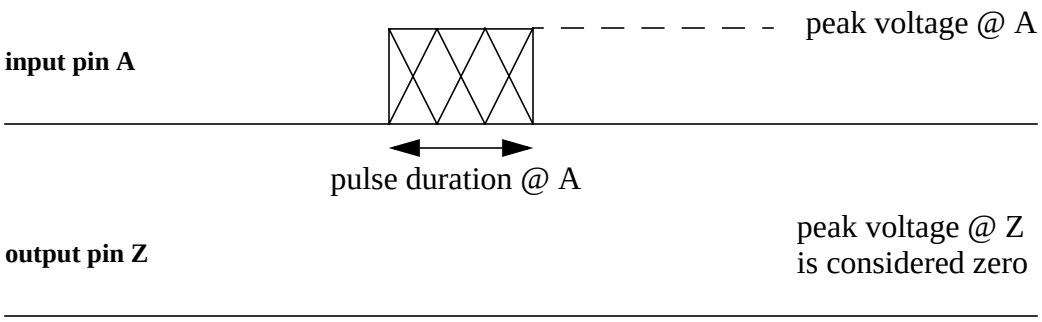



**Figure 8-24: Example of noise rejection**

The peak voltage limit for noise rejection can depend on the duration of the noise pulse at pin A and other electrical quantities, e.g., the load capacitance at pin `Z`. If the peak voltage limit does not depend on the duration of the noise pulse, the `NOISE_MARGIN` statement shall be used rather than the vector-specific `LIMIT` construct for noise rejection.

# 8.15 Interconnect parasitics and analysis

This section defines interconnect parasitics and analysis.

## 8.15.1 Principles of the WIRE statement

Parasitic descriptions shall be in the context of a `WIRE` statement. The following fundamental modeling styles are supported.

- Statistical wireload models
- Boundary parasitics

Statistical wireload models as well as interconnect analysis calculation models can be used within the context of a `LIBRARY`, `SUBLIBRARY`, or `CELL` statement. The latter applies only for cells with `CELLTYPE=block`, i.e., hierarchical cells. Boundary parasitics apply exclusively for hierarchical cells. Statistical wireload models can be mixed with boundary parasitics within the same `WIRE` statement.

Interconnect analysis models shall also be defined within a WIRE statement. However, they shall not be mixed with statistical wireload models or boundary parasitic descriptions.

The purpose of interconnect analysis is to calculate electrical quantities such as DELAY, SLEWRATE, and noise VOLTAGE in the context of a netlist consisting of electrical components, such as CAPACITANCE, RESISTANCE, and INDUCTANCE.

As opposed to boundary parasitics, where the components are connected to physical nodes and pins of a cell, the components represent an abstract network targeted for analysis. The interconnect analysis model specifies a directive for reducing the parasitic extraction/delay calculation tool to an arbitrary network. In addition, the model specifies the calculation models for delay, noise, etc. in the context of the reduced network.

### 8.15.2 Statistical wireload models

A statistical wireload model is a collection of arithmetic models for estimated the electrical quantities CAPACITANCE, RESISTANCE, and INDUCTANCE, representing the interconnect load and estimated AREA and SIZE of the interconnect nets.

These arithmetic models shall have no PIN annotation. Only environmental quantities such as PROCESS, DERATE_CASE, and TEMPERATURE shall be allowed as arguments in the HEADER.

In addition, the quantities AREA, SIZE, FANOUT, FANIN, and CONNECTIONS are allowed as arguments in the HEADER.

FANOUT and FANIN represent the number of receiver pins and driver pins, respectively, connected to the net. CONNECTIONS is the total number of pins connected to the net. CONNECTIONS equals to the sum of FANOUT and FANIN.

AREA represents a physically measurable area of an object, whereas SIZE represents an abstract symbolic quantity or cost function for area. When AREA or SIZE is used as argument within the HEADER, it shall represent the total area or size, respectively, allocated for place and route of the block for which the wireload model applies. An arithmetic model given for AREA or SIZE itself shall represent the estimated or actual area or size, respectively, of the object in the context of which the model appears. CELL and WIRE are applicable objects for AREA or SIZE models.

In order to convert SIZE to AREA (analogous to converting DRIVE_STRENGTH to RESISTANCE; see Section 8.8.1), an arithmetic model for SIZE with AREA as an argument can be used outside the WIRE statement. Arithmetic models for SIZE inside the WIRE statement shall be interpreted as a calculation model rather than a conversion model.

The total area or size of a block shall be larger or equal to the area or size, respectively, of all objects within the block, i.e., cells and wires.

Note: The area or size of a block is design-specific data, whereas the area or size of cells and wires is given in the library.

Example:

```
LIBRARY my_library {
   WIRE my_wlm {
      CAPACITANCE {
```

```
            HEADER {
               CONNECTIONS { TABLE { 2 3 4 5 10 20 } }
               AREA { TABLE { 1000 10000 100000 } }
            }
            TABLE {
               0.03 0.06 0.08 0.10 0.15 0.25
               0.05 0.10 0.15 0.18 0.25 0.35
               0.10 0.18 0.25 0.32 0.50 0.65
            }
         }
         AREA {
            HEADER {
               CONNECTIONS { TABLE { 2 3 4 5 10 20 } }
               AREA { TABLE { 1000 10000 100000 } }
            }
            TABLE {
               0.3 0.6 0.8 1.0 1.5 2.5
               0.5 1.0 1.5 1.8 2.5 3.5
               1.0 1.8 2.5 3.2 5.0 6.5
            }
         }
      }
      CELL my_cell {
         AREA = 1.5;
         PIN my_input  { DIRECTION = input;  CAPACITANCE = 0.1; }
         PIN my_output { DIRECTION = output; CAPACITANCE = 0.0; }
      }
   }
```

A net routed in a block of `AREA=10000`, driven by an instance of `my_cell` connecting to five receivers (i.e., `CONNECTIONS=5`), each of which is an instance of `my_cell`, shall have an estimated capacitance of 0.18+4*0.1 = `0.58` and wire area of `1.8`. The five cell instances together shall have an area of `7.5`.

Note: `CAPACITANCE`, `RESISTANCE`, and `AREA` can each be independent arithmetic models within the `WIRE` statement. No multiplication factor between area and capacitance or between area and resistance is assumed.

### 8.15.3 Boundary parasitics

Boundary parasitics for a `CELL` can be given within a `WIRE` statement in the context of the `CELL`. The parasitics shall be identified by arithmetic models for `CAPACITANCE`, `RESISTANCE`, and `INDUCTANCE` containing a `NODE` annotation.

The syntax is as follows:

```
two_node_multi_value_assignment ::=
          NODE { node_identifier node_identifier }

four_node_multi_value_assignment ::=
          NODE { node_identifier node_identifier node_identifier
          node_identifier }
```

where *node*_identifier is one of the following:

a simple identifier, referring to a declared PIN of the CELL.

a hierarchical_identifier, referring to a declared PORT of a PIN of the CELL (see Section 9.10.4)

a simple identifier, referring to a declared NODE of the WIRE (see Section 8.15.4)

a simple identifier, not referring to a declared object. This can be used for connectivity inside the WIRE only.

The *two_node*_multi_value_assignment applies for capacitance, resistance, and self-inductance. These components imply the following relationship between voltage and current across the nodes:

$$\text{VOLTAGE(node1, node2)} = \text{RESISTANCE(node1, node2)} \cdot \text{CURRENT(node1, node2)}$$

$$\text{CURRENT(node1, node2)} = \text{CAPACITANCE(node1, node2)} \cdot \frac{d}{dt}\text{VOLTAGE(node1, node2)}$$

$$\text{VOLTAGE(node1, node2)} = \text{INDUCTANCE(node1, node2)} \cdot \frac{d}{dt}\text{CURRENT(node1, node2)}$$

The *four_node*_multi_value_assignment applies for mutual inductance. This component implies the following relationship between voltage and current across the nodes:

$$\text{VOLTAGE(node1, node2)} = \text{INDUCTANCE(node1, node2, node3, node4)} \cdot \frac{d}{dt}\text{CURRENT(node3, node4)}$$

Note: Both PIN assignments (e.g., PIN=A;) and NODE assignments (e.g., NODE { A B }) can refer to PINs or PORTs. The fundamental semantic difference between a PIN assignment and a NODE assignment is the PIN assignment within an object defines the object is *applied* or *measured* at the PIN or PORT. (e.g., DELAY and SLEWRATE); the NODE assignment within an object defines the object is fundamentally *connected* with the PIN or PORT in the same way an object inside a PIN is also fundamentally connected with the PIN. Therefore, the CAPACITANCE with NODE assignment is a more detailed way of describing a CAPACITANCE of a PIN, whereas a CAPACITANCE with PIN assignment describes a load capacitance, which is applied externally to the pin.

A CELL can contain a WIRE statement describing boundary parasitics as well as PIN statements containing arithmetic models for CAPACITANCE, RESISTANCE, or INDUCTANCE. In this case the latter shall be considered as a reduced form of the former. An analysis tool shall either use the set of components inside the PIN or inside the WIRE, but not a combination of both.

Example:

```
CELL my_cell {
   PIN A { PINTYPE = digital; CAPACITANCE = 4.8; RESISTANCE = 37.9;
      PORT p1 { VIEW = physical; } // see Section 9.10
      PORT p2 { VIEW = none; } // see Section 9.10
   }
```

```
    PIN B { PINTYPE = digital; CAPACITANCE = 2.6; }
    PIN gnd { PINTYPE = supply; SUPPLYTYPE = ground; }
    WIRE my_boundary_parasitics {
       CAPACITANCE = 1.3 { NODE { A.p1 gnd } }
       CAPACITANCE = 2.8 { NODE { A.p2 gnd } }
       RESISTANCE = 65 { NODE { A.p1 A.p2 } }
       CAPACITANCE = 0.7 { NODE { A.p1 B } }
       CAPACITANCE = 1.9 { NODE { B gnd } }
    }
}
```

This example corresponds to the netlist shown in Figure 8-25.

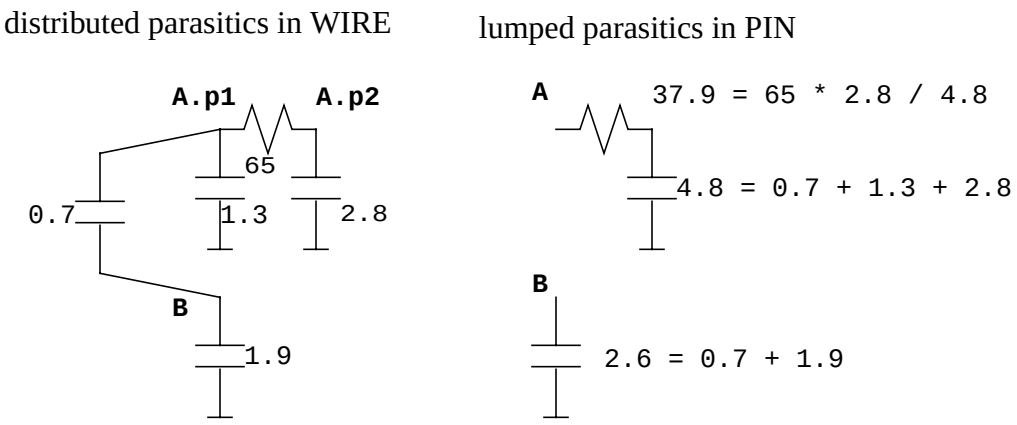



**Figure 8-25: Example of boundary parasitic description**

The distributed parasitics in the `WIRE` statement can be reduced to the lumped parasitics in the `PIN` statement.

### 8.15.4 NODE declaration

The nodes used for interconnect analysis shall be declared within the `WIRE` statement, using the following syntax.

```
node ::=
          NODE node_identifier { all_purpose_items }
```

The `NODETYPE` annotation and the `NODE_CLASS` annotation also specifically apply to a `NODE`.

```
nodetype_annotation ::=
          NODETYPE = nodetype_identifier ;

nodetype_identifier ::=
          ground
   |      power
   |      source
```

```
|     sink
|     driver
|     receiver
```

- A *driver node* is the interface between a cell output pin and interconnect
- A *receiver node* is the interface between interconnect and a cell input pin
- A *source node* is a virtual start point of signal propagation; it can be collapsed with a driver node
- A *sink node* is a virtual end point of signal propagation; it can be collapsed with a receiver node
- A *power node* provides the current for rising signals at the source/driver side and a reference for logic high signals at the sink/receiver side
- A *ground node* provides the current for falling signals at the source/driver side and a reference for logic low signals at the sink/receiver side

The arithmetic models for electrical components which are part of the network shall have names and `NODE` annotations, referring either to the pre-declared nodes or to internal nodes which need not be declared.

Example:

```
WIRE my_interconnect_model {
   NODE n0  { NODETYPE = source;   }
   NODE n2  { NODETYPE = driver;   }
   NODE n4  { NODETYPE = receiver; }
   NODE n5  { NODETYPE = sink;     }
   NODE vdd { NODETYPE = power;    }
   NODE vss { NODETYPE = ground;   }
   RESISTANCE R1 { NODE { n0 n1 } }
   RESISTANCE R2 { NODE { n1 n2 } }
   RESISTANCE R3 { NODE { n2 n3 } }
   RESISTANCE R4 { NODE { n3 n4 } }
   RESISTANCE R5 { NODE { n4 n5 } }
   CAPACITANCE C1 { NODE { n1 vss } }
   CAPACITANCE C2 { NODE { n2 vss } }
   CAPACITANCE C3 { NODE { n3 vss } }
   CAPACITANCE C4 { NODE { n4 vss } }
   CAPACITANCE C5 { NODE { n5 vss } }
}
```

This example is illustrated in Figure 8-26.

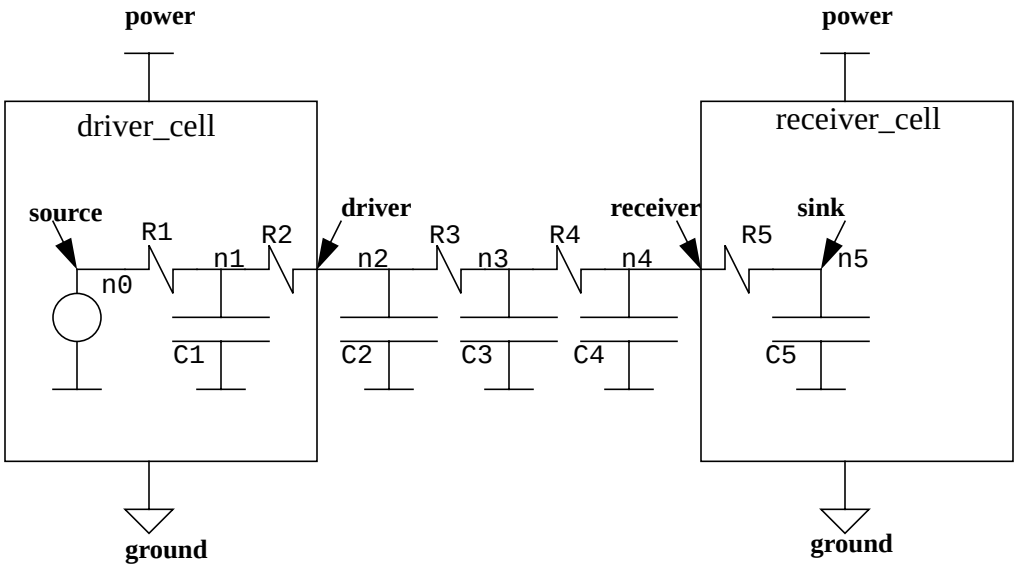



**Figure 8-26: Example for interconnect description**

The NODE_CLASS annotation is optional and orthogonal to the NODETYPE annotation.

*node_class*_annotation ::=
**NODE_CLASS =** *node_class*_identifier **;**

The NODE_CLASS annotation shall refer to a pre-declared CLASS within the WIRE statement to indicate which node belongs to which device in the case of separate power supplies.

Example:

```
WIRE my_interconnect_model {
   CLASS driver_cell;
   CLASS receiver_cell;
   NODE n0   { NODETYPE = source; NODE_CLASS = driver_cell; }
   NODE n2   { NODETYPE = driver; NODE_CLASS = driver_cell; }
   NODE n4   { NODETYPE = receiver; NODE_CLASS = receiver_cell; }
   NODE n5   { NODETYPE = sink;     NODE_CLASS = receiver_cell; }

   NODE vdd1 { NODETYPE = power;  NODE_CLASS = driver_cell; }
   NODE vss1 { NODETYPE = ground; NODE_CLASS = driver_cell; }
   NODE vdd2 { NODETYPE = power;  NODE_CLASS = receiver_cell; }
   NODE vss2 { NODETYPE = ground; NODE_CLASS = receiver_cell; }
}
```

If NODE_CLASS is not specified, the nodes with NODETYPE=power|ground are supposed to be global. The DC-connected nodes with NODETYPE=driver|source and NODETYPE=receiver|sink are supposed to belong to the same device.

### 8.15.5 Interconnect delay and noise calculation

Calculation models for DELAY and SLEWRATE can be described in the context of a VECTOR inside a WIRE. The PIN assignments in these models shall refer to pre-declared NODEs inside the WIRE.

Example:

```
WIRE my_interconnect_model {
   /* node declarations */
   /* electrical component declarations */
   VECTOR ( (01 n0 ~> 01 n5) | (10 n0 ~> 10 n5) ) {
      /* DELAY model */
      /* SLEWRATE model */
   }
}
```

The pre-declared electrical components which are part of the network can be used within an EQUATION without being re-declared in the HEADER of the model.

Example:

```
DELAY {
   FROM { PIN = n0; } TO { PIN = n5; }
   EQUATION {
      R1*(C1+C2+C3+C4+C5) + R2*(C2+C3+C4+C5)
      + R3*(C3+C4+C5) + R4*(C4+C5) + R5*C5
   }
}
```

External components or stimuli which are not part of the network shall be declared in the HEADER. Also, all arguments for TABLE-based models shall be in the HEADER. To avoid re-declaration of pre-declared components, an EQUATION shall also be used for those arguments in the HEADER which refer to pre-declared components.

Example:

```
SLEWRATE {
   PIN = n5;
   HEADER {
      SLEWRATE { PIN = n0; TABLE {/* numbers */} }
      RESISTANCE  { EQUATION { R1+R2+R3+R4+R5 } TABLE {/* numbers */} }
      CAPACITANCE { EQUATION { C1+C2+C3+C4+C5 } TABLE {/* numbers */} }
   }
   TABLE { /* numbers */ }
}
```

In order to model crosstalk delay and noise, at least two driver and receiver nodes are required. The symbolic state "*" (see Section 5.4.13) shall be used to indicate the signal subjected to noise.

Example:

```
WIRE interconnect_model_with_coupling {
   NODE aggressor_source { NODETYPE = driver; }
   NODE victim_source    { NODETYPE = driver; }
   NODE aggressor_sink   { NODETYPE = receiver; }
   NODE victim_sink      { NODETYPE = receiver; }
```

```
      NODE vdd { NODETYPE = power;  }
      NODE gnd { NODETYPE = ground; }
      CAPACITANCE cc { NODE {aggressor_sink victim_sink}}
      CAPACITANCE cv { NODE {victim_sink gnd }}
      RESISTANCE  rv { NODE {victim_source victim_sink}}
      VECTOR ( 01 aggressor_sink -> ?* victim_sink -> *? victim_sink ) {
         /* xtalk noise model */
      }
      VECTOR (
         ( 01 aggressor_source <&> 01 victim_source )
          -> 01 aggressor_sink -> 01 victim_sink
      ) {
      /* xtalk DELAY model */
      }
   }
```

Example for noise model:

```
   VOLTAGE {
         PIN = victim_sink;
         MEASUREMENT = peak;
         CALCULATION = incremental;
         HEADER {
            SLEWRATE   tra { PIN = aggressor_sink; }
            VOLTAGE     va { NODE {vdd gnd} }
         }
         EQUATION { (1-EXP(-tra/(rv*cv)))*va*rv*cc/tra }
      }
   }
```

Example for delay model:

```
   DELAY {
      FROM { PIN = victim_source; } TO { PIN = victim_sink; }
      CALCULATION = incremental;
      HEADER {
         SLEWRATE tra { PIN = aggressor_sink; }
         SLEWRATE trv { PIN = victim_source; }
      }
      EQUATION { (1-EXP(-tra/(rv*cv)))*rv*cc*trv/tra }
   }
```

The `VOLTAGE` model applies for a rising aggressor signal while the victim signal is stable. The `DELAY` model applies for rising victim signal simultaneous with or followed by a rising aggressor signal at the coupling point. The `VECTOR` implicitly defines the time window of interaction between aggressor and victim; interaction occurs only if the aggressor signal at the coupling point intervenes during the propagation of the victim signal from its source to the coupling point. Both `VOLTAGE` and `DELAY` represent incremental numbers.

### 8.15.6 SELECT_CLASS annotation for WIRE statement

A sophisticated tool can support more than one interconnect model. Each calculation model can have its “netlist” with the appropriate validity range of the RC components. For instance, a lumped model can be used for short nets and a distributed model can be used for longer nets.

Also, models with different accuracy for the same net can be defined. For instance, the lumped model can be used for estimation purpose and the distributed model for signoff.

For this purpose, classes can be defined to select a set of models. The selection shall be defined by the user, in a similar way as a user can select wireload models for pre-layout parasitic estimation. The selected class shall be indicated by the `SELECT_CLASS` annotation within the `WIRE` statement.

Example:

```
LIBRARY my_library {
   CLASS estimation;
   CLASS verification;
   WIRE rough_model_for_short_nets {
      SELECT_CLASS = estimation; /* etc.*/
   }
   WIRE detailed_model_for_short_nets {
      SELECT_CLASS = verification; /* etc.*/
   }
   WIRE rough_model_for_long_nets {
      SELECT_CLASS = estimation; /* etc.*/
   }
   WIRE detailed_model_for_long_nets {
      SELECT_CLASS = verification; /* etc.*/
   }
}
```

# Section 9
# Physical Modeling

## 9.1 Overview

Table 9-1 summarizes the ALF statements for physical modeling.

**Table 9-1 Statements in ALF describing physical objects**

| Statement | Scope | Comment |
|---|---|---|
| LAYER | LIBRARY, SUBLIBRARY | description of a plane provided for physical objects consisting of electrically conducting material |
| VIA | LIBRARY, SUBLIBRARY | description of a physical object for electrical connection between layers |
| SITE | LIBRARY, SUBLIBRARY | placement grid for a class of physically placeable objects |
| BLOCKAGE | CELL | physical object on a layer, forming an obstruction against placing or routing other objects |
| PORT | PIN | physical object on a layer, providing electrical connections to a pin |
| PATTERN | VIA, RULE, BLOCKAGE, PORT | physical object on a layer, described for the purpose of defining relationships with other physical objects |
| RULE | LIBRARY, SUBLIBRARY, CELL, PIN | set of rules defining calculable relationships between physical objects |
| ANTENNA | LIBRARY, SUBLIBRARY, CELL | set of rules defining restrictions for physical size of electrically connected objects for the purpose of manufacturing |
| ARTWORK | VIA, CELL | reference to an imported object from GDS2 |
| ARRAY | LIBRARY, SUBLIBRARY | description of a regular grid for placement, global and detailed routing |
| geometric model | PATTERN | description of the geometric form of a physical object |
| REPEAT | physical object | algorithm to replicate a physical object in a regular way |
| SHIFT | physical object | specification to shift a physical object in x/y direction |
| FLIP | physical object | specification to flip a physical object around an axis |
| ROTATE | physical object | specification to rotate a physical object around an axis |
| BETWEEN | CONNECTIVITY, DISTANCE | reference to objects with a relation to each other |

## 9.2 Arithmetic models in the context of layout

Table 9-2 shows keywords for arithmetic models in the context of layout.

**Table 9-2 Arithmetic models for layout data**

| Keyword | Value type | Base units | Default units | Description |
|---|---|---|---|---|
| SIZE | non-negative number | N/A | 1 | abstract, unitless measurement for the size of a physical object |
| AREA | non-negative number | Square Meter | p (pico) | area in square microns (pico = micro$^2$) |
| DISTANCE | non-negative number | Meter | u (micro) | distance between two points in microns |
| HEIGHT | positive number | Meter | u (micro) | y- dimension of a placeable object (e.g., cell or block)<br>z- dimension of a routeable object (e.g., pattern on routing layer), representing the absolute height above substrate |
| LENGTH | positive number | Meter | u (micro) | x-, or y- dimension of a routeable object (e.g., pattern on routing layer) measured in routing direction |
| WIDTH | positive number | Meter | u (micro) | x-dimension of a placeable object (e.g., cell or block)<br>x- or y- dimension of a routeable object (e.g., pattern on routing layer) measured in orthogonal direction to the route |
| PERIMETER | positive number | Meter | u (micro) | circumference of a physical object |
| THICKNESS | positive number | Meter | u (micro) | z- dimension of a manufacturable physical object, representing the distance between the bottom of the object above and the top of the object below |
| OVERHANG | non-negative number | Meter | u (micro) | distance between the edges of two overlapping physical objects |
| EXTENSION | non-negative number | Meter | u (micro) | distance between the center and the outer edge of a physical object |

Table 9-3 through Table 9-12 summarize the semantic meanings of arithmetic model keywords in the context of layout.

**Table 9-3 Semantic meaning of SIZE**

| Context | Meaning |
|---|---|
| CELL | abstract measure for size of the cell, cost function for design implementation |
| WIRE | - as a model (TABLE or EQUATION):<br>abstract measure for the size of the wire itself<br>- as argument of a model (HEADER):<br>abstract measure for size of the block for which the wireload model applies,<br>can be calculated by combining the size of all cells and all wires in the block |
| ANTENNA | abstract measure for size of the antenna for which the antenna rule applies |

**Table 9-4 Semantic meaning of WIDTH**

| Context | Meaning |
|---|---|
| CELL, SITE | horizontal distance between cell or site boundaries, respectively |
| WIRE | as argument of a model (HEADER):<br>horizontal distance between block boundaries for which wireload model applies |
| LAYER, ANTENNA | width of a wire, orthogonal to routing direction |

**Table 9-5 Semantic meaning of HEIGHT**

| Context | Meaning |
|---|---|
| CELL, SITE | vertical distance between cell or site boundaries, respectively |
| WIRE | as argument of a model (HEADER):<br>vertical distance between block boundaries for which wireload model applies |
| LAYER | distance from top of ground plane to bottom of wire |

**Table 9-6 Semantic meaning of LENGTH**

| Context | Meaning |
|---|---|
| WIRE | estimated routing length of a wire in a wireload model |
| LAYER, ANTENNA | actual routing length of a wire in layout |

**Table 9-7 Semantic meaning of AREA**

| Context | Meaning |
|---|---|
| CELL | physical area of the cell, product of width and height of a rectangular cell |
| WIRE | - as a model (TABLE or EQUATION):<br>physical area of the wire itself<br>- as argument of a model (HEADER):<br>physical area of the block for which wireload model applies,<br>product of width and height of rectangular block |
| LAYER, VIA, ANTENNA | physical area of a placeable or routeable object, measured in the x-y plane |

**Table 9-8 Semantic meaning of PERIMETER**

| Context | Meaning |
|---|---|
| CELL | perimeter of the cell, twice the sum of height and width for rectangular cell |
| WIRE | - as a model (TABLE or EQUATION):<br>perimeter the wire itself<br>- as argument of a model (HEADER):<br>perimeter of the block for which wireload model applies,<br>twice the sum of height and width for rectangular block |
| LAYER, VIA, ANTENNA | perimeter of a placeable or routeable object, measured in the x-y plane |

**Table 9-9 Semantic meaning of DISTANCE**

| Context | Meaning |
|---|---|
| RULE | distance between objects for which the rule applies |

**Table 9-10 Semantic meaning of THICKNESS**

| Context | Meaning |
|---|---|
| LAYER, ANTENNA | distance between top and bottom of a physical object, orthogonal to the x-y plane |

**Table 9-11 Semantic meaning of OVERHANG**

| Context | Meaning |
|---|---|
| RULE | distance between the outer border of an object and the outer border of another object inside the first one |

**Table 9-12 Semantic meaning of EXTENSION**

| Context | Meaning |
|---|---|
| LAYER, VIA, RULE, geometric model | distance between the border of the original object and the border of the same object after enlargement |

# 9.3 Statements for geometric transformation

This section defines SHIFT, ROTATE, FLIP, and REPEAT.

## 9.3.1 SHIFT statement

The SHIFT statement defines the horizontal and vertical offset measured between the coordinates of the geometric model and the actual placement of the object. Eventually, a layout tool only supports integer numbers. The numbers are in units of DISTANCE.

```
shift_annotation_container ::=
          SHIFT { horizontal_or_vertical_annotations }

horizontal_or_vertical_annotations ::=
          horizontal_annotation
   |      vertical_annotation
   |      horizontal_annotation vertical_annotation

horizontal_annotation ::=
          HORIZONTAL = number ;

vertical_annotation ::=
          VERTICAL = number ;
```

If only one annotation is given, the default value for the other one is 0. If the SHIFT statement is not given, both values default to 0.

## 9.3.2 ROTATE statement

The *rotate_annotation* statement defines the angle of rotation in degrees measured between the orientation of the object described by the coordinates of the geometric model and the actual placement of the object measured in counter-clockwise direction, specified by a number between 0 and 360. Eventually, a layout tool can only support angles which are multiple of 90 degrees. The default is 0.

```
rotate_annotation ::=
          ROTATE = number ;
```

The object shall rotate around its origin.

### 9.3.3 FLIP statement

The *flip*_annotation describes a transformation of the specified coordinates by flipping the object around an axis specified by a number between 0 and 180. The number represents the angle of the flipping direction in degrees. Eventually, a layout tool can only support angles which are multiple of 90 degrees. The axis is orthogonal to the flipping direction. The axis shall go through the origin of the object.

```
flip_annotation ::=
          FLIP = number ;
```

Example:

FLIP = 0 means flip in horizontal direction, axis is vertical.
FLIP = 90 means flip in vertical direction, axis is horizontal.

### 9.3.4 REPEAT statement

The REPEAT statement shall be defined as follows:

```
repeat ::=
          REPEAT [ = unsigned ] {
                shift_annotation_container
                [ repeat ]
          }
```

The purpose of the REPEAT statement is to describe the replication of a physical object in a regular way, for example SITE (see Section 9.12). The REPEAT statement can also appear within a geometric_model.

The unsigned number defines the total number of replications. The number 1 means, the object appears just once. If this number is not given, the REPEAT statement defines a rule for an arbitrary number of replications.

REPEAT statements can also be nested.

Examples:

The following example replicates an object three times along the horizontal axis in a distance of 7 units.

```
REPEAT = 3 {
   SHIFT { HORIZONTAL = 7; }
}
```

The following example replicates an object five times along a 45-degree axis.

```
REPEAT = 5 {
   SHIFT { HORIZONTAL = 4; VERTICAL = 4; }
}
```

The following example replicates an object two times along the horizontal axis and four times along the vertical axis.

```
REPEAT = 2 {
   SHIFT { HORIZONTAL = 5; }
   REPEAT = 4 {
      SHIFT { VERTICAL = 6; }
   }
}
```

Note: The order of nested `REPEAT` statements does not matter. The following example gives the same result as the previous example.

```
REPEAT = 4 {
   SHIFT { VERTICAL = 6; }
   REPEAT = 2 {
      SHIFT { HORIZONTAL = 5; }
   }
}
```

### 9.3.5 Summary of geometric transformations

```
geometric_transformations ::=
          geometric_transformation { geometric_transformation }

geometric_transformation ::=
          shift_annotation_container
   |      rotate_annotation
   |      flip_annotation
   |      repeat
```

Rules and restrictions:

- A physical object can contain a `geometric_transformation` statement of any kind, but no more than one of a specific kind.
- The `geometric_transformation` statements shall apply to all `geometric_models` within the context of the object.
- The `geometric_transformation` statements shall refer to the origin of the object, i.e., the point with coordinates `{ 0 0 }`. Therefore, the result of a combined transformation shall be independent of the order in which each individual transformation is applied.

These are demonstrated in Figure 9-1.

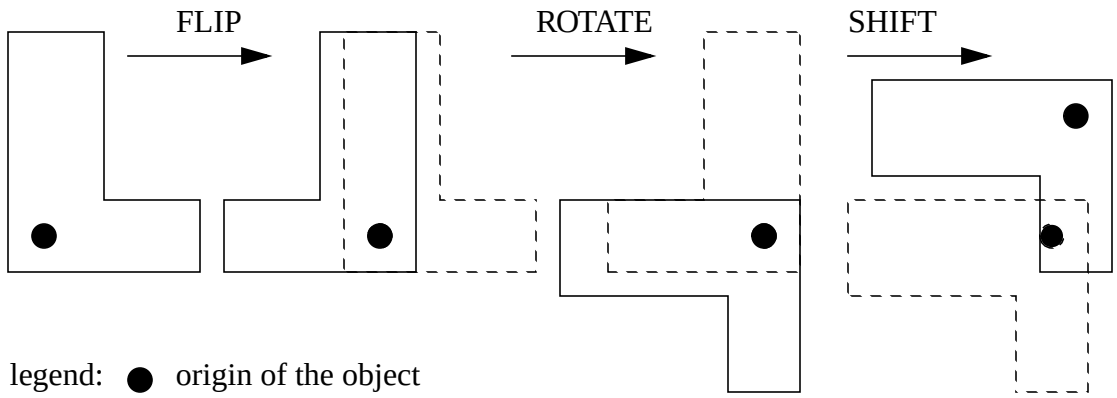



**Figure 9-1: Illustration of FLIP, ROTATE, and SHIFT**

## 9.4 ARTWORK statement

The ARTWORK statement shall be defined as follows:

```
artwork ::=
          ARTWORK = artwork_identifier {
                [ geometric_transformations ]
                { pin_assignments }
          }
```

The ARTWORK statement creates a reference between the cell in the library and the original cell imported from a physical layout database (e.g., GDS2).

