OpenSPARC Slide-Cast

In Twelve Chapters
Presented by OpenSPARC designers, developers, and programmers
• to guide users as they develop their own OpenSPARC designs and
• to assist professors as they teach the next generation

Sun microsystems

OpenSPARC

Chapter Ten

DEVELOPING APPLICATIONS FOR CMT PROCESSORS

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Sun Microsystems

Agenda

- Compiler and tools options
- Compiling applications
- Profiling applications
- Writing parallel applications
- System utilisation
- Other resources

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Sun Studio 12

100% standards compliant
- ANSI C/C++
- C99, IEEE-754 ...
- OpenMP 2.5
- **GCC compatibility with gccfss**

Rapid debugging
- Set breakpoints
- Single-step
- NetBeans IDE
- **Thread analyzer**

Performance tuning
- Advanced optimizations
- Sun Performance Analyzer
- SPOT
- **Automatic Tuning System**

Extensive libraries
- Media and graphics
- Science and math
- Portable performance
- **Parallelized for CMT**

Sun Studio 12

- IDE (based on NetBeans.org)
- Compilers
  - C/C++/Fortran
- Debugger
  - dbx
- Performance Analyzer
- Thread Analyzer
- Solaris and Linux
- SPARC and x86/64

GCC for Sun Systems

- Enables GCC to use Sun Studio code generator
  - GCC compatibility
  - Sun Studio optimisations
  - Compatibility with Sun Studio tools

Mapping Tools to the Development or Migration Lifecycle

- **Application Selection**
  - cooltst
- **Observing**
  - SPOT
  - Corestat

http://cooltools.sunsource.net/ Over 15k

- **Development**
  - GCC4SS (Compiling)
  - ATS (tuning)
  - BIT (instrumenting)
  - Discover (checking)
  - Thread Analyzer (checking)
- **Deployment**
  - CoolTuner
  - Cool Stack
  - Consolidation Tool

Sun Studio Express

- Preview of next Sun Studio release
- July 2008 release includes
  - CMT Developer Tools
  - Initial support for OpenMP 3.0

CMT Developer Tools

- Automatic Tuning and Troubleshooting System (`ats`)
- Binary Improvement Tool (`bit`)
- Sun Memory Error Discovery Tool (`discover`)
- Simple Performance Optimisation Tool (`spot`)
- Free download from:
  http://cooltools.sunsource.net/

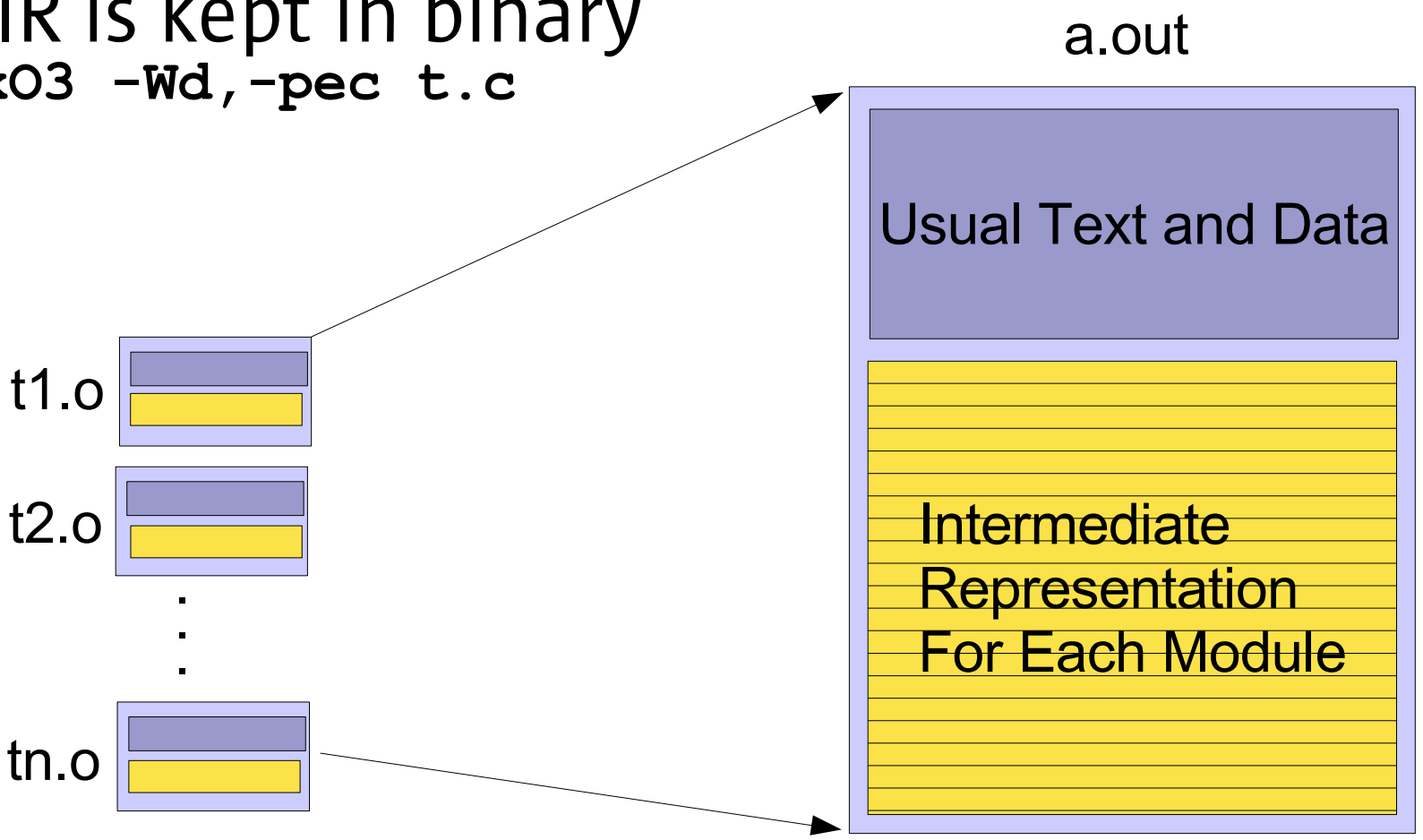
ATS

- Recompile application without access to source
- Automated performance tuning
  - Find the best compiler flags
- Automated application debug
  - Find problem compiler flag
  - Find problem module
- SPARC & x86
- http://cooltools.sunsource.net/ats/

