

Correction: Remove the last limitation of Section 12.4.4 (Distribution)
(The correct limitation is listed at the end of Section 12.4) – BNF is also modified

Limitations:

- A **dist** operation shall not be applied to **randc** variables.
- A **dist** expression requires that expression contain at least one **rand** variable.
- ~~— A **dist** expression can only be a top level constraint (not a predicated constraint).~~

Updated Constraints BNF – Replace in section 12.4 and Annex A.1.9

```
constraint_declaration ::= //from Annex A.1.9  
    [ static ] constraint constraint_identifier { { constraint_block } }  
  
constraint_block ::=  
    solve [ priority ] identifier_list before identifier_list ;  
    | constraint_expression  
  
constraint_expression ::=  
    expression ;  
    | expression => constraint_set  
    | if ( expression ) constraint_set [ else constraint_set ]  
    | expression dist { dist_list } ;  
    | foreach ( array_identifier [ loop_variables ] ) constraint_set  
  
constraint_set ::=  
    constraint_expression  
    | { { constraint_expression } }  
  
dist_list ::= dist_item { , dist_item }  
  
dist_item ::=  
    value_range := expression  
    | value_range :/ expression  
  
constraint_prototype ::= [ static ] constraint constraint_identifier  
  
extern_constraint_declaration ::=  
    [ static ] constraint class_identifier :: constraint_identifier { { constraint_block } }  
  
identifier_list ::= identifier { , identifier }  
  
loop_variables ::= [ index_indentifier ] { , [ index_indentifier ] }
```

Insert as new 12.10.4 sub-section -- After \$srandom()

12.10.4 \$get_randstate()

The system function `$get_randstate()` retrieves the current state of the Random Number Generator (RNG) of an object or process.

The prototype for the `$get_randstate()` function is:

```
function string $get_randstate( [class_identifier obj] );
```

The `$get_randstate()` system function returns a copy of the internal state of a RNG. With no arguments, this function returns the state of the RNG associated with the calling process. If the optional object argument is specified, the function returns the state of the RNG associated with the given object.

The RNG state is a string of unspecified length and format.

12.10.5 \$set_randstate()

The system function `$set_randstate()` allows setting the current state to of the Random Number Generator (RNG) of an objects or process.

The prototype for the `$set_randstate()` function is:

```
function void $set_randstate( string state [, class_identifier obj] );
```

The `$set_randstate()` system function copies the given state into the internal state of a RNG. With one argument, this function sets the state of the RNG associated with the calling process. If the optional object argument is specified, the function sets the state of the RNG associated with the given object.

The RNG state is a string of unspecified length and format.

Calling `$set_randstate()` with a string value that was not obtained from `$get_randstate()` is undefined.

Insert as Section 12.4.10

12.4.10 Functions in Constraints

Some properties are unwieldy or impossible to express in a single expression. For example, the natural way to compute the number of 1's in a packed array uses a loop:

```
function int count_ones ( bit [9:0] w );
    for( count_ones = 0; w != 0; w = w >> 1 )
        count_ones += w & 1'b1;
endfunction
```

Such a function could be used to constrain other random variables to the number of 1 bits:

```
constraint C1 { length == count_ones( v ) };
```

Without the ability to call a function, this constraint requires the loop to be unrolled and expressed as a sum of the individual bits:

```
constraint C2
{
    length == ((w>>9)&1) + ((w>>8)&1) + ((w>>7)&1) + ((w>>6)&1) + ((w>>5)&1) +
              ((w>>4)&1) + ((w>>3)&1) + ((w>>2)&1) + ((w>>1)&1) + ((w>>0)&1);
}
```

Unlike the `count_ones` function, more complex properties, which require temporary state or unbounded loops, may be impossible to convert into a single expression. The ability to call functions, thus, enhances the expressive power of the constraint language and reduces the likelihood of errors.

To handle these common cases, SystemVerilog allows constraint expressions to include function calls, but it imposes certain semantic restrictions.