The geometric_transformations define the operations for transformation from the artwork geometry to the actual cell geometry. In other words, the artwork is considered as the original object whereas the cell is the transformed object.

The imported cell can have pins with different names. The LHS of the pin_assignments describes the pin names of the original cell, the RHS describes the pin names of the cell in this library. See Section 11.4 for the syntax of pin_assignments .

Example:

```
CELL my_cell {
   PIN A { /* fill in pin items */ }
   PIN Z { /* fill in pin items */ }
   ARTWORK = \GDS2$!@#$ {
      SHIFT { HORIZONTAL = 0; VERTICAL = 0; }
      ROTATE = 0;
      \GDS2$!@#$A = A;
      \GDS2$!@#$B = B;
   }
}
```

# 9.5 LAYER statement

This section defines the LAYER statements.

## 9.5.1 Definition

The LAYER statement shall be defined as follows:

```
layer ::=
          LAYER identifier { layer_items }

layer_items ::=
          layer_item { layer_item }

layer_item ::=
          all_purpose_item
   |      arithmetic_model
   |      arithmetic_model_container
```

The syntax and semantics of all_purpose_item, arithmetic_model_container, and arithmetic_model are defined in Section 11.7 and Section 11.16.

Specific items applicable for LAYER are listed in Table 9-3.

**Table 9-13 Items for LAYER description**

| Item | Applies for layer | Usable ALF statement | Comment |
|---|---|---|---|
| purpose | all | PURPOSE = <identifier> ; | see Section 9.5.2 |
| property | routing, cut, master | PROPERTY { ... } | see Section 3.2.7 |
| current density limit | routing, cut | LIMIT { CURRENT { ... MAX { ... } } } | see Section 7.5, Section 8.1.2, Section 7.6.1, Section 8.9.1, and Section 9.5.5 |
| resistance | routing, cut | RESISTANCE { ... } | see Section 8.7.2 and Section 9.5.5 |
| capacitance | routing | CAPACITANCE {... } | see Section 8.7.2 and Section 9.5.5 |
| default width or minimum width | routing | WIDTH { DEFAULT = <number>; } | see Section 7.1.4., Section 9.2, and Section 9.5.5 |
| manufacturing tolerance for width | routing | WIDTH { MIN = <number>; TYP = <number>; MAX = <number>; } | see Section 7.6.1, Section 8.9.1, and Section 9.5.5 |

**Table 9-13 Items for LAYER description, *continued***

| Item | Applies for layer | Usable ALF statement | Comment |
|---|---|---|---|
| default wire extension | routing | `EXTENSION { DEFAULT = <number>; }` | see Section 9.7.4 and Section 9.5.5 |
| height | routing, cut, master | `HEIGHT = <number>;` | see Section 9.2 |
| thickness | routing, cut, master | `THICKNESS = <number>;` | see Section 9.2 |
| preferred routing direction | routing | `PREFERENCE` | see Section 9.5.4 |

Note: Rules involving relationships between objects within one or several layers is described in the RULE statement (see Section 9.11).

### 9.5.2 PURPOSE annotation

The purpose of each layer shall be identified using the `PURPOSE` annotation.

```
layer_purpose_assignment ::=
        PURPOSE = layer_purpose_identifier ;

layer_purpose_identifier ::=
        routing
    |   cut
    |   substrate
    |   dielectric
    |   reserved
    |   abstract
```

The identifiers have the following definitions:

- *routing*: layer provides electrical connections within one plane
- *cut*: layer provides electrical connections between planes
- *substrate*: layer(s) at the bottom
- *dielectric*: provides electrical isolation between planes
- *reserved*: layer is for proprietary use only
- *abstract*: not a manufacturable layer, used for description of boundaries between objects

`LAYER` statements shall be in sequential order defined by the manufacturing process, starting bottom-up in the following sequence: one or multiple substrate layers, followed by alternating cut and routing layers, then the dielectric layer. Abstract layers can appear at the end of the sequence.

### 9.5.3 PITCH annotation

The PITCH annotation identifies the routing pitch for a layer with PURPOSE=routing.

```
pitch_annotation ::=
   PITCH = non_negative_number ;
```

The pitch is measured between the center of two adjacent parallel wires routed on the layer.

### 9.5.4 PREFERENCE annotation

The PREFERENCE annotation for LAYER shall have the following form:

```
routing_preference_annotation ::=
         PREFERENCE = routing_preference_identifier ;

routing_preference_identifier ::=
         horizontal
   |     vertical
```

The purpose is to indicate the preferred routing direction.

### 9.5.5 Example

This example contains a default width (the syntax is all_purpose_item), resistance, capacitance, and current limits (the syntax is arithmetic_model) for arbitrary wires in a routing layer. Since width and thickness are arguments of the models, special wires and fat wires are also taken into account.

```
LAYER metal1 {
   PURPOSE = routing;
   PREFERENCE { HORIZONTAL = 0.75; VERTICAL = 0.25; }
   WIDTH { DEFAULT = 0.4; MIN = 0.39; TYP = 0.40; MAX = 0.41; }
   THICKNESS { DEFAULT = 0.2; MIN = 0.19; TYP = 0.20; MAX = 0.21; }
   EXTENSION { DEFAULT = 0; }
   RESISTANCE {
      HEADER { LENGTH WIDTH THICKNESS TEMPERATURE }
      EQUATION {

         0.5*(LENGTH/(WIDTH*THICKNESS))
         *(1.0+0.01*(TEMPERATURE-25))
      }
   }
   CAPACITANCE {
      HEADER { AREA PERIMETER }
      EQUATION { 0.48*AREA + 0.13*PERIMETER*THICKNESS }
   }
   LIMIT {
      CURRENT ac_limit_for_avg {
         UNIT = mAmp ;
         MEASUREMENT = average ;
         HEADER {
```

```
            WIDTH { UNIT = uM; TABLE { 0.4 0.8 } }
            FREQUENCY { UNIT = megHz; { 1 100 } }
            THICKNESS { UNIT = uM; TABLE { 0.2 0.4 } }
         }
         TABLE {
            2.0e-6 4.0e-6 1.5e-6 3.0e-6
            4.0e-6 8.0e-6 3.0e-6 6.0e-6
         }
      }
      CURRENT ac_limit_for_rms {
         UNIT = mAmp ;
         MEASUREMENT = rms ;
         HEADER {
            WIDTH { UNIT = uM; TABLE { 0.4 0.8 } }
            FREQUENCY { UNIT = megHz; { 1 100 } }
            THICKNESS { UNIT = uM; TABLE { 0.2 0.4 } }
         }
         TABLE {
            4.0e-6 7.0e-6 4.5e-6 7.5e-6
            8.0e-6 14.0e-6 9.0e-6 15.0e-6
         }
      }
      CURRENT ac_limit_for_peak {
         UNIT = mAmp ;
         MEASUREMENT = peak ;
         HEADER {
            WIDTH { UNIT = uM; TABLE { 0.4 0.8 } }
            FREQUENCY { UNIT = megHz; { 1 100 } }
            THICKNESS { UNIT = uM; TABLE { 0.2 0.4 } }
         }
         TABLE {
            6.0e-6 10.0e-6 5.9e-6 9.9e-6
            12.0e-6 20.0e-6 11.8e-6 19.8e-6
         }
      }
      CURRENT dc_limit {
         UNIT = mAmp ;
         MEASUREMENT = static ;
         HEADER {
            WIDTH { UNIT = uM; TABLE { 0.4 0.8 } }
            THICKNESS { UNIT = uM; TABLE { 0.2 0.4 } }
         }
         TABLE { 2.0e-6 4.0e-6 4.0e-6 8.0e-6 }
      }
   }
}
```

## 9.6 Geometric model statement

This section defines the geometric model statement and how to predefine commonly used objects (using `TEMPLATE`).

### 9.6.1 Definition

The geometric model statement shall be defined as follows:

```
geometric_model ::=
          geometric_model_identifier
             [ geometric_model_name_identifier ] {
                all_purpose_items
                coordinates
          }
   |      geometric_model_template_instantiation

geometric_models ::=
          geometric_model { geometric_model }

geometric_model_identifier ::=
          DOT
   |      POLYLINE
   |      RING
   |      POLYGON

coordinates ::=
          COORDINATES { x_number y_number { x_number y_number } }
```

A point is a pair of `x_number` and `y_number`.

A **`DOT`** is `1` point.

A **`POLYLINE`** is defined by `N>1` connected points, forming an open object.

A **`RING`** is defined by `N>1` connected points, forming a closed object, i.e., the last point is connected with first point. The object occupies the edges of the enclosed space.

A **`POLYGON`** is defined by `N>1` connected points, forming a closed object, i.e., the last point is connected with first point. The object occupies the entire enclosed space.

All of these are depicted in Figure 9-2.

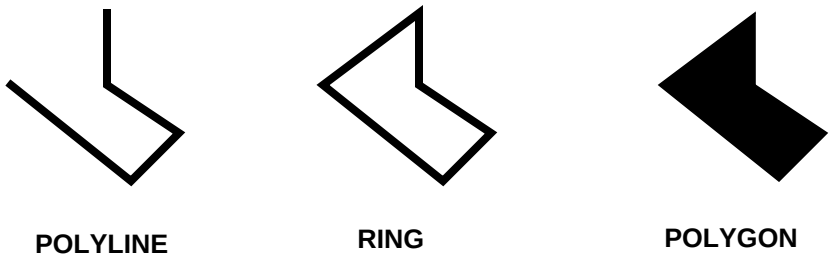



**Figure 9-2: Illustration of geometric models**

See Section 9.3.4 for the definition of the `repeat` statement.

The *point_to_point*_annotation applies for **POLYLINE**, **RING**, and **POLYGON**. It specifies how the connections between points is made. The default is `straight`, which defines a straight connection (see Figure 9-3). The value `rectilinear` specifies a connection by moving in the x-direction first and then moving in the y-direction (see Figure 9-4). This enables a non-redundant specification of rectilinear objects using `N/2` points instead of `N` points.

```
point_to_point_annotation ::=
        POINT_TO_POINT = point_to_point_identifier ;

point_to_point_identifier ::=
        straight
   |    rectilinear
```

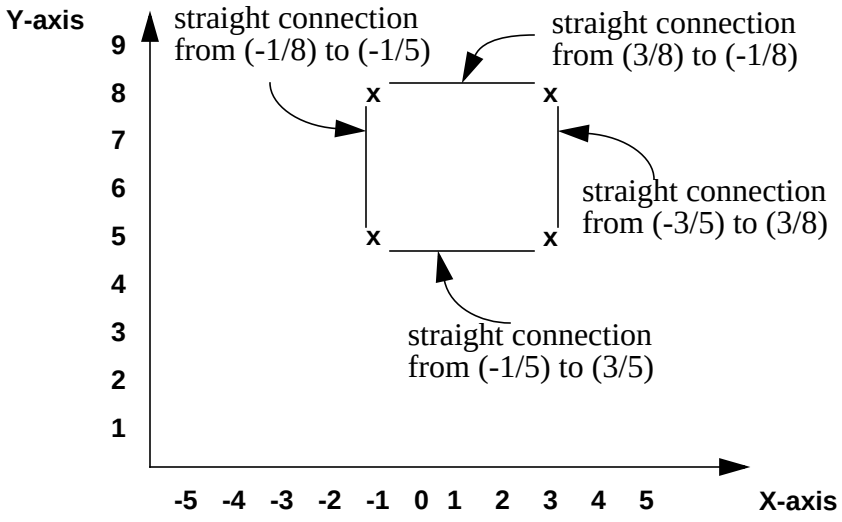



**Figure 9-3: Illustration of straight point-to-point connection**

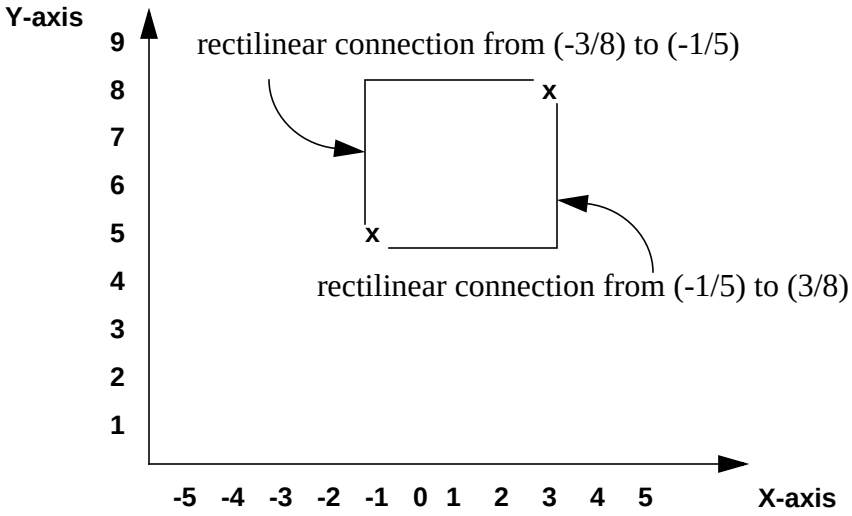



**Figure 9-4: Illustration of rectilinear point-to-point connection**

Example:

```
POLYGON {
   POINT_TO_POINT = straight;
   COORDINATES { -1 5 3 5 3 8 -1 8 }
}
POLYGON {
   POINT_TO_POINT = rectilinear;
   COORDINATES { -1 5 3 8 }
}
```

Both objects describe the same rectangle.

### 9.6.2 Predefined geometric models using TEMPLATE

The `TEMPLATE` construct (see Section 3.2.6) can be used to predefine some commonly used objects.

The templates RECTANGLE and LINE shall be predefined as follows:

```
TEMPLATE RECTANGLE {
   POLYGON {
      POINT_TO_POINT = rectilinear;
      COORDINATES { <left> <bottom> <right> <top> }
   }
}
```

```
TEMPLATE LINE {
   POLYLINE {
      POINT_TO_POINT = straight;
      COORDINATES { <x_start> <y_start> <x_end> <y_end> }
   }
}
```

The following example shows the instantiation of predefined templates.

```
// same rectangle as in previous example
RECTANGLE {left = -1; bottom = 5; right = 3; top = 8; }
//or
RECTANGLE {-1 5 3 8 }

// diagonals through the rectangle
LINE {x_start = -1; y_start = 5; x_end = 3; y_end = 8; }
LINE {x_start = 3; y_start = 5; x_end = -1; y_end = 8; }
//or
LINE { -1 5 3 8 }
LINE { 3 5 -1 8 }
```

The definitions for predefined templates are fixed. Therefore the keywords RECTANGLE and LINE are reserved. On the other hand, the definitions for user-defined templates are only known by the library supplied by the user.

The following example shows some user-defined templates.

```
TEMPLATE HORIZONTAL_LINE {
   POLYLINE {
      POINT_TO_POINT = straight;
      COORDINATES { <left> <y> <right> <y> }
   }
}

TEMPLATE VERTICAL_LINE {
   POLYLINE {
      POINT_TO_POINT = straight;
      COORDINATES { <x> <bottom> <x> <top> }
   }
}
```

The following example shows the instantiation of user-defined templates.

```
// lines bounding the rectangle
HORIZONTAL_LINE { y = 5; left = -1; right = 3; }
HORIZONTAL_LINE { y = 8; left = -1; right = 3; }
VERTICAL_LINE { x = -1; bottom = 5; top = 8; }
VERTICAL_LINE { x = 3; bottom = 5; top = 8; }
//or
HORIZONTAL_LINE { 5 -1 3 }
HORIZONTAL_LINE { 8 -1 3 }
VERTICAL_LINE { -1 5 8 }
VERTICAL_LINE { 3 5 8 }
```

# 9.7 PATTERN statement

This section defines the PATTERN statement and its annotations.

## 9.7.1 Definition

The PATTERN statement shall be defined as follows:

```
pattern ::=
   PATTERN [ identifier ] {
         [ all purpose_items ]
         [ geometric_models ]
         [ geometric_transformations ]
   }
```

## 9.7.2 SHAPE annotation

The SHAPE annotation is defined as follows

```
shape_assignment ::=
         SHAPE = shape_identifier ;

shape_identifier ::=
         line
   |     tee
   |     cross
   |     jog
   |     corner
   |     end
```

SHAPE applies only for a PATTERN in a routing layer, as shown in Figure 9-5. The default is **line**.

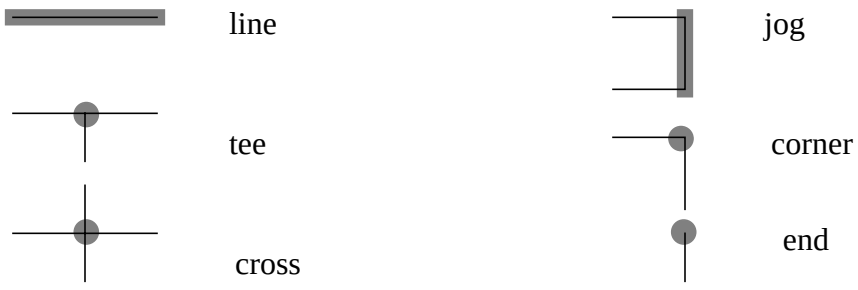



**Figure 9-5: Routing layer shapes**

line and jog represent routing segments, which can have an individual LENGTH and WIDTH. The LENGTH *between* routing segments is defined as the common run length. The DISTANCE *between* routing segments is measured orthogonal to the routing direction.

tee, cross, and corner represent intersections between routing segments. end represents the end of a routing segment. Therefore, they have points rather than lines as references. The points can have an EXTENSION. The DISTANCE between points can be measured straight or by using HORIZONTAL and VERTICAL.

### 9.7.3 LAYER annotation

The *layer*_annotation defines the layer where the object resides. The layer shall have been declared before.

```
layer_annotation ::=
        LAYER = layer_identifier ;
```

### 9.7.4 EXTENSION annotation

The *extension*_annotation specifies the value by which the drawn object is extended at all sides.

```
extension_annotation ::=
        EXTENSION = non_negative_number ;
```

The default value of *extension*_annotation is 0.

### 9.7.5 VERTEX annotation

The *vertex*_annotation shall appear only in conjunction with the *extension*_annotation. It specifies the form of the extended object, as shown in Figure 9-6.

```
vertex_annotation ::=
        VERTEX = vertex_identifier ;

vertex_identifier ::=
        round
   |    straight
```

The default value of *vertex*_annotation is **straight**.

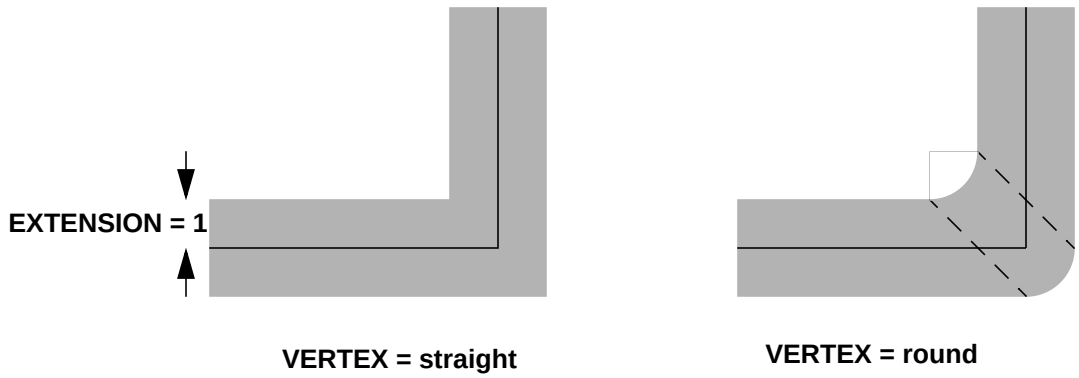



**Figure 9-6: Illustration of VERTEX annotation**

### 9.7.6 PATTERN with geometric model

A `geometric_model` describes the form of a physical object; it does not describe a physical object itself. The `geometric_model` shall be in the context of a `PATTERN`.

A pattern can contain `geometric_model` statements, geometric transformation statements (see Section 9.3.5), and `all_purpose_items` (see Section 11.7).

### 9.7.7 Example

```
PATTERN {
   LAYER = metal1;
   EXTENSION = 1;
   DOT { COORDINATES { 5 10 } }
}
```

This object is effectively a square, with a lower left corner (`x=4, y=9`) and upper right corner (`x=6, y=11`).

## 9.8 VIA statement

This section defines the `VIA` statement and its annotations.

### 9.8.1 Definition

The `VIA` statement shall be defined as follows:

```
via ::=
          VIA [ identifier ] { via_items }
```

```
via_items ::=
          via_item { via_item }

via_item ::=
          all_purpose_item
    |     pattern
    |     arithmetic_model
```

The VIA statement shall contain at least three patterns, referring to the cut layer and two adjacent routing layers. Stacked vias can contain more than three patterns.

The all_purpose_items and arithmetic_models for VIA are listed in Table 9-14.

**Table 9-14 Items for VIA description**

| Item | Usable ALF statement | Comment |
|---|---|---|
| property | PROPERTY | see Section 3.2.7 |
| resistance | RESISTANCE | see Section 8.7.2 |
| GDS2 reference | ARTWORK | see Section 9.4 and Section 9.8.3 |
| usage | USAGE | see Section 9.8.2 and Section 9.8.3 |

### 9.8.2 USAGE annotation

The USAGE annotation for a VIA shall have one of the following mutually exclusive values.

```
usage_annotation ::=
          USAGE = usage_identifier ;

usage_identifier ::=
          default
    |     non_default
    |     partial_stack
    |     full_stack
```

The identifiers have the following definitions:

- *default*: via can be used per default
- *non_default*: via can only be used if authorized by a RULE
- *partial_stack*: via contains 3 patterns: lower and upper routing layer and cut layer in-between. It can only be used to build stacked vias. The bottom of a stack can be a default or a non_default via.
- *full_stack*: via contains 2N+1 patterns (N>1). It describes the full stack from bottom to top.

### 9.8.3 Example

```
VIA via_with_two_contacts_in_x_direction {
   ARTWORK = GDS2_name_of_my_via {
      SHIFT { HORIZONTAL = -2; VERTICAL = -3; }
      ROTATE = 180;
   }
   PATTERN via_contacts {
      LAYER = cut_1_2 ;
      RECTANGLE { 1 1 3 3 }
      REPEAT = 2 {
         SHIFT{ HORIZONTAL = 4; }
         REPEAT = 1 {
            SHIFT { VERTICAL = 4; }
   }  }  }
   PATTERN lower_metal {
      LAYER = metal_1 ;
      RECTANGLE { 0 0 8 4 }
   }
   PATTERN upper_metal {
      LAYER = metal_2 ;
      RECTANGLE { 0 0 8 4 }
   }
}
```

A TEMPLATE (see Section 3.2.6) can be used to define a construction rule for a via.

```
TEMPLATE my_via_rule
   VIA <via_rule_name> {
      PATTERN via_contacts {
         LAYER = cut_1_2 ;
         RECTANGLE { 1 1 3 3 }
         REPEAT = <x_repeat> {
            SHIFT{ HORIZONTAL = 4; }
            REPEAT = <y_repeat> {
               SHIFT { VERTICAL = 4; }
      }  }  }
      PATTERN lower_metal {
         LAYER = metal_1 ;
         RECTANGLE { 0 0 <x_cover> <y_cover> }
      }
      PATTERN upper_metal {
         LAYER = metal_2 ;
         RECTANGLE { 0 0 <x_cover> <y_cover> }
      }
   }
}
```

A static instance of the TEMPLATE can be used to create the same via as in the first example (except for the reference to GDS2):

```
my_via_rule {
   via_rule_name = via_with_two_contacts_in_x_direction;
   x_cover = 8;
   y_cover = 4;
   x_repeat = 2;
   y_repeat = 1;
}
```

A dynamic instance of the TEMPLATE (see Section 5.6.8) can be used to create a via rule.

```
my_via_rule = dynamic {
   via_rule_name = via_with_NxM_contacts;
   x_cover = 8;
   y_cover = 4;
   x_repeat {
      HEADER { x_cover { TABLE { 4 8 12 16 } } }
      TABLE { 1 2 3 4 }
   }
   y_repeat {
      HEADER { y_cover { TABLE { 4 8 12 16 } } }
      TABLE { 1 2 3 4 }
   }
}
```

Instead of defining fixed values for the placeholders, here the mathematical relationships between the placeholders are defined, which can generate a via rule for any set of values.

### 9.8.4 VIA reference

Certain physical objects can contain a reference to one or more vias, using the following statement.

```
via_reference ::=
          VIA { via_instantiations }

via_instantiations ::=
          via_instantiation { via_instantiation }

via_instantiation ::=
          via_identifier { geometric_transformations }
```

The *via*_identifier shall be the name of an already defined VIA.

Example for a via reference in a PORT, see Section 9.10.

## 9.9 BLOCKAGE statement

This section defines the BLOCKAGE statement and its use.

### 9.9.1 Definition

The BLOCKAGE statement shall be defined as follows:

```
blockage ::=
          BLOCKAGE [ identifier ] {
                     [ all_purpose_items ]
                     [ patterns ]
          }
```

See Section 11.7 for applicable `all_purpose_items`.

### 9.9.2 Example

```
CELL my_cell {
   BLOCKAGE my_blockage {
      PATTERN p1 {
         LAYER = metal1;
         RECTANGLE { -1 5 3 8 }
         RECTANGLE { 6 12 3 8 }
      }
      PATTERN p2 {
         LAYER = metal2;
         RECTANGLE { -1 5 3 8 }
      }
   }
}
```

The `BLOCKAGE` consists of two rectangles covering `metal1` and one rectangle covering `metal2`.

## 9.10 PORT statement

This section defines the `PORT` statement and its use.

### 9.10.1 Definition

A port is a collection of geometries within a pin, representing electrically equivalent points.

The `PORT` statement shall be defined as follows:

```
port ::=
          PORT port_identifier ;
     |    PORT [ port_identifier ] {
                     [ all_purpose_items ]
                     [ patterns ]
                     [ via_reference ]
          }
```

A numerical digit can be used as the first character in *`port`*`_identifier`. In this case the number shall be proceeded by the escape character (see Section 10.3.8) in the declaration of the `PORT`.

The PORT statement is legal within the context of a PIN statement. For this purpose, the syntax for pin_item (see Section 11.11) shall be augmented as follows:

```
pin_item ::=
          all_purpose_item
   |      arithmetic_model
   |      port
```

A pin can have either no PORT statement, an arbitrary number of PORT statements with a *port*_identifier, or exactly one PORT statement without a *port*_identifier.

### 9.10.2 VIA reference

A PORT can contain a reference to one or more vias by using the via_reference statement (see Section 9.8.4).

Example:

```
VIA my_via { /* put via definition here */ }

// later in the same library
CELL my_cell {
   PIN my_pin {
      PORT my_port {
         VIA {
            my_via { SHIFT { HORIZONTAL = 1.0 ; VERTICAL = 2.0 ; } }
            my_via { SHIFT { HORIZONTAL = 5.0 ; VERTICAL = 8.0 ; } }
         }
      }
   }
}
```

The VIA my_via is instantiated twice in the PORT my_port within the PIN my_pin of the CELL my_cell. The origin of the instantiated vias is shifted with respect to the origin of the cell, as specified by the SHIFT statements.

### 9.10.3 CONNECTIVITY rules for PORT and PIN

By default, all connections to a pin shall be made to the same port. Different ports of a pin shall not be connected externally. Those defaults can be overridden by using connectivity rules for ports within a pin.

Pins of the same cell shall not be shorted externally by default. This default can also be overridden by using connectivity rules for pins within a cell.

Example:

```
PIN A {
   PORT P1 { VIEW=physical; }
}
```

```
PIN B {
   PORT Q1 { VIEW=physical; }
   PORT Q2 { VIEW=physical; }
   PORT Q3 { VIEW=physical; }
   CONNECTIVITY {
      CONNECT_RULE = can_short;
      BETWEEN { Q1 Q3 }
   }
   CONNECTIVITY {
      CONNECT_RULE = cannot_short;
      BETWEEN { Q1 Q2 }
   }
   CONNECTIVITY {
      CONNECT_RULE = cannot_short;
      BETWEEN { Q2 Q3 }
   }
}
CONNECTIVITY {
   CONNECT_RULE = must_short;
   BETWEEN { A B }
}
```

The router can make external connections between `Q1` and `Q3`, but not between `Q1` and `Q2` or between `Q2` and `Q3`, respectively. The router shall make an external connection between `A.P1` and any port of `B` (`B.Q1`, `B.Q2`, or `B.Q3`).

### 9.10.4 Reference of a declared PORT in a PIN annotation

In the context of timing modeling, a `PORT` can have the semantic meaning of a `PIN`. For examples, `PORT`s can be used as `FROM` and/or `TO` points of delay measurements -- use a reference by a `hierarchical_identifier`.

Example:

```
CELL my_cell {
   PIN A {
      DIRECTION = input;
      PORT p1;
      PORT p2;
   }
   PIN Z {
      DIRECTION = output;
   }
   VECTOR ( 01 A -> 01 Z ) {
      DELAY {
         FROM { PIN = A.p1; }
         TO { PIN = Z; }
```

```
        }
        DELAY {
           FROM { PIN = A.p2; }
           TO { PIN = Z; }
        }
     }
  }
```

### 9.10.5 VIEW annotation

A subset of values for the VIEW annotation inside a PIN (see Section 6.4.1) shall be applicable for a PORT as well.

```
port_view_annotation ::=
          VIEW = port_view_identifier ;

port_view_identifier ::=
          physical
   |      none
```

VIEW=physical shall qualify the PORT as a real port with the possibility to connect a routing wire to it.

VIEW=none shall qualify the PORT as a virtual port for modeling purpose only.

### 9.10.6 LAYER annotation

The *layer*_annotation can appear inside a PORT (see Section 9.10).

### 9.10.7 ROUTING_TYPE

A PORT can inherit the ROUTING_TYPE from its PIN or it can have its own ROUTING_TYPE annotation.

## 9.11 RULE statement

This section defines the RULE statement and its use.

### 9.11.1 Definition

The RULE statement shall be defined as follows:

```
rule ::=
          RULE [ identifier ] { rule_items }

rule_items ::=
          rule_item { rule_item }
```

```
rule_item ::=
          pattern
   |      all_purpose_item
   |      arithmetic_model
```

The `all_purpose_items` for RULE are listed in Table 9-15.

**Table 9-15 Items for RULE description**

| Item | Usable ALF statement | Comment |
|---|---|---|
| rule is for same net or different nets | `CONNECTIVITY` | see Section 9.10.3 and Section 9.15 |
| spacing rule | `LIMIT { DISTANCE ... }` | see Section 7.5 and Section 9.11.2 |
| overhang rule | `LIMIT { OVERHANG ... }` | see Section 7.5 and Section 9.11.3 |

The rules for spacing and overlap, respectively, shall be expressed using the `LIMIT` construct with `DISTANCE` and `OVERHANG`, respectively, as keywords for the arithmetic models (see Section 7.5 and Section 7.6.1). The keywords `HORIZONTAL` and `VERTICAL` shall be introduced as qualifiers for arithmetic submodels (see Section 7.6) to distinguish rules for different routing directions. If these qualifiers are not used, the rule shall apply in any routing direction.