ATS uses PEC

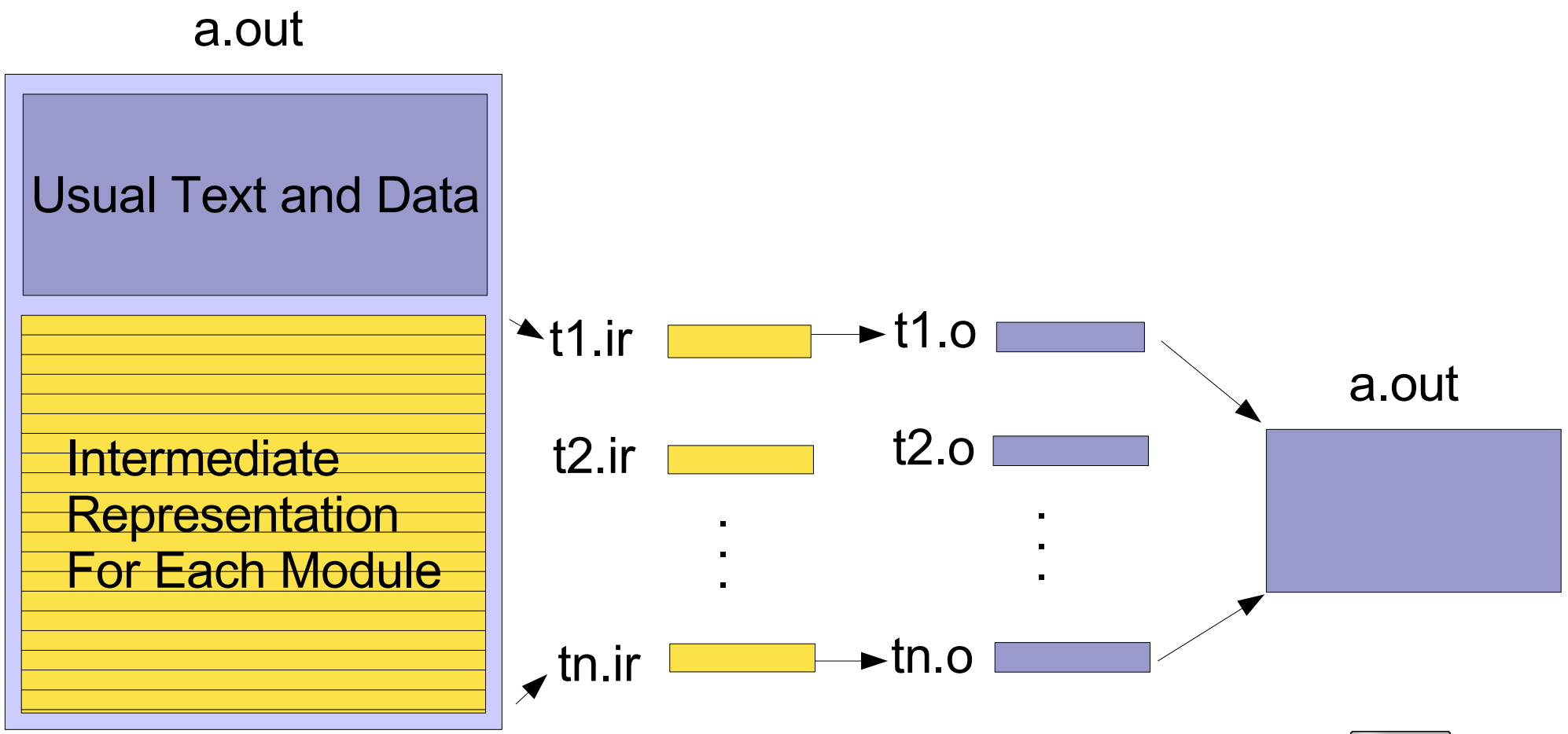
- Portable Executable Code
  - Sun IR is kept in binary