- Functions that appear in constraint expressions may not contain output or **ref** arguments (**const ref** are allowed). They should be automatic (or preserve no state information) and have no side effects.
- Function that appear in constraints may not modify the constraints, for example, calling `rand_mode` or `constraint_mode` methods.
- Functions shall be called before constraints are solved, and their return values shall be treated as state variables.
- Random variables used as function arguments shall establish an implicit variable ordering or priority. Constraints that include only variables with higher priority are solved before other, lower priority, constraints. Random variables solved as part of a higher priority set of constraints become state variables to the remaining set of constraints. For example:

```
class B;
    rand int x, y;
    constraint C { x <= F(y); };
    constraint D { y inside { 2, 4, 8 } };
endclass
```

Forces `y` to be solved before `x`. Thus, constraint `D` is solved separately before constraint `C`, which uses the values of `y` and `F(y)` as a state variables.

- Circular dependencies created by the implicit variable ordering shall result in an error.

- Function calls in active constraints are executed an unspecified number of times (at least once), in an unspecified order.

And change Section 12.4 to:

The declarative nature of constraints imposes the following restrictions on constraint expressions:

- Functions are allowed with certain limitations (see Section 12.4.10).
- Operators with side effects, such as `++` and `--` are not allowed.
- **randc** variables cannot be specified in ordering constraints (see **solve...before** in Section 12.4.8).
- **dist** expressions cannot appear in other expressions (unlike **inside**); they can only be top-level expressions.

Insert as Section 12.4.11

12.4.11 Object handle guards

Constraint expressions may involve multiple objects. Typically, inter-object constraints are written under the assumption that all constrained objects already exist. If, however, some of the objects do not exist, the constraint will fail. Predicate expressions involving object handles can be useful to guard against expressions that should not be evaluated. For example:

```
class SList;  
  rand int n;  
  rand Slist next;  
  
  constraint sort { n < next.n; }  
endclass
```

The sort constraint of the singly-linked list, SList, is intended to assign a random sequence of numbers that is sorted in ascending order. However, the constraint expression will fail on the last element. That condition can be avoided by writing a handle guard using a predicate expression:

```
constraint sort { if( next != null ) n < next.n; }
```

Predicate expressions involving object handle comparisons prevent the creation of constraints and do not behave as logical relations. Thus, in the sort constraint above, the **if** prevents the creation of a constraint when `next == null`, which in this case prevents accessing a non-existent object.

Predicate expressions (implication and `if...else`) involving comparisons between two object handles or an object handle and the constant **null** shall behave as a guard, that is, if the condition expression is false then that constraint expression is disregarded.

Insert as new 12.4.7 section (After if..else constraints)

12.4.7 Iterative Constraints

Iterative constraints allow arrayed variables to be constrained in a parameterized manner using loop variables and indexing expressions.

The syntax to define an iterative constraint is:

```
constraint_expression ::= //from Annex A.1.9  
    ...  
    | foreach (array_identifier [ loop_variables ]) constraint_set  
loop_variables ::= [ index_identifier ] { , [ index_identifier ] }
```

The **foreach** construct specifies iteration over the elements of an array. Its argument is an identifier that designates any type of array (fixed-size, dynamic, or associative) followed by a list of loop variables enclosed in square brackets. Each loop variable corresponds to one of the dimensions of the array.

For example:

```
class C;  
    rand byte A[] ;  
  
    constraint C1 { foreach ( A [ i ] ) A[i] inside {2,4,8,16}; }  
    constraint C2 { foreach ( A [ j ] ) A[j] > 2 * j; }  
endclass
```

C1 constrains each element of the array *A* to be in the set $\{2, 4, 8, 16\}$. C2 constrains each element of the array *A* to be greater than twice its index.