### 9.11.2 Width-dependent spacing

An example of width-dependent spacing is:

```
RULE width_and_length_dependent_spacing {
   PATTERN segment1 { LAYER = metal_1; SHAPE = line; }
   PATTERN segment2 { LAYER = metal_1; SHAPE = line; }
   CONNECTIVITY {
      CONNECT_RULE = cannot_short;
      BETWEEN { segment1 segment2 }
   }
   LIMIT {
      DISTANCE { BETWEEN { segment1 segment2 }
         MIN {
            HEADER {
               WIDTH w1 {
                  PATTERN = segment1;
                  /* TABLE, if applicable */
               }
               WIDTH w2 {
                  PATTERN = segment2;
                  /* TABLE, if applicable */
               }
               LENGTH common_run {
                  BETWEEN { segment1 segment2 }
                  /* TABLE, if applicable */
               }
```

```
            }
            /* EQUATION or TABLE */
         }
         MAX { /* some technology have MAX spacing rules */ }
      }
   }
}
```

Spacing rules dependent on routing direction can be expressed as follows:

```
LIMIT {
   DISTANCE { BETWEEN { segment1 segment2 }
      HORIZONTAL {
         MIN { /* HEADER, EQUATION or TABLE */ }
      }
      VERTICAL {
         MIN { /* HEADER, EQUATION or TABLE */ }
      }
   }
}
```

### 9.11.3 End-of-line rule

End-of-line rules can be expressed as follows:

```
RULE lonely_via {
   PATTERN via_lower { LAYER = metal_1; SHAPE = line; }
   PATTERN via_cut   { LAYER = cut_1_2; }
   PATTERN via_upper { LAYER = metal_2; SHAPE = end;  }
   PATTERN adjacent  { LAYER = metal_2; SHAPE = line; }
   CONNECTIVITY {
      CONNECT_RULE = must_short;
      BETWEEN { via_lower via_cut via_upper }
   }
   CONNECTIVITY {
      CONNECT_RULE = cannot_short;
      BETWEEN { via_upper adjacent }
   }
   LIMIT {
      OVERHANG {
         BETWEEN { via_cut via_upper }
         MIN {
            HEADER {
               DISTANCE {
                  BETWEEN { via_cut adjacent }
                  /* TABLE, if applicable */
               }
            }
            /* TABLE or EQUATION */
         }
      }
   }
}
```

Overhang dependent on routing direction can be expressed as follows:

```
LIMIT {
   OVERHANG { BETWEEN { via_cut via_upper }
      HORIZONTAL {
         MIN { /* HEADER, EQUATION or TABLE */ }
      }
      VERTICAL {
         MIN { /* HEADER, EQUATION or TABLE */ }
      }
   }
}
```

### 9.11.4 Redundant vias

Rules for redundant vias can be expressed as follows:

```
RULE constraint_for_redundant_vias {
   PATTERN via_lower { LAYER = metal_1; }
   PATTERN via_cut   { LAYER = cut_1_2; }
   PATTERN via_upper { LAYER = metal_2; }
   CONNECTIVITY {
      CONNECT_RULE = must_short;
      BETWEEN { via_lower via_cut via_upper }
   }
   LIMIT {
      WIDTH {
         PATTERN = via_cut;
         MIN = 3; MAX = 5;
      }
      DISTANCE {
         BETWEEN { via_cut }
         MIN = 1; MAX = 2;
      }
      OVERHANG {
         BETWEEN { via_lower via_cut }
         MIN = 2; MAX = 4;
      }
      OVERHANG {
         BETWEEN { via_upper via_cut }
         MIN = 2; MAX = 4;
      }
   }
}
```

### 9.11.5 Extraction rules

Extraction rules can be expressed as follows:

```
RULE parallel_lines_same_layer {
   PATTERN segment1 { LAYER = metal_1; SHAPE = line; }
   PATTERN segment2 { LAYER = metal_1; SHAPE = line; }
   CAPACITANCE {
      BETWEEN { segment1 segment2 }
      HEADER {
         DISTANCE {
            BETWEEN { segment1 segment2 }
            /* TABLE, if applicable */
         }
         LENGTH {
            BETWEEN { segment1 segment2 }
            /* TABLE, if applicable */
         }
      }
      /* EQUATION or TABLE */
   }
}
```

### 9.11.6 RULES within BLOCKAGE or PORT

General width-dependent spacing rules can not apply to blockages which are abstractions of smaller blockages collapsed together. The spacing rule between the constituents of the blockage and their neighboring objects shall be applied instead.

For example, a blockage can consist of two parallel wires in vertical direction of `width=1` and `distance=1`. They can be collapsed to form a blockage of `width=3`. Left and right of the blockage, the spacing rule shall be based on the width of the constituent wires (i.e., `1`) instead of the width of the blockage (i.e., `3`).

Therefore, it shall be legal within a `RULE` statement to appear within the context of a `BLOCKAGE` or `PORT` and reference a `PATTERN` which has been defined within the context of the `BLOCKAGE` or `PORT`.

Example:

```
CELL my_cell {
   BLOCKAGE my_blockage {
      PATTERN my_pattern {
         LAYER = metal1;
         RECTANGLE { 5 0 8 10 }
      }
      RULE for_my_pattern {
         PATTERN my_metal1 { LAYER = metal1; }
         LIMIT {
            DISTANCE {
               BETWEEN { my_metal1 my_pattern }
               MIN = 1;
            }
         }
      }
   }
}
```

It shall also be legal to define the spacing rule, which normally would be inside the RULE statement, directly within the context of a PATTERN using the LIMIT construct and the arithmetic model for DISTANCE. This arithmetic model shall not contain a BETWEEN statement. The spacing rule shall apply between the PATTERN and any external object on the same layer.

Example:

```
CELL my_cell {
   BLOCKAGE my_blockage {
      PATTERN p1 {
         LAYER = metal1;
         RECTANGLE { 5 0 8 10 }
         LIMIT { DISTANCE { MIN = 1; } }
      }
   }
}
```

### 9.11.7 VIA reference

A RULE can contain a reference to one or more vias, using the via_reference statement (see Section 9.8.4).

## 9.12 SITE statement

This section defines the SITE statement and its use.

### 9.12.1 Definition

The SITE statement shall be defined as follows:

```
site ::=
          SITE site_identifier { all_purpose_items }
```

The *width*_annotation and *height*_annotation (see Section 9.2) are mandatory.

### 9.12.2 ORIENTATION_CLASS and SYMMETRY_CLASS

A set of CLASS statements shall be used to define a set of legal orientations applicable to a SITE. Both the CLASS and the SITE statements shall be within the context of the same LIBRARY or SUBLIBRARY.

```
orientation_class ::=
          CLASS orientation_class_identifier {
                [ geometric_transformations ]
          }
```

To refer to a predefined orientation class, use the ORIENTATION_CLASS statement within a SITE and/or a CELL. ORIENTATION of a CELL means the orientation of the cell itself. ORIENTATION of a SITE means the orientation of rows that can be created using that site.

```
orientation_class_multivalue_annotation ::=
        ORIENTATION { orientation_class_identifiers }
```

The SYMMETRY_CLASS statement shall be used for a SITE to indicate symmetry between legal orientations. Multiple SYMMETRY statements shall be legal to enumerate all possible combinations in case they cannot be described within a single SYMMETRY statement.

```
symmetry_class_multivalue_annotation ::=
        SYMMETRY_CLASS { orientation_class_identifiers }
```

Legal orientation of a cell within a site shall be defined as the intersection of legal cell orientation and legal site orientation. If there is a set of common legal orientations for both cell and site without symmetry, the orientation of cell instance and site instance shall match.

If there is a set of common legal orientations for both cell and site with symmetry, the cell can be placed on the side using any orientation within that set.

*Case 1*: no symmetry

Site has legal orientations A and B. Cell has legal orientations A and B. When the site is instantiated in the A orientation, the cell shall be placed in the A orientation.

*Case 2*: symmetry

Site has legal orientations A and B and symmetry between A and B. Cell has legal orientations A and B. When the site is instantiated in the A orientation, the cell can be placed in the A or B orientation.

### 9.12.3 Example

```
LIBRARY my_library {
   CLASS north { ROTATE = 0; }
   CLASS flip_north { ROTATE = 0; FLIP = 0; }
   CLASS south { ROTATE = 180; }
   CLASS flip_south { FLIP = 90; }

   SITE Site1 {
      ORIENTATION_CLASS { north flip_north }
   }

   SITE Site2 {
      ORIENTATION_CLASS { north flip_north south flip_south}
      SYMMETRY_CLASS { north flip_north }
      SYMMETRY_CLASS { south flip_south }
   }
   CELL Cell1 {
      SITE { Site1 Site2 }
```

```
         ORIENTATION_CLASS { north flip_north }
      }
      CELL Cell2 {
         SITE { Site2 }
         ORIENTATION_CLASS { north south }
      }
   }
```

`Cell1` can be placed on `site1`. The orientation of `Site1` and `Cell1` shall match because there is no symmetry between `north` and `flip_north` in `Site1`.

`Cell1` can be placed on `Site2`, provided `Site2` is instantiated in the `north` or `flip_north` orientation. The orientation of `site2` and `cell1` need not match because of the symmetry between `north` and `flip_north` in `Site2`.

`Cell2` can be placed on `Site2`, provided `Site2` is instantiated in the `north` or `south` orientation. The orientation of `Site2` and `Cell2` shall match because there is no symmetry between `north` and `south` in `Site2`.

## 9.13 ANTENNA statement

This section defines the `SITE` statement and its use.

### 9.13.1 Definition

The `ANTENNA` statement shall be defined as follows:

```
antenna ::=
          ANTENNA [ antenna_identifier ] { antenna_items }

antenna_items ::=
          antenna_item { antenna_item }

antenna_item ::=
          all_purpose_item
   |      arithmetic_model
   |      arithmetic_model_container
```

The syntax and semantics of `all_purpose_item`, `arithmetic_model_container`, and `arithmetic_model` are already defined in defined in Section 11.7 and Section 11.16.

The items applicable for ANTENNA are shown in Table 9-16.

**Table 9-16 Items for ANTENNA description**

| Item | Usable ALF statement | Scope | Comment |
|---|---|---|---|
| maximum allowed antenna size | `LIMIT { SIZE { MAX { ... } } }` | `LIBRARY, SUBLIBRARY CELL, PIN` | see Section 7.5, Section 8.1.2, Section 7.6.1, Section 8.9.1, and Section 9.13.2 |
| calculation method for antenna size | `SIZE { HEADER { ... } TABLE { ...}` or `SIZE [id] { HEADER { ... } EQUATION { ...}` | `LIBRARY, SUBLIBRARY` | see Section 8.1.3, and Section 9.13.2 |
| argument values for antenna size calculation | *argument* = *value* ; or *argument* = *value* { ... } | `CELL, PIN` | see Section 11.2and Section 9.13.2 |

The use of the keyword SIZE (see Section 8.1.3) in the context of ANTENNA is proposed to represent an abstract, dimensionless model of the antenna size. It is related to the area of the net which forms the antenna, but it is not necessary a measure of area. It can be a measure of area ratio as well. However, the arguments of the calculation function for antenna SIZE shall be measurable data, such as AREA, PERIMETER, LENGTH, THICKNESS, WIDTH, and HEIGHT of metal segments connected to the net. The argument also need an annotation defining the applicable LAYER for the metal segments.

A process technology can have more than one antenna rule calculation method. In this case, the *antenna*_identifier is mandatory for each rule.

Antenna rules apply for routing and cut layers connected to poly silicon and eventually to diffusion. The CONNECT_RULE statement in conjunction with the BETWEEN statement shall be used to specify the connected layers. Connectivity shall only be checked up to the highest layer appearing in the CONNECT_RULE statement. Connectivity through higher layers shall not be taken into account, since such connectivity does not yet exist in the state of manufacturing process when the antenna effect occurs.

### 9.13.2 Layer-specific antenna rules

Antenna rules can be checked individually for each layer. In this case, the SIZE model contains only two or three arguments: AREA of the layer or perimeter (calculated from the LENGTH and WIDTH) of the layer causing the antenna effect, the area of poly silicon, and, eventually, the area of diffusion.

Example:

```
ANTENNA individual_m1 {
   LIMIT { SIZE { MAX = 1000; } }
   SIZE {
      CONNECTIVITY {
         CONNECT_RULE = must_short; BETWEEN { metal1 poly }
      }
      CONNECTIVITY {
         CONNECT_RULE = cannot_short; BETWEEN { metal1 diffusion }
      }
      HEADER {
         AREA a1 { LAYER = metal1; }
         AREA a0 { LAYER = poly; }
      }
      EQUATION { a1 / a0 }
   }
ANTENNA individual_m2 {
   LIMIT { SIZE { MAX = 1000; } }
   SIZE {
      CONNECTIVITY {
         CONNECT_RULE = must_short; BETWEEN { metal2 poly }
      }
      CONNECTIVITY {
         CONNECT_RULE = cannot_short; BETWEEN { metal2 diffusion }
      }
      HEADER {
         AREA a2 { LAYER = metal2; }
         AREA a0 { LAYER = poly; }
      }
      EQUATION { a2 / a0 }
   }
}
```

### 9.13.3 All-layer antenna rules

Antenna rules can also be checked globally for all layers. In that case, the `SIZE` model contains area or perimeter of all layers as additional arguments.

Example:

```
ANTENNA global_m2_m1 {
   LIMIT { SIZE { MAX = 2000; } }
   SIZE {
      CONNECTIVITY {
         CONNECT_RULE = must_short;
         BETWEEN { metal2 metal1 poly }
      }
      CONNECTIVITY {
         CONNECT_RULE = cannot_short;
         BETWEEN { metal2 diffusion }
      }
      HEADER {
```

```
            AREA a2 { LAYER = metal1; }
            AREA a1 { LAYER = metal1; }
            AREA a0 { LAYER = poly; }
         }
         EQUATION { (a2 + a1) / a0 }
      }
   }
```

### 9.13.4 Cumulative antenna rules

Antenna rules can also be checked by accumulating the individual effect. In that case, the `SIZE` model can be represented as a nested arithmetic model, each of which contain the model of the individual effect.

Example:

```
ANTENNA accumulate_m2_m1 {
   LIMIT { SIZE { MAX = 3000; } }
   SIZE {
      HEADER {
         SIZE ratio1 {
            CONNECTIVITY {
               CONNECT_RULE = must_short;
               BETWEEN { metal1 poly }
            }
            CONNECTIVITY {
               CONNECT_RULE = cannot_short;
               BETWEEN { metal1 diffusion }
            }
            HEADER {
               AREA a1 { LAYER = metal1; }
               AREA a0 { LAYER = poly; }
            }
            EQUATION { a1 / a0 }
         }
         SIZE ratio2 {
            CONNECTIVITY {
               CONNECT_RULE = must_short;
               BETWEEN { metal2 poly }
            }
            CONNECTIVITY {
               CONNECT_RULE = cannot_short;
               BETWEEN { metal2 diffusion }
            }
            HEADER {
               AREA a2 { LAYER = metal2; }
               AREA a0 { LAYER = poly; }
```

```
            }
            EQUATION { a2 / a0 }
         }
      }
      EQUATION { ratio1 + ratio2 }
   }
}
```

Note the arguments `a0` in `ratio1` and `ratio2` can are not the same. In `ratio1`, `a0` represents the area of poly silicon connected to `metal1` in a net. In `ratio2`, `a0` represents the area of poly silicon connected to `metal2` in a net, where the connection can be established through more than one subnet in `metal1`.

### 9.13.5 Illustration

Consider the structure shown in Figure 9-7.

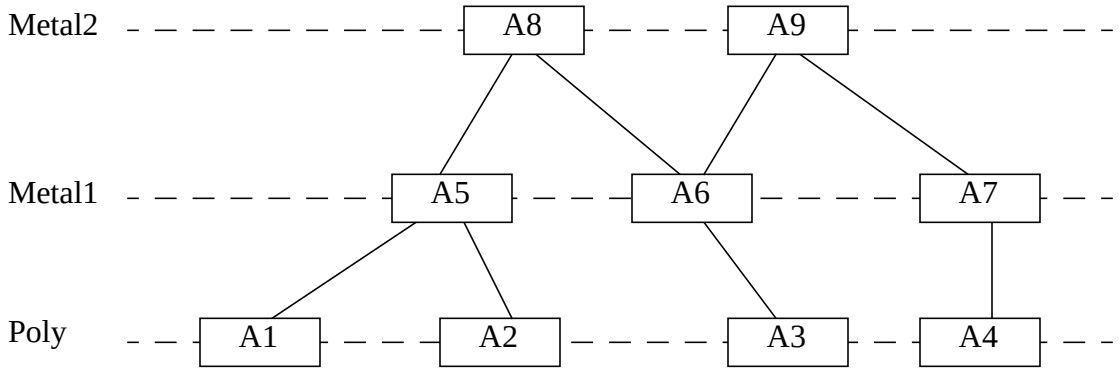



**Figure 9-7: Metal-poly illustration**

Checking this structure against the rules in the examples yields the following results:

```
individual_m1:
   1000 > A5 / (A1+A2)
   1000 > A6 / A3
   1000 > A7 / A4
individual_m2:
   1000 > (A8+A9) / (A1+A2+A3+A4)

global_m2_m1:
   2000 > (A8+A9+A5+A6+A7) / (A1+A2+A3+A4)

accumulate_m2_m1:
   3000 > (A8+A9) / (A1+A2+A3+A4) + A5 / (A1+A2)
   3000 > (A8+A9) / (A1+A2+A3+A4) + A6 / A3
   3000 > (A8+A9) / (A1+A2+A3+A4) + A7 / A4
```

# 9.14 ARRAY Statement

This section defines the ARRAY statement and its use.

## 9.14.1 Definition

The ARRAY statement shall be defined as follows:

```
array ::=
          ARRAY identifier {
                all_purpose_items
                geometric_transformations
          }
```

The geometric_transformations define the locations of the starting points within the array and the number of repetitions of the components of the array. Details are defined in the next section.

## 9.14.2 PURPOSE annotation

Each array shall have a PURPOSE assignment.

```
array_purpose_assignment ::=
          PURPOSE = array_purpose_identifier ;

array_purpose_identifier :: =
          floorplan
    |     placement
    |     global
    |     routing
```

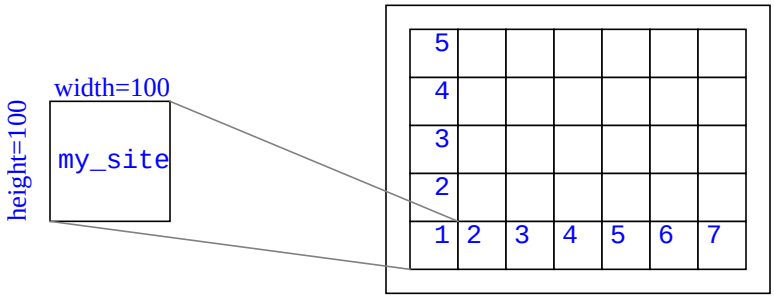
An array with purpose **floorplan** or **placement** shall have a reference to a SITE and a *shift*_annotation_container , *rotate*_annotation, and eventually a *flip*_annotation to define the location and orientation of the SITE in the context of the array.

An array with purpose **routing** shall have a reference to one or more routing LAYERs and a *shift*_annotation_container to define the location of the starting point.

An array with purpose **global** shall have a *shift*_annotation_container to define the location of the starting point.

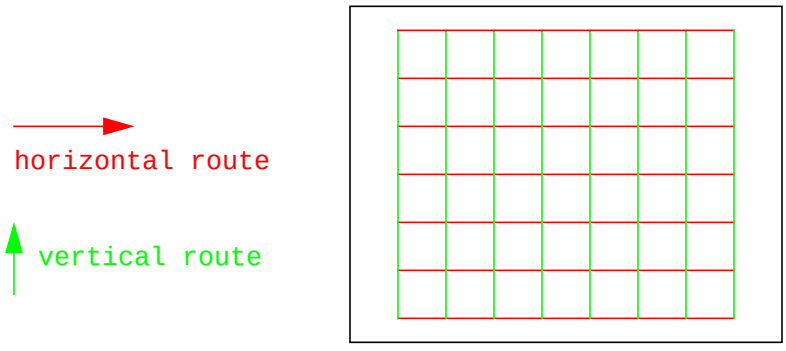
### 9.14.3 Examples

Example 1:



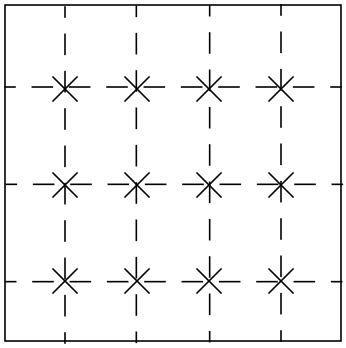

```
ARRAY grid_for_my_site {
   PURPOSE = placement;
   SITE = my_site;
   SHIFT { HORIZONTAL = 50; VERTICAL = 50; }
   REPEAT = 7 {
      SHIFT { HORIZONTAL = 100; }
      REPEAT = 5 {
         SHIFT { VERTICAL = 5; }
      }
   }
}
```

Example 2:



```
ARRAY grid_for_detailed_routing {
   PURPOSE = routing;
   LAYER { metal1 metal2 metal3 }
   SHIFT { HORIZONTAL = 100; VERTICAL = 50; }
   REPEAT = 7 {
      SHIFT { VERTICAL = 100; }
      REPEAT = 8 {
         SHIFT { HORIZONTAL = 100; }
      }
   }
}
```

Example 3:



```
ARRAY grid_for_global_routing {
   PURPOSE = global;
   SHIFT { HORIZONTAL = 100; VERTICAL = 100; }
   REPEAT = 3 {
      SHIFT { VERTICAL = 150; }
      REPEAT = 4 {
         SHIFT { HORIZONTAL = 100; }
      }
   }
}
```

# 9.15 CONNECTIVITY statement

This section defines the `CONNECTIVITY` statement and its use.

## 9.15.1 Definition

A `CONNECTIVITY` statement shall have the following form:

```
connectivity ::=
        CONNECTIVITY [ identifier ] {
              connect_rule_annotation
              between_multi_value_assignment
        }
```

```
|       CONNECTIVITY [ identifier ] {
              connect_rule_annotation
              table_based_model
        }
```

### 9.15.2 CONNECT_RULE annotation

The *connect_rule annotation* can be only inside a `CONNECTIVITY` object. It specifies the connectivity requirement.

```
CONNECT_RULE = string ;
```

which can take the values shown in Table 9-17.

**Table 9-17 : CONNECT_RULE annotation**

| Annotation string | Description |
|---|---|
| `must_short` | electrical connection required |
| `can_short` | electrical connection allowed |
| `cannot_short` | electrical connection disallowed |

It is not necessary to specify more than one rule between a given set of objects. If one rule is specified to be *True*, the logical value of the other rules can be implied shown in Table 9-18.

**Table 9-18 : Implications between connect rules**

| must_short | cannot_short | can_short |
|---|---|---|
| *False* | *False* | *True* |
| *False* | *True* | *False* |
| *True* | *False* | N/A |

### 9.15.3 CONNECTIVITY modeled with BETWEEN statement

The `BETWEEN` statement specifies the objects for which the connectivity applies.

```
between_multi_value_assignment ::=
        BETWEEN { identifiers }
```

If the `BETWEEN` statement contains only one identifier, than the `CONNECTIVITY` shall apply between multiple instances of the same object.

Example:

```
CLASS analog_power;
CLASS analog_ground;
CLASS digital_power;
CLASS digital_ground;
CONNECTIVITY Aground { // connect all members of CLASS analog_ground
   CONNECT_RULE = must_short;
   BETWEEN { analog_ground }
}
CONNECTIVITY Dground { // connect all members of CLASS digital_ground
   CONNECT_RULE = must_short;
   BETWEEN { digital_ground }
}
CONNECTIVITY Apower { // connect all members of CLASS analog_power
   CONNECT_RULE = must_short;
   BETWEEN { analog_power }
}
CONNECTIVITY Dpower { // connect all members of CLASS digital_power
   CONNECT_RULE = must_short;
   BETWEEN { digital_power }
}
CONNECTIVITY Aground2Dground {
   CONNECT_RULE = must_short;
   BETWEEN { analog_ground digital_ground }
}
CONNECTIVITY Apower2Dpower {
   CONNECT_RULE = can_short;
   BETWEEN { analog_power digital_power }
}
CONNECTIVITY Apower2Aground {
   CONNECT_RULE = cannot_short;
   BETWEEN { analog_power analog_ground }
}
CONNECTIVITY Apower2Dground {
   CONNECT_RULE = cannot_short;
   BETWEEN { analog_power digital_ground }
}
CONNECTIVITY Dpower2Aground {
   CONNECT_RULE = cannot_short;
   BETWEEN { digital_power analog_ground }
}
CONNECTIVITY Dpower2Dground {
   CONNECT_RULE = cannot_short;
   BETWEEN { digital_power digital_ground }
}
```

### 9.15.4 CONNECTIVITY modeled as lookup TABLE

Connectivity can also be described as a lookup table model. This description is usually more compact than the description using the `BETWEEN` statements.

The connectivity model can have the arguments shown in Table 9-19 in the HEADER.

**Table 9-19 : Arguments for connectivity**

| Argument | Value type | Description |
|---|---|---|
| DRIVER | string | argument of connectivity function |
| RECEIVER | string | argument of connectivity function |

Each argument shall contain a TABLE.

The connectivity model specifies the allowed and disallowed connections amongst drivers or receivers in one-dimensional tables or between drivers and receivers in two-dimensional tables.The boolean literals in the table refer to the CONNECT_RULE as shown in Table 9-20.

**Table 9-20 : Boolean literals in non-interpolateable tables**

| Boolean literal | Description |
|---|---|
| 1 | CONNECT_RULE is *True* |
| 0 | CONNECT_RULE is *False* |
| ? | CONNECT_RULE does not apply |

Example:

```
CLASS analog_power;
CLASS analog_ground;
CLASS digital_power;
CLASS digital_ground;
CONNECTIVITY all_must_short {
   CONNECT_RULE = must_short;
   HEADER {
      RECEIVER r1 {
         TABLE {analog_ground analog_power digital_ground digital_power}
      }
      RECEIVER r2 {
         TABLE {analog_ground analog_power digital_ground digital_power}
      }
   }
   TABLE {
      1 0 1 0
      0 1 0 0
      1 0 1 0
      0 0 0 1
   }
/*
The following table would apply, if the CONNECT_RULE was "cannot_short":
   TABLE {
      0 1 0 1
      1 0 1 0
      0 1 0 1
      1 0 1 0
   }
The following table would apply, if the CONNECT_RULE was "can_short":
```

```
    TABLE {
       ? 0 ? 0
       0 ? 0 ?
       ? 0 ? 0
       0 ? 0 ?
    }
*/
}
```

# 9.16 Physical annotations for CELL

This section defines the physical annotations for a CELL.

## 9.16.1 PLACEMENT_TYPE annotation

A CELL can contain the following PLACEMENT_TYPE statement:

```
placement_type_assignment ::=
          PLACEMENT_TYPE = placement_type_identifier ;

placement_type_identifier ::=
          pad
    |     core
    |     ring
    |     block
    |     connector
```

The identifiers have the following definitions:

- *pad*: I/O pad, to be placed in the I/O rows
- *core*: regular macro, to be placed in the core rows
- *block*: hierarchical block with regular power structure
- *ring*: macro with built-in power structure
- *connector*: macro at the end of core rows connecting with power or ground

## 9.16.2 Reference of a SITE by a CELL

A CELL can point to one or more legal placement SITEs.

Example:

```
CELL my_cell {
   SITE { my_site /* fill in other sites, if applicable */ }
   /* fill in contents of cell definition */
}
```

# 9.17 Physical annotations for PIN

This section defines the physical annotations for a PIN.

## 9.17.1 CONNECT_CLASS annotation

**CONNECT_CLASS { *class*_identifiers }**

annotates a declared class object for connectivity determination.

Connectivity rules involving those classes shall apply for the pin.

## 9.17.2 SIDE annotation

**SIDE = string ;**

which can take the values shown in Table 9-21.

**Table 9-21 : SIDE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| left | pin is on the left side |
| right | pin is on the right side |
| top | pin is at the top |
| bottom | pin is at the bottom |

## 9.17.3 ROW and COLUMN annotation

The following annotation shall be used for a pin in order to indicate the location of the pin within a placement row or column:

```
row_assignment ::=
        ROW = unsigned ;

column_assignment ::=
        COLUMN = unsigned ;
```

where *row*_assignment applies for pins with **SIDE = right | left** and *column*_assignment applies for pins with **SIDE = top | bottom**.

For bus pins, *row*_assignment and *column*_assignment shall have the form of multi_value_assignments.

```
row_multi_value_assignment ::=
        ROW { unsigned { unsigned } }

column_multi_value_assignment ::=
        COLUMN { unsigned { unsigned } }
```

### 9.17.4 ROUTING_TYPE annotation

A PIN can contain the following ROUTING_TYPE statement:

```
routing_type_assignment ::=
        ROUTING_TYPE = routing_type_identifier ;

routing_type_identifier ::=
        regular
   |    abutment
   |    ring
   |    feedthrough
```

The identifiers have the following definitions:

- *regular*: connection by regular routing
- *abutment*: connection by abutment, no routing
- *ring*: pin forms a ring around the block with connection allowed to any point of the ring
- *feedthrough*: both ends of the pin align and can be used for connection

## 9.18 Physical annotations for arithmetic models

This section defines the physical annotations for arithmetic models.

### 9.18.1 BETWEEN statement within DISTANCE

The BETWEEN statement within DISTANCE shall identify the objects for which the distance measurement applies.

```
between_multi_value_assignment ::=
        BETWEEN { identifiers }
```

If the BETWEEN statement contains only one identifier, than the DISTANCE shall apply between multiple instances of the same object.

### 9.18.2 MEASUREMENT annotation for DISTANCE

The following statement shall specify how the distance between objects is measured.

```
distance_measurement_assignment ::=
        MEASUREMENT = distance_measurement_identifier ;

distance_measurement_identifier ::=
        straight
   |    horizontal
   |    vertical
   |    manhattan
```

The default is **straight**.

The mathematical definitions for distance measurements between two points with differential coordinates Δx and Δy are:

- straight distance = $(\Delta x^2 + \Delta y^2)^{1/2}$
- horizontal distance = Δx
- vertical distance = Δy
- manhattan distance = Δx + Δy

### 9.18.3 REFERENCE annotation for DISTANCE

The *reference*_annotation shall specify the reference for distance measurements between objects, as shown in Figure 9-8.

```
reference_annotation ::=
        REFERENCE = reference_identifier ;

reference_identifier ::=
        center
    |   origin
    |   edge
```

The default shall be **edge**. The value **center** is only applicable for objects with EXTENSION, whereas the value **edge** is applicable for any physical object. The value **origin** is only applicable for objects with specified coordinates.

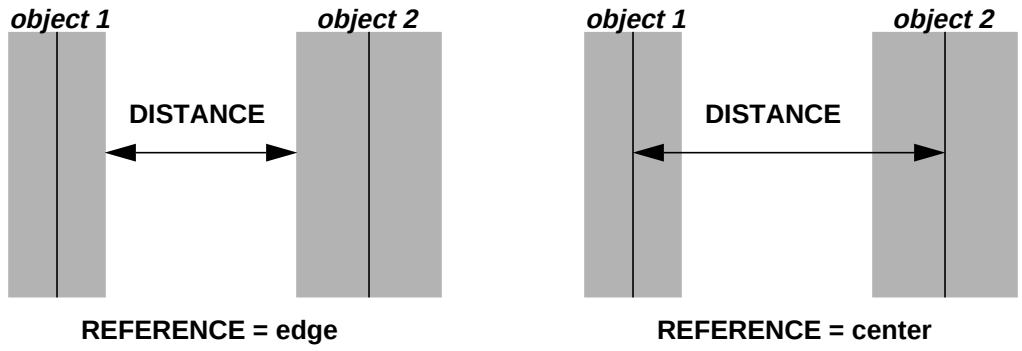



**Figure 9-8: Illustration of REFERENCE for DISTANCE**

### 9.18.4 Reference to ANTENNA within SIZE, AREA, and PERIMETER

In hierarchical design, a PIN with physical PORTs can be abstracted. Therefore, an arithmetic model for size, area, perimeter etc. relevant for certain antenna rules can be precalculated. The

following statement within the arithmetic model enables references to the set of antenna rules for which the arithmetic model applies.

```
antenna_reference_multi_value_assignment ::=
          ANTENNA { antenna_identifiers }
```

Example:

```
CELL cell1 {
   PIN pin1 {
      AREA poly_area = 1.5 {
         LAYER = poly;
         ANTENNA { individual_m1 individual_via1 }
      }
      AREA m1_area = 1.0 {
         LAYER = metal1;
         ANTENNA { individual_m1 }
      }
      AREA via1_area = 0.5 {
         LAYER = via1;
         ANTENNA { individual_via1 }
      }
   }
}
```

The area `poly_area` is used in the rules `individual_m1` and `individual_via1`.
The area `m1_area` is used in the rule`individual_m1` only.
The area `via1_area` is used in the rule `individual_via1` only.

The case with diffusion is illustrated in the following example:

```
CELL my_diode {
   CELLTYPE = special; ATTRIBUTE { DIODE }
   PIN my_diode_pin {
      AREA = 3.75 {
         LAYER = diffusion;
         ANTENNA { rule1_for_diffusion rule2_for_diffusion }
      }
   }
}
```

# Section 10
# Lexical Rules

This section discusses the lexical rules.

## 10.1 Cross-reference of lexical tokens

Table 10-1 cross-references the lexical tokens used in ALF.

**Table 10-1 : Cross-reference of lexical tokens**

| Lexical token | Section | | Lexical token | Section |
|---|---|---|---|---|
| alphabetic_bit_literal | 10.3.4 | | integer | 10.3.3 |
| any_character | 10.2.3 | | nonescaped_identifier | 10.3.8 |
| based_literal | 10.3.5 | | non_negative_number | 10.3.3 |
| binary_base | 10.3.5 | | nonreserved_character | 10.2.3 |
| binary_digit | 10.3.5 | | number | 10.3.3 |
| bit_edge_literal | 10.3.6 | | numeric_bit_literal | 10.3.4 |
| bit_literal | 10.3.4 | | octal_base | 10.3.5 |
| block_comment | 10.3.2 | | octal_digit | 10.3.5 |
| comment | 10.3.2 | | placeholder_identifier | 10.3.8 |
| decimal_base | 10.3.5 | | quoted_string | 10.3.7 |
| delimiter | 10.3.1 | | reserved_character | 10.2.3 |
| digit | 10.2.3 | | sign | 10.3.3 |
| dont_care_literal | 10.3.4 | | single_line_comment | 10.3.2 |
| edge_literal | 10.3.6 | | symbolic_edge_literal | 10.3.6 |
| escape_character | 10.2.3 | | unsigned | 10.3.3 |
| escaped_identifier | 10.3.8 | | whitespace | 10.2.2 |
| hex_base | 10.3.5 | | word_edge_literal | 10.3.6 |
| hex_digit | 10.3.5 | | | |

## 10.2 Characters

This section defines the use of characters in ALF.