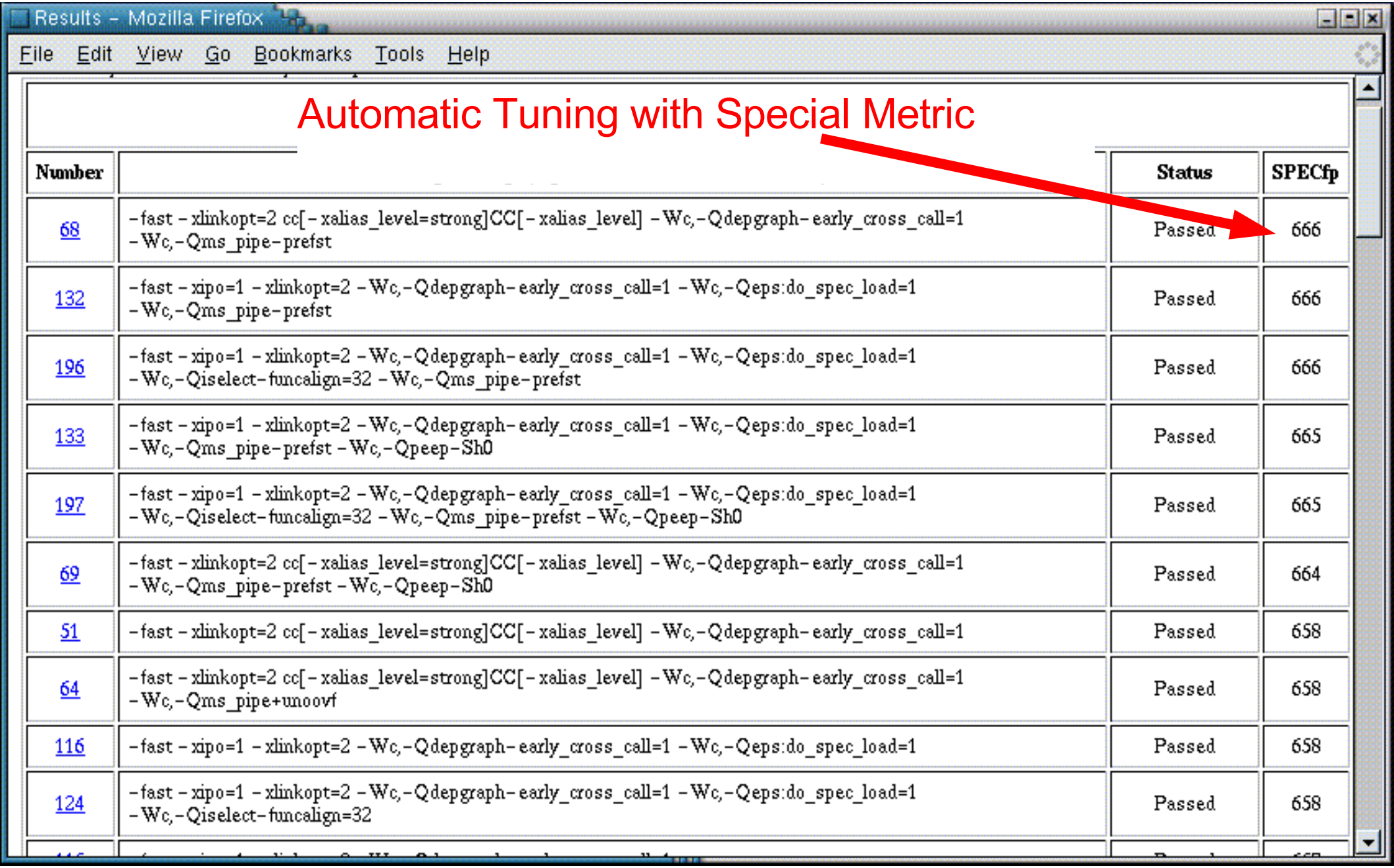
```
cc -xO3 -Wd,-pec t.c
```



Recompiling Binaries

- IR is extracted and reprocessed



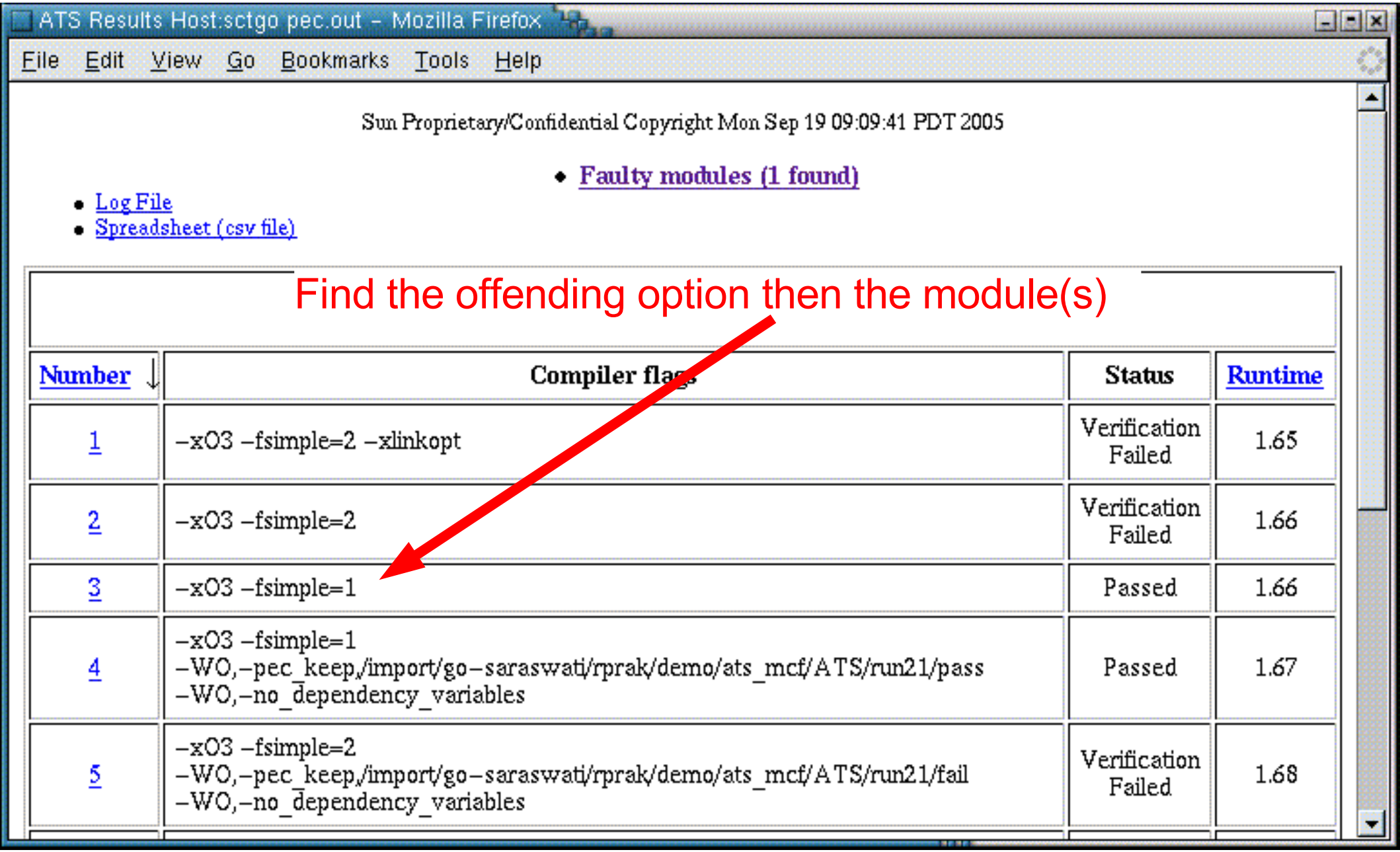


Number		Status	SPECfp
68	-fast -xlinkopt=2 cc[-xalias_level=strong]CC[-xalias_level] -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qms_pipe-prefst	Passed	666
132	-fast -xipo=1 -xlinkopt=2 -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qeps:do_spec_load=1 -Wc,-Qms_pipe-prefst	Passed	666
196	-fast -xipo=1 -xlinkopt=2 -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qeps:do_spec_load=1 -Wc,-Qiselect-funcalign=32 -Wc,-Qms_pipe-prefst	Passed	666
133	-fast -xipo=1 -xlinkopt=2 -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qeps:do_spec_load=1 -Wc,-Qms_pipe-prefst -Wc,-Qpeep-Sh0	Passed	665
197	-fast -xipo=1 -xlinkopt=2 -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qeps:do_spec_load=1 -Wc,-Qiselect-funcalign=32 -Wc,-Qms_pipe-prefst -Wc,-Qpeep-Sh0	Passed	665
69	-fast -xlinkopt=2 cc[-xalias_level=strong]CC[-xalias_level] -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qms_pipe-prefst -Wc,-Qpeep-Sh0	Passed	664
51	-fast -xlinkopt=2 cc[-xalias_level=strong]CC[-xalias_level] -Wc,-Qdepgraph-early_cross_call=1	Passed	658
64	-fast -xlinkopt=2 cc[-xalias_level=strong]CC[-xalias_level] -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qms_pipe+unoovf	Passed	658
116	-fast -xipo=1 -xlinkopt=2 -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qeps:do_spec_load=1	Passed	658
124	-fast -xipo=1 -xlinkopt=2 -Wc,-Qdepgraph-early_cross_call=1 -Wc,-Qeps:do_spec_load=1 -Wc,-Qiselect-funcalign=32	Passed	658

Find bug

- Locate problem flags and problem module

```
% ats -i 'script:findbug -xO3 -fsimple=2 -xlinkopt'
      a.out
```



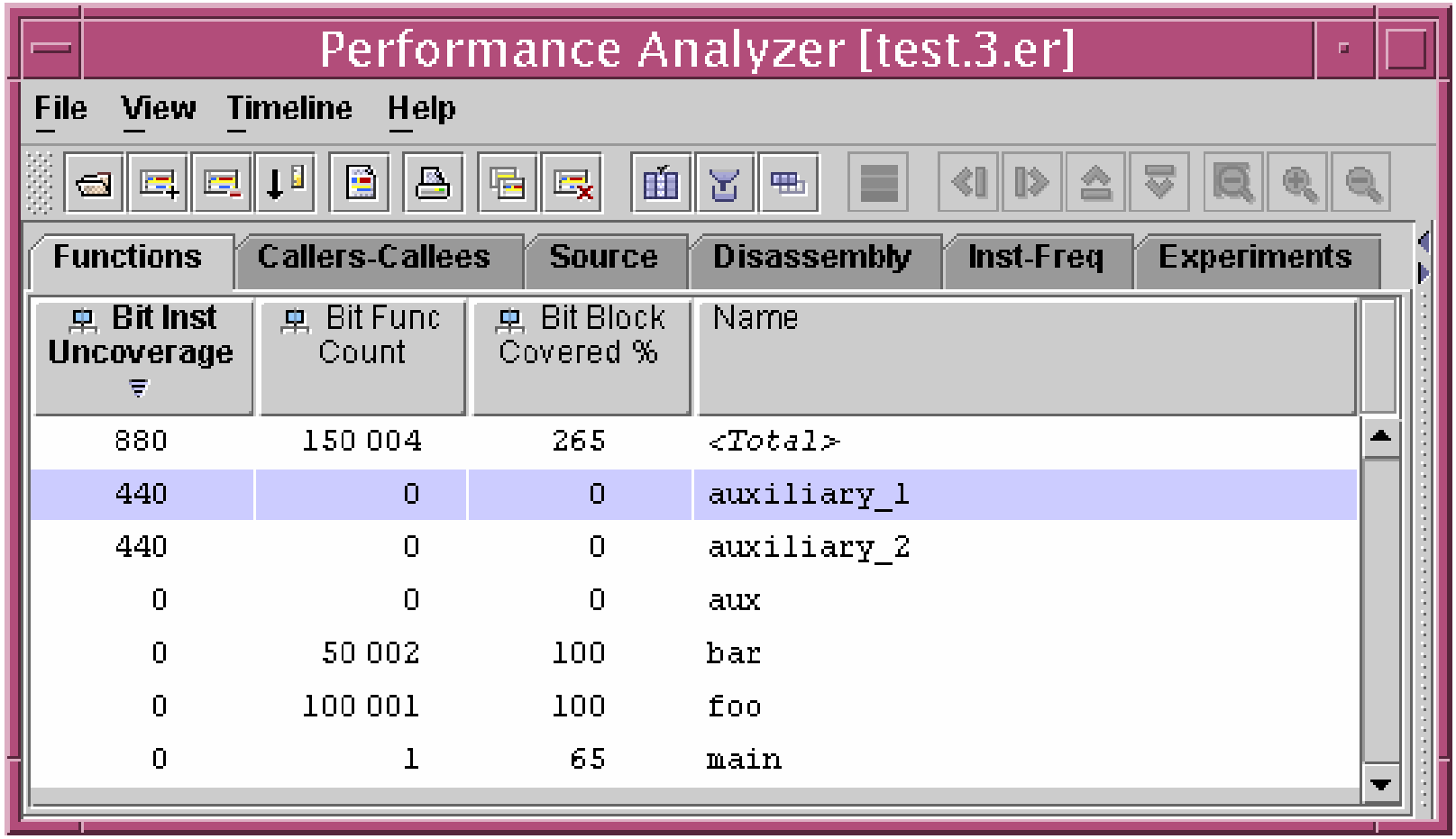
BIT

- Gathers runtime information
  - Instruction execution count
  - Branch taken probabilities
    - Compare behaviour of different workloads
- Generates coverage information
- SPARC only
- http://cooltools.sunsource.net/bit/