The number of loop variables must not exceed the number of dimensions of the array variable. The scope of each loop variable is the **foreach** constraint construct, including its constraint_set. The type of each loop variable is implicitly declared to be consistent with the type of array index. An empty loop variable indicates no iteration over that dimension of the array. As with default arguments, a list of commas at the end may be omitted, thus, **foreach**(arr [j]) is a shorthand for **foreach**(arr [j, , ,]). It shall be an error for any loop variable to have the same identifier as the array.

The mapping of loop variables to array indexes is determined by the dimension cardinality, as described in Section 22.4.

```
//      1  2  3          3  4          1  2      -> Dimension numbers  
int A [2][3][4];    bit [3:0][2:1] B [5:1][4];  
  
foreach( A [ i,j,k ] ) ...  
foreach( B [ q,r,s ] ) ...
```

The first **foreach** causes *i* to iterate from 0 to 1, *j* from 0 to 2, and *k* from 0 to 3. The second **foreach** causes *q* to iterate from 5 to 1, *r* from 0 to 3, and *s* from 2 to 1.

Iterative constraints may include predicates. For example:

```
class C;  
    rand int A[] ;  
  
    constraint c1 { arr.size inside {1:10}; }
```

```

    constraint c2 { foreach ( A[ k ] ) (k < A.size - 1) => A[k + 1] > A[k]; }
endclass

```

The first constraint, c1, constrains the size of the array A to be between 1 and 10. The second constraint, c2, constrains each array value to be greater than the preceding one, i.e., an array sorted in ascending order.

Predicate expressions within a **foreach** behave as guards against the creation of constraints, and not as logical relations. For example, the implication in constraint c2 above prevents the creation of a constraint when $k < A.size() - 1$, which in this case prevents an out of bounds access in the constraint.

Index expressions may include loop variables, constants, and state variables. Invalid or out of bound array indexes are not automatically eliminated; users must explicitly exclude these indexes using predicates.

The size method of a dynamic or associative array can be used to constrain the size of the array (see constraint c1 above). If an array is constrained by both size constraints and iterative constraints, the size constraints are solved first, and the iterative constraints next. As a result of this implicit ordering between size constraints and iterative constraints, the size method shall be treated as a state variable within the **foreach** block of the corresponding array. For example, the expression `A.size` is treated as a random variable in constraint c1, and as a state variable in constraint c2. This implicit ordering can cause the solver to fail in some situations.

12.10 In-line random variable control

The `randomize()` method can be used to temporarily control the set of random and state variables within a class instance or object. When the `randomize` method is called with no arguments, it behaves as described in the previous sections, that is, it assigns new values to all random variables in an object --- those declared as **rand** or **randc** --- such that all of the constraints are satisfied. When `randomize` is called with arguments, those arguments designate the complete set of random variables within that object; all other variables in the object are considered state variables. For example, consider the following class and calls to `randomize`:

```
class CA;
    rand byte x, y;
    byte v, w;

    constraint c1 { x < v && y > w };
endclass

CA a = new;

a.randomize();           // random variables: x, y    state variables: v, w
a.randomize( x );       // random variables: x      state variables: y, v, w
a.randomize( v, w );    // random variables: v, w    state variables: x, y
a.randomize( w, x );    // random variables: w, x    state variables: y, v
```

This mechanism controls the set of active random variables for the duration of the call to `randomize`, which is conceptually equivalent to making a set of calls to the `rand_mode()` method to disable or enable the corresponding random variables. Calling `randomize()` with arguments allows changing the random mode of any class property, even those not declared as **rand** or **randc**. This mechanism, however, does not affect the cyclical random mode: it cannot change a non-random variable into a cyclical random variable (**randc**), and cannot change a cyclical random variable into a non-cyclical random variables (change from **randc** to **rand**).

The scope of the arguments to the `randomize` method is the object class. Arguments are limited to the names of properties of the calling object; expressions are not allowed. The random mode of **local** class members can only be changed when the call to `randomize` has access to those properties, that is, within the scope of the class in which the local members are declared.