### 10.2.1 Character set

Each graphic character corresponds to a unique code of the ISO eight-bit coded character set [ISO 8859-1 : 1987(E)] and is represented (visually) by a graphical symbol.

### 10.2.2 Whitespace characters

The characters shown in Figure 10-1shall be considered *whitespace characters*:

| Character | ASCII code (hex) |
|---|---|
| space | 20 |
| vertical tab | 0B |
| horizontal tab | 09 |
| line feed (new line) | 0A |
| carriage return | 0D |
| form feed | 0C |

**Figure 10-1: List of whitespace characters**

Comments are also considered white space (see Section 10.3.2).

A whitespace character shall be ignored except when it separates other lexical tokens or when it appears in a quoted string.

### 10.2.3 Reserved and non-reserved characters

The ASCII character set shall be divided in three categories: whitespace (see Section 10.2.2), reserved characters, and non-reserved characters. The reserved characters are symbols that make up punctuation marks and operators. The non-reserved characters shall be used for creating identifiers and numbers. Both are shown in Figure 10-2.

```
reserved_character ::=
     & | | | ^ | ~ | + | - | * | / | % | ? | ! | = | < | > | :
   | ( | ) | [ | ] | { | } | @ | ; | , | . | ” | ’

nonreserved_character ::=
     letter | digit | _ | $ | #

letter ::=
     a | b | c | d | e | f | g | h | i | j | k | l | m
   | n | o | p | q | r | s | t | u | v | w | x | y | z
   | A | B | C | D | E | F | G | H | I | J | K | L | M
   | N | O | P | Q | R | S | T | U | V | W | X | Y | Z

digit ::=
   0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

escape_character ::=
     \

any_character ::=
     reserved_character
   | nonreserved_character
```

```
    | escape_character
    | whitespace
```

**Figure 10-2: Reserved and non-reserved characters**

ALF treats uppercase and lowercase characters as the same characters. In other words, ALF is a *case-insensitive language*.

Notes:

The characters $ and # can be reserved in other languages, such as VERILOG. Therefore, if translation from ALF into VERILOG is required, these characters shall not be used for items which need to be translated, e.g., the names of cells and pins.
Other languages can be case-sensitive, such as VERILOG. Therefore, if translation from ALF into VERILOG is required, the case of the name used in the declaration of the object, e.g., the name of a cell or a pin, shall always be preserved as a reference. For example, if the name of a cell is declared as "**MyCell**", reference to the cell can be made as "**MYCELL**" or "**mycell**". However, it shall always be translated into VERILOG as "**MyCell**".

## 10.3 Lexical tokens

The ALF source text files shall be a stream of lexical tokens. Each lexical token is either a *delimiter*, a *comment*, a *bit literal*, a *based literal*, an *edge literal*, a *number*, a *quoted string* or an *identifier*.

### 10.3.1 Delimiters

A *delimiter* is either a reserved character or one of the following compound operators, each composed of two or three adjacent reserved characters, as shown in Figure 10-3.

```
delimiter ::=
     reserved_character
   | && | ~& | || | ~| | ~^ | == | != | ** | >= | <=
   | ?! | ?~ | ?- | ?? | -> | <-> | &> | <&> | >> | <<
```

**Figure 10-3: Tokens that make up delimiters**

Each special character in a single character delimiter list shall be a single delimiter unless this character is used as a character in a compound operator or as a character in a quoted string.

### 10.3.2 Comments

ALF has two forms to introduce comments, as shown in Figure 10-4.

A *single-line comment* shall start with the two characters `//` and end with a new line.

A *block comment* shall start with `/*` and end with `*/`. Comments shall not be nested. The single-line comment token `//` shall not have any special meaning in a block comment.

```
comment ::=
     single_line_comment
   | block_comment
```

**Figure 10-4: Single-line and block comments**

### 10.3.3 Numbers

Constant numbers can be specified as integer or real, as shown in Figure 10-5.

```
sign      ::=  +  |  -

unsigned ::=  digit  { _ | digit }

integer ::=  [ sign ] unsigned

non_negative_number ::=
     unsigned [ .  unsigned ]
   | unsigned [ . unsigned ] E [ sign ] unsigned

number ::=
     [ sign ] non_negative_number
```

**Figure 10-5: Integer and real numbers**

The *integer* is a decimal integer constant.

### 10.3.4 Bit literals

A *bit* literal shall represent a single bit constant, as shown in Table 10-2.

```
bit_literal ::=
     numeric_bit_literal
   | alphabetic_bit_literal
   | dont_care_literal
   | random_literal

numeric_bit_literal ::= 0 | 1

alphabetic_bit_literal ::=
     X | Z | L | H | U | W
   | x | z | l | h | u | w

dont_care_literal ::= ?

random_literal ::= *
```

**Table 10-2 : Single bit constants**

| Literal | Description |
|---|---|
| **0** | value is logic zero |
| **1** | value is logic one |
| **X** or **x** | value is unknown |
| **L** or **l** | value is logic zero with weak drive strength |
| **H** or **h** | value is logic one with weak drive strength |
| **W** or **w** | value is unknown with weak drive strength |
| **Z** or **z** | value is high-impedance |
| **U** or **u** | value is uninitialized |
| **?** | value is any of the above, yet stable |
| ***** | value may randomly change |

### 10.3.5 Based literals

A *based literal* is a constant expressed in a form that specifies the base explicitly. The base can be specified in *binary*, *octal*, *decimal* or *hexadecimal* format, as shown in Figure 10-6.

```
based_literal ::=
     binary_base { _ | binary_digit }
   | octal_base { _ | octal_digit }
   | decimal_base { _ | digit }
   | hex_base { _ | hex_digit }

binary_base ::=
   'B | 'b

octal_base ::=
   'O | 'o

decimal_base ::=
   'D | 'd

hex_base ::=
   'H | 'h

binary_digit ::=
   bit_literal

octal_digit ::=
   binary_digit | 2 | 3 | 4 | 5 | 6 | 7
```

```
hex_digit ::=
   octal_digit | 8 | 9 | A | B | C | D | E | F | a | b | c | d | e | f
```

**Figure 10-6: Based constants**

The underscore **(**_) shall be legal anywhere in the number, except as the first character and this character is ignored. This feature can be used to break up long numbers for readability purposes. No white space shall be allowed between base and digit token in a based literal.

When an alphabetic bit literal is used as an octal digit, it shall represent three repeated bits with the same literal. When an alphabetic bit literal is used as a hex digit, it shall represent four repeated bits with the same literal.

For example,

| | | |
|---|---|---|
| `'o2xw0u` | is same as | `'b010_xxx_www_000_uuu` |
| `'hLux` | is same as | `'bLLLL_uuuu_xxxx` |

### 10.3.6 Edge literals

An *edge literal* shall be constructed by two bit literals or two based literals, as shown in Figure 10-7. It shall describe the transition of a signal from one discrete value to another. No white space shall be allowed within (between) the two literals. An underscore can be used.

```
edge_literal ::=
     bit_edge_literal
   | word_edge_literal
   | symbolic_edge_literal

bit_edge_literal ::=
     bit_literal  bit_literal

word_edge_literal ::=
     based_literal  based_literal

symbolic_edge_literal::=
     ?? | ?~ | ?! | ?-
```

**Figure 10-7: Edge literals**

### 10.3.7 Quoted strings

The *quoted string* shall be a sequence of zero or more characters enclosed between two quotation marks (") and contained on a single line, as shown in Figure 10-8. Character *escape codes* are used inside the string literal to represent some common special characters.

```
quoted_string ::=
     " { any_character } "
```

**Figure 10-8: A quoted string**

The characters which can follow the backslash (\) and their meanings are listed in Table 10-3.

**Table 10-3 : Special characters in quoted strings**

| Symbol | ASCII Code (octal) | Meaning |
|---|---|---|
| \g | 007 | alert/bell |
| \h | 010 | backspace |
| \t | 011 | horizontal tab |
| \n | 012 | new line |
| \v | 013 | vertical tab |
| \f | 014 | form feed |
| \r | 015 | carriage return |
| \" | 042 | double quotation mark |
| \\ | 134 | backslash |
| **\ddd** | | 3-digit octal value of ASCII character |

A non-quoted string can not contain any reserved character. Therefore, use of a quoted string is necessary when referencing file names (which typically contain a dot (.) character).

### 10.3.8 Identifiers

*Identifiers* are used in ALF as names of objects, reserved words, and context-sensitive keywords. An identifier shall be any sequence of letters, digits, underscore (_), and dollar sign ($) character. If an identifier is constructed from one or more non-reserved characters, it is called *non-escaped identifier*. A `digit` shall not be allowed as first character of a non-escaped identifier.

```
nonescaped_identifier ::=
   nonreserved_character { nonreserved_character }
```

A sequence of characters starting with an `escape_character` is called an *escaped identifier*. The escaped identifier legalizes the use of a `digit` as first character of an identifier and the use of `reserved_character` anywhere in an identifier. Or it can be used to prevent the misinterpretation of an identifier as a keyword. The escape character shall be followed by at least one non-white space character to form an escaped identifier. The escaped identifier shall contain all characters up to first white space character.

```
escaped_identifier ::=
   escape_character escaped_characters

escaped_characters ::=
   escaped_character { escaped_character }
```

```
escaped_character ::=
     nonreserved_character
   | reserved_character
   | escape_character
```

A *placeholder identifier* shall be a non-escaped identifier between the less-than character (<) and the greater-than character (>). No whitespace or delimiters are allowed between the non-escaped identifier and the placeholder characters (< and >). The placeholder identifier is used in template objects as a formal parameter, which is replaced by the actual parameter in template instantiation.

```
placeholder_identifier ::=
   < nonescaped_identifier >
```

Identifiers are treated in a case-insensitive way. They can be used in the definition of objects and in reference to already defined objects. A parser should preserve the case of an identifier in the definition of an object, since a downstream application could be case-sensitive.

### 10.3.9 Hierarchical identifier

A hierarchical identifier shall be defined as follows:

```
hierarchical_identifier ::=
          identifier . { identifier . } identifier
```

with no whitespace in-between.

The dot (`.`) shall take precedence over the `escape_character`. In order to escape the dot, the `escape_character` shall be placed directly in front of it.

Examples:

```
\id1.id2 //Only id1 is escaped.
id1\.id2 //Only the dot is escaped.
id1.\id2 //Only id2 is escaped.
```

## 10.4 Keywords

Keywords are case-insensitive non-escaped identifiers. For clarity, this document uses uppercase letters for keywords and lowercase letters elsewhere, unless otherwise mentioned.

Keywords are reserved for use as object identifiers, not for general symbols. To use an identifier that conflicts with the list of keywords, use the escape character, e.g., to declare a pin that is called `PIN`, use the form:

```
PIN \PIN {..}
```

A keyword can either be a *reserved keyword* (also called *hard keyword*) or a *context-sensitive keyword* (also called *soft keyword*). The hard keywords have fixed meaning and shall be understood by any parser of ALF. The soft keywords might be understood only by specific

applications. For example, a parser for a timing analysis application can ignore objects that contain power related information described using soft keywords.

### 10.4.1 Keywords for objects

The keywords shown in Figure 10-9 are used to identify object types:

```
ALIAS          ATTRIBUTE      BEHAVIOR       CELL
CLASS          CONSTANT       EQUATION       FUNCTION
GROUP          HEADER         INCLUDE        LIBRARY
PIN            PRIMITIVE      PROPERTY       STATETABLE
SUBLIBRARY     TABLE          TEMPLATE       VECTOR
WIRE
```

**Figure 10-9: Keywords for objects**

### 10.4.2 Keywords for operators

The keywords shown in Figure 10-10 are used for built-in arithmetic functions:

```
ABS            absolute value
EXP            natural exponential function
LOG            natural logarithm
MIN            minimum
MAX            maximum
```

**Figure 10-10: Keywords for built-in arithmetic functions**

### 10.4.3 Context-sensitive keywords

In order to address the need of extensible modeling, ALF provides a predefined set of *public* context-sensitive keywords. Additional private context-sensitive keywords can be introduced as long as they do not have the same name as any existing public keyword.

The public context-sensitive keywords and their semantic meaning are defined in Section 5.6. This set can be extended to include private context-sensitive keywords.

## 10.5 Rules against parser ambiguity

The following rules shall apply when resolving ambiguity in parsing ALF source:

- In a context where both `bit_literal` and `identifier` are legal syntax items, `nonescaped_identifier` shall take priority over `alphabetic_bit_literal`.
- In a context where both `bit_literal` and `number` are legal syntax items, `number` shall take priority over `numeric_bit_literal`.
- In a context where both `edge_literal` and `identifier` are legal syntax items, `identifier` shall take priority over `bit_edge_literal`.

- In a context where both `edge_literal` and `number` are legal syntax items, `number` shall take priority over `bit_edge_literal`.

In such contexts, `based_literal` shall be used instead of `bit_literal`.

# Section 11

# Syntax Rules

This section discusses the syntactical rules. The formal syntax of ALF language is described using Backus-Naur Form (BNF).

## 11.1 Cross-reference of BNF items

A BNF item with a singular name is defined in the same section as the BNF item with the plural name. A plural item name implies one or more items with the corresponding singular name.

**Table 11-1 : Cross-reference of BNF items with short semantic explanation**

| BNF item | Section | Short semantic explanation |
|---|---|---|
| alias | 11.8 | statement defining an alias |
| all_purpose_item(s) | 11.7 | item(s) that can appear inside any hierarchical object |
| annotation | 11.7 | parameter-value assignment inside an object, may be nested |
| annotation_container | 11.7 | unnamed object containing annotations |
| antenna | new | statement describing a set of process antenna rules |
| arithmetic_assignment | 11.2 | statement assigning an arithmetic expression to a variable |
| arithmetic_binary | 11.6 | arithmetic operator requiring two operands |
| arithmetic_expression | 11.3 | expression involving arithmetic operations |
| arithmetic_function_operator | 11.6 | arithmetic operator prefixing a list of arguments |
| arithmetic_model(s) | 11.16 | statement(s) for description of characterization data using single numbers, tables or equations |
| arithmetic_model_item(s) | 11.16 | statement(s) inside arithmetic model statement |
| arithmetic_model_container | 11.16 | unnamed object containing arithmetic models |
| arithmetic_submodel(s) | 11.16 | statement(s) inside an arithmetic model statement for categorizing the characterization data |
| arithmetic_submodel_item(s) | 11.16 | statement(s) inside arithmetic submodel statement |
| arithmetic_unary | 11.6 | arithmetic operator requiring one operand |
| array | new | statement describing a regular floorplan or placement or routing definition in gate array technology |
| artwork | new | statement making reference between the cell in the library and the corresponding description in the layout (GDS2) database |
| assignment | 11.2 | terminated statement for single value assignment to an object |

**Table 11-1 : Cross-reference of BNF items with short semantic explanation, *continued***

| BNF item | Section | Short semantic explanation |
|---|---|---|
| assignment_base | 11.2 | unterminated statement for single value assignment to an object |
| attribute | 11.8 | statement associating attributes to an object |
| attribute_item(s) | 11.8 | item(s) inside an attribute statement |
| behavior | 11.17 | statement describing the logic function of a digital circuit in a behavioral language |
| behavior_item(s) | 11.17 | item(s) inside a behavior statement |
| blockage | new | statement describing routing obstructions |
| boolean_and | 11.6 | boolean AND operator |
| boolean_arithmetic | 11.6 | operator for boolean arithmetic |
| boolean_binary | 11.6 | boolean operator requiring two operands |
| boolean_case_compare | 11.6 | binary boolean operator for magnitude comparison |
| boolean_cond | 11.6 | boolean postfix operator evaluating the preceding boolean expression (if-clause) |
| boolean_else | 11.6 | boolean infix operator separating if-and else-clauses |
| boolean_expression | 11.3 | expression involving boolean operations |
| boolean_logic_compare | 11.6 | binary boolean operator for logic comparison |
| boolean_or | 11.6 | boolean OR operator |
| boolean_unary | 11.6 | boolean operator requiring one operand |
| cell(s) | 11.9 | statement(s) describing the entire model of a digital or analog circuit |
| cell_item(s) | 11.9 | item(s) inside a cell statement |
| cell_instantiation | 11.4 | statement inside a cell, describing a reference to another cell with pin-to-pin correspondence |
| class | 11.8 | statement describing a class for the use of reference and inheritance by other objects |
| class_item(s) | 11.7 | item(s) inside a class statement, which will be inherited by any object referring to the class |
| connectivity | new | statement describing a set of electrical connectivity rules |
| constant | 11.8 | statement defining a numeric constant |
| context_sensitive_keyword | 11.5 | identifier of an object for which the semantic meaning is established by its context |
| coordinates | new | statement containing numbers representing the coordinates of reference points within a physical object |
| dynamic_instantiation_item(s) | 11.4 | item(s) inside a dynamic instantiation of a template |
| edge_literal(s) | 11.5 | symbol(s) describing a transition between two states |
| equation | 11.16 | statement inside arithmetic model containing an arithmetic expression for the calculation of characterization data |
| from | 11.16 | statement inside arithmetic model defining start point of timing measurement |

**Table 11-1 : Cross-reference of BNF items with short semantic explanation, *continued***

| BNF item | Section | Short semantic explanation |
|---|---|---|
| function | 11.17 | statement describing the logic function of a circuit in a canonical way, using behavior and/or statetable statement |
| generic_object | 11.7 | statement with the sole purpose of being used by other objects |
| geometric_model(s) | new | statement(s) describing the form of a physical object |
| geometric_transformation(s) | new | statement(s) describing shift, rotate, flip or repetition of a physical object |
| group | 11.8 | statement allowing expansion of one object to multiple objects |
| header | 11.16 | statement inside arithmetic model containing a list of parameters of the arithmetic model |
| identifier(s) | 11.5 | literal(s) defining a keyword or a name or a string value |
| include | 11.8 | statement defining the inclusion of a file |
| index | 11.5 | symbol defining an integer or a range of integers for the use as indices |
| layer | new | statement describing the stackup information in a technology library |
| layer_item(s) | new | statement(s) inside a layer statement |
| library (libraries) | 11.10 | statement(s) describing the entire contents of a library |
| keyword_declaration | 11.8 | statement declaring a new keyword |
| library_item(s) | 11.10 | item(s) inside a library statement |
| library_specific_object | 11.7 | statement describing an object which is part of the library. Multiple statements describing objects of the same type may appear within a given context. |
| library_specific_singular_object | 11.7 | statement describing an object which is part of the library. Only one statement describing an object of a specific type may appear within a given context. |
| logic_assignment(s) | 11.2 | statement(s) assigning a logic expression to a logic variable |
| logic_value(s) | 11.5 | variable(s) or constant logic value(s) |
| logic_constant(s) | 11.5 | constant logic value(s) |
| logic_variable(s) | 11.5 | variable(s) containing a logic value |
| multi_value_assignment | 11.2 | statement for assignment of multiple values to an object |
| named_assignment | 11.2 | terminated statement for single value assignment to a named object |
| named_assignment_base | 11.2 | unterminated statement for single value assignment to a named object |
| named_cell_instantiation(s) | 11.4 | statement(s) describing a reference to another cell with pin-to-pin correspondence and instance name |
| node | 11.15 | statement declaring a node in a physical wireload model |
| non_scan_cell | 11.9 | statement inside a cell making reference to the cells which can be replaced by this cell in scan design |

**Table 11-1 : Cross-reference of BNF items with short semantic explanation, *continued***

| BNF item | Section | Short semantic explanation |
|---|---|---|
| number(s) | 11.5 | integer or floating point number(s) |
| pattern | new | statement describing a physical object |
| pin(s) | 11.11 | statement(s) describing a pin inside a cell |
| pin_group | 11.11 | statement defining a group of pins that can henceforth be referenced as a bus |
| pin_instantiation | 11.4 | statement inside a bus pin containing information pertinent to a subset of pins within that bus |
| pin_item(s) | 11.11 | item(s) inside a pin statement |
| pin_assignment(s) | 11.2 | statement(s) defining a correspondence between two pins or between a pin and a constant logic value |
| port | new | statement describing an electrical connection point within a pin |
| primitive(s) | 11.12 | statement(s) describing a technology-independent cell |
| primitive_instantiation | 11.4 | statement inside a behavior statement for logic function description by reference to a primitive |
| primitive_item(s) | 11.12 | item(s) inside a primitive statement |
| property | 11.8 | statement describing private properties without standardized semantics |
| property_item(s) | 11.8 | item(s) inside a property statement |
| range | 11.7 | definition of a contiguous range of integer numbers |
| repeat | new | statement defining how to duplicate a physical object several times |
| rule(s) | new | statement describing a set of design or extraction rules |
| sequential_else_if | 11.6 | operator indicating a lower-priority logic state or event sequence |
| sequential_if | 11.6 | operator indicating a top-priority logic state or event sequence |
| sequential_logic_statement | 11.2 | statement inside a behavior statement for logic function description with storage elements |
| site | new | statement describing a legal placement site |
| source_text | 11.7 | contents of a self-sufficient file in ALF |
| statetable(s) | 11.17 | statement(s) describing the logic function of a digital circuit in table format |
| statetable_body | 11.17 | list of values inside a statetable |
| statetable_header | 11.17 | list of variables inside a statetable |
| statetable_value(s) | 11.5 | literal(s) inside a statetable |
| string | 11.5 | identifier consisting of a restricted set of characters or quoted string containing arbitrary characters |
| structure | 11.17 | statement(s) describing the structure of a cell in form of a netlist |

**Table 11-1 : Cross-reference of BNF items with short semantic explanation, *continued***

| BNF item | Section | Short semantic explanation |
|---|---|---|
| sublibrary (sublibraries) | 11.13 | statement(s) describing the contents of a sub-library inside a library |
| table | 11.16 | statement inside arithmetic model containing a list of characterization data |
| table_item(s) | 11.16 | item(s) inside a table statement |
| template | 11.8 | statement defining an object with placeholders |
| template_instantiation | 11.4 | statement referring to a template and filling the place-holders |
| template_item(s) | 11.8 | statement(s) inside a template statement |
| test | new | statement containing information about the test algorithm applicable to the cell |
| to | 11.16 | statement inside arithmetic model defining end point of timing measurement |
| unnamed_assignment(s) | 11.2 | terminated statement(s) for single value assignment to an unnamed object |
| unnamed_assignment_base | 11.2 | unterminated statement for single value assignment to an unnamed object |
| unnamed_cell_instantiation(s) | 11.4 | statement(s) describing a reference to another cell with pin-to-pin correspondence without instance name |
| value(s) | lexical | number(s) or string(s) or logic value(s) |
| vector(s) | 11.14 | statement(s) describing event sequence and data for char-acterization of a circuit |
| vector_and | 11.6 | operator used for description of simultaneous events or simultaneous event sequences |
| vector_binary | 11.6 | operator requiring two operands used for description of event sequences |
| vector_boolean_and | 11.6 | operator used for description of event sequences with condition, one operand is an expression describing a complex event, other operand is a boolean expression |
| vector_boolean_cond | 11.6 | condition operator indicating if-clause |
| vector_boolean_else | 11.6 | condition operator separating if-and else-clauses |
| vector_complex_event | 11.3 | expression describing complex event sequences without condition |
| vector_conditional_event | 11.3 | expression describing complex event sequences with condition |
| vector_event_sequence | 11.3 | expression describing one event sequence |
| vector_expression | 11.3 | expression describing complex event sequences |
| vector_followed_by | 11.6 | operator used for description of subsequent events |
| vector_item(s) | 11.14 | item(s) inside a vector statement |
| vector_event | 11.3 | expression describing one single event or multiple simul-taneous events |
| vector_or_boolean_expression | 11.3 | a vector expression or a boolean expression |

**Table 11-1 : Cross-reference of BNF items with short semantic explanation, *continued***

| BNF item | Section | Short semantic explanation |
|---|---|---|
| vector_expression | 11.3 | expression describing complex event sequences |
| vector_single_event | 11.3 | expression describing one single event |
| vector_unary | 11.6 | operator requiring one operand used for description of event sequences |
| via | new | statement describing the construction of electrical connections accross layers |
| via_instantiation(s) | 11.4 | statement describing the correspondence between the model and the instance of a via |
| via_item(s) | new | statement(s) inside a via statement |
| via_reference | new | statement containing via_instantiation statements |
| violation | 11.16 | statement inside arithmetic model defining consequences of timing violation or illegal operation |
| wire(s) | 11.15 | statement(s) describing a wireload model |
| wire_item(s) | 11.15 | item(s) inside a wire statement |

## 11.2 Assignments

```
unnamed_assignment_base ::=
     context_sensitive_keyword = value

unnamed_assignment ::=
     unnamed_assignment_base  ;

unnamed_assignments ::=
     unnamed_assignment { unnamed_assignment }

named_assignment_base ::=
     context_sensitive_keyword identifier = value

named_assignment ::=
     named_assignment_base  ;

named_assignments ::=
     named_assignment { named_assignment }

assignment_base ::=
     named_assignment_base
   | unnamed_assignment_base

multi_value_assignment ::=
     identifier { values }

assignment ::=
     named_assignment
```

```
   | unnamed_assignment
   | multi_value_assignment

pin_assignment ::=
     pin_identifier [index] = pin_identifier [index] ;
   | pin_identifier [index] = logic_constant ;
   | logic_constant  = pin_identifier [index] ;

pin_assignments ::=
     pin_assignment { pin_assignment }

arithmetic_assignment ::=
     identifier = arithmetic_expression ;
```

## 11.3 Expressions

```
arithmetic_expression ::=
     ( arithmetic_expression )
   | number
   | [ arithmetic_unary ] identifier

   | arithmetic_expression arithmetic_binary
          arithmetic_expression
   |arithmetic_function_operator
          ( arithmetic_expression { , arithmetic_expression } )
   | boolean_expression ? arithmetic_expression :
          { boolean_expression ? arithmetic_expression : }
          arithmetic_expression

boolean_expression ::=
     ( boolean_expression )
   | logic_constant
   | logic_variable
   | boolean_unary boolean_expression
   | boolean_expression boolean_binary boolean_expression
   | boolean_expression
          boolean_cond boolean_expression boolean_else
          { boolean_expression boolean_cond boolean_else }
          boolean_expression

vector_single_event ::=
     ( vector_single_event )
   | vector_unary boolean_expression

vector_event ::=
     ( vector_event )
   | vector_single_event
   | vector_event vector_and vector_event

vector_event_sequence ::=
     ( vector_event_sequence )
```

```
   | vector_event
   | vector_event_sequence vector_followed_by vector_event_sequence

vector_complex_event ::=
     ( vector_complex_event )
   | vector_event_sequence
   | vector_complex_event vector_binary vector_complex_event

vector_conditional_event ::=
     vector_expression vector_boolean_and boolean_expression
   | boolean_expression vector_boolean_and vector_expression
   | boolean_expression vector_boolean_cond vector_expression
          vector_boolean_else
          { boolean_expression vector_boolean_cond vector_expression
          vector_boolean_else } vector_expression

vector_expression ::=
     ( vector_expression )
   | vector_complex_event
   | vector_conditional_event
   | vector_expression vector_binary vector_expression

vector_or_boolean_expression ::=
     vector_expression
   | boolean_expression
```

## 11.4 Instantiations

```
cell_instantiation ::=
          unnamed_cell_instantiation
   |      named_cell_instantiation

unnamed_cell_instantiations ::=
          unnamed_cell_instantiation { unnamed_cell_instantiation }

unnamed_cell_instantiation ::=
     cell_identifier { logic_values }
   | cell_identifier { pin_assignments }

named_cell_instantiations ::=
          named_cell_instantiation { named_cell_instantiation }

named_cell_instantiation ::=
          cell_identifier instance_identifier { logic_values }
   |      cell_identifier instance_identifier { pin_assignments }

pin_instantiation ::=
          pin_identifier [ index ] {
                pin_items
          }
```

```
primitive_instantiation ::=
     primitive_identifier [ identifier ] { logic_values }
   | primitive_identifier [ identifier ] { logic_assignments }
   | primitive_identifier [ identifier ] { pin_assignments }

template_instantiation ::=
     template_identifier ;
   | template_identifier [ = static ] { values }
   | template_identifier [ = static ] { all_purpose_items }
   | template_identifier  = dynamic { values }
   | template_identifier  = dynamic { dynamic_instantiation_items }

dynamic_instantiation_items ::=
   dynamic_instantiation_item { dynamic_instantiation_item }

dynamic_instantiation_item ::=
     all_purpose_item
   | arithmetic_model
   | arithmetic_assignment

via_instantiations ::=
          via_instantiation { via_instantiation }

via_instantiation ::=
          via_identifier { geometric_transformations }
```

## 11.5 Literals

```
context_sensitive_keyword ::=
     nonescaped_identifier

edge_literal ::=
     bit_edge_literal
   | word_edge_literal
   | symbolic_edge_literal

edge_literals::=
     edge_literal { edge_literal }

identifier ::=
     nonescaped_identifier
   | escaped_identifier
   | placeholder_identifier

identifiers ::=
   identifier { identifier }

index ::=
     [ unsigned ]
   | [ unsigned : unsigned ]
```

```
   | [ identifier ]
   | [ identifier : identifier ]

logic_value ::=
     logic_constant
   | logic_variable

logic_values ::=
     logic_value { logic_value }

logic_constant ::=
     bit_literal
   | based_literal

logic_constants ::=
    logic_constant { logic_constant }

statetable_value ::=
     logic_constant
   | edge_literal
   | ( [!] logic_variable )

statetable_values ::=
     statetable_value { statetable_value }

logic_variable ::=
     pin_identifier [ index ]

logic_variables ::=
     logic_variable { logic_variable }

numbers ::=
     number { number }

string ::=
     quoted_string
   | identifier

value ::=
     number
   | string
   | logic_value

values ::=
     value { value }
```

## 11.6 Operators

```
arithmetic_unary ::=
     + | -
```

```
arithmetic_binary ::=
   + | - | * | / | ** | %

arithmetic_function_operator ::=
     abs
   | exp
   | log
   | min
   | max

boolean_unary ::=
     ! | ~ | & | ~& | | | ~| | ^ | ~^

boolean_and ::=
     & | &&

boolean_or ::=
     | | ||

boolean_logic_compare ::=
     ^ | ~^

boolean_case_compare ::=
     != | == | >= | <= | > | <

boolean_arithmetic ::=
     + | - | * | / | % | >> | <<

boolean_binary ::=
     boolean_and
   | boolean_or
   | boolean_logic_compare
   | boolean_case_compare
   | boolean_arithmetic

boolean_cond ::=
     ?

boolean_else ::=
     :

vector_unary ::=
     edge_literal

vector_and ::=
     & | &&

vector_or ::=
     | | ||
```

```
vector_followed_by ::=
   -> | ~>

vector_binary ::=
     vector_and
   | vector_or
   | vector_followed_by
   | <->
   | &>
   | <&>

vector_boolean_and ::=
   & | &&

vector_boolean_cond ::=
   ?

vector_boolean_else ::=
   :

sequential_if ::=
     @

sequential_else_if ::=
     :
```

See Section 5.1.2, Section 5.1.3, and Section 7.2.2 for the semantics of these operators.