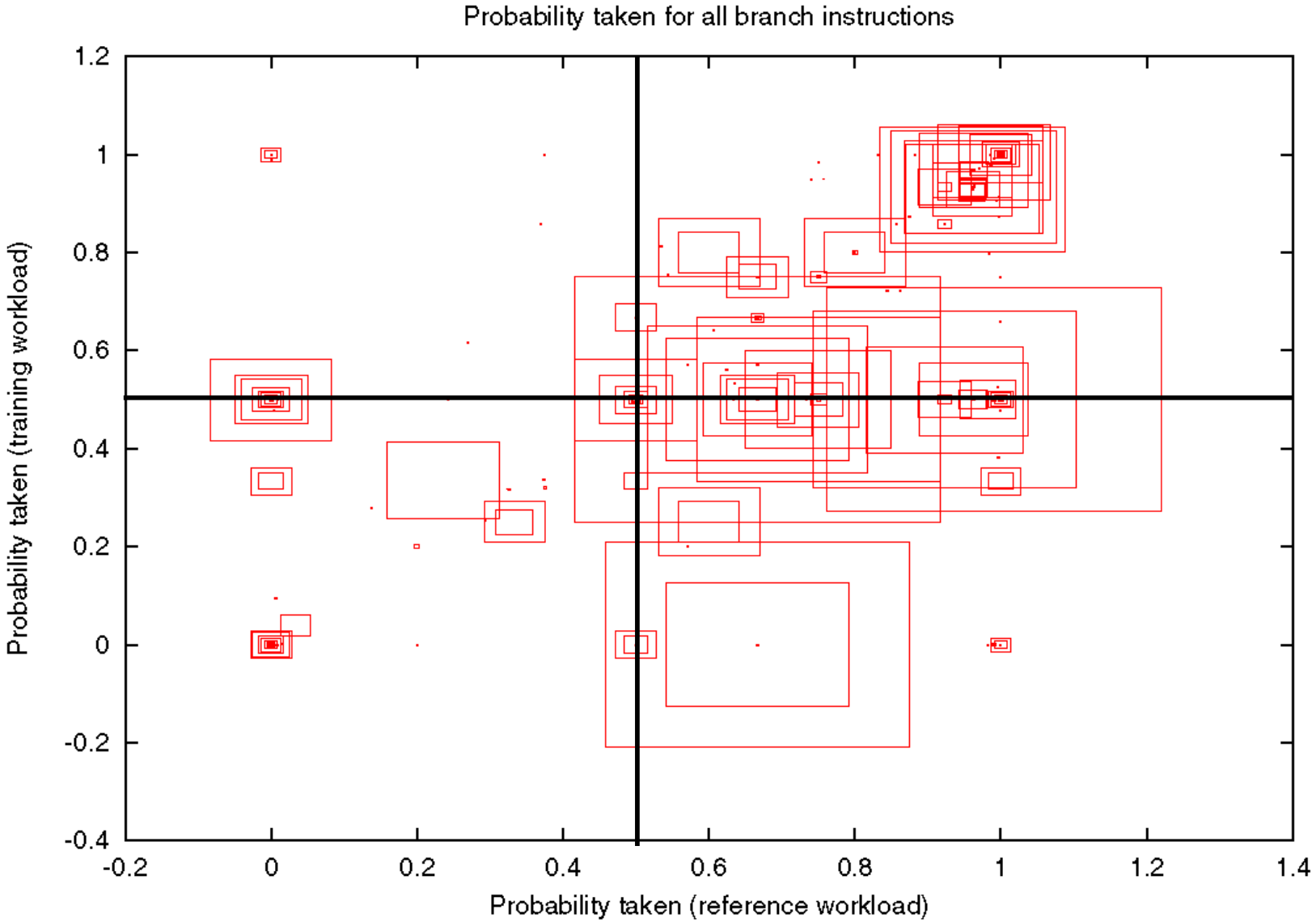
BIT coverage results

```
bit coverage -R -d  nmBasic.t.exe
...
BIT Code Coverage
Total Functions: 179
Covered Functions: 19
Function Coverage: 10.6%
Total Basic Blocks: 775
Covered Basic Blocks: 508
Basic Block Coverage: 65.5%
Total Basic Block Executions: 1,296
Average Executions per Basic Block: 1.67
Total Instructions: 3,168
Covered Instructions: 1,719
Instruction Coverage: 54.3%
Total Instruction Executions: 6,373
Average Executions per Instruction: 2.01
Creating experiment database test.1.er
```

BIT Uncoverage



Branch probability analysis



http://www.spec.org/workshops/2006/papers/10_Darryl_Gove.pdf

Discover

- Memory access error detection
  - Write past end of array
  - Read of uninitialised data
  - Use of freed memory
- SPARC only
- http://cooltools.sunsource.net/discover/

Discover - example

```
$ more memerr.c
#include<stdlib.h>
void main()
{
  int* a;
  int i;
  a=(int*)malloc(sizeof(int)*5);
  for (i=0; i<6; i++)
  {
   a[i]=0;
  }
}

$ cc -O -xbinopt=prepare memerr.c

% discover a.out
```

```
% a.out

ERROR (ABW):
writing to memory beyond array
bounds at:

        main() + 0x158
        _start() + 0x108

block  was allocated at:

        malloc() + 0x144
        main() + 0x1c
        _start() + 0x108

DISCOVER SUMMARY:
 unique errors   : 1 (1 total)
 unique warnings : 0 (0 total)
```

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Optimisation flags

- No optimisation flags = no optimisation
- `-O` = good degree of optimisation
- `-fast` = aggressive optimisation

Debug flags

- `-g` for C/Fortran
- `-g0` for C++
  - `-g` disables front-end inlining in C++
- Minor changes to code at low optimisation
  - Tail call optimisation
- Allows attribution of time to lines of source

32-bit or 64-bit

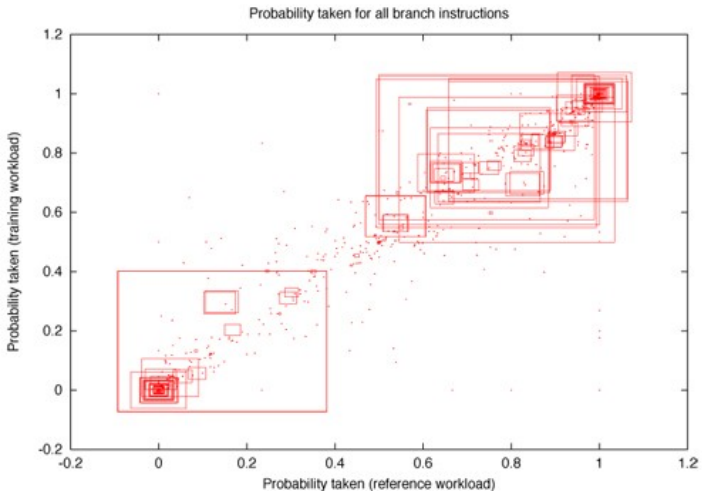
- Compiler flags: `-m32 | -m64`
- 64-bit:
  - Larger address space
  - Pointers and longs 64-bits
  - Larger memory footprint
  - Potentially lower performance

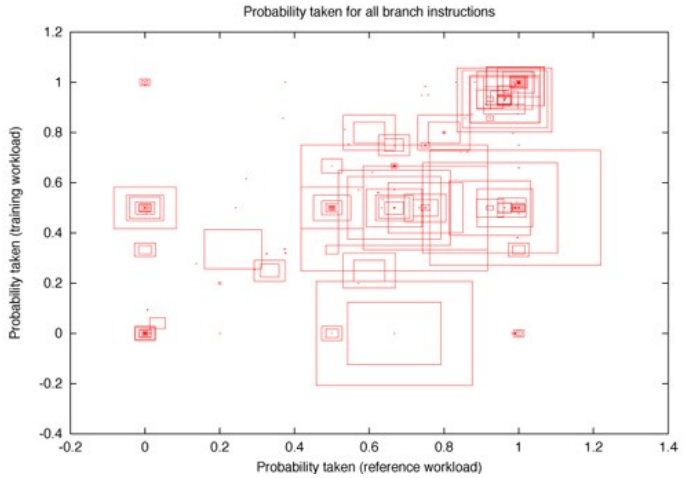
Inlining

- Inlining
  - Within file **`-xO4`**
  - Across files **`-xipo`**
- Avoid cost of calling routine
- Expose further performance opportunities