12.10.1 In-line constraint checker

Normally, calling the `randomize` method of a class that has no random variables causes the method to behave as a checker, that is, it assigns no random values, and only returns a status: one if all constraints are satisfied and zero otherwise. The in-line random variable control mechanism can also be used to force the `randomize()` method to behave as a checker.

The `randomize` method accepts the special argument **null** to indicate no random variables for the duration of the call, that is, all class members behave as state variables. This causes the `randomize` method to behave as a checker instead of a generator. A checker evaluates all constraints and simply returns one if all constraints are satisfied, and zero otherwise. For example, if class `CA` defined above executes the following call:

```
success = a.randomize( null );           // no random variables
```

Then the solver considers all variables as state variables and only checks whether the constraint is satisfied, namely, that the relation $(x < v \ \&\& \ y > w)$ is true using the current values of x , y , v , and w .

12.11 Randomization of scope variables - `::randomize()`

The built-in class `randomize` method operates exclusively on class member variables. Using classes to model the data to be randomized is a powerful mechanism that enables the creation of generic, reusable objects containing random variables and constraints that can be later extended, inherited, constrained, overridden, enabled, disabled, merged with or separated from other objects. The ease with which classes and their associated random variables and constraints can be manipulated make classes an ideal vehicle for describing and manipulating random data and constraints. However, some less-demanding problems that do not require the full flexibility of classes, can use a simpler mechanism to randomize data that does not belong to a class. The scope `randomize` method, `::randomize()`, enables users to randomize data in the current scope, without the need to define a class or instantiate a class object.

The syntax of the scope `randomize` method is:

```
scope_randomize ::=                                     // from Annex A.?
    [::] randomize ( [ variable_identifier_list ] ) [ with { constraint_block } ]

variable_identifier_list ::= variable_identifier {, variable_identifier }
```

The scope `randomize` method behaves exactly the same as a class `randomize` method, except that it operates on the variables of the current scope instead of class member variables. Arguments to this method specify the variables that are to be assigned random values. i.e., the random variables.

For example:

```
module stim;
    bit[15:0] addr;
    bit[31:0] data;

    function bit gen_stim();
        bit success, rd_wr;

        success = ::randomize( addr, data, rd_wr );
        return rd_wr ;
    endfunction

    ...
endmodule
```

The function `gen_stim` calls `::randomize()` with three variables as arguments: `addr`, `data`, and `rd_wr`. Thus, `::randomize()` assigns new random variables to those variables that are visible in the scope of the `gen_stim` function. Note that `addr` and `data` have module scope, whereas `rd_wr` has scope local to the function. The above example can also be written using a class:

```
class stimc;
    rand bit[15:0] addr;
    rand bit[31:0] data;
    rand bit rd_wr;
endclass

function bit gen_stim( stimc p );
    bit success;
    success = p.randomize();
    addr = p.addr;
    data = p.data;
    return p.rd_wr;
endfunction
```

However, for this simple application, the scope `randomize` method leads to a straightforward implementation.

The scope randomize method returns 1 if it successfully sets all the random variables to valid values, otherwise it returns 0. If the scope randomize method is called with no arguments then it behaves as a checker, and simply returns status.

12.11.1 Adding constraints to scope variables - ::randomize() with

The ::randomize() **with** form of the scope randomize method allows users to specify random constraints to be applied to the local scope variables. When specifying constraints, the arguments to the scope randomize method become random variables, all other variables are considered state variables.

```
task stimulus( int length );  
    int a, b, c, success;  
  
    success = ::randomize( a, b, c ) with { a < b ; a + b < length };  
    ...  
    success = ::randomize( a, b ) with { b - a > length };  
    ...  
endfunction
```

The task stimulus above calls ::randomize twice resulting in two sets of random values for its local variables a, b, and c. In the first call variables a and b are constrained such that variable a is less than b, and their sum is less than the task argument length, which is designated as a state variable. In the second call, variables a and b are constrained such that their difference is greater than state variable length.