## 11.7 Auxiliary objects

```
all_purpose_item ::=
     annotation
   | annotation_container
   | generic_object
   | template_instantiation

all_purpose_items ::=
     all_purpose_item { all_purpose_item }

annotation ::=
     assignment
   | assignment_base { all_purpose_items }

annotation_container ::=
     context_sensitive_keyword { all_purpose_items }

generic_object ::=
     alias
   | attribute
   | constant
   | class
```

```
      | group
      | include
      | keyword_declaration
      | property
      | template

library_specific_object ::=
   cell
      | library
      | node
      | pin
      | pin_group
      | primitive
      | sublibrary
      | vector
      | wire
      | antenna
      | array
      | blockage
      | connectivity
      | layer
      | port
      | rule
      | site
      | via

library_specific_singular_object ::=
      | function
      | non_scan_cell
      | test
      | range
      | artwork

source_text ::=
     ALF_REVISION version_string library
```

## 11.8 Generic objects

```
alias ::=
     ALIAS identifier = identifier ;

attribute ::=
     ATTRIBUTE { attribute_items }

attribute_item ::=
     identifier [ { unnamed_assignments } ]

attribute_items ::=
     attribute_item { attribute_item }
```

```
class ::=
     CLASS identifier  ;
   | CLASS identifier  { class_items }

class_item ::=
     all_purpose_item
   | logic_assignment
   | vector_assignment

class_items ::=
   class_item { class_item }

constant ::=
     CONSTANT identifier = number ;
   | CONSTANT identifier = logic_constant ;

group ::=
     GROUP group_identifier { identifiers }
   | GROUP group_identifier { numbers }
   | GROUP group_identifier { edge_literals }
   | GROUP group_identifier { logic_constants }
   | GROUP group_identifier { logic_variables }
   | GROUP group_identifier { integer : integer }

include ::=
     INCLUDE quoted_string ;

keyword_declaration ::=
   KEYWORD context_sensitive_keyword = syntax_item_identifier ;

property ::=
     PROPERTY [ identifier ] { property_items }

property_items ::=
   property_item { property_item }

property_item ::=
          unnamed_assignment
   |      multi_value_assignment

template ::=
     TEMPLATE template_identifier { template_items }

template_item ::=
     all_purpose_item
   | library_specific_object
   | library_specific_singular_object
   | arithmetic_model
   | arithmetic_model_container
   | header
   | table
```

```
   | equation
   | behavior_item
   | geometric_model

template_items ::=
     template_item { template_item }
```

## 11.9 CELL

```
cell ::=
     CELL cell_identifier { cell_items }
   | CELL cell_identifier ;
   | cell_template_instantiation

cell_item ::=
     all_purpose_item
   | pin
   | pin_group
   | primitive
   | function
   | non_scan_cell
   | test
   | arithmetic_model
   | vector
   | wire
   | blockage
   | artwork
   | connectivity

cell_items ::=
     cell_item {cell_item}

non_scan_cell ::=
   NON_SCAN_CELL { unnamed_cell_instantiations }
```

## 11.10 LIBRARY

```
library ::=
     LIBRARY library_identifier { library_items [sublibraries] }
   | LIBRARY library_identifier ;
   | library_template_instantiation

libraries ::=
     library { library }

library_item ::=
     all_purpose_item
   | arithmetic_model
   | cell
   | primitive
```

```
   | wire
   | layer
   | via
   | rule
   | antenna
   | array
   | site
   | connectivity

library_items ::=
     library_item { library_item }
```

## 11.11 PIN

```
pin ::=
     PIN [ index ] pin_identifier { pin_items }
     PIN [ index ] pin_identifier ;
   | pin_template_instantiation

pins ::=
     pin { pin }

pin_item ::=
     all_purpose_item
   | arithmetic_model
   | range
   | port
   | connectivity

pin_items ::=
     pin_item { pin_item }

pin_group ::=
          PIN_GROUP [ index ] pin_group_identifier {
               pin_items
               MEMBERS { pins }
          }

range ::=
          RANGE { unsigned : unsigned }
```

## 11.12 PRIMITIVE

```
primitive ::=
     PRIMITIVE primitive_identifier { primitive_items }
   | PRIMITIVE primitive_identifier ;
   | primitive_template_instantiation

primitives ::=
     primitive { primitive }
```

```
primitive_item ::=
     all_purpose_item
   | pin
   | function

primitive_items ::=
     primitive_item { primitive_item }
```

## 11.13 SUBLIBRARY

```
sublibrary ::=
     SUBLIBRARY library_identifier { library_items }
   | SUBLIBRARY library_identifier ;
   | sublibrary_template_instantiation

sublibraries ::=
     sublibrary { sublibrary }
```

## 11.14 VECTOR

```
vector ::=
     VECTOR ( vector_or_boolean_expression )  { vector_items }
   | VECTOR ( vector_or_boolean_expression ) ;
   | vector_template_instantiation

vector_item ::=
     all_purpose_item
   | arithmetic_model
   | logic_assignment
   | vector_assignment

vector_items ::=
     vector_item { vector_item }

vector_assignment ::=
   context_sensitive_keyword = ( vector_expression ) ;
```

## 11.15 WIRE

```
wire ::=
     WIRE wire_identifier { wire_items }
   | WIRE wire_identifier ;
   | wire_template_instantiation

wire_item ::=
     all_purpose_item
   | node
   | arithmetic_model
```

```
wire_items ::=
     wire_item { wire_item }

node ::=
          NODE node_identifier { all_purpose_items }
```

## 11.16 Arithmetic model

```
arithmetic_model ::=
          context_sensitive_keyword [ identifier ] = value ;
   |      context_sensitive_keyword [ identifier ]
                [ = value ] { arithmetic_model_items }
   |      arithmetic_model_template_instantiation

arithmetic_models ::=
     arithmetic_model { arithmetic_model }

arithmetic_model_item ::=
          all_purpose_item
   |      from
   |      to
   |      violation
   |      header
   |      table
   |      equation
   |      arithmetic_submodel

arithmetic_model_items ::=
   arithmetic_model_item { arithmetic_model_item }

arithmetic_model_container ::=
     context_sensitive_keyword { arithmetic_models }

arithmetic_submodel ::=
          context_sensitive_keyword = value ;
   |      context_sensitive_keyword
                [ = value ] { arithmetic_submodel_items }
   |       context_sensitive_keyword { arithmetic_submodels }
   |      arithmetic_submodel_template_instantiation

arithmetic_submodels ::=
     arithmetic_submodel { arithmetic_submodel }

arithmetic_submodel_item ::=
          all_purpose_item
   |      header
   |      table
   |      equation

arithmetic_submodel_items ::=
   arithmetic_submodel_item { arithmetic_submodel_item }
```

```
header ::=
     HEADER { [ all_purpose_items ] arithmetic_models }
   | header_template_instantiation

table ::=
     TABLE { table_items }
   | table_template_instantiation

table_item ::=
     number
   | identifier

table_items ::=
     table_item { table_item }

equation ::=
     EQUATION { arithmetic_expression }
   | equation_template_instantiation

violation ::=
          VIOLATION {
                [ message_type_assignment ]
                [ message_assignment ]
                [ behavior ]
          }

from ::=
          FROM {
                [ pin_assignment ]
                [ edge_assignment ]
                [ threshold_arithmetic_model ]
          }

to ::=
          TO {
                [ pin_assignment ]
                [ edge_assignment ]
                [ threshold_arithmetic_model ]
          }
```

# 11.17 FUNCTION

```
function ::=
     FUNCTION [ identifier ]
      { [all_purpose_items] [behavior] [structure] [statetables] } }
   | function_template_instantiation

structure ::=
          STRUCTURE { named_cell_instantiations }
```

```
statetable ::=
     STATETABLE [ identifier ]  { statetable_header statetable_body }

statetables ::=
      statetable { statetable }

statetable_body ::=
     statetable_values : statetable_values ;
     { statetable_values : statetable_values ; }

statetable_header ::=
     logic_variables : logic_variables ;

behavior ::=
     BEHAVIOR [ identifier ]  { behavior_items }

behavior_item ::=
     logic_assignment
   | sequential_logic_statement
   | primitive_instantiation

behavior_items ::=
   behavior_item { behavior_item }

logic_assignment ::=
     identifier [index] = boolean_expression ;

logic_assignments ::=
     logic_assignment  { logic_assignment }

sequential_logic_statement ::=
      sequential_if ( vector_or_boolean_expression )
           { logic_assignments }
     { sequential_else_if ( vector_or_boolean_expression )
           { logic_assignments } }
```

## 11.18 TEST

```
test ::=
          TEST { behavior }
```

## 11.19 Geometric Model

```
geometric_model ::=
          geometric_model_identifier
             [ geometric_model_name_identifier ] {
                all_purpose_items
                coordinates }
   |      geometric_model_template_instantiation
```

```
geometric_models ::=
          geometric_model { geometric_model }

coordinates ::=
          COORDINATES { x_number y_number { x_number y_number } }

geometric_transformations ::=
          geometric_transformation { geometric_transformation }

geometric_transformation ::=
          shift_annotation_container
   |      rotate_assignment
   |      flip_assignment
   |      repeat

repeat ::=
          REPEAT [ = unsigned ] {
                shift_annotation_container
                [ repeat ]
          }
```

## 11.20 ARTWORK

```
artwork ::=
          ARTWORK = artwork_identifier {
                [ geometric_transformations ]
                { pin_assignments }
          }
```

## 11.21 LAYER

```
layer ::=
          LAYER identifier { layer_items }

layer_items ::=
          layer_item { layer_item }

layer_item ::=
          all_purpose_item
   |      arithmetic_model
   |      arithmetic_model_container
```

## 11.22 PATTERN

```
pattern ::=
   PATTERN [ identifier ] {
          [ all purpose_items ]
          [ geometric_models ]
```

```
            [ geometric_transformations ]
      }
```

## 11.23 VIA

```
via ::=
            VIA [ identifier ] { via_items }

via_items ::=
            via_item { via_item }

via_item ::=
            all_purpose_item
      |     pattern
      |     artwork
      |     arithmetic_model

via_reference ::=
            VIA { via_instantiations }
```

## 11.24 BLOCKAGE

```
blockage ::=
            BLOCKAGE [ identifier ] {
                        [ all_purpose_items ]
                        [ patterns ]
                        [ rules ]
            }
```

## 11.25 PORT

```
port ::=
            PORT port_identifier ;
      |     PORT [ port_identifier ] {
                        [ all_purpose_items ]
                        [ patterns ]
                        [ rules ]
                        [ via_reference ]
            }
```

## 11.26 RULE

```
rule ::=
            RULE [ identifier ] { rule_items }

rules ::= rule { rule }
```

```
rule_items ::=
          rule_item { rule_item }

rule_item ::=
          pattern
   |      all_purpose_item
   |      arithmetic_model
```

## 11.27 SITE

```
site ::=
          SITE site_identifier { all_purpose_items }
```

## 11.28 ANTENNA

```
antenna ::=
          ANTENNA [ antenna_identifier ] { antenna_items }

antenna_items ::=
          antenna_item { antenna_item }

antenna_item ::=
          all_purpose_item
   |      arithmetic_model
   |      arithmetic_model_container
```

## 11.29 ARRAY

```
array ::=
          ARRAY identifier {
                all_purpose_items
                geometric_transformations
          }
```

## 11.30 Connectivity

```
connectivity ::=
          CONNECTIVITY [ identifier ] {
                connect_rule_assignment
                between_multi_value_assignment
          }

   |      CONNECTIVITY [ identifier ] {
                connect_rule_assignment
                header table
          }
```

# Appendix A
# Sample Applications

This section shows various examples of library elements modeled using ALF.

## A.1 Truth table versus boolean equation

A combinational logic cell and a sequential logic cell are shown below using two different constructs - truth table and boolean equation.

### A.1.1 NAND gate

A two-input `NAND` gate library cell can be modeled as shown below. The `FUNCTION` of the cell can be modeled either as a `STATETABLE` or as `BEHAVIOR` using a boolean equation.

Modeling a `NAND` gate using a truth table:

```
CELL ND2 { /* 2 input NAND gate */
   PIN a {DIRECTION=input;}
   PIN b {DIRECTION=input;}
   PIN z {DIRECTION=output;}

   FUNCTION {
      STATETABLE {
         a b : z ;
         0 ? : 1 ;
         1 ? : (!b);
      }
   }
)
```

Modeling a `NAND` gate using a boolean expression:

```
CELL ND2 { /* 2 input NAND gate */
   PIN a {DIRECTION=input;}
   PIN b {DIRECTION=input;}
   PIN z {DIRECTION=output;}

   FUNCTION {
      BEHAVIOR {
         z = !(a & b);
      }
   }
)
```

### A.1.2 Flip-flop

A flip-flop with asynchronous set and clear signals is shown below using a truth table.

```
CELL FLIPFLOP {
   PIN CLEAR {DIRECTION=input; SIGNALTYPE=clear; POLARITY=low;}
   PIN SET {DIRECTION=input; SIGNALTYPE=set; POLARITY=low;}
   PIN CLOCK {DIRECTION=input; SIGNALTYPE=clock; POLARITY=rising_edge;}
   PIN D {DIRECTION=input;}
   PIN Q {DIRECTION=output;}
   FUNCTION {
   .../* One of the descriptions below go here */
   }
}
```

```
   STATETABLE {
      CLEAR SET CLOCK  D Q : Q;
      0     ?     ??   ? ? : 0;
      1     0     ??   ? ? : 1;
      1     1     01   ? ? : (d);
      1     1     1?   ? ? : (q);
      1     1     ?0   ? ? : (q);
   }
```

Modeling a flip-flop with asynchronous set and clear using a boolean expression:

```
   BEHAVIOR {
      @(!CLEAR) {Q = 0;} : (!SET) {Q = 1;} : (01 CLOCK) {Q = D;}
   }
```

## A.2 Use of primitives

The functionality of a cell can be described using instances of other cells.

### A.2.1 D-flip-flop with asynchronous clear

Modeling aD- flip-flop with asynchronous clear:

```
CELL d_flipflop_clr {
   PIN cd {DIRECTION=input; SIGNALTYPE=clear; POLARITY=low;}
   PIN cp {DIRECTION=input; SIGNALTYPE=clock; POLARITY=rising_edge;}
   PIN d {DIRECTION=input;}
   PIN q {DIRECTION=output;}
   FUNCTION {
   .../* One of the descriptions below go here */
   }
}
```

Explicit description does not use instances of other cells defined in the library:

```
   BEHAVIOR {
      @(01 cp & cd) {q = d;}
      @(!cd) {q = 0;}
   }
```

Use of primitives permit the derivation of new cells from other cells. Below, a D-flip-flop with asynchronous clear is derived from a predefined `ALF_FLIPFLOP` with asynchronous set and clear (see Section A.1.2):

```
BEHAVIOR {
   ALF_FLIPFLOP {CLOCK=cp; D=d; Q=q; SET='b0; CLEAR=!cd;}
}
```

### A.2.2 JK-flipflop

This example shows three ways of modeling a JK-flip-flop.

```
CELL jk_flipflop {
   PIN cp {DIRECTION=input; SIGNALTYPE=clock; POLARITY=rising_edge;}
   PIN j {DIRECTION=input;}
   PIN k {DIRECTION=input;}
   PIN q {DIRECTION=output;}
   FUNCTION {
   .../* One of the descriptions below go here */
   }
}
```

Explicit description:

```
   BEHAVIOR {
      d =
         (!j & k) ? 0 :
         ( j & !k) ? 1 :
         ( j & k) ? !(q) :
         (!j & !k) ? (q) :
                     'bx ;
      @(01 cp) {q = d;}
   }
```

Use of primitives (using predefined `ALF_MUX` and `ALF_FLIPFLOP`):

```
   BEHAVIOR {
      ALF_MUX {Q=d; D[0]=j; D[1]=!k; S=q;}
      ALF_FLIPFLOP {CLOCK=cp; D=d; Q=q; SET='b0; CLEAR='b0;}
   }
```

Use of a hybrid form (boolean expressions within primitive instantiation):

```
   BEHAVIOR {
      ALF_FLIPFLOP {CLOCK=cp; D=(q ? !k : j); Q=q; SET='b0; CLEAR='b0;}
   }
```

Use of a truth table:

```
   STATETABLE {
      cp j k q : (q) ;
      01 0 0 ? : (q) ;
      01 0 1 ? : 0 ;
      01 1 0 ? : 1 ;
      01 1 1 ? : (!q);
      1? ? ? ? : (q) ;
      ?0 ? ? ? : (q) ;
   }
```

### A.2.3 D-flip-flop with synchronous load and clear

This example shows two different models of a synchronous D-flip-flop.

```
CELL d_flipflop_ld_clr {
   PIN cs {DIRECTION=input; SIGNALTYPE=clear;
          POLARITY=low; ACTION=synchronous;}
   PIN ls {DIRECTION=input;}
   PIN cp {DIRECTION=input; SIGNALTYPE=clock; POLARITY=rising_edge;}
   PIN d {DIRECTION=input;}
   PIN q {DIRECTION=output;}
   FUNCTION { ... }
}
```

Explicit description:

```
   BEHAVIOR {
      d1 = (ls)? d : q;
      d2 = d1 & cs;
      @(01 cp) {q = d2;}
   }
```

Use of primitives:

```
   BEHAVIOR {
      ALF_MUX {Q=d1; D0=q; D1=d; SELECT=ls;}/* Connection by pin name */
      ALF_AND {d2 d1 cs}                  /* Connection by pin order */
      ALF_FLIPFLOP {CLOCK=cp; D=d2; Q=q; SET='b0; CLEAR='b0; }
   }
```

### A.2.4 D-flip-flop with input multiplexor

This example shows three different modeling styles for a D-flip-flop with input multiplexor, asynchronous set, and asynchronous clear:

```
CELL d_flipflop_mux_set_clr {
   PIN sel {DIRECTION=input;}
   PIN sd {DIRECTION=input; SIGNALTYPE=set; POLARITY=low;}
   PIN cd {DIRECTION=input; SIGNALTYPE=clear; POLARITY=low;}
   PIN cp {DIRECTION=input; SIGNALTYPE=clock; POLARITY=rising_edge;}
   PIN d1 {DIRECTION=input;}
   PIN d2 {DIRECTION=input;}
   PIN q {DIRECTION=output;}
   FUNCTION { /* fill in BEHAVIOR */ }
}
```

Explicit description:

```
   BEHAVIOR {
      @(!cd) {q = 0;}
      @(!sd & cd) {q = 1;}
      @(01 cp & cd & sd) {q = sel? d1 : d2;}
   }
```

More efficient description can be created using an *if-then-else* style:

```
   BEHAVIOR {
      @(!cd) {q = 0;}
      :(!sd) {q = 1;}
      :(01 cp){q = sel ? d1 : d2;}
   }
```

Use of primitive:

```
BEHAVIOR {
   ALF_FLIPFLOP {CLOCK=cp; D=(sel ? d1: d2); Q=q; SET=!sd; CLEAR=!cd;}
}
```

The use of an `ALF_MUX` primitive is eliminated here by using an assignment expression to the `D` input in the `ALF_FLIPFLOP` instance.

### A.2.5 D-latch

This example shows a level-sensitive cell in two different styles.

```
CELL d_latch {
   PIN g {DIRECTION=input; SIGNALTYPE=clock; POLARITY=high;}
   PIN d {DIRECTION=input;}
   PIN q {DIRECTION=output;}
   FUNCTION { ... }
}
```

Explicit description:

```
BEHAVIOR {
   @(g) {q = d;}
}
```

Use of primitive:

```
BEHAVIOR {
   ALF_LATCH {ENABLE=g; D=d; Q=q; SET='b0; CLEAR='b0;}
}
```

### A.2.6 SR-latch

The example below shows how some of the input pins can be left unconnected if they represent a “don’t care” situation.

```
CELL sr_latch {
   PIN sn {DIRECTION=input; SIGNALTYPE=set; POLARITY=low;}
   PIN rn {DIRECTION=input; SIGNALTYPE=clear; POLARITY=low;}
   PIN q {DIRECTION = output;}
   PIN qn {DIRECTION = output;}
   FUNCTION { ... }
}
```

Explicit description:

```
BEHAVIOR {
   @ (!sn) {q = 'b1; qn = !rn;}
   @ (!rn) {qn = 'b1; q = !sn;}
}
```

Use of primitive:

```
BEHAVIOR {
   ALF_LATCH {ENABLE='b0; Q=q; SET=!sn; CLEAR=!rn;}
}
```

Since the ENABLE pin is always set to 0, the connection of D pin is irrelevant. Even if D is considered 'bX or 'bZ, the behavior shall not change.

### A.2.7 JTAG BSR

The following example shows a JTAG BSR cell with built-in scan chain.

```
CELL F10_18 {
   PIN SysOut {DIRECTION = output;}
   PIN TDO {DIRECTION = output; SIGNALTYPE = scan_data;}
   PIN SysIn {DIRECTION = input;}
   PIN TDI {DIRECTION = input; SIGNALTYPE = scan_data;}
   PIN Shift {DIRECTION = input; SIGNALTYPE = scan_enable;}
   PIN Clk {DIRECTION = input; POLARITY = rising_edge;
            SIGNALTYPE = master_clock;}
   PIN Update {DIRECTION = input; POLARITY = rising_edge;
            SIGNALTYPE = slave_clock;}
   PIN Mode {DIRECTION = input; SIGNALTYPE = select;}
   PIN STATE0 { // This state is on the scan chain
            SCAN_POSITION = 1; DIRECTION = output; VIEW = none;}
   PIN STATE1 { // NOT on scan chain (just update latch)
            DIRECTION = output; VIEW = none;}
   FUNCTION {
      BEHAVIOR {
         @(01 Clk) {STATE0 = Shift ? TDI : SysIn;}
         @(01 Update) {STATE1 = STATE0;}
         TDO = STATE0;
         SysOut = Mode ? STATE1 : SysIn;
      }
   }
}
```

### A.2.8 Combinational scan cell

The following example shows a combinational scan cell with a reused primitive.

```
LIBRARY major_ASIC_vendor {
   INFORMATION {
      version = v2.1.0;
      title = “0.35 standard cell”;
      product = p35sc;
      author = “Major Asic Vendor, Inc.”;
      datetime = “Wed Jul 23 13:50:12 MST 1997”;
   }
   ..
   CELL ND3A {
      INFORMATION {
         version = v6.0;
         title = “3 input nand”;
         product = p35sc_lib;
         author = “Joe Cell Designer”;
         datetime = “Tue Apr 1 01:39:47 PST 1997”;
      }
```

```
      PIN Z {DIRECTION=output;}
      PIN A {DIRECTION=input;}
      PIN B {DIRECTION=input;}
      PIN C {DIRECTION=input;}
      FUNCTION {
         BEHAVIOR {
            ALF_NAND {Z A B C}
         }
      }
      /* fill in timing and power data for ND3A cell */
   }
   ..
   CELL ND3B {
      PIN Z {DIRECTION=output;}
      PIN A {DIRECTION=input;}
      PIN B {DIRECTION=input;}
      PIN C {DIRECTION=input;}
      FUNCTION {
         BEHAVIOR {
            ALF_NAND {Z A B C}
         }
      }
      /* fill in timing and power data for ND3B cell */
   }
   ..
   CELL SCAN_ND4 {
      PIN Z {DIRECTION=output;}
      PIN A {DIRECTION=input;}
      PIN B {DIRECTION=input;}
      PIN C {DIRECTION=input;}
      PIN D {DIRECTION=input; SIGNALTYPE=scan_enable;}

      SCAN_TYPE = control_0;
      NON_SCAN_CELL = ALF_NAND {Z A B C}
      FUNCTION {
         BEHAVIOR {
            Z = !(A & B & C & D);
         }
      }
   }
   ..
}
```

### A.2.9 Scan flip-flop

The following example shows a scan flip-flop using the generic `ALF_FLIPFLOP` primitive.

```
LIBRARY major_ASIC_vendor {
   ...
   CELL F614 {
      PIN H01 {DIRECTION = input; SIGNALTYPE = data;}
      PIN H02 {DIRECTION = input; SIGNALTYPE = clock;}
      PIN H03 {DIRECTION = input; SIGNALTYPE = clear; POLARITY = high;}
      PIN H04 {DIRECTION = input; SIGNALTYPE = set; POLARITY = high;}
```

```
         PIN N01 {DIRECTION = output;}
         PIN N02 {DIRECTION = output;}
         FUNCTION {
            BEHAVIOR {
               ALF_FLIPFLOP {
                  D=H01; CLOCK=H02; CLEAR=H03; SET=H04;
                  Q=N01; QN=N02; Q_CONFLICT='bX; QN_CONFLICT='bX;
               }
            }
         }
      }
      ...
      CELL S000 {
         PIN H01 {DIRECTION = input; SIGNALTYPE = scan_data;}
         PIN H02 {DIRECTION = input; SIGNALTYPE = clock;}
         PIN H03 {DIRECTION = input; SIGNALTYPE = scan_enable;
                  POLARITY = low;}
         PIN H04 (DIRECTION = input; SIGNALTYPE = set; POLARITY = high;}
         PIN H05 {DIRECTION = input; SIGNALTYPE = clear; POLARITY = high;}
         PIN H06 {DIRECTION = input; SIGNALTYPE = data;}
         PIN N01 {DIRECTION = output; SIGNALTYPE = data;}
         PIN N02 {DIRECTION = output;}
         FUNCTION{
            BEHAVIOR{ // flipflop_d is an implicitly defined internal pin
               ALF_MUX {Q=flipflop_d; D0=H06; D1=H01; SELECT=H03;}
               ALF_FLIPFLOP {
                  D=flipflop_d; CLOCK=H02; CLEAR=H05; SET=H04;
                  Q=N01; QN=N02; Q_CONFLICT='bX; QN_CONFLICT='bX;
               }
            }
         }
         SCAN_TYPE = muxscan;
         NON_SCAN_CELL = ALF_FLIPFLOP {D=H06; CLOCK=H02; CLEAR=H05; SET=H04;
                                       Q=N01; QN=N02; Q_CONFLICT='bX;
                                       QN_CONFLICT='bX; 'b0=H03; 'b0=H01;}
      } // H03 and H01 have no corresponding pin in ALF_FLIPFLOP
      ...
   }
```

### A.2.10 Quad D-flip-flop

The following example shows a quad D-flip-flop with and without built-in scan chain.

```
   LIBRARY major_ASIC_vendor {
      PRIMITIVE FFX4 {
         PIN CK { DIRECTION = input; }
         PIN D0 { DIRECTION = input; }
         PIN D1 { DIRECTION = input; }
         PIN D2 { DIRECTION = input; }
         PIN D3 { DIRECTION = input; }
         PIN Q0 { DIRECTION = output; }
         PIN Q1 { DIRECTION = output; }
         PIN Q2 { DIRECTION = output; }
         PIN Q3 { DIRECTION = output; }
```

```
      FUNCTION {
         BEHAVIOR {
            ALF_FLIPFLOP {Q=Q0; D=D0; CLOCK=CK; SET='b0; CLEAR='b0;}
            ALF_FLIPFLOP {Q=Q1; D=D1; CLOCK=CK; SET='b0; CLEAR='b0;}
            ALF_FLIPFLOP {Q=Q2; D=D2; CLOCK=CK; SET='b0; CLEAR='b0;}
            ALF_FLIPFLOP {Q=Q3; D=D3; CLOCK=CK; SET='b0; CLEAR='b0;}
         }
      }
   }
   CELL SCAN_FFX4 {
      PIN OUT0 {DIRECTION = output;}
      PIN OUT1 {DIRECTION = output;}
      PIN OUT2 {DIRECTION = output;}
      PIN OUT3 {DIRECTION = output;}
      PIN SO {DIRECTION = output; SIGNALTYPE = scan_data;}
      PIN IN0 {DIRECTION = input; SIGNALTYPE = data;}
      PIN IN1 {DIRECTION = input; SIGNALTYPE = data;}
      PIN IN2 {DIRECTION = input; SIGNALTYPE = data;}
      PIN IN3 {DIRECTION = input; SIGNALTYPE = data;}
      PIN CLK {DIRECTION = input; SIGNALTYPE = clock;}
      PIN SI {DIRECTION = input; SIGNALTYPE = scan_data;}
      PIN SE {DIRECTION = input; SIGNALTYPE = scan_enable;}
      PIN STATE0 {SCAN_POSITION = 1; DIRECTION = output; VIEW = none;}
      PIN STATE1 {SCAN_POSITION = 2; DIRECTION = output; VIEW = none;}
      PIN STATE2 {SCAN_POSITION = 3; DIRECTION = output; VIEW = none;}
      PIN STATE3 {SCAN_POSITION = 4; DIRECTION = output; VIEW = none;}
      FUNCTION {
         BEHAVIOR {
            OUT0 = STATE0; OUT1 = STATE1; OUT2 = STATE2; OUT3 = STATE3;
            SO = !STATE3;
            @(01 CLK) {
               STATE0 = SE ? !SI : IN0;
               STATE1 = SE ? !STATE0 : IN1;
               STATE2 = SE ? !STATE1 : IN2;
               STATE3 = SE ? !STATE2 : IN3;
            }
         }
      }
      SCAN_TYPE = muxscan;
      NON_SCAN_CELL = FFX4 {CLK IN0 IN1 IN2 IN3 OUT0 OUT1 OUT2 OUT3}
      } // this example shows referencing by order
   }
}
```

## A.3 Templates and vector-specific models

This section describes how to use templates and vector-specific models.

### A.3.1 Vector-specific delay and power tables

In this example, the use of vector specific models for input-to-output delay, output slewrate, and switching energy is shown.

```
CELL nand2 {
   PIN a {DIRECTION = input; CAPACITANCE = 0.02 {UNIT = pF;}}
   PIN b {DIRECTION = input; CAPACITANCE = 0.02 {UNIT = pF;}}
   PIN z {DIRECTION = output;}
   FUNCTION {
      BEHAVIOR {z = !(a & b); }
   }
   VECTOR (10 a -> 01 z){    /* Vector specific characterization */
      DELAY {
         UNIT = ns;
         FROM {PIN = a; THRESHOLD = 0.4;}
         TO {PIN = z; THRESHOLD = 0.6;}
         HEADER {
            CAPACITANCE {
               PIN = z; UNIT = pF;
               TABLE {0.01 0.02 0.04 0.08 0.16}
            }
            SLEWRATE {
               PIN = a; UNIT = ns;
               FROM {THRESHOLD = 0.5;}
               TO {THRESHOLD = 0.3;}
               TABLE {0.1 0.3 0.9}
            }
         }
         TABLE {
               0.1 0.2 0.4 0.8 1.6
               0.2 0.3 0.5 0.9 1.7
               0.4 0.5 0.7 1.1 1.9
         }
      }
      SLEWRATE {
         PIN = z; UNIT = ns;
         FROM {THRESHOLD = 0.3;}
         TO {THRESHOLD = 0.5;}
         HEADER {
            CAPACITANCE {
               PIN = z; UNIT = pF;
               TABLE {0.01 0.02 0.04 0.08 0.16}
            }
            SLEWRATE {
               PIN = a; UNIT = ns;
               FROM {THRESHOLD = 0.5;}
               TO {THRESHOLD = 0.3;}
               TABLE {0.1 0.3 0.9}
            }
         }
         TABLE {
               0.1 0.2 0.4 0.8 1.6
               0.1 0.2 0.4 0.8 1.6
               0.2 0.4 0.6 1.0 1.8
         }
      }
      ENERGY {
         UNIT = pJ;
```

```
            HEADER {
               CAPACITANCE {
                  PIN = z; UNIT = pF;
                  TABLE {0.01 0.02 0.04 0.08 0.16}
               }
               SLEWRATE {
                  PIN = a; UNIT = ns;
                  FROM {THRESHOLD = 0.5;}
                  TO {THRESHOLD = 0.3;}
                  TABLE {0.1 0.3 0.9}
               }
            }
            TABLE {
                  0.21 0.32 0.64 0.98 1.96
                  0.22 0.33 0.65 0.99 1.97
                  0.31 0.42 0.74 1.08 2.06
            }
         }
      }
      VECTOR (01 a -> 10 z){
         DELAY { ... }
         SLEWRATE { ... }
         ENERGY { ... }
      }
      VECTOR (10 b -> 01 z){
         DELAY { ... }
         SLEWRATE { ... }
         ENERGY { ... }
      }
      VECTOR (01 b -> 10 z){
         DELAY { ... }
         SLEWRATE { ... }
         ENERGY { ... }
      }
   }
```

### A.3.2 Use of TEMPLATE

Notice the header for the delay, ramptime, and energy models was the same in the example in Section A.3.1. Therefore, creating a template definition can eliminate duplicate information, reduce the possibility of inadvertent errors, and make the models compact. For example, a header template can be created as shown below:

```
TEMPLATE std_header_2d {
   HEADER {
      CAPACITANCE {
         PIN = <out_pin>; UNIT = pF;
         TABLE {0.01 0.02 0.04 0.08 0.16}
      }
      SLEWRATE {
         PIN = <in_pin>; UNIT = ns;
         FROM {THRESHOLD {RISE = 0.3; FALL = 0.5;} }
         TO {THRESHOLD {RISE = 0.5; FALL = 0.3;} }
         TABLE {0.1 0.3 0.9}
      }
   }
```

The use of `TEMPLATE` eliminates the repetition of header information by rewriting the previous example (only the first vector) as shown below.

```
      DELAY {
         UNIT = ns;
         THRESHOLD {RISE=0.4; FALL=0.6;}
         FROM {PIN = a;}
         TO {PIN = z;}
         std_header_2d {     /* Template is used */
            in_pin = a;
            out_pin = z;
         }
         TABLE {
               0.1 0.2 0.4 0.8 1.6
               0.2 0.3 0.5 0.9 1.7
               0.4 0.5 0.7 1.1 1.9
         }
      }
      SLEWRATE {
         PIN = z; UNIT = ns;
         FROM {THRESHOLD {RISE = 0.3; FALL = 0.5;} }
         TO {THRESHOLD {RISE = 0.5; FALL = 0.3;} }
         std_header_2d {     /* Template is used */
            in_pin = a;
            out_pin = z;
         }
         TABLE {
               0.1 0.2 0.4 0.8 1.6
               0.1 0.2 0.4 0.8 1.6
               0.2 0.4 0.6 1.0 1.8
         }
      }
      ENERGY {
         UNIT = pJ;
         std_header_2d {     /* Template is used */
            in_pin = a;
            out_pin = z;
         }
```

```
            TABLE {
                  0.21 0.32 0.64 0.98 1.96
                  0.22 0.33 0.65 0.99 1.97
                  0.31 0.42 0.74 1.08 2.06
            }
         }
      }
```

Note the entire characterization model for CELL `nand2` is the same for each vector (i.e., pair of input and output pins), so further efficiency can be achieved by defining the characterization model itself as a template. This template definition uses the instantiation of the previously defined header template.

```
TEMPLATE std_char_2d {
   DELAY {
      UNIT = ns;
      THRESHOLD {RISE=0.4; FALL=0.6;}
      FROM {PIN = <in_pin>; }
      TO {PIN = <out_pin>; }
      std_header_2d {
         in_pin = <input_pin>;
         out_pin = <output_pin>;
      }
      TABLE <delay_data>
   }
   SLEWRATE {
      PIN = <out_pin>; UNIT = ns;
      FROM {THRESHOLD {RISE = 0.3; FALL = 0.5;} }
      TO {THRESHOLD {RISE = 0.5; FALL = 0.3;} }
      std_header_2d {
         in_pin = <input_pin>;
         out_pin = <output_pin>;
      }
      TABLE <slewrate_data>
   }
   ENERGY {
      UNIT = pJ;
      std_header_2d {
         in_pin = <input_pin>;
         out_pin = <output_pin>;
      }
      TABLE <energy_data>
   }
}
```

Now only the delay, slewrate, and energy models contain specific data that is different for each vector. All repetitive information is in the template definition. The characterization model can be rewritten compactly as shown below:

```
std_char_2d {
   in_pin = a;
   out_pin = z;
   delay_data {
         0.1 0.2 0.4 0.8 1.6
         0.2 0.3 0.5 0.9 1.7
         0.4 0.5 0.7 1.1 1.9
   }
   slewrate_data {
         0.1 0.2 0.4 0.8 1.6
         0.1 0.2 0.4 0.8 1.6
         0.2 0.4 0.6 1.0 1.8
   }
   energy_data {
         0.21 0.32 0.64 0.98 1.96
         0.22 0.33 0.65 0.99 1.97
         0.31 0.42 0.74 1.08 2.06
   }
}
```

### A.3.3 Vector description styles for timing arcs

In previous examples, the vectors were specified as timing arcs. This is not ambiguous, since the sequence of transitions can only happen under one test condition.

```
VECTOR (10 a -> 01 z){
   std_char_2d { ... }
}
VECTOR (01 a -> 10 z){
   std_char_2d { ... }
}
VECTOR (10 b -> 01 z){
   std_char_2d { ... }
}
VECTOR (01 b -> 10 z){
   std_char_2d { ... }
}
```

An alternate way of describing the above vectors is to specify the input transition and the state of the other input(s) which control the output transition.

```
VECTOR (10 a && b){
   std_char_2d { ... }
}
VECTOR (01 a && b){
   std_char_2d { ... }
}
VECTOR (10 b && a){
   std_char_2d { ... }
}
VECTOR (01 b && a){
   std_char_2d { ... }
}
```

A more concise method of vector description is to specify both output transition and input state(s) together with the input transition.

```
VECTOR (10 a -> 01 z && b){
   std_char_2d { ... }
}
VECTOR (01 a -> 10 z && b){
   std_char_2d { ... }
}
VECTOR (10 b -> 01 z && a){
   std_char_2d { ... }
}
VECTOR (01 b -> 10 z && a){
   std_char_2d { ... }
}
```

In the non-redundant specification, either the input state or the output transition can be derived from the functional description.