Profile feedback

- Two compile passes (complicates build)
- Good for “branchy” code
- Helps inlining decisions
- Profile feedback
  - `-xprofile=[collect:|use:]`





http://developers.sun.com/solaris/articles/coverage.h

Target architecture

- General
  - `-xtarget=generic`
- If build and run machine is the same
  - `-xtarget=native`
- Flags evaluated from left to right

Aliasing

- Compiler has to assume pointers alias
  - Unless it can prove otherwise
  - Or it is told to assume otherwise
- Specify degree of aliasing
  - `-xalias_level=<level>`
- Specify pointers passed into functions don't alias
  - `-xrestrict`
- Restrict qualify pointers
  - `int * restrict p`

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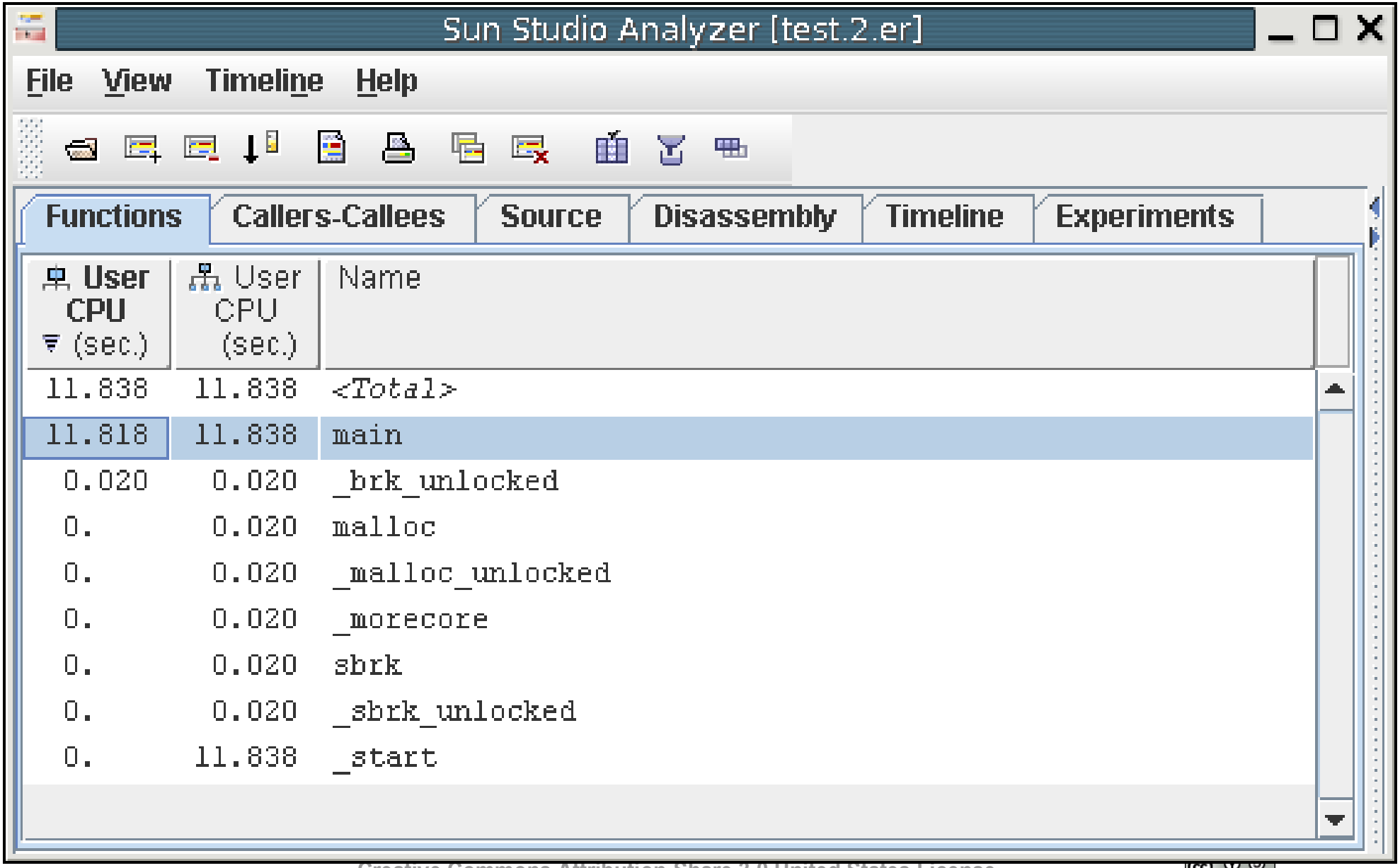
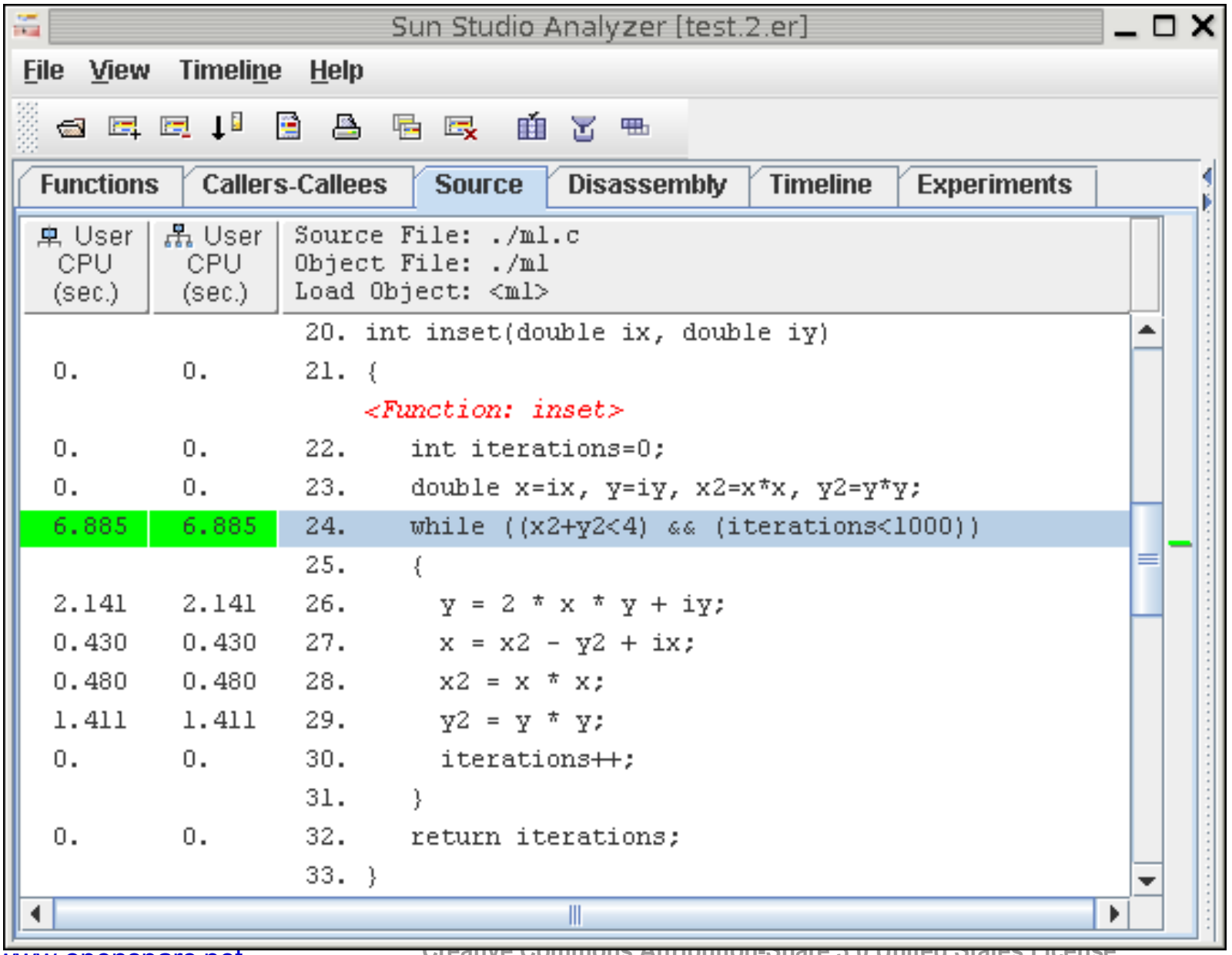
Profiling