### A.3.4 Vectors for delay, power, and timing constraints

A D-flip-flop model without the set and clear signals is shown below. This model has vectors with a specific purpose: some for delay and power, some for power only (output is not switching), and some for timing constraints. However, each vector has the same structure, although the input variables change. The vectors for delay and power model require two-dimensional tables with load capacitance and input ramptime as variables, the vectors for power model require one-dimensional tables with input ramptime as variable, and the vectors for time constraints require 2-dimensional tables with ramptime on two inputs as variables.

```
CELL d_flipflop {
   PIN cp {DIRECTION = input; SIGNALTYPE = clock; POLARITY = rising_edge;}
   PIN d {DIRECTION = input;}
   PIN q {DIRECTION = output;}
   FUNCTION {
      BEHAVIOR { @(01 cp) {q = d; } }
   }
   VECTOR (01 cp -> 01 q) {
      /* fill in arithmetic models for delay and power */
   }
   VECTOR (01 cp -> 10 q) {
      /* fill in arithmetic models for delay and power */
   }
   VECTOR (01 cp && d == q) {
      /* fill in arithmetic model for power */
   }
   VECTOR (10 cp && d == q) {
      /* fill in arithmetic model for power */
   }
   VECTOR (10 cp && d != q) {
      /* fill in arithmetic model for power */
   }
   VECTOR (01 d && !cp) {
      /* fill in arithmetic model for power */
   }
```

```
      VECTOR (10 d && !cp) {
         /* fill in arithmetic model for power */
      }
      VECTOR (01 d && cp) {
         /* fill in arithmetic model for power */
      }
      VECTOR (10 d && cp) {
         /* fill in arithmetic model for power */
      }
      VECTOR (01 d <&> 01 cp)
         SETUP {
            /* fill in arithmetic model for setup time constraint */
            VIOLATION {
               BEHAVIOR {q = 'bx;}
               MESSAGE_TYPE = error;
               MESSAGE = “setup violation 01 d <-> 01 cp“;
            }
         }
         HOLD {
            /* fill in arithmetic model for hold time constraint */
            VIOLATION {
               BEHAVIOR {q = 'bx;}
               MESSAGE_TYPE = error;
               MESSAGE = “hold violation 01 d <-> 01 cp“;
            }
         }
      VECTOR (10 d <&> 01 cp)
         SETUP {
            /* fill in arithmetic model for setup time constraint */
            VIOLATION {
               BEHAVIOR {q = 'bx;}
               MESSAGE_TYPE = error;
               MESSAGE = “setup violation 10 d <-> 01 cp“;
            }
         }
         HOLD {
            /* fill in arithmetic model for hold time constraint */
            VIOLATION {
               BEHAVIOR {q = 'bx;}
               MESSAGE_TYPE = error;
               MESSAGE = “hold violation 10 d <-> 01 cp“;
            }
         }
      }
   }
```

## A.4 Combining tables and equations

This section describes how to combine tables and equations.

### A.4.1 Table versus equation

The following examples show the usage of TABLE and EQUATION in the model.

An example using a table:

```
CURRENT {
   PIN = VDD;
   UNIT = mA;
   TIME = 30 {UNIT = ns;}
   MEASUREMENT = average;
   HEADER {
      CAPACITANCE {
         PIN = z; UNIT = pF;
         TABLE {0.02 0.04 0.08 0.16}
      }
      SLEWRATE {
         PIN = a; UNIT = ns;
         TABLE {0.1 0.3 0.9}
      }
   }
   TABLE {
      0.0011 0.0021 0.0041 0.0081
      0.0013 0.0023 0.0043 0.0083
      0.0019 0.0029 0.0049 0.0089
   }
}
```

The equivalent example using an equation:

```
CURRENT {
   PIN = VDD; UNIT = mA;
   TIME = 30 {UNIT = ns;}
   MEASUREMENT = average;
   HEADER {
      CAPACITANCE {PIN = z; UNIT = pF;}
      SLEWRATE {PIN = a; UNIT = ns;}
   }
   EQUATION { 0.05*CAPACITANCE + 0.001*SLEWRATE }
}
```

If the model uses an EQUATION, then each argument shall appear in the HEADER. If the model uses a TABLE, then the HEADER shall contain a TABLE for each argument. The number of values in the main table and the indexing scheme is defined by the order and the number of values in each table inside the header.

### A.4.2 Cell with multiple output pins

The following example shows how to use combinations of tables and equations for efficient modeling of energy consumption of a cell with two (buffered) outputs. When two outputs are switching and are triggered by the same input, the dynamic energy consumption depends on ramptime of the input signal and load capacitance on each output.

Instead of creating a three-dimensional table, 2 two-dimensional tables are used, which varies the load capacitance at one output and keeps zero load at the other output. The equation

calculates the energy for both outputs switching by adding the values from each table together for the applicable load capacitance and by subtracting a corresponding correction term. The result is exact for cells with buffered outputs.

As shown in the example below, an arithmetic model becomes a named object if several objects of the same type occur within the same scope (e.g,. `ENERGY`). For named objects, the equation uses the object name instead of the object type.

```
VECTOR (01 ci -> (01 co <-> 10 s) & a) {
   ENERGY {
      UNIT = pJ;
      HEADER {
         ENERGY energy_co {        // named object
            UNIT = pJ;
            HEADER {
               CAPACITANCE {
                  PIN = co; UNIT = pF;
                  TABLE { ... }
               }
               SLEWRATE {
                  PIN = ci; UNIT = ns;
                  TABLE { ... }
               }
            }
            TABLE { ... }
         }
         ENERGY energy_s {         // named object
            UNIT = pJ;
            HEADER {
               CAPACITANCE {
                  PIN = s; UNIT = pF;
                  TABLE { ... }
               }
               SLEWRATE {
                  PIN = ci; UNIT = ns;
                  TABLE { ... }
               }
            }
            TABLE { ... }
         }
         ENERGY energy_noload {    // named object
            UNIT = pJ;
            HEADER {
               SLEWRATE {
                  PIN = ci; UNIT = ns;
                  TABLE { ... }
               }
            }
            TABLE { ... }
         }
      }
      EQUATION { energy_co + energy_s - energy_noload }
   }
}
```

### A.4.3 PVT derating

Combinations of tables and equations can also be used for derating with respect to voltage and temperature, since those variables can add more dimensions to a purely table-based model.

In this example, the DELAY objects are named, since there is both a nominal and a derated DELAY.

```
DELAY rise_out{
   HEADER {
      PROCESS {
         HEADER {nom snsp snwp wnsp wnwp}
         TABLE {0.0 -0.1 -0.2 +0.3 +0.2}
      }
      VOLTAGE {//fill in any annotations
      }
      TEMPERATURE {//fill in any annotations
      }
      DELAY nom_rise_out {
         HEADER {
            CAPACITANCE {
               TABLE {0.03 0.06 0.12 0.24}
            }
            SLEWRATE {
               TABLE {0.1 0.3 0.9}
            }
         }
         TABLE {
            0.07 0.10 0.14 0.22
            0.09 0.13 0.19 0.30
            0.10 0.15 0.25 0.41
         }
      }
   }
   EQUATION {
      nom_rise_out
      * (1 + PROCESS)
      * (1 + (TEMPERATURE - 25)*0.001)
      * (1 + (VOLTAGE - 3.3)*(-0.3))
   }
}
```

The HEADER in the process object contains exclusively named variables (nom, snsp, etc.), similar to the truth table of a FUNCTION containing only pin names. Therefore, the TABLE is expected to have as many entries as the HEADER. The TABLE inside nom_rise_out shall follow the format defined by each TABLE inside the declarations of load and ramptime. Other declared objects in the HEADER are ignored for the TABLE format if they do not have a TABLE inside themselves.

For convenience, the derating equation can be defined as a template for future reuse.

```
TEMPLATE std_derating {
   EQUATION {
      <variable>
      * (1 + <Kp>)
      * (1 + (TEMPERATURE - 25)*<Kt>)
      * (1 + (VOLTAGE - 3.3)*<Kv>)
   }
}
```

Instantiation of the template in the model:

```
DELAY rise_out{
   HEADER {
      PROCESS {
         HEADER {nom snsp snwp wnsp wnwp}
         TABLE {0.0 -0.1 -0.2 +0.3 +0.2}
      }
      VOLTAGE { ... }
      TEMPERATURE { ... }
      DELAY nom_rise_out {
         HEADER {
            CAPACITANCE {TABLE { ... }}
            SLEWRATE {TABLE { ... }}
         }
         TABLE { ... }
   }
   std_derating {
      variable = nom_rise_out ;
      Kp = PROCESS ;
      Kt = 0.001 ;
      Kv = -0.3 ;
   }
}
```

It is possible to express voltage, temperature and delay with the derating case as an independent variable:

```
VOLTAGE {
   HEADER { DERATE_CASE { TABLE {nom bccom wcmil} } }
   TABLE {3.3 3.5 2.8}
}
TEMPERATURE {
   HEADER { DERATE_CASE { TABLE {nom bccom wcmil} } }
   TABLE {25 0 125}
}
DELAY {
   HEADER {
      DERATE_CASE {
         HEADER {nom bccom wcmil}
         TABLE {0 -0.0835 0.265}
      }
      PROCESS
         HEADER {nom snsp snwp wnsp wnwp}
         TABLE {0.0 -0.1 -0.2 +0.3 +0.2}
      }
```

```
      DELAY nom_rise_out { ... }
   }
   EQUATION {
      nom_rise_out
      * (1 + PROCESS)
      * (1 + DERATE_CASE)
   }
```

Another possibility is a completely tabulated model, where the process and derating identifiers can be directly used as table items.

```
DELAY {
   HEADER {
      DERATE_CASE {
         TABLE {nom bccom wcmil}
      }
      PROCESS
         TABLE {nom snsp snwp wnsp wnwp}
      }
   TABLE {
      // 3*5 = 15 values
   }
```

## A.5 Use of annotations

This section describes how to use annotations.

### A.5.1 Annotations for a PIN

Direct annotation:

```
PIN data_in {DIRECTION = input; THRESHOLD = 0.35; CAPACITANCE = 0.010;}
```

Using annotation containers:

```
PIN data_in {
   DIRECTION = input;
   THRESHOLD = 0.35;
   CAPACITANCE = 0.010; {
         UNIT = pF; MEASUREMENT = average;
         MIN = 0.009; TYP = 0.010; MAX = 0.012;
   }
   LIMIT {
      SLEWRATE {UNIT=ns;MAX=3.0;}
      VOLTAGE {MAX=3.5; MIN=-0.2;}
   }
}
```

The input pin `data_in` has a non-linear capacitance that was characterized by using an average measurement (as opposed to RMS or peak measurements). Different measurements yield average capacitances between `0.009 pF` and `0.012 pF`; the typical average capacitance is `0.010 pF`. The slewrate applied to the pin shall not exceed `3.0 ns`. The voltage swing shall not exceed the lower bound of `-0.2 V` and the upper bound of `3.5 V`.

```
CAPACITANCE {UNIT = pF;}
PIN data_out {
   DIRECTION = output; CAPACITANCE = 0.002;
   LIMIT {CAPACITANCE {MAX = 0.96;} }
}
```

The output pin `data_out` has a capacitance of `0.002 pF`. The maximum load capacitance that can be applied to the pin is `0.96 pF`.

### A.5.2 Annotations for a timing arc

Specifications for a particular timing arc which references specific pins:

```
DELAY {
   UNIT = ns;
   FROM {PIN = data_in; THRESHOLD = 0.4;}
   TO {PIN = data_out; THRESHOLD = 0.6;}
}

SLEWRATE {
   PIN = data_out; UNIT = ns;
   FROM {THRESHOLD = 0.3;}
   TO {THRESHOLD = 0.5;}
}
```

Specifications for a generic timing arc which does not reference specific pins, but where the values for both switching directions are defined:

```
DELAY {
   UNIT = ns;
   THRESHOLD {RISE=0.4; FALL=0.6;}
}

SLEWRATE {
   UNIT = ns;
   FROM {THRESHOLD {RISE=0.3; FALL=0.5;}}
   TO  {THRESHOLD {RISE=0.5; FALL=0.3;}}
}
```

### A.5.3 Creating self-explaining annotations

Self-explaining annotations can be created using `TEMPLATE`.

Example:

The number of connections allowed for each pin:

```
TEMPLATE must_connect {
      LIMIT {CONNECTION {MIN = 1;}}
}

TEMPLATE can_float {
      LIMIT {CONNECTION {MIN = 0;}}
}

TEMPLATE no_connection {
      LIMIT {CONNECTION {MAX = 0;}}
}
```

```
CELL a_flipflop {
   PIN q {must_connect DIRECTION=output;}
   PIN qn {can_float DIRECTION=output;}
   PIN qi {no_connection DIRECTION=output;}
   ...
}
```

## A.6 Providing a fall-back position for applications (using DEFAULT)

ALF's modeling capabilities address the needs for all types of applications. However, ALF shall also work for applications that use only a subset of information. Capability can be modelled using a `DEFAULT` to control a subset of information. The information provided by `DEFAULT` can be strictly ignored by applications that understand the full information.

A particular application might not be able to use three-dimensional tables or understand certain models. `DEFAULT` values can be provided for each model.

Example:

```
      DELAY {
         HEADER {
            SLEWRATE {
               PIN = a; UNIT = 1e-9;
               TABLE {0.5 1.0 1.5}
               DEFAULT = 1.0;
            }
            CAPACITANCE {
               PIN = z; UNIT = 1e-12;
               TABLE {0.1 0.2 0.3 0.4}
               DEFAULT = 0.1;
            }
            VOLTAGE {
               PIN = vdd; UNIT = 1;
               TABLE {3.0 3.3 3.6}
               DEFAULT = 3.3;
            }
         }
         TABLE {
            // arrangement of whitespaces and comments
            // is only for readability
            // parser sees just a sequence of 3x4x3=36 numbers

//slewrate  0.5 1.0 1.5 capacitance      voltage
//          --------------+--------------+-------
            0.2 0.8 1.1   // 0.1            3.0
            0.4 1.0 1.2   // 0.2
            0.7 1.2 1.4   // 0.3
            0.9 1.5 1.8   // 0.4

            0.1 0.7 1.2   // 0.1            3.3
            0.3 0.9 1.3   // 0.2
```

```
            0.6 1.1 1.5   // 0.3
            0.8 1.3 1.7   // 0.4

            0.1 0.6 1.0   // 0.1           3.6
            0.2 0.8 1.1   // 0.2
            0.4 1.0 1.3   // 0.3
            0.7 1.2 1.6   // 0.4
         }
      }
```

An application that does not understand `VOLTAGE` shall extract the following information from this example:

```
      DELAY {
         HEADER {
            SLEWRATE {
               PIN = a; UNIT = 1e-9;
               TABLE {0.5 1.0 1.5}
            }
            CAPACITANCE {
               PIN = z; UNIT = 1e-12;
               TABLE {0.1 0.2 0.3 0.4}
            }
         }
         TABLE {

   //slewrate  0.5 1.0 1.5       capacitance   voltage
   //             --------------+--------------+-------
               0.1 0.7 1.2      // 0.1            3.3
               0.3 0.9 1.3      // 0.2
               0.6 1.1 1.5      // 0.3
               0.8 1.3 1.7      // 0.4
         }
      }
```

An application that does not understand `SLEWRATE` shall extract only the following information:

```
      DELAY {
         HEADER {
            CAPACITANCE {
               UNIT = 1e-12;
               PIN = z;
               TABLE {0.1 0.2 0.3 0.4}
            }
         }
         TABLE {

   //slewrate  1.0 capacitance     voltage
   //          ----+--------------+-------
               0.7   // 0.1           3.3
               0.9   // 0.2
               1.1   // 0.3
               1.3   // 0.4
         }
      }
```

# A.7 Bus modeling

This section describes how to model buses.

## A.7.1 Tristate driver

Bus drivers are usually tristate buffers, which have straightforward functional models. If both the input signal and enable signal have well-defined logic states, the output is driven to `'b1`, `'b0`, or `'bz`; otherwise it is driven to `'bx`.

```
CELL tristate_buffer {
   PIN a {DIRECTION = input; SIGNALTYPE = data;}
   PIN e {DIRECTION = input; SIGNALTYPE = out_enable;}
   PIN z {DIRECTION = output; SIGNALTYPE = data;
          ATTRIBUTE {tristate} }
   FUNCTION {
      BEHAVIOR {
         z =
          (e & a) ? 'b1:
          (e & !a) ? 'b0:
          (!e)     ? 'bz:
                     'bx;
      }
   }
}
```

A different model can be used for transmission-gate type of buffers, which also passes the high impedance state from input to output.

```
      BEHAVIOR {
         z =
          ( e) ? a :
          (!e) ? 'bz:
                 'bx;
      }
   }
```

The drive strength information of tristate buffers is also needed to model a bus contention. This is easily achieved by annotating a pin property, using a context-sensitive keyword.

```
CELL tristate_buffer {
   ...
   PIN z {DIRECTION = output; DRIVE_STRENGTH = 4;ATTRIBUTE {tristate}}
   ...
}
```

The pin-property `DRIVE_STRENGTH` can take an arbitrary positive integer or a real number. In general, greater values override smaller values and here `DRIVE_STRENGTH=0` is equivalent to

```
BEHAVIOR {z='bz;}.
```

ALF does not assume a particular set of legal drive strengths. The scale and granularity is left to the discretion of the ASIC vendor (user).

Modeling of state-dependent drive strength is achieved by annotating drive strength within a vector rather than within a pin. The following example shows a buffer with a `strong-0` and `weak-1` drive.

```
CELL tristate_buffer {
   ...
   PIN z {DIRECTION = output; ATTRIBUTE {tristate}}
   ...
   VECTOR (z==0) {
      DRIVE_STRENGTH = 4; {PIN = z;}
   }
   VECTOR (z==1) {
      DRIVE_STRENGTH = 2; {PIN = z;}
   }
}
```

The bus itself is not described by an ALF model, since the bus is a design construct rather than a library cell. A simulation model (Verilog or VHDL) can handle the bus contention. However, since buses can also be embedded within a core cell, the functional model of the core need a functional model of that bus as well.

### A.7.2 Bus with multiple drivers

The following example shows a bus with three drivers of equal strength. The output is the resolved value of the bus.

```
CELL bus3 {
   PIN z1 {DIRECTION = input;}
   PIN z2 {DIRECTION = input;}
   PIN z3 {DIRECTION = input;}
   PIN z {DIRECTION = output;}
   FUNCTION {
      BEHAVIOR {
         z =
          ((z2=='bz || z2==z1) && z3=='bz)? z1:
          ((z3=='bz || z3==z2) && z1=='bz)? z2:
          ((z1=='bz || z1==z3) && z2=='bz)? z3:
          (z1=='b1 && z2=='b1 && z3=='b1)? 'b1:
          (z1=='b0 && z2=='b0 && z3=='b0)? 'b0:
                                          'bx;
      }
   }
}
```

The following example shows a bus with two drivers of equal strength and one driver with weaker strength (e.g., a busholder).

```
CELL bus2s1w {
   PIN z_strong1 {DIRECTION = input;}
   PIN z_strong2 {DIRECTION = input;}
   PIN z_weak   {DIRECTION = input;}
   PIN z        {DIRECTION = output;}
   FUNCTION {
      BEHAVIOR {
```

```
            z =
             (z_strong1=='b1 && z_strong2=='b1)? 'b1:
             (z_strong1=='b0 && z_strong2=='b0)? 'b0:
             (z_strong1=='bz && z_strong2=='bz)? z_weak:
                                                'bx;
         }
      }
   }
```

### A.7.3 Busholder

A *busholder* is a cell that retains the previous value of a tristate bus when all drivers go to high impedance. This device has only one external pin, which is bidirectional. The input to this bidirectional pin is the resolved value of the bus.

```
CELL busholder {
   PIN a {DIRECTION = both;}
   PIN z {DIRECTION = output; VIEW = none;}
   FUNCTION {
      BEHAVIOR {
         a = !z;
         @(a==0) {z = 1;}
         @(a==1) {z = 0;}
         @(a=='bx) {z = 'bx;}
      }
   }
}
```

In order to understand the functionality of a bidirectional pin, the pin can be split conceptually into an input pin and an output pin as shown below.

```
CELL busholder_explicit {
   PIN a_in {DIRECTION = input;}
   PIN a_out {DIRECTION = output;}
   PIN z {DIRECTION = output; VIEW = none;}
   FUNCTION {
      BEHAVIOR {
         a_out = !z;
         @(a_in==0) {z = 1;}
         @(a_in==1) {z = 0;}
         @(a_in=='bx) {z = 'bx;}
      }
   }
}
```

The function of this device is well defined if `a_out==a_in` for all cases where `a_in!='bz`. In the case of `a_in=='bz, a_out` can take any value. This is a general modeling rule for functions with bidirectional pins.

## A.8 Wire models

This section describes how to model wire models.

### A.8.1 Basic wire model

This example shows two wire models, using tables and equations. The equation is used outside the defined table range. If no equation was defined, the table is extrapolated.

```
WIRE small_wire {
   CAPACITANCE {
      UNIT = pF;
      HEADER {
         CONNECTIONS {
            TABLE {2 3 4 5}
         }
      }
      TABLE {0.05 0.09 0.13 0.17}
      EQUATION {CONNECTIONS * 0.04 - 0.03}
   }
   RESISTANCE {
      UNIT = mOHM;
      HEADER {
         CONNECTIONS {
            TABLE {2 3 4 5}
         }
      }
      TABLE {7.5 10.0 12.5 15.0}
      EQUATION {CONNECTIONS * 2.5 + 2.5}
   }
}
WIRE large_wire {
   CAPACITANCE {
      UNIT = pF;
      HEADER {
         CONNECTIONS {
            TABLE {2 3 4}
         }
      }
      TABLE {0.10 0.16 0.22}
      EQUATION {CONNECTIONS * 0.06 - 0.2}
   }
   RESISTANCE {
      UNIT = mOhm;
      HEADER {
         CONNECTIONS {
            TABLE {2 3 4}
         }
      }
      TABLE {10.0 12.5 15.0}
      EQUATION {CONNECTIONS * 2.5 + 5.0}
   }
}
```

### A.8.2 Wire select model

Since a library can contain multiple wire models, it is necessary to specify which model needs to be selected for an application. The annotations inside each wire model can be used for this purpose.

```
WIRE small_wire {
   LIMIT {AREA {UNIT=1e-6; MAX=25;}}
   ...
}
WIRE large_wire {
   LIMIT {AREA {UNIT=1e-6; MIN=25; MAX=100;}}
   ...
}
```

If the area covering the routing space is smaller than 25mm$^2$, the `small_wire` model shall be chosen. If the area covering the routing space is between 25mm$^2$ and 100mm$^2$, the `large_wire` model is chosen. The unit for area is 1mm$^2$.

To enable customized wire model selection, more annotations using the `USAGE` keyword can also be introduced.

## A.9 Megacell modeling

This section describes how to model megacells.

### A.9.1 Expansion of timing arcs

`GROUP` can be used for sets of numbers or for a continuous range of numbers. This can be useful for defining timing arcs between all bits of two vectors. For example,

```
GROUP adr_bits {1 2 3}
GROUP data_bits {1 2}
VECTOR (01 adr[adr_bits] -> 01 dout[data_bits]) { ... }
```

replaces the following statements:

```
VECTOR (01 adr[1] -> 01 dout[1]) { ... }
VECTOR (01 adr[2] -> 01 dout[1]) { ... }
VECTOR (01 adr[3] -> 01 dout[1]) { ... }
VECTOR (01 adr[1] -> 01 dout[2]) { ... }
VECTOR (01 adr[2] -> 01 dout[2]) { ... }
VECTOR (01 adr[3] -> 01 dout[2]) { ... }
```

The following example shows bit-wise expansion of two vectors:

```
GROUP data_bits {1 2}
VECTOR (01 din[data_bits] -> 01 dout[data_bits]) { ... }
```

This replaces the following statements:

```
VECTOR (01 din[1] -> 01 dout[1]) { ... }
VECTOR (01 din[2] -> 01 dout[2]) { ... }
```

Example for byte-wise (or sub-word-wise) expansion:

```
GROUP low_byte {1 2}
GROUP high_byte {3 4}
VECTOR (01 we[0] -> 01 din[low_byte]) { ... }
VECTOR (01 we[1] -> 01 din[high_byte]) { ... }
```

This replaces the following statements:

```
VECTOR (01 we[0] -> 01 din[1]) { ... }
VECTOR (01 we[0] -> 01 din[2]) { ... }
VECTOR (01 we[1] -> 01 din[3]) { ... }
VECTOR (01 we[1] -> 01 din[4]) { ... }
```

### A.9.2 Two-port memory

The memory model example below shows the use of abstract transition operators on words in various vectors. This example also contains some vectors with distinction between events on row and column address lines.

```
CELL async_1write_1read_ram {
   GROUP col {1:0}
   GROUP row {4:2}
   GROUP all {row col}
   GROUP byte{7:0}
   GROUP \* {0:31}
   PIN enable_write {DIRECTION = input}
   PIN [4:0] adr_write {DIRECTION = input}
   PIN [4:0] adr_read {DIRECTION = input}
   PIN [7:0] data_write {DIRECTION = input}
   PIN [7:0] data_read {DIRECTION = output}
   PIN [7:0] data_store [0:31] {DIRECTION = output VIEW = none}

   FUNCTION {
      BEHAVIOR {
         data_read = data_store[adr_read];
         @(enable_write) {data_store[adr_write] = data_write;}
      }
   }

   VECTOR
   (?! adr_read[col] -> ?? data_read[byte]) {
      /* fill in arithmetic models for delay and power */
   }
   VECTOR
   (?! adr_read[row] -> ?? data_read[byte]) {
      /* fill in arithmetic models for delay and power */
   }
   VECTOR
   ((?!adr_read[col] && ?!adr_read[row]) -> ?? data_read[byte]){
      /* fill in arithmetic models for delay and power */
   }
   VECTOR (01 enable_write -> ?? data_read[byte]) {
      /* fill in arithmetic models for delay and power */
   }
```

```
   VECTOR (?! data_write[byte] -> ?? data_read[byte]) {
      /* fill in arithmetic models for delay and power */
   }

   VECTOR (?! adr_write[col]) {
      /* fill in arithmetic models for power */
   }
   VECTOR (?! adr_write[row]) {
      /* fill in arithmetic models for power */
   }
   VECTOR (?! adr_write[row] && ?! adr_write[col]) {
      /* fill in arithmetic models for power */
   }
   VECTOR (01 enable_write) {
      /* fill in arithmetic models for power */
   }
   VECTOR (10 enable_write) {
      /* fill in arithmetic models for power */
   }
   VECTOR (?! data_write[byte] && !enable_write) {
      /* fill in arithmetic models for power */
   }
   VECTOR (?! data_write[byte] && enable_write) {
      /* fill in arithmetic models for power */
   }
}
   VECTOR (?! adr_write[all] <-> 01 enable_write) {
      SETUP {
         VIOLATION {
            BEHAVIOR { data_store[\*] = 'bxxxxxxxx; }
            MESSAGE_TYPE = error;
            MESSAGE =
"setup violation: changing 'adr_write' -> rising 'enable_write',
memory -> 'X'";
         }
         FROM { pin = adr_write; }
         TO { pin = enable_write; }
         /* fill in header, table or equation */
      }
   }
   VECTOR (10 enable_write <-> ?! adr_write[all]) {
      HOLD {
         VIOLATION {
            BEHAVIOR { data_store[\*] = 'bxxxxxxxx; }
            MESSAGE_TYPE = error;
            MESSAGE =
"hold violation: falling 'enable_write' -> changing 'adr_write',
memory -> 'X'";
         }
         FROM { pin = enable_write; }
         TO { pin = adr_write; }
         /* fill in header, table or equation */
      }
   }
```

```
   VECTOR (?! data_write[byte] <-> 10 enable_write) {
      SETUP {
         VIOLATION {
            BEHAVIOR { data_store[adr_write] = 'bxxxxxxxx; }
            MESSAGE_TYPE = error;
            MESSAGE =
"setup violation: changing 'data_write' -> falling 'enable_write',
memory[adr_write] -> 'X'";
         }
         FROM { pin = data_write; }
         TO { pin = enable_write; }
         /* fill in header, table or equation */
      }
      HOLD {
         VIOLATION {
            BEHAVIOR { data_store[adr_write] = 'bxxxxxxxx; }
            MESSAGE_TYPE = error;
            MESSAGE =
"hold violation: falling 'enable_write' -> changing 'data_write',
memory[adr_write] -> 'X'";
         }
         FROM { pin = enable_write; }
         TO { pin = data_write; }
         /* fill in header, table or equation */
      }
   }
   VECTOR (01 enable_write -> 10 enable_write) {
      PULSEWIDTH {
         VIOLATION {
            MESSAGE_TYPE = error;
            MESSAGE = "pulsewidth violation: high 'enable_write'";
         }
         PIN = enable_write;
         /* fill in header, table or equation */
      }
   }
   VECTOR (10 enable_write -> 01 enable_write) {
      PULSEWIDTH {
         VIOLATION {
            MESSAGE_TYPE = error;
            MESSAGE = "pulsewidth violation: low 'enable_write'";
         }
         PIN = enable_write;
         /* fill in header, table or equation */
      }
   }
}
```

The energy consumption for each operation depends on the number of switching bits of the bus. Therefore, the model for power inside a particular vector can look like this:

```
VECTOR (?! data_write && enable_write) {
   ENERGY {
      UNIT = pJ;
      HEADER {switching_bits {PIN = data_write;}}
      EQUATION {1.3 * switching_bits}
   }
}
```

The rule that the address on a write port shall not change during write-enable high can be incorporated easily in the functional model. A pessimistic model assumes the whole memory content shall become unknown if such an illegal address change occurs.

```
BEHAVIOR {
   data_read = data_store[adr_read];
   @(enable_write) {data_store[adr_write] = data_write;}
   @(!?adr_write && enable_write)
      {data_store[\*] = 'bxxxxxxxx;}
}
```

### A.9.3 Three-port memory

Functional models of more complex memories are also straightforward. The conflicts of writing to one memory location simultaneously from different ports can be modeled in a pessimistic way as follows:

```
CELL async_2write_1read_ram {
   PIN enb_write1 {DIRECTION = input;}
   PIN enb_write2 {DIRECTION = input;}
   PIN [4:0] adr_write1 {DIRECTION = input;}
   PIN [4:0] adr_write2 {DIRECTION = input;}
   PIN [4:0] adr_read {DIRECTION = input;}
   PIN [7:0] data_write1 {DIRECTION = input;}
   PIN [7:0] data_write2 {DIRECTION = input;}
   PIN [7:0] data_read {DIRECTION = output;}
   PIN [7:0] data_store [0:31] {DIRECTION = output; VIEW = none;}
   FUNCTION {
      BEHAVIOR {
         data_read = data_store[adr_read];
         @(enb_write1 & !enb_write2)
            {data_store[adr_write1] = data_write1;}
         @(enb_write2 & !enb_write1)
            {data_store[adr_write2] = data_write2;}
         @(enb_write1 & enb_write2 && adr_write1!=adr_write2) {
            data_store[adr_write1] = data_write1;
            data_store[adr_write2] = data_write2;
         }
         @(enb_write1 & enb_write2 && adr_write1==adr_write2) {
            data_store[adr_write1] =
               (data_write1==data_write2)? data_write1:8'bx;
            data_store[adr_write2]
               (data_write2==data_write1)? data_write2:8'bx;
```

```
            }
         }
      }
   }
```

### A.9.4 Annotation for pins of a bus

Annotations of numeric values to a bus apply to the total bus, not to each individual pin.

Example:

```
PIN [1:4] my_bus_pin {
   CAPACITANCE = 0.04 ;
}
```

The total bus pin capacitance is `0.4`; the capacitance values on each individual pin are not defined.

The individual pin capacitance can be defined as follows:

```
PIN [1:4] my_bus_pin {
   CAPACITANCE c1 = 0.01 { PIN = my_bus_pin[1]; }
   CAPACITANCE c2 = 0.01 { PIN = my_bus_pin[2]; }
   CAPACITANCE c3 = 0.01 { PIN = my_bus_pin[3]; }
   CAPACITANCE c4 = 0.01 { PIN = my_bus_pin[4]; }
}
```

### A.9.5 Skew for simultaneously switching signals on a bus

Vectors with simultaneously switching bits on a bus can contain a specification of the allowed skew in order to be still considered as simultaneously switching bits.

Example:

```
PIN [1:3] address;
VECTOR (?! address )
   SKEW {
         PIN = address;
         /* fill in data */
   }
}
```

`SKEW` applied to a bus pin is the maximal allowed time window between the earliest and latest edge within simultaneously switching signals of a bus.

The multiple value annotation feature allows the definition of a group of pins equivalent to a bus for `SKEW` modeling in the following way:

```
PIN A;
PIN [1:4] B;
VECTOR (?! A && ?! B)
   SKEW { PIN { A B[2:3] } }
}
```

SKEW applies to the group of pins A, B[2], and B[3]. The following example is semantically different, since this results in expansion of each object where the group is instantiated:

```
PIN A;
PIN [1:4] B;
GROUP my_group { A B[2] B[3] }
VECTOR (?! my_group)
   SKEW { PIN = my_group; }
}
```

The expansion yields the following:

```
PIN A;
PIN [1:4] B;
VECTOR (?! A)
   SKEW { PIN = A ; }
}
VECTOR (?! B[2])
   SKEW { PIN = B[2] ; }
}
VECTOR (?! B[3])
   SKEW { PIN = B[3] ; }
}
```

See Section B.2.7 for the definition of SKEW for scalar pins.

# A.10 Special cells

This section describes how to model special cells.

## A.10.1 Pulse generator

The following cell generates a one-shot pulse of 1 ns duration when enable goes high.

```
CELL one_shot {
   PIN enable {DIRECTION = input;}
   PIN q {DIRECTION = output;}
   FUNCTION {
      BEHAVIOR {
         @(01 enable) {q = 1;}
         @(q) {q = 0;}
      }
   }
   VECTOR (01 q -> 10 q) {
      DELAY = 1.0 {UNIT = ns;}
   }
}
```

## A.10.2 VCO

The following cell is a voltage controlled oscillator with 50% duty cycle and enable.

```
CELL vco {
   PIN enable {DIRECTION = input; PINTYPE = digital;}
   PIN v_in {DIRECTION = input; PINTYPE = analog;}
   PIN q {DIRECTION = output; PINTYPE = digital;}
   FUNCTION {
      BEHAVIOR {
         @(!enable) {q = 0;}
         @(!q & enable) {q = 1;}
         @( q & enable) {q = 0;}
      }
   }
   TEMPLATE voltage_controlled_delay {
      DELAY {
         UNIT = ns;
         HEADER {
            voltage {
               PIN = v_in;
               TABLE {0.5 1.0 1.5 2.0 2.5 3.0}
            }
         }
         TABLE {10.00 5.00 3.33 2.50 2.00 1.67}
      }
   }
   VECTOR (01 q -> 10 q)
      voltage_controlled_delay
   }
   VECTOR (10 q -> 01 q)
      voltage_controlled_delay
   }
}
```

The template shown above can also be written as an equation to map voltage to frequency:

```
   TEMPLATE voltage_controlled_delay {
      DELAY {
         UNIT = ns;
         HEADER {voltage {PIN = v_in;}}
         EQUATION {5.0 / voltage}
      }
   }
```

## A.11 Core modeling (using a digital filter)

This example illustrates the potential of ALF for modeling complex blocks. It shows a digital filter performing the following operation

```
dout(t) = state(t) + b1 * state(t-1) + b2 * state(t-2)
state(t) = din(t) - a1 * state(t-1) - a2 * state(t-2)
```

This second order infinite impulse response (*IIR*) filter is implemented with a single multiplier and a single adder/subtractor in a way that a new `dout` is produced every four clock cycles. The variable coefficients `a1, a2, b1,` and `b2` are stored in a dual port RAM.

The model uses templates for the functional blocks of a two-bit counter used as the controller for memory access and I/O operation, a RAM for coefficient storage, and the filter itself. They are instantiated as a structural netlist in the top module.

The use of templates is more general than the use of primitives, since not all basic blocks of the core might be supported as primitives.

```
LIBRARY core_lib {
   TEMPLATE CNT2 {
      BEHAVIOR {
         @ (!<cd>) {<cnt> = 2'b0;}
         : (01 <cp>) {<cnt> = <start> ? 2'b0 : <cnt> + 1;}
      }
   }

   TEMPLATE RAM16X4 {
      BEHAVIOR {
         <dout> = <dmem>[<r_adr>];
         @ (<we>) {<dmem>[<w_adr>] = <din>;}
      }
   }

   TEMPLATE IIR2 {
      BEHAVIOR {
         sum =
            (<cntrl>=='d0)? <din> - product :
            (<cntrl>=='d1)? accu - product :
            (<cntrl>=='d2)? accu + product :
            (<cntrl>=='d3)? accu + product;
         @ (!<cd>) {
            product = 16'b0;
            accu = 16'b0;
         }
         : (01 <cp>){
            product =
               (<cntrl>=='d0)? coeff * state2 :
               (<cntrl>=='d1)? coeff * state1 :
               (<cntrl>=='d2)? coeff * state2 :
               (<cntrl>=='d3)? coeff * state1 :
               16'bX;
            accu = sum;
         }
         @ (!<cd>) {
            <dout> = 16'b0;
            state1 = 16'b0;
            state2 = 16'b0;
         }
         : (01 <cp> && <cntrl>=='d0){
            state2 = state1;
            state1 = accu;
            <dout> = accu;
         }
      }
   }
```

```
      CELL digital_filter {
         PIN [15:0] data_out {DIRECTION = output;}
         PIN [15:0] data_in {DIRECTION = input;}
         PIN [1:0] index_coeff {DIRECTION = input;}
         PIN write_coeff {DIRECTION = input;}
         PIN [15:0] coeff_in {DIRECTION = input;}
         PIN [15:0] coeff_out {DIRECTION = output; VIEW = none;}
         PIN [15:0] coeff_array [1:4] {DIRECTION = output; VIEW = none;}
         PIN data_strobe {DIRECTION = input;}
         PIN [1:0] count {DIRECTION = output VIEW = none;}
         PIN clock {DIRECTION = input;}
         PIN reset {DIRECTION = input;}
         FUNCTION {
            IIR2 {   din=data_in; dout=data_out; coeff=coeff_out;
                     cp=clock; cd=reset; cntrl = count;}
            CNT2 {   start=data_strobe; cnt=count; ck=clock; cd=reset;}
            RAM16X4{ we=write_coeff; din=coeff_in; dout=coeff_out;
                     dmem=coeff_array; r_adr=count; w_adr=index_coeff;}
         }
      }
   }
```

# A.12 Connectivity

Connectivity information can be specified within the definition of the ALF language format as described below. A connectivity object always contains a rule specifying the type of connections (e.g., must short, can short, or cannot short) and a table. If no header is given, then the table contains the pins or pin classes subject to the connectivity rule. If a header is given, then the table contains the values of the connectivity function between arguments in the header. There shall be a table inside each connectivity argument, containing the pins or pin classes subject to the connectivity rule. Valid arguments are `DRIVER` and/or `RECEIVER`. Valid values are the boolean digits `0`, `1`, and `?`. The value `1` implies the connection rule is *True*, the value `0` implies the connection rule is *False*, the value `?` implies a “don’t care” situation with the connection rule.

## A.12.1 External connections between pins of a cell

The following example shows how to specify required and disallowed interconnections external to a cell.

```
   CELL pll {
      PIN vdd_ana {PINTYPE=supply;}
      PIN vdd_dig {PINTYPE=supply;}
      PIN vss_ana {PINTYPE=supply;}
      PIN vss_dig {PINTYPE=supply;}
      CONNECTIVITY common_ground {
         CONNECT_RULE = must_short;
         TABLE {vss_ana vss_dig}
```

```
      CONNECTIVITY separate_supply {
         CONNECT_RULE = cannot_short;
         TABLE {vdd_ana vdd_dig}
      }
   }
```

### A.12.2 Allowed connections for classes of pins

The following example defines allowable pin interconnections. The constants for the desired connectivity classes, the grouping of these classes, and the allowable class connectivity table are first defined at the library level. The non-zero values within the matrix specify allowable connectivity of indexed classes. The connectivity classes for pins are then specified with the pin annotation sections.

```
   LIBRARY example_library {
      ...
      CLASS default_class;
      CLASS clock_class;
      CLASS enable_class;
      CLASS reset_class;
      CLASS tristate_class;
      ...
      TEMPLATE drivers {
         default_class
         clock_class
         enable_class
         reset_class
         tristate_class
      }
      TEMPLATE receivers {
         default_class
         clock_class
         enable_class
         reset_class
      }
      CONNECTIVITY driver_to_driver {
         CONNECT_RULE = can_short;
         HEADER {
            DRIVER {TABLE {drivers}}
         }
         TABLE {// def clk enb rst tri
                0 0 0 0 1
         }
      }
      CONNECTIVITY receiver_to_receiver {
         CONNECT_RULE = can_short;
         HEADER {
            RECEIVER {TABLE {receivers}}
         }
         TABLE {// def clk enb rst
                1 1 1 1
         }
      }
```

```
CONNECTIVITY driver_to_receiver {
   CONNECT_RULE = can_short;
   HEADER {
      DRIVER {TABLE {drivers}}
      RECEIVER {TABLE {receivers}}
   }
   TABLE {// def clk enb rst tri // driver/receiver
               1    1   1    1   0    // def
               0    1   0    0     0  // clk
               0    0   1    0     0  // enb
               0    0   0    1     0  // rst
   }
}
```

The above table specifies the applicable connectivity from each class to itself, as well as from each class to the `default_class`, except for the `tristate_class` class (which can only connect to itself). While any class can connect to the `default_class`, the `default_class` can only connect to itself.

Once the library level connectivity is defined, connection class specifications are defined for each pin within cells. The default integer value for the `CLASS` annotation is `0`, which corresponds to the constant declaration value for `default_class`.

```
CELL d_flipflop_clr {
   PIN cd {DIRECTION = input; SIGNALTYPE = clear;
          POLARITY = low; CONNECT_CLASS = reset_class;}
   PIN cp {DIRECTION = input; SIGNALTYPE = clock;
          POLARITY = rising_edge; CONNECT_CLASS = clock_class;}
   PIN d {DIRECTION = input;}
   PIN q {DIRECTION = output; CONNECT_CLASS = default_class;}
}

CELL d_latch {
   PIN g {DIRECTION = input; SIGNALTYPE = enable;
          POLARITY = high; CONNECT_CLASS = enable_class;}
   PIN d {DIRECTION = input; CONNECT_CLASS = default_class;}
   PIN q {DIRECTION = output; CONNECT_CLASS = default_class;}
}

CELL tristate_buffer {
   PIN a {DIRECTION = input;}
   PIN enable {DIRECTION = input; CONNECT_CLASS = enable_class;}
   PIN z {DIRECTION = output; CONNECT_CLASS = tristate_class;}
   ...
}
```

Net-specific connectivity, as opposed to the pin-specific connectivity shown above, is also possible within the syntax of the language, since a `CLASS` is not restricted to pins. Specific applications can assign all pins of a specific type, as well as nets like power and ground rails to a defined class. This class can be used within the connectivity tables to allow or disallow certain connectivity.

For example, if `vddrail_class` is defined as a net-specific connectivity class, then a specific pin can be disallowed from connecting to any net in the `vddrail_class` connectivity class.

```
CLASS vddrail_class
...
CELL inverter {
   PIN in_pin {DIRECTION = input; SIGNALTYPE = clear;
               POLARITY = low; CONNECT_CLASS = reset_class;}
   CONNECTIVITY dont_tie {
      CONNECT_RULE = cannot_short;
      TABLE {in_pin vddrail_class}
   }
   ...
}
```

# A.13 Signal integrity

This section describes how to model signal integrity.

## A.13.1 I/V curves

I/V curves describe the driven or drawn current at a pin as a function of the voltage at one or several pins. The following example describes the output current of a buffer as a function of the input and output voltage with a two-dimensional lookup table.

```
CELL simple_buffer {
   PIN z { DIRECTION = output; }
   PIN a { DIRECTION = input; }
   // current @ z dependent on voltage @ z and @ a
   CURRENT {
      PIN = z;
      UNIT = ma;
      HEADER {
         VOLTAGE vout {
            PIN = z;
            TABLE { 0.0 0.5 1.0 1.5 2.0 2.5 3.0 }
         }
         VOLTAGE vin {
            PIN = a;
            TABLE { 0.0  1.0  2.0  3.0 }
         }
      }
      TABLE {
         5.0 5.0 4.8 4.2 3.2 1.6 0.0
         2.5 1.5 0.2 -0.4 -1.8 -2.7 -3.5
         1.2 0.1 -1.3 -1.9 -2.5 -3.8 -4.6
         0.0 -2.0 -3.8 -4.7 -5.5 -6.2 -6.3
      }
   }
   // fill in function, vector and other stuff
}
```

An equation can also be used instead of a lookup table, for example:

```
      CURRENT {
         PIN = z;
         UNIT = ma;
         HEADER {
            VOLTAGE vout {
               PIN = z;
            }
            VOLTAGE vin {
               PIN = a;
            }
         }
         EQUATION {
            (1 - exp(6.3 - 2.4*vout))*exp(0.9 - 0.3*vin)
            - (1 - exp(3.2*vout))*exp(0.3*vin)
         }
      }
```

A buffer can have programmable drive strength controlled by the state of additional input pins. State-dependent I/V curves can be described by vector-specific `CURRENT` models.

```
   CELL programmable_drive_strength_buffer {
      PIN z { DIRECTION = output; }
      PIN a { DIRECTION = input; }
      // control pins for drive strength
      PIN p1 { DIRECTION = input; }
      PIN p2 { DIRECTION = input; }
      VECTOR (!p1 & !p2) {
         CURRENT {
            // fill in the model
         }
      }
      VECTOR (!p1 & p2) {
         CURRENT {
            // fill in the model
         }
      }
      VECTOR ( p1 & !p2) {
         CURRENT {
            // fill in the model
         }
      }
      VECTOR ( p1 & p2) {
         CURRENT {
            // fill in the model
         }
      }
   }
```

It is also possible to describe other analog cell characteristics (state-dependent or state-independent), for instance, voltage versus voltage, frequency versus voltage, or current versus temperature, in the same way.

### A.13.2 Driver resistance

Driver resistance is used to model the transient behavior of signals especially for crosstalk. The drivers are modeled by voltage sources and driver resistances, as illustrated in Figure A-1.

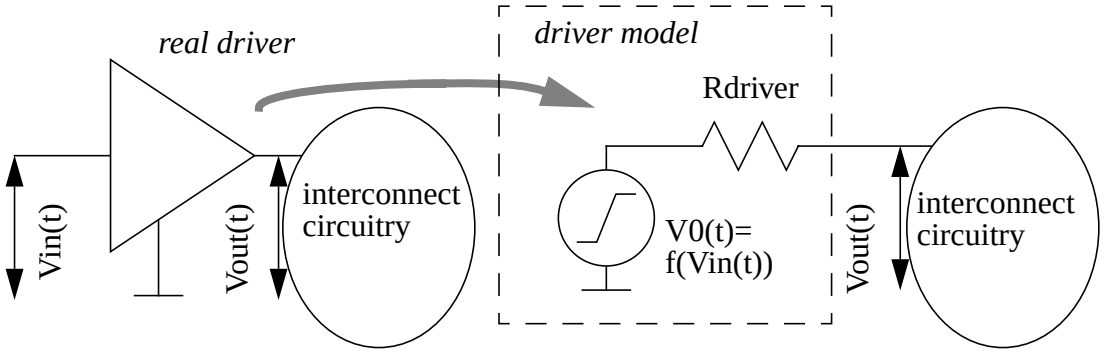



**Figure A-1: Modeling driver resistance**

The idea here is to use linear circuit theory for the analysis of multiple drivers interacting with coupled RC-interconnect networks. In reality, the drivers have non-linear resistance. The linear resistance is a model of the non-linear resistance with the best-fitting linear resistance. Therefore, the driver resistance is state-dependent and eventually also load- and slewrate dependent, because the best-fitting value for driver resistance is different for different states and different ranges of load and slewrates.

The following example shows a buffer featuring different driver resistance values for static low and high states, and tables of slewrate- and load-dependent transient driver resistance values for rise and fall transitions.

```
cell simple_buffer {
   PIN z { DIRECTION = output; }
   PIN a { DIRECTION = input; }
   // state-dependent static driver resistance
   VECTOR (!z) {
      RESISTANCE = 0.7k { PIN = z; }
   }
   VECTOR (z) {
      RESISTANCE = 1.2k { PIN = z; }
   }

   // slew & load dependent transient driver resistance
   VECTOR (01 a -> 01 z) {
      RESISTANCE {
         PIN = z;
         UNIT = kohm;
         HEADER {
            CAPACITANCE {
               PIN = z;
               UNIT = pfarad;
```

```
                TABLE { 0.1  0.4  1.6 }
             }
             SLEWRATE {
                PIN = a;
                UNIT = nsec;
                TABLE { 0.5  1.5}
             }
          TABLE { 1.4  1.3  1.3  1.6  1.4  1.3 }
       }
    }
    VECTOR (10 a -> 10 z) {
       RESISTANCE {
          PIN = z;
          UNIT = kohm;
          HEADER {
             CAPACITANCE {
                PIN = z;
                UNIT = pfarad;
                TABLE { 0.1  0.4  1.6 }
             }
             SLEWRATE {
                PIN = a;
                UNIT = nsec;
                TABLE { 0.5  1.5}
             }
          TABLE { 0.9  0.8  0.8  1.1  0.9  0.8 }
       }
    }
}
```

The transient driver resistance can also be state-dependent, for example, in the case of a buffer with programmable drive-strength.

```
CELL programmable_drive_strength_buffer {
   PIN z { DIRECTION = output; }
   PIN a { DIRECTION = input; }
   // control pins for drive strength
   PIN p1 { DIRECTION = input; }
   PIN p2 { DIRECTION = input; }
   // state-dependent static driver resistance
   VECTOR (!z & !p1 & !p2) {
      RESISTANCE = 0.7k { PIN = z; }
   }
   VECTOR (!z & !p1 & p2) {
      RESISTANCE = 0.6k { PIN = z; }
   }
   VECTOR (!z & p1 & !p2) {
      RESISTANCE = 0.5k { PIN = z; }
   }
   VECTOR (!z & p1 & !p2) {
      RESISTANCE = 0.4k { PIN = z; }
   }
   VECTOR (z & !p1 & !p2) {
      RESISTANCE = 1.2k { PIN = z; }
   }
```

```
        VECTOR (z & !p1 & p2) {
           RESISTANCE = 1.0k { PIN = z; }
        }
        VECTOR (z & p1 & !p2) {
           RESISTANCE = 0.8k { PIN = z; }
        }
        VECTOR (z & p1 & p2) {
           RESISTANCE = 0.6k { PIN = z; }
        }
        // slew & load and state dependent transient driver resistance
        VECTOR (01 a -> 01 z && !p1 & !p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (01 a -> 01 z && !p1 & p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (01 a -> 01 z &&  p1 & !p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (01 a -> 01 z &&  p1 & p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (10 a -> 10 z && !p1 & !p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (10 a -> 10 z && !p1 & p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (10 a -> 10 z &&  p1 & !p2) {
           RESISTANCE {
              // fill in the model
        }
        VECTOR (10 a -> 10 z &&  p1 & p2) {
           RESISTANCE {
              // fill in the model
        }
    }
```

The model for transient driver resistance has the same form as a slewrate- and load-dependent model for delay. Voltage-, process-, and temperature-dependent driver resistance can also be modeled in the same way as voltage-, process-, and temperature-dependent delay.

## A.14 Resistance and capacitance on a pin

This section describes how to model pin resistance and capacitance.

### A.14.1 Self-resistance and capacitance on input pin

A *pin resistance* is a resistance inside a `PIN` object.

```
PIN <pin_identifier> {
   DIRECTION = input;
   RESISTANCE = <resistance_value>;
   CAPACITANCE = <capacitance_value>;
}
```

The pin resistance is in series with the pin capacitance, as shown in Figure A-2.

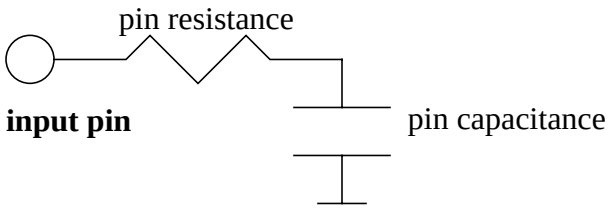



**Figure A-2: Resistance and capacitance on a pin**

### A.14.2 Pullup and pulldown resistance on input pin

A pullup or pulldown resistance, or a combination of both on an input pin, can be described as follows:

```
PIN <pin_identifier> {
   DIRECTION = input;
   PULL = < up | down | both > {
      VOLTAGE = <voltage_value>;
      RESISTANCE = <resistance_value>;
   }
}
```

The pullup/pulldown resistance is in series with a clamp voltage, as shown in Figure A-3.

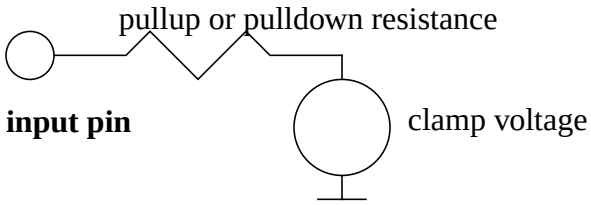



**Figure A-3: Pullup or pulldown resistance**

In the case of a pullup/pulldown combination, the resistance and voltage represent the Thevenin equivalent resistance and voltage, respectively, as shown in Figure A-4.

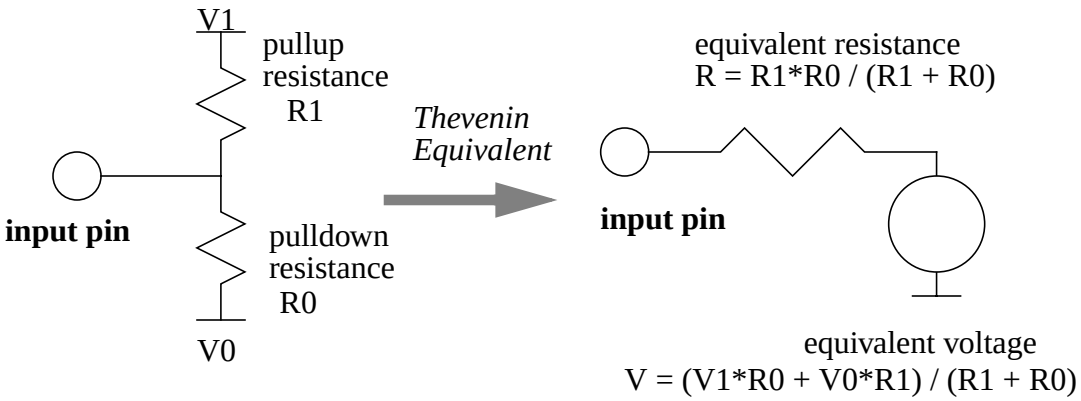



**Figure A-4: Thevenin equivalent resistance**

### A.14.3 Pin and load resistance and capacitance on an output pin

The driver resistance (see Section A.13.2) can also be represented as a pin capacitance of an output pin, where there is no state dependency.

```
PIN <pin_identifier> {
   DIRECTION = output;
   CAPACITANCE = <capacitance_value>;
   RESISTANCE {
      RISE = <rise_resistance_value>;
      FALL = <rise_resistance_value>;
   }
}
```

Note the distinction of capacitance and resistance of the pin itself and capacitance and resistance applied as load to the pin in Figure A-5. The load capacitance and resistance are specified in a characterization vector (see Section A.3).

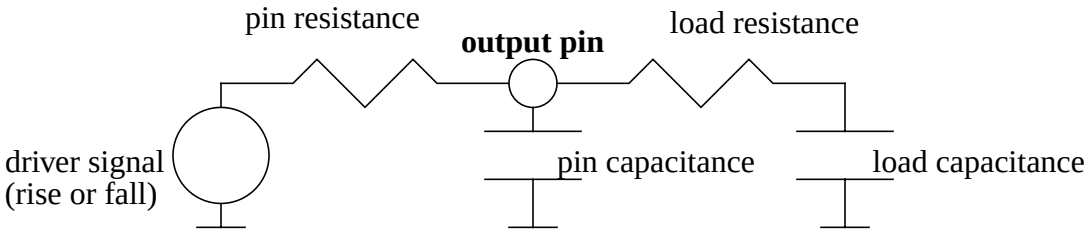



**Figure A-5: Resistance and capacitance on an output pin**

# Appendix B
# ALF/SDF Cross Reference

This section provides a cross reference between the representation of timing data in ALF and SDF. In general, ALF is used as a characterization library, which is the input to a delay calculator, whereas SDF is the output from a delay calculator. Therefore, ALF typically contains tables or equations (i.e., arithmetic models) for timing data, whereas SDF contains a discrete set of data in fixed format. However, in an ALF representation of timing shells for cores, which are typically represented in SDF today, the ALF library contains the same data as the SDF.

The specification of the stimulus for a particular timing measurement (i.e., the timing diagram) is pertinent to both ALF and SDF. In ALF, timing diagrams are directly described in the vector expression language, and the timing measurements are always specified in relation to a particular timing vector. In SDF, timing diagrams are partly described in the language and partly implied by the keyword for timing measurements. Therefore SDF needs a larger set of keywords than ALF for the same description capability.

## B.1 SDF delays

This section details the different types of SDF delays.

### B.1.1 SDF DELAY for IOPATH and INTERCONNECT

A `DELAY` is a measurement of the time needed for a signal to travel from one port to another port. In ALF, delay measurements are described in a uniform language, independent of whether `A` and `Z` are the input and output port of the same cell, respectively, the driver and receiver connected to the same net, or both outputs of a cell. Therefore, the SDF keywords `IOPATH` and `INTERCONNECT` have no counterpart in ALF.

```
VECTOR (01 A -> 01 Z) {
   DELAY {
      FROM {PIN = A;}
      TO {PIN = Z;}
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-1:

*rising edge at A followed by rising edge at Z.*

The `FROM` and `TO` pin annotations define the sense of measurement for `DELAY`.

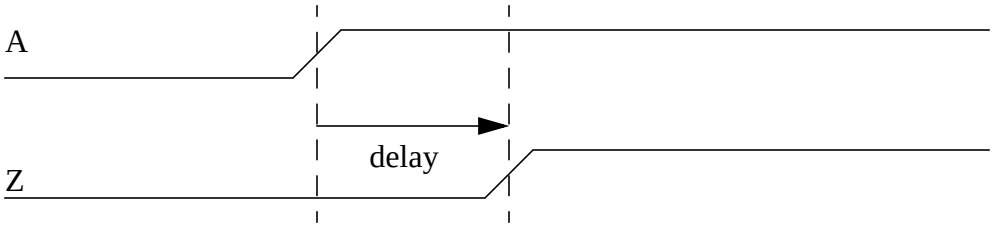



**Figure B-1: Measurement of SDF IOPATH or INTERCONNECT delay**

As opposed to SDF, where input ports of an `IOPATH` can have an edge specification and output ports can not, the vector expression language in ALF always contains the specification of the edge:

```
rising edge = “01”, falling edge = “10”, any edge = “?!”.
```

### B.1.2 SDF PATHPULSE

A `PATHPULSE` in SDF defines the smallest pulse that can appear at a port in form of

- a full-swing pulse
- a pulse to X.

The equivalent model in ALF uses two vectors in conjunction with the keyword `PULSEWIDTH`.[1]

The ALF keywords are of more general use than the SDF `PATHPULSE` keyword, which is just for one specific use.

```
VECTOR (01 Z -> 10 Z) {
   PULSEWIDTH {
      PIN = Z;
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-2:

*rising edge at Z followed by falling edge at Z.*

The smallest possible full-swing pulse applies at pin `Z`.

1. The same keyword `PULSEWIDTH` is also used for a timing constraint in ALF. The semantic meaning in both usage cases is consistent: `PULSEWIDTH =` *smallest possible pulse at output or smallest allowed pulse at input*. Therefore, the usage of the same keyword is justified.


**Figure B-2: Measurement of SDF PATHPULSE full-swing**

```
VECTOR ('b0'bX Z -> 'bX'b0 Z) {
   PULSEWIDTH {
      PIN = Z;
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-3:

*rising edge at Z from 0 to X followed by falling edge at Z from X to 0.*

The smallest possible pulse to `X` applies at pin `Z`.

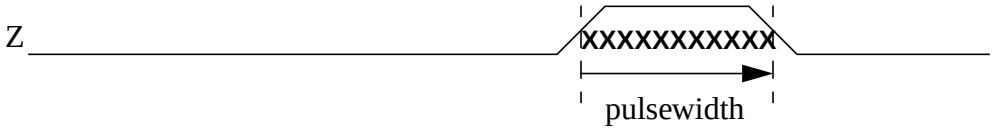



**Figure B-3: Measurement of SDF PATHPULSE to X**

```
VECTOR (01 A -> 10 B -> 01 Z -> 10 Z) {
   PULSEWIDTH {
      PIN = Z;
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-4:

*rising edge at A followed by falling edge at B followed by rising edge at Z followed by falling edge at Z.*

This is a detailed specification of the pulse itself at pin `Z`, as well as of the triggering input signals `A` and `B`.

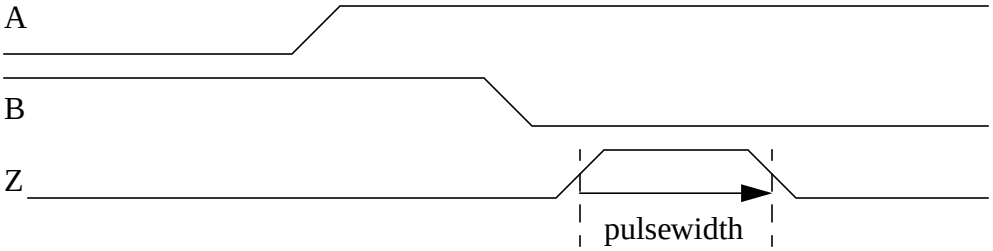



**Figure B-4: Measurement of SDF PATHPULSE with triggering inputs**

### B.1.3 SDF RETAIN delays

A `RETAIN` delay in SDF is a measurement of the time when an output signal shall retain its value after a change at a related input signal occurs. It appears always in conjunction with a `IOPATH` delay, which is the time for which an output shall stabilize after changing its value.

`RETAIN` is mainly used for asynchronous memories, where decoder glitches can appear at the data output port.

```
VECTOR (01 A -> ?! Z) {
   RETAIN {
      FROM {PIN = A;}
      TO {PIN = Z;}
      /* fill in data */
   }
   DELAY {
      FROM {PIN = A;}
      TO {PIN = Z;}
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-5:

*rising edge at A followed by any edge at Z.*

The intermediate events at `Z`, occurring eventually between retain and delay time, are not specified.

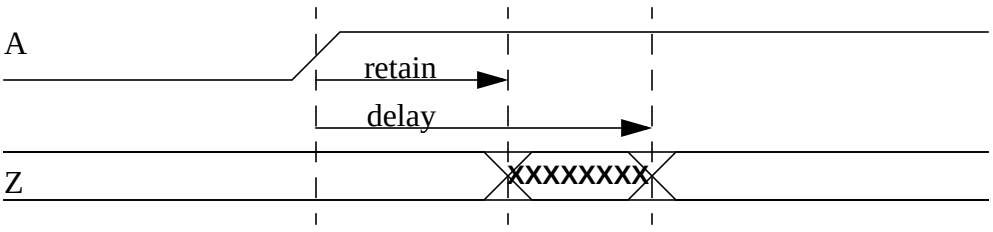



**Figure B-5: RETAIN and IOPATH delay**

### B.1.4 SDF PORT delays

A **PORT** delay in SDF is a delay measurement with unspecified start point, since the start point is going to be established by a connection to a driver in the design and not in the library.