- Performance Analyzer
  - Gather
    - `collect <app> <params>`
    - `collect -P <pid>`
  - Analyse
    - `analyzer <test.N.er>`
    - `er_print <test.N.er>`
- spot
  - Generate html report
    - `spot <app> <params>`
    - `spot -P <pid>`
- Compiler flags: `-g -g0 -xbinopt=prepare`

Application profile

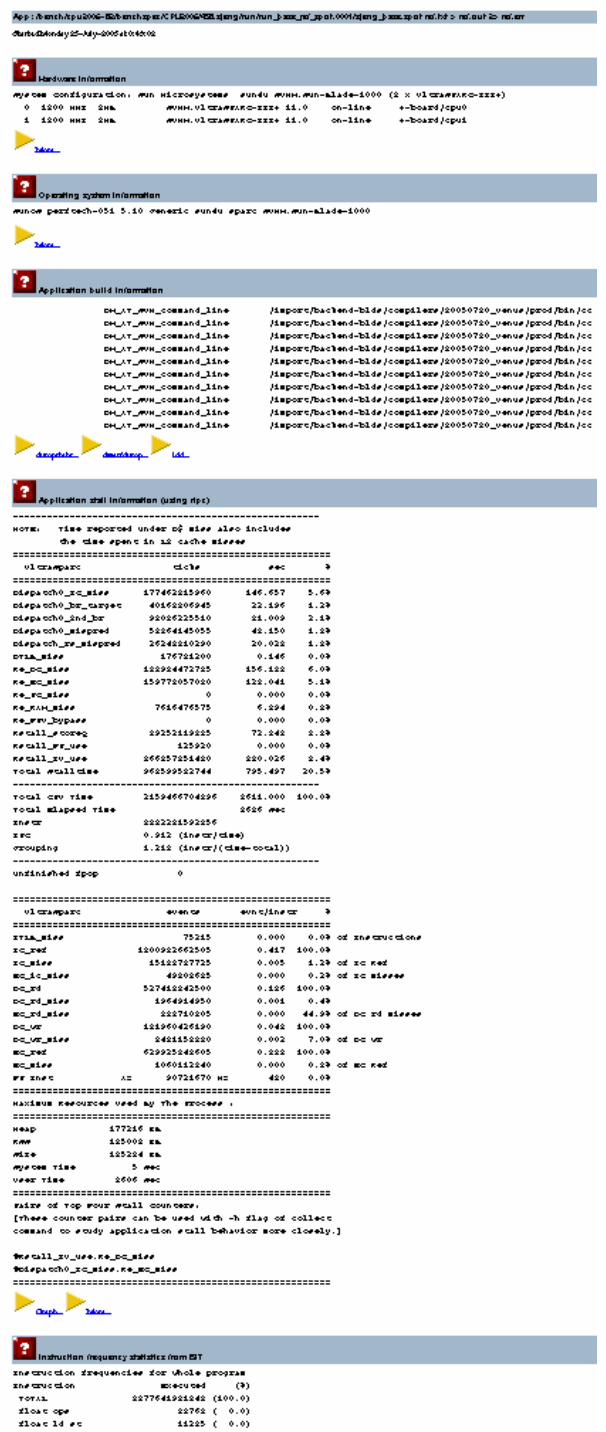


Source level profile



spot

- Generates detailed html report on application performance
  - > Performance counter events
  - > Time-based profile
  - > Event-based profile
  - > Bandwidth utilisation
  - > ...
- Helps in remote collaboration
- SPARC & x86
- http://cooltools.sunsource.net/spot/



Instruction Frequency (from BIT)

Instruction frequency statistics from BIT

Instruction frequencies for whole program

Instruction	Executed	(%)
TOTAL	7963939485	(100.0)
float ops	4194304000	(52.7)
float ld st	3145728000	(39.5)
load store	3502243842	(44.0)
load	2432696322	(30.5)
store	1069547520	(13.4)

Instruction	Executed	(%)	Annulled	In Delay Slot
TOTAL	7963939485	(100.0)		
lddf	2097152000	(26.3)	100	0
add	1415578342	(17.8)	0	5242882
stdf	1048576000	(13.2)	0	262143900
faddd	1048576000	(13.2)	0	0
prefetch	791674576	(9.9)	0	0
br	602931826	(7.6)	0	0
subcc	602931628	(7.6)	0	2
lduw	335544322	(4.2)	0	335544320
stw	20971520	(0.3)	1	0

Whole program... Functions...

Instruction frequency summary information

Instruction frequency detail

Performance Counters

Time lost due to various processor stall conditions

Memory consumption & system time

Graph of events over time

Application stall information (using ripc)

```
NOTE:   Time reported under D$ miss also includes
        the time spent in L2 cache misses
```

UltraSparc	ticks	sec	%
Dispatch0_IC_miss	177463815960	146.657	5.6%
Dispatch0_br_target	40168806945	33.196	1.3%
Dispatch0_2nd_br	98026335510	81.009	3.1%
Dispatch0_mispred	58264145055	48.150	1.8%
Dispatch_rs_mispred	36348210390	30.038	1.2%
DTLB_miss	176781300	0.146	0.0%
Re_DC_miss	188924473785	156.128	6.0%
Re_EC_miss	159778057020	132.041	5.1%
Re_PC_miss	0	0.000	0.0%
Re_RAW_miss	7616476575	6.294	0.2%
Re_FPU_bypass	0	0.000	0.0%
Rstall_storeQ	89353119825	73.842	2.8%
Rstall_FP_use	135930	0.000	0.0%
Rstall_IU_use	266257251420	220.036	8.4%
Total Stalltime	962599583744	795.497	30.5%
Total CPU Time	3159466704896	2611.000	100.0%
Total Elapsed Time		2686 Sec	
Instr	2883321593856		
IPC	0.913 (instr/time)		
Grouping	1.312 (instr/(time-total))		
unfinished fpop	0		

UltraSparc	events	evnt/instr	%
ITLB_miss	75315	0.000	0.0% o
IC_ref	1200923663505	0.417	100.0%
IC_miss	15128727735	0.005	1.3% o
EC_ic_miss	49808625	0.000	0.3% o
DC_rd	537418243500	0.186	100.0%
DC_rd_miss	1964914950	0.001	0.4%
EC_rd_miss	882710205	0.000	44.9% o
DC_wr	121960436190	0.042	100.0%
DC_wr_miss	8481158880	0.003	7.0% o
EC_ref	639935242605	0.222	100.0%
EC_miss	1060113240	0.000	0.2% o
FP Inst A=	90731670 M=	480	0.0%

```
Maximum Resources Used By The Process :
Heap            177316 KB
RSS             185008 KB
Size            185384 KB
System Time          5 Sec
User Time         2606 Sec
Pairs Of Top Four Stall Counters:
[These counter pairs can be used with -h flag of collect
command to study application stall behavior more closely.]

#Rstall_IU_use,Re_DC_miss
#Dispatch0_IC_miss,Re_EC_miss
```

Graph ... More ...