```
VECTOR (01 A) {
   DELAY {
      TO {PIN = A;}
      /* fill in data */
   }
}
```

This ALF VECTOR describes the event shown in Figure B-6:

*rising edge at A.*

The absence of a FROM pin defines the absence of a start point, which corresponds to the exact meaning of PORT in SDF.

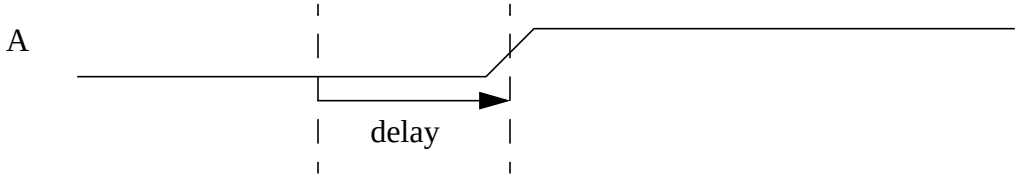



**Figure B-6: SDF PORT delay**

ALF also has the capability of describing a delay measurement with unspecified end point.

```
VECTOR (01 Z) {
   DELAY {
      FROM {PIN = Z;}
      /* fill in data */
   }
}
```

Hence, ALF provides the description capability for both a delay from unspecified driver to specified receiver and a delay from specified driver to unspecified receiver.

### B.1.5 SDF DEVICE delays

A DEVICE delay in SDF is a delay that applies from all input ports of a device to one specific output port or to all output ports by default.

The ALF vector expression language has no notion of “all input ports of a device”. ALF has a more general capability of declaring groups of pins and define delays from group to group or from group to pin or from pin to group.

```
GROUP any_input { A B }
GROUP any_output { Y Z }
VECTOR (01 any_input -> 01 any_output) {
   DELAY {
      FROM {PIN = any_input;}
      TO {PIN = any_output;}
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the event

*rising edge at any_input (i.e., A or B) followed by rising edge at any_output (i.e., Y or Z).*

This construct is equivalent to the following four vectors:

```
VECTOR (01 A -> 01 Y) {
   DELAY {
      FROM {PIN = A;}
      TO {PIN = Y;}
      /* fill in data */
   }
}
VECTOR (01 B -> 01 Y) {
   DELAY {
      FROM {PIN = B;}
      TO {PIN = Y;}
      /* same data */
   }
}
VECTOR (01 A -> 01 Z) {
   DELAY {
      FROM {PIN = A;}
      TO {PIN = Z;}
      /* same data */
   }
}
VECTOR (01 B -> 01 Z) {
   DELAY {
      FROM {PIN = B;}
      TO {PIN = Z;}
      /* same data */
   }
}
```

## B.2 SDF timing constraints

This section details the different types of SDF timing constraints.

### B.2.1 SDF SETUP

A `SETUP` in SDF is the minimal time required for a data signal to arrive before the sampling edge of a clock signal in order to be sampled correctly.

```
VECTOR (?! din -> 01 clk) {
   SETUP {
      FROM {PIN = din;}
      TO {PIN = clk;}
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-7:

*any edge at din followed by rising edge at clk.*

The `FROM` and `TO` pin annotations define the sense of measurement for `SETUP`. Since setup time is measured in positive sense from data to clock, `din` is the data pin and `clk` is the clock pin.

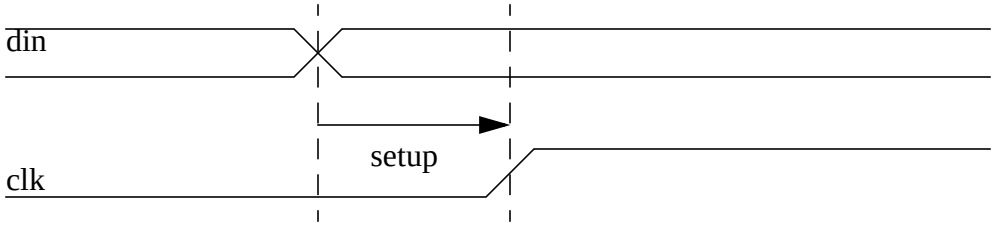



**Figure B-7: Measurement of SDF SETUP**

### B.2.2 SDF HOLD

A `HOLD` in SDF is the minimal non-negative time required for a data signal to stay at its value after the sampling edge of a clock signal in order to be sampled correctly.

```
VECTOR (01 clk -> ?! din) {
   HOLD {
      FROM {PIN = clk;}
      TO {PIN = din;}
      /* fill in data */
}
```

This ALF `VECTOR` describes the sequence of events as shown in Figure B-8:

*rising edge at clk followed by any edge at din.*

The `FROM` and `TO` pin annotations define the sense of measurement for `HOLD`. Since hold time is measured in positive sense from clock to data, `clk` is the clock pin and `din` is the data pin.

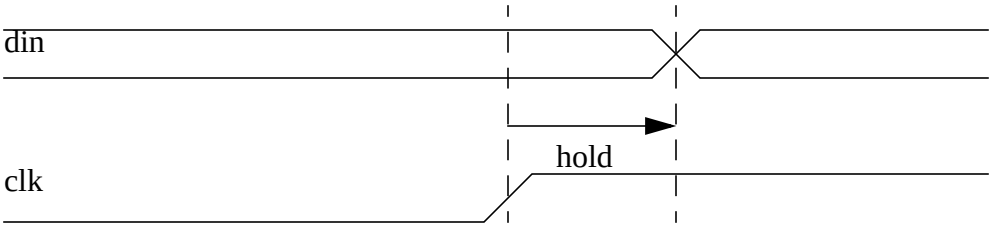



**Figure B-8: Measurement of SDF HOLD**

### B.2.3 SDF SETUPHOLD

A SETUPHOLD in SDF is a combination of SETUP and HOLD. In this combination, either SETUP or HOLD can be a negative value, but the sum of both values, which represents the minimal pulsewidth of the data in order to be sampled correctly, shall be non-negative. The time from the leading data edge to the sampling clock edge is SETUP. The time from the sampling clock edge to the trailing data edge is HOLD.

```
VECTOR // for SETUPHOLD
   (  ?! din -> 01 clk -> ?! din   //setup & hold both positive
   |  01 clk -> ?! din -> ?! din   //negative setup, positive hold
   |  ?! din -> ?! din -> 01 clk   //positive setup, negative hold
   ) {
   SETUP {
      FROM {PIN = din;
       TO {PIN = clk;}
      /* fill in data */
   }
   HOLD {
      FROM {PIN = clk;}
      TO {PIN = din;}
      /* fill in data */
   }
}
```

These ALF VECTORs describe the alternative sequences of events shown in Figure B-9:

*any edge at din followed by rising edge at clk followed by any edge at din*
or *rising edge at clk followed by any edge at din followed by any edge at din*
or *any edge at din followed by any edge at din followed by rising edge at clk.*

The FROM and TO pin annotations define the sense of measurement for SETUP and HOLD, respectively, in the same way as if they were specified in separate vectors.

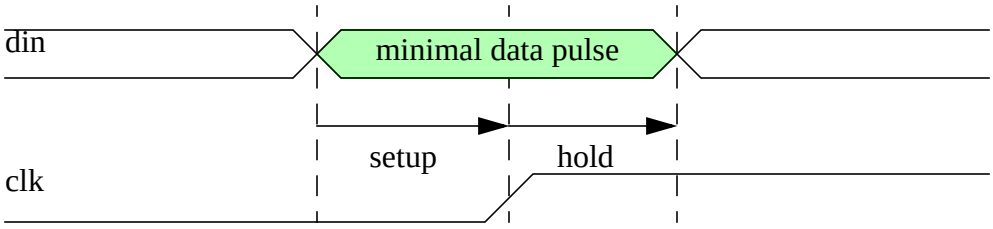



**Figure B-9: Measurement of SDF SETUPHOLD**

### B.2.4 SDF RECOVERY

A `RECOVERY` in SDF is the minimal time required for a higher priority asynchronous control signal to be released before a lower priority clock signal in order to allow the clock to be in control.

```
VECTOR (01 clearbar -> 01 clk) {
   RECOVERY {
      FROM {PIN = clearbar;}
      TO {PIN = clk;}
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-10:

*rising edge at clearbar followed by rising edge at clk.*

The `FROM` and `TO` pin annotations define the sense of measurement for `RECOVERY`. Since recovery time is measured in positive sense from the higher priority asynchronous control signal to the lower priority clock, `clearbar` is the asynchronous control pin and `clk` is the clock pin.

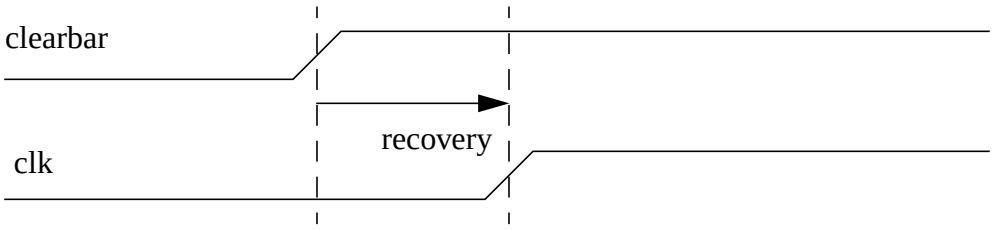



**Figure B-10: Measurement of SDF RECOVERY**

### B.2.5 SDF REMOVAL

A `REMOVAL` in SDF is the minimal time required for a higher priority asynchronous control signal to stay active after a lower priority clock signal in order to keep overriding the clock.

```
VECTOR (01 clk -> 01 clearbar) {
   REMOVAL {
      FROM {PIN = clk;}
      TO {PIN = clearbar;}
}
```

This ALF VECTOR describes the sequence of events shown in Figure B-11:

*rising edge at clk followed by rising edge at clearbar.*

The FROM and TO pin annotations define the sense of measurement for REMOVAL. Since removal time is measured in positive sense from the lower priority clock to the higher priority asynchronous control signal, clk is the clock pin and clearbar is the asynchronous control pin.

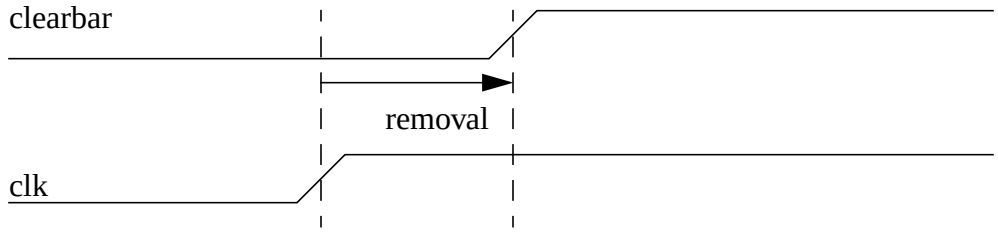



**Figure B-11: Measurement of SDF REMOVAL**

### B.2.6 SDF RECREM

A RECREM in SDF is a combination of RECOVERY and REMOVAL. In this combination, either RECOVERY or REMOVAL can be negative, but the sum of both shall be non-negative. The sum of RECOVERY and REMOVAL represents the width of the "forbidden zone" for the phase between the higher priority and the lower priority signal. The boundary to the left is RECOVERY and the boundary to the right is REMOVAL.

In a characterization vector for RECREM, either the RECOVERY or the REMOVAL effect can be observed, depending on the phase relationship between the signals. This is different from SETUPHOLD, where the effects of both SETUP and HOLD can be observed in the same characterization vector.

```
VECTOR // for RECREM
   (  01 clearbar -> 01 clk// pos. recovery or neg. removal
   |  01 clk -> 01 clearbar// neg. recovery or pos. removal
   ) {
   RECOVERY{
      FROM {PIN = clearbar;}
      TO {PIN = clk;}
      /* fill in data */
   }
```

```
   REMOVAL {
      FROM {PIN = clk;}
      TO {PIN = clearbar;}
      /* fill in data */
   }
}
```

These ALF VECTORs describe the alternative sequences of events shown in Figure B-12:

*rising edge at clearbar followed by rising edge at clk*
or *rising edge at clk followed by rising edge at clearbar*

The FROM and TO pin annotations define the sense of measurement for RECOVERY and REMOVAL, respectively, in the same way as if they were specified in separate vectors.

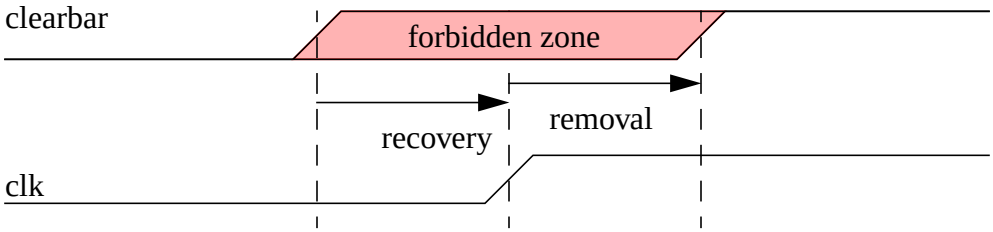



**Figure B-12: Measurement of SDF RECREM**

### B.2.7 SDF SKEW

A SKEW in SDF is the maximum allowed difference in arrival time between signals. The allowed region for the phase between signals is bound by zero (0) to the left and SKEW to the right for positive SKEW, or by SKEW to the left and zero (0) to the right for negative SKEW.

```
VECTOR (01 clk1 <&> 01 clk2) {// pos. or neg. or zero skew
   SKEW {
      FROM {PIN = clk1;}
      TO {PIN = clk2;}
      /* fill in data */
   }
}
```

These ALF VECTORs describe the alternative sequences of events shown in Figure B-13:

*rising edge at clk1 followed by rising edge at clk2*
or *rising edge at clk2 followed by rising edge at clk1*
or *rising edge at clk2 simultaneously with rising edge at clk1*

This is the most general case, where the skew can be positive, negative, or zero (0) across the characterization space. The FROM and TO pin annotations define the sense of measurement for SKEW.

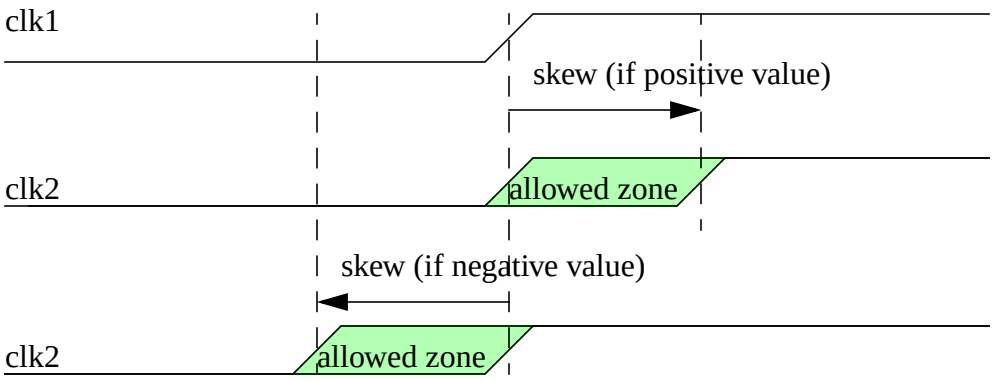



**Figure B-13: Measurement of SDF SKEW**

### B.2.8 SDF WIDTH

```
VECTOR (01 clk -> 10 clk) {// high pulse
   PULSEWIDTH {
      PIN = clk;
      /* fill in data */
   }
}
```

This ALF `VECTOR` describes the sequence of events shown in Figure B-14:

*rising edge at clk followed by falling edge at clk.*

The pulsewidth applies to the positive phase of the signal `clk`.

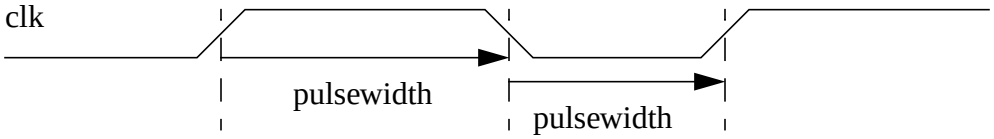



**Figure B-14: Measurement of SDF WIDTH**

```
VECTOR (10 clk -> 01 clk) {// low pulse
   PULSEWIDTH {
      PIN = clk;
      /* fill in data */
   }
}
```

This ALF `VECTOR` describe the sequence of events:

*falling edge at clk followed by rising edge at clk.*

The pulsewidth applies to the negative phase of the signal `clk`.

```
VECTOR (01 clk -> 10 clk | 10 clk -> 01 clk) {// high or low pulse
   PULSEWIDTH {
      PIN = clk;
      /* fill in data */
   }
}
```

These ALF VECTORs describe the alternative sequences of events shown in Figure B-14:

*rising edge at clk followed by falling edge at clk*
or *falling edge at clk followed by rising edge at clk.*

The pulsewidth applies to both phases of the signal clk.

### B.2.9 SDF PERIOD

```
VECTOR (01 clk -> 10 clk -> 01 clk) {
   PERIOD {
      PIN = clk;
      /* fill in data */
   }
}
```

This ALF VECTOR describes the sequence of events shown in Figure B-15:

*rising edge at clk followed by falling edge at clk followed by rising edge at clk.*

Thus the period is measured between two consecutive rising edges at the signal clk.

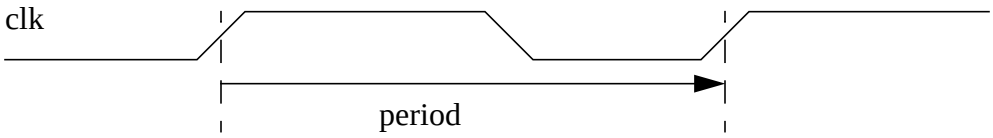



**Figure B-15: Measurement of SDF PERIOD**

### B.2.10 SDF NOCHANGE

```
VECTOR (?! addr -> 10 write -> 01 write -> ?! addr) {
   SETUP {
      FROM {PIN = addr;}
      TO {PIN = write;}
      /* fill in data */
   HOLD {
      FROM {PIN = write;}
      TO {PIN = addr;}
      /* fill in data */ }
   NOCHANGE {
      PIN = addr;
      /* fill in optional data */
   }
}
```

This ALF VECTOR describes the sequence of events shown in Figure B-16:

*any edge at addr followed by falling edge at write followed by rising edge at write followed by any edge at addr.*

The SETUP time is measured from the first edge at addr to the first edge at write. The HOLD time is measured from the second edge at write to the second edge at addr. The signal addr can not change between the start time of the setup measurement until the end time of the hold measurement. ALF allows the specification of an additional measurement between the first and second edge of the signal (subject to NOCHANGE). However, this additional measurement can not be directly translated into SDF (and would be for characterization and future purposes only).

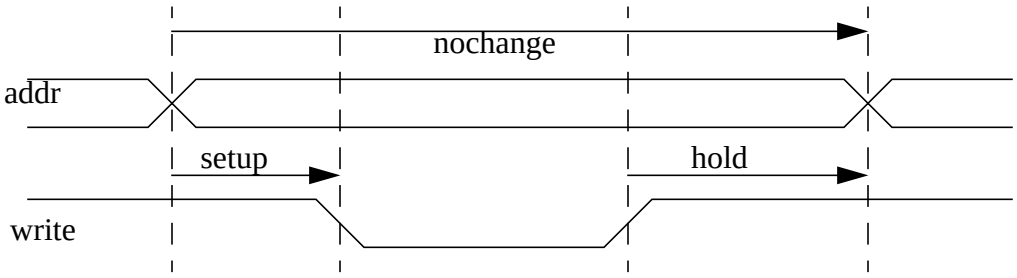



**Figure B-16: Detection of SDF NOCHANGE**

## B.3 SDF conditions and labels for delays and timing constraints

Conditions for IOPATH timing arcs in SDF apply to the entire timing arc. The condition is evaluated during the event on the “from” port (i.e., an input pin) and the event on the “to” port (i.e., an output pin) is scheduled consequently.

Conditions for timing constraints in SDF can be defined individually for each port. The condition associated with the *start point* of the timing constraint (i.e., data for SETUP and clock for HOLD) is called the *stamp condition*. The condition associated with the *end point* of the timing constraint (i.e., clock for SETUP and data for HOLD) is called the *check condition*.

The use of SETUPHOLD instead of the combination of SETUP and HOLD or RECREM instead the combination of RECOVERY and REMOVAL in SDF imposes restrictions in the definition of conditions. Where the use of two individual timing constraints allows the definition of four conditions (two stamp and two check), the use of one combined timing constraint allows only the definition of two conditions (one stamp and one check).

The ALF vector expression language can be used to specify conditions during the sequence of events in a more general way than SDF.

Some more examples in ALF:

```
VECTOR ( C & ( 01 A -> 01 B) )
```

Some alternative specification options:

```
VECTOR ( ?1 C -> 01 A -> 01 B -> 1? C ) // verbose
```

```
VECTOR ( ?1 C -> 01 A -> 01 B ) // C shall be true before start
VECTOR ( 01 A -> 01 B -> 1? C ) // C shall be true until the end
```

These ALF VECTORs describe the sequence of events shown in Figure B-17:

*rising edge at A is followed by rising edge at B, C is true before rising edge of A until after rising edge of B.*

Either of the pseudo-events (?1 C or 1? C) at the boundary can be omitted, since either one of them is sufficient to specify the condition C shall be *True* during the entire event sequence.

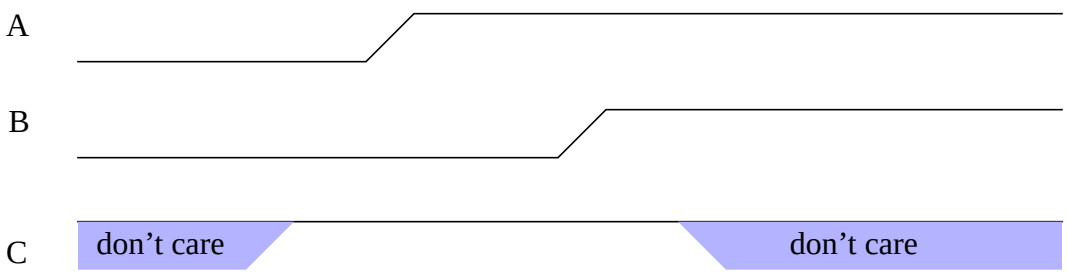



**Figure B-17: Condition during sequence of two events**

```
VECTOR ( ( C & 01 A ) -> 01 B )
```

alternative specification options:

```
VECTOR ( ?1 C -> 01 A -> 1? C -> 01 B )
VECTOR ( 01 A -> 1? C -> 01 B )
```

These ALF VECTORs describe the sequence of events shown in Figure B-18:

*rising edge at A is followed by rising edge at B, C is true before rising edge of A until after rising edge of A.*

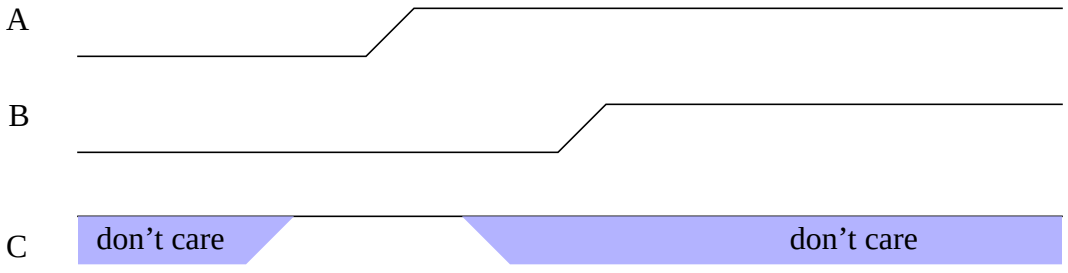



**Figure B-18: Condition during leading event**

```
VECTOR ( 01 A -> (C & 01 B) )
```

alternative syntax:

```
VECTOR ( 01 A -> ?1 C -> 01 B -> 1? C )
```

This ALF VECTOR describes the sequence of events shown in Figure B-19:

*rising edge at A is followed by rising edge at B, C is true before rising edge of B until after rising edge of B.*

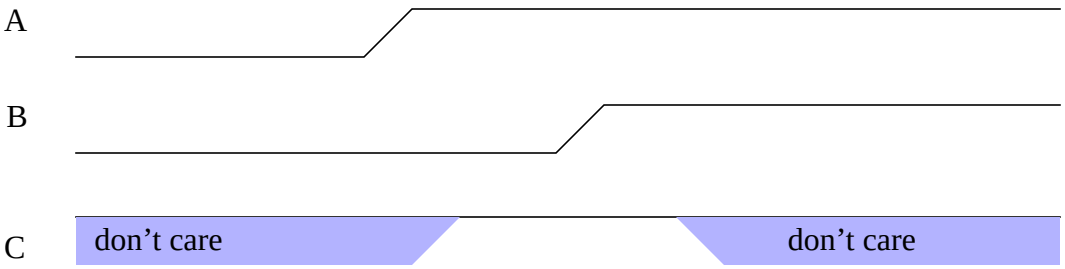



**Figure B-19: Condition during trailing event**

A SETUPHOLD with SCOND (stamp condition) and CCOND (check condition) in SDF can be described in ALF in the following way (and depicted in Figure B-20):

```
VECTOR ( ?! din -> ?1 ccond -> 01 clk -> 1? scond -> ?! din ) {
   SETUP {
      FROM {PIN = din;
       TO {PIN = clk;}
      /* fill in data */
   }
   HOLD {
      FROM {PIN = clk;}
      TO {PIN = din;}
      /* fill in data */
   }
}
```

A more verbose specification of the vector is:

```
VECTOR (
   ?1 scond // scond shall be true at the beginning
-> ?! din   // din toggles
-> ?1 ccond // last chance for ccond to become true
-> 01 clk   // rising edge at clk
-> 1? scond // scond gets a break
-> ?! din   // din toggles
-> 1? ccond // ccond gets a break at last
)
```

The optional condition label in SDF has its counterpart in ALF (see Section 6.6.3). As in SDF, the use and interpretation of this label is defined by the application tool and not by the standard.

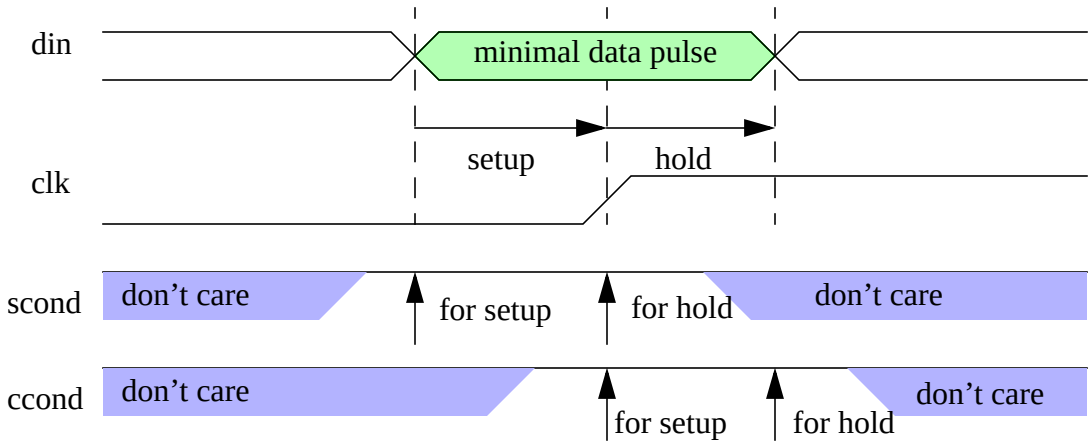



**Figure B-20: SETUPHOLD with SCOND and CCOND**

# Appendix C
# Phased-out Items

This section contains all items from the ALF 1.1 spec., which are phased out for ALF 2.0, because they are considered obsolete.

## C.1 Polarity for output pin

The polarity of an output pin (i.e. `DIRECTION = output;`) can take the following values:

**Table C-1 POLARITY (output) annotations for a PIN object (phased out)**

| Annotation string | Description |
|---|---|
| `inverted` | polarity change between input and output |
| `non_inverted` | no polarity change between input and output |
| `both` | polarity may change or not (e.g. XOR) (default) |
| `none` | polarity has no meaning(e.g. analog signal) |

**Reason for phase-out**

Not required by any tool today. Applies to very few signals in a library (e.g. inverted and non-inverted output of flipflop). Different semantics than polarity for input signal, therefore potentially confusing.

**Substitution**

Use attributes for pins representing double-rail signals (see Section 6.4.18).

**Example**

old style:

```
CELL my_flipflop {
   PIN Q { DIRECTION = output; POLARITY = non_inverted; }
   PIN Qbar { DIRECTION = output; POLARITY = inverted; }
   /* other pins and stuff */
}
```

new style:

```
CELL my_flipflop {
   PIN Q { DIRECTION = output; ATTRIBUTE { non_inverted } }
   PIN Qbar { DIRECTION = output; ATTRIBUTE { inverted } }
   /* optional */ PIN_GROUP [0:1] Q_double_rail { MEMBERS Q Qbar } }
   /* other pins and stuff */
}
```

## C.2 ENABLE_PIN annotation

```
ENABLE_PIN = string ;
```

references an output enable pin (i.e., a pin with `SIGNALTYPE = out_enable;` ).

**Reason for phase-out**

This annotation is phased out, since it provides only a very limited capability to describe a relationship between two pins, which is normally not described as an annotation for a pin. Relationships between pins can be described using VECTOR, supplemented by new features in ALF 2.0. Also, the ENABLE_PIN can make a reference to a pin which may not yet be declared. This clashes with the general rule: an object shall not be referenced before it is declared.

**Substitution**

The ENABLE_PIN can be infered according to the following rules:

For cells with `CELLTYPE = buffer | combinational | latch | flipflop` the following rule applies:

- For a PIN with `SIGNALTYPE = data` and `DIRECTION = output | both`, the PIN with `SIGNALTYPE = out_enable` is the enable-pin.
- For a PIN with `SIGNALTYPE = scan_data` and `DIRECTION = output | both`, the PIN with `SIGNALTYPE = scan_out_enable` is the enable-pin.

For cells with `CELLTYPE = memory` the following rule applies:

- For a PIN with `SIGNALTYPE = data` and `DIRECTION = output | both`, the PIN with `SIGNALTYPE = read_enable` is the enable-pin.
- For a PIN with `SIGNALTYPE = test_data` and `DIRECTION = output | both`, the PIN with `SIGNALTYPE = test_read_enable` is the enable-pin.

Port-specific enable-pins in multiport memories must have the same `SIGNAL_CLASS` as the related output pin.

## C.3 ATTRIBUTE with POLARITY annotation

The following attributes within a PIN object can also have `POLARITY` annotation:

**Table C-2 Attributes with POLARITY annotation (phased out)**

| Attribute item | Description |
|---|---|
| `TIE` | signal that needs to be tied to a fixed value |
| `READ` | read enable mode |
| `WRITE` | write enable mode |

**Reason for phase-out**

Not a very concise modeling style. Also, this is the only case where ATTRIBUTE contains non-atomic objects. By removing this special case, ATTRIBUTE will contain only atomic objects, which simplifies the datamodel.

**Substitution**

Use mode-specific polarity for signal with composite signaltype based on fundamental signaltype “control” (see Section 6.4.4 and Section 6.4.6).

**Example**

old style:

```
PIN rw {
   ATTRIBUTE {
      WRITE { POLARITY = high; }
      READ  { POLARITY = low ; }
   }
}
```

new style:

```
PIN rw {
   SIGNALTYPE = read_write_control;
   POLARITY {
      WRITE = high;
      READ = low;
   }
}
```

## C.4 OFF_STATE annotation

```
OFF_STATE = string ;
```

which can be:

**Table C-3 OFF_STATE annotations for a PIN object**

| Annotation string | Description |
|---|---|
| `inverted` | pin is inverted when in off state |
| `non_inverted` | pin is not inverted when in off state |

**Reason for phase-out**

The purpose of this feature is not clear. No practical example could be found.

## C.5 SCAN annotation container

A SCAN container may be used inside a CELL or a PIN object and may contain annotations which are allowed inside a CELL or a PIN object for limiting the scope of those annotations.

Example:

```
PIN clk1 { signaltype = master_clock; SCAN {signaltype = slave_clock;} }
PIN clk2 { SCAN {signaltype = master_clock;} }
```

In normal mode, `clk1` is master clock, `clk2` is unused. In scan mode, `clk2` is master clock, `clk1` is slave clock.

**Reason for phase-out**

This feature is not required for DFT, since all DFT items are already identified by dedicated keywords. The example above is unrealistic.

## C.6 PRIMITIVE definition in FUNCTION

BNF in ALF 1.1, chapter 3.4.16

```
function ::=
   FUNCTION [ identifier ] {
         [ all_purpose_items ]
         [ primitives ]
         [ behavior ]
         [ statetables ]
   }
```

Proposed change: remove `[ primitives ]`.

**Reason for change**

PRIMITIVE definitions must contain a FUNCTION statement themselves. Therefore, the possibility of having PRIMITIVE inside FUNCTION and FUNCTION inside PRIMITIVE bears the potential risk of circular reference in the datamodel.

**Substitution**

use PRIMITIVE definitions inside the CELL which contains the FUNCTION.

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