Profile - hardware events (-x flag)

Application HW counter profile output

```
./spot_run4/test.Dispatch0_br_target_Re_DC_miss.er: Experiment has warnings, see header for details
Current metrics: e.Dispatch0_br_target:e.Re_DC_miss:e.bit_fcount:e.bit_instx:e.bit_annul:name
Current Sort Metric: Exclusive Dispatch0_br_target Events ( e.Dispatch0_br_target )
Functions sorted by metric: Exclusive Dispatch0_br_target Events

Excl.                 Excl.          Excl. Bit   Excl. Bit    Excl. Bit    Name
Dispatch0_br_target   Re_DC_miss     Func        Inst Exec    Inst
Events sec.           Events sec.    Count                    Annul
0.826                 80.459         103         7963939485   204          <Total>
0.522                 30.965           1          705167424     2          tlb_miss
0.177                 31.270           1          705167424     2          cache_miss
0.127                 18.224         100         6553603800   200          fp_routine
0.                     0.              1                 837    0          main
0.                     0.              0                   0    0          _start
```

Time lost to Data Cache miss events

Profile - time

```
current filename for subsequent output: ./spot_run4/html/functions.func
Functions sorted by metric: Exclusive User CPU Time

Excl.      Incl.      Excl.      Excl.      Excl. Bit  Excl. Bit    Excl. Bit   Name
User CPU   User CPU   Sys. CPU   Wall       Func       Inst Exec    Inst
   sec.       sec.       sec.       sec.   Count                   Annul
119.834    119.834    0.570      120.684    103        7963939485   204         <Total>
 56.059     56.059    0.          56.219      1         705167424      2        [trimmed] tlb_miss src Caller-callee
 35.815     35.815    0.          35.915      1         705167424      2        [trimmed] cache_miss src Caller-callee
 27.709     27.709    0.          27.729    100        6553603800   200         [trimmed] fp_routine src Caller-callee
  0.250      0.250    0.570        0.821      0                0      0        _memset
  0.       119.834    0.           0.         1              837      0        [trimmed] main src Caller-callee
  0.       119.834    0.           0.         0                0      0        _start
```

Instructions executed in each routine

Number of times routine was called

Time spent in each routine

Assembly level profile

Source code for loop

```
                                                    26.    for (int i=0; i<size*16; i++) {cp= (int**)*cp;}
    0.          0                1    0                  [26]     12418:  sll          %o1, 4, %o1
    0.          0                1    0                  [26]     1241c:  cmp          %o1, 0
    0.          0                1    0                  [26]     12420:  ble,pn       %icc, 0x1246c
    0.          0                1    0                  [26]     12424:  add          %o1, -1, %o3
    0.          0                1    0                  [26]     12428:  add          %o3, 1, %g3
    0.          0                1    0                  [26]     1242c:  clr          %o1
    0.          0                1    0                  [26]     12430:  cmp          %g3, 1
    0.          0                1    0                  [26]     12434:  bl,pn        %icc, 0x12458
    0.          0                1    0                  [26]     12438:  mov          %o3, %o5
## 55.209       0        167772160    0                  [26]     1243c:  inc          %o1
    0.          0        167772160    0                  [26]     12440:  cmp          %o1, %o5
    0.          0        167772160    0                  [26]     12444:  ble,pt       %icc, 0x1243c
    0.801       0        167772160    0                  [26]     12448:  ld           [%o2], %o2
    0.          0                1    0                  [26]     1244c:  cmp          %o1, %o3
    0.          0                1    0                  [26]     12450:  bg,pn        %icc, 0x1246c
    0.          0                1    0                  [26]     12454:  nop
    0.          0                0    0                  [26]     12458:  ld           [%o2], %o2
    0.          0                0    0                  [26]     1245c:  inc          %o1
    0.          0                0    0                  [26]     12460:  cmp          %o1, %o3
    0.          0                0    0                  [26]     12464:  ble,a,pt     %icc, 0x1245c
    0.          0                0    0                  [26]     12468:  ld           [%o2], %o2
                                                    27.    return cp;
```

Loop entered once, trip count = ~170M

Load instruction that takes the time

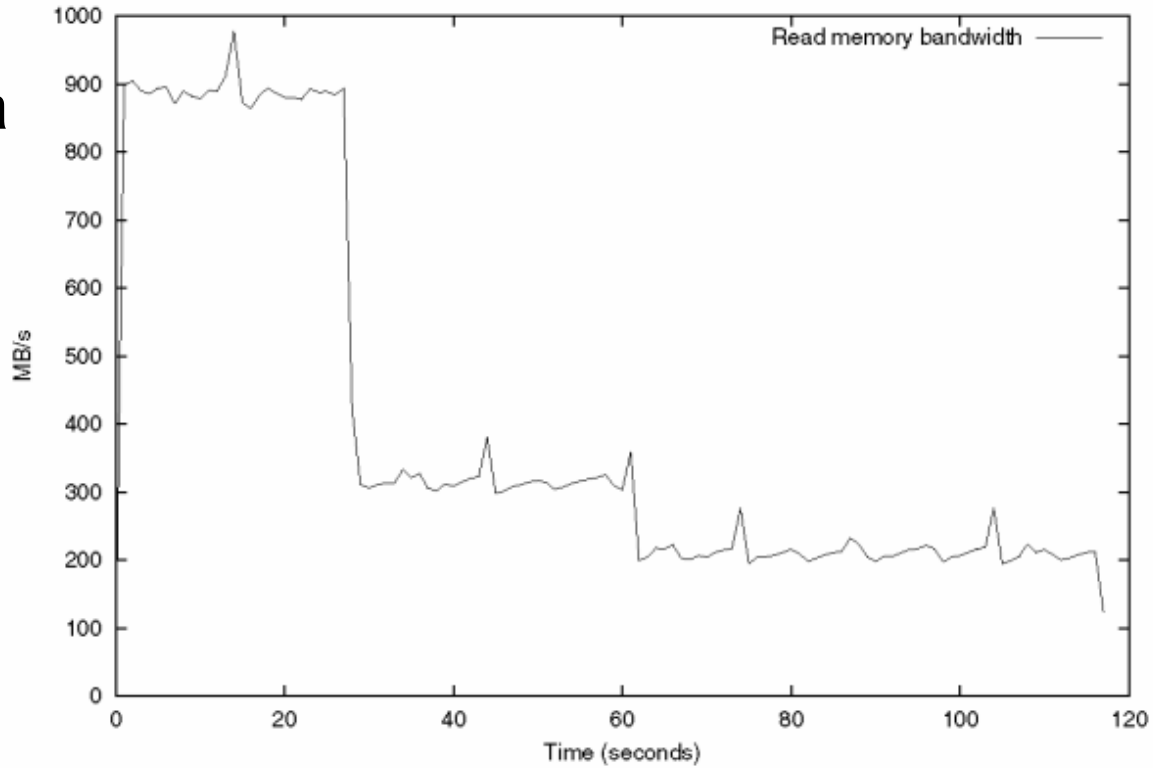
System-wide bandwidth data (-x flag)

Bandwidth data

```
Graph ./spot_run4/bandwidth.ps produced
Output graph ./spot_run4/bandwidth.ps generated.
------------------------------------------------------------
 Read  memory bandwidth: 399.613289596688 MB/sec (total bytes = 49025913856)
 Write memory bandwidth: 72.3323692908654 MB/sec (total bytes = 8873980416)
 Total memory bandwidth: 471.945658887553 MB/sec (total bytes = 57899894272)
 Elapsed time    : 117 secs
------------------------------------------------------------
```

Graph ... More ...

System-wide bandwidth data collected with -X flag and root permissions.

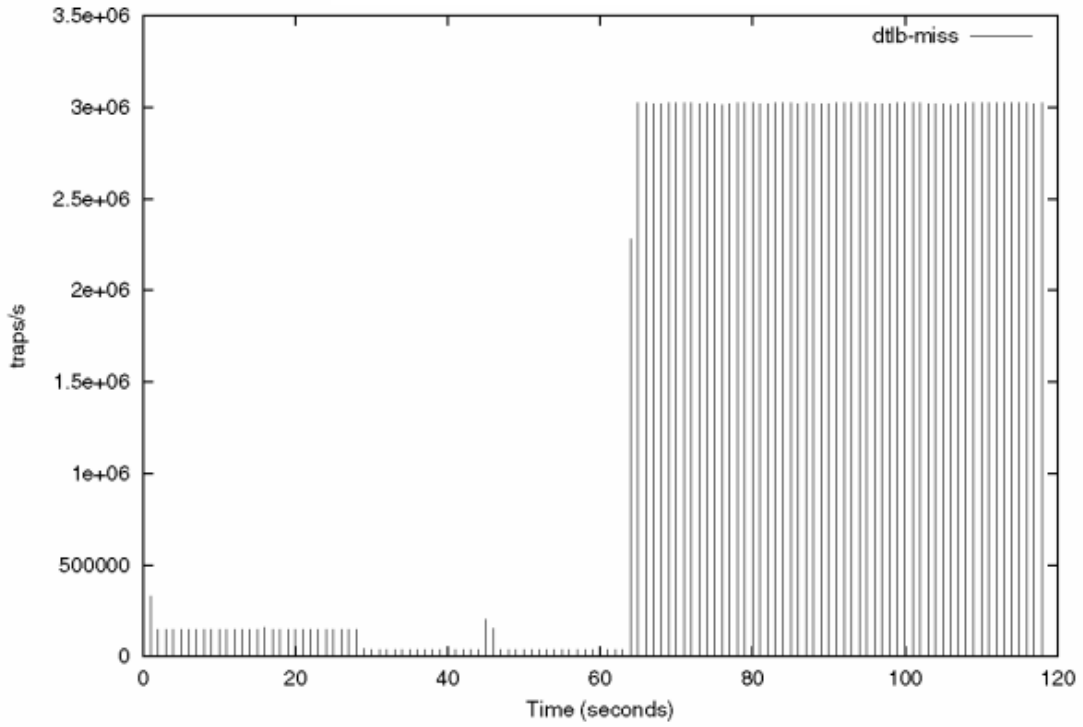


System-wide trap data (-x flag)

traps data

```
Output graph ./spot_run4/traps.ps generated.
Graph ./spot_run4/traps.ps produced
------------------------------------------------------------
               cleanwin          5.7 (traps/sec)
              dtlb-miss     182986.8 (traps/sec)
              dtlb-prot         54.3 (traps/sec)
            fill-kern-64        297.8 (traps/sec)
            fill-user-32          0.2 (traps/sec)
          fill-user-32-cln       19.3 (traps/sec)
             flush-wins           0.0 (traps/sec)
           fp-disabled            0.0 (traps/sec)
               get-psr            0.0 (traps/sec)
             gethrtime            0.1 (traps/sec)
               int-vec           18.6 (traps/sec)
              itlb-miss           2.2 (traps/sec)
               level-1            2.3 (traps/sec)
              level-10          100.0 (traps/sec)
              level-13            1.6 (traps/sec)
              level-14           16.0 (traps/sec)
               level-4           27.0 (traps/sec)
               level-6            0.3 (traps/sec)
               level-9            0.0 (traps/sec)
         spill-asuser-32          6.6 (traps/sec)
       spill-asuser-32-cln       51.6 (traps/sec)
            spill-kern-64       313.6 (traps/sec)
            spill-user-32          2.3 (traps/sec)
          spill-user-32-cln        1.2 (traps/sec)
             syscall-32            7.9 (traps/sec)
------------------------------------------------------------
```

Graph ... More ...



System-wide trap information collected with -x flag and root permissions

Agenda

- Compiler and tools options
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Multithreading code

- Profile application
  - Identify hot code
- Estimate performance gain
  - Amdahl's law
  - Performance gain depends on time spent in region to be parallelised
- Parallelise code

Expected gains from parallelisation

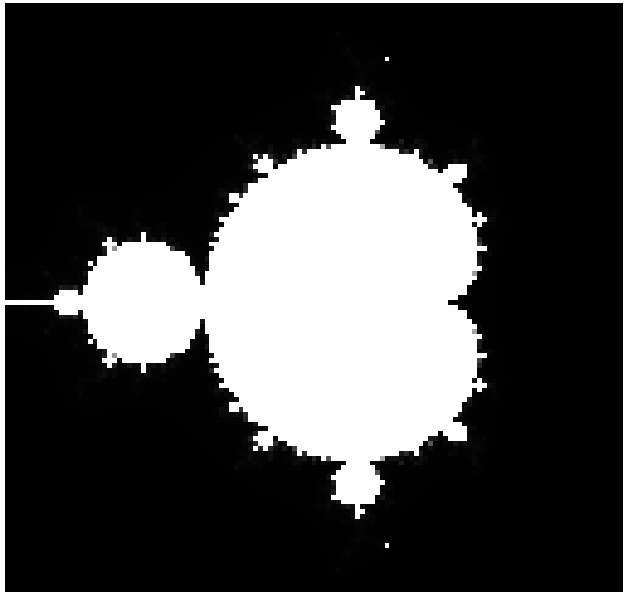
Maximum performance gain from parallelisation is determined by the time spent in code that can be parallelised.
Amdahl's Law.

Approaches to multithreaded coding

- Automatic parallelisation
  - Easy to use
  - Relies on compiler extracting parallelism
- OpenMP
  - Easy to use
  - Relies on developer to add directives to source code
- POSIX Threads
  - Very flexible
  - Potentially hard to use

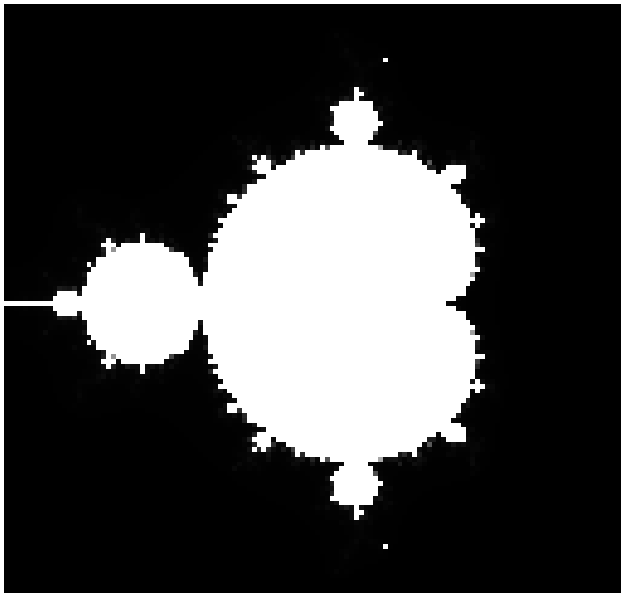
Mandlebrot set example

```
void calculate()
{
  int x,y;
  double xv,yv;
  for (x=0; x<SIZE; x++)
  {
    for (y=0; y<SIZE; y++)
    {
      xv = ((double)(x-SIZE/2))
           /(double)(SIZE/4);
      yv = ((double)(y-SIZE/2))
           /(double)(SIZE/4);
      data[x][y]=inset(xv,yv);
    }
  }
}
```



Determining if a point is in the set

```
int inset(double ix, double iy)
{
    int iterations=0;
    double x=ix, y=iy, x2=x*x, y2=y*y;
    while ((x2+y2<4) && (iterations<1000))
    {
      y = 2 * x * y + iy;
      x = x2 - y2 + ix;
      x2 = x * x;
      y2 = y * y;
      iterations++;
    }
    return iterations;
}
```



Running the application

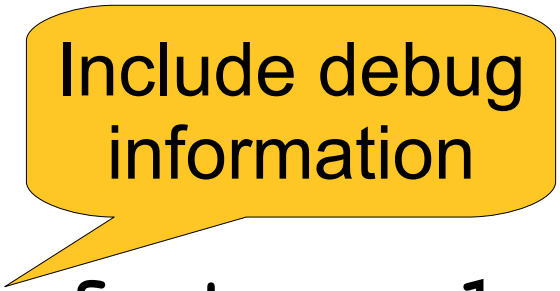
Include debug information

High optimisation

```
$ cc -g -fast -o m1 m1.c
$ timex m1
real        11.78
user        11.70
sys          0.05
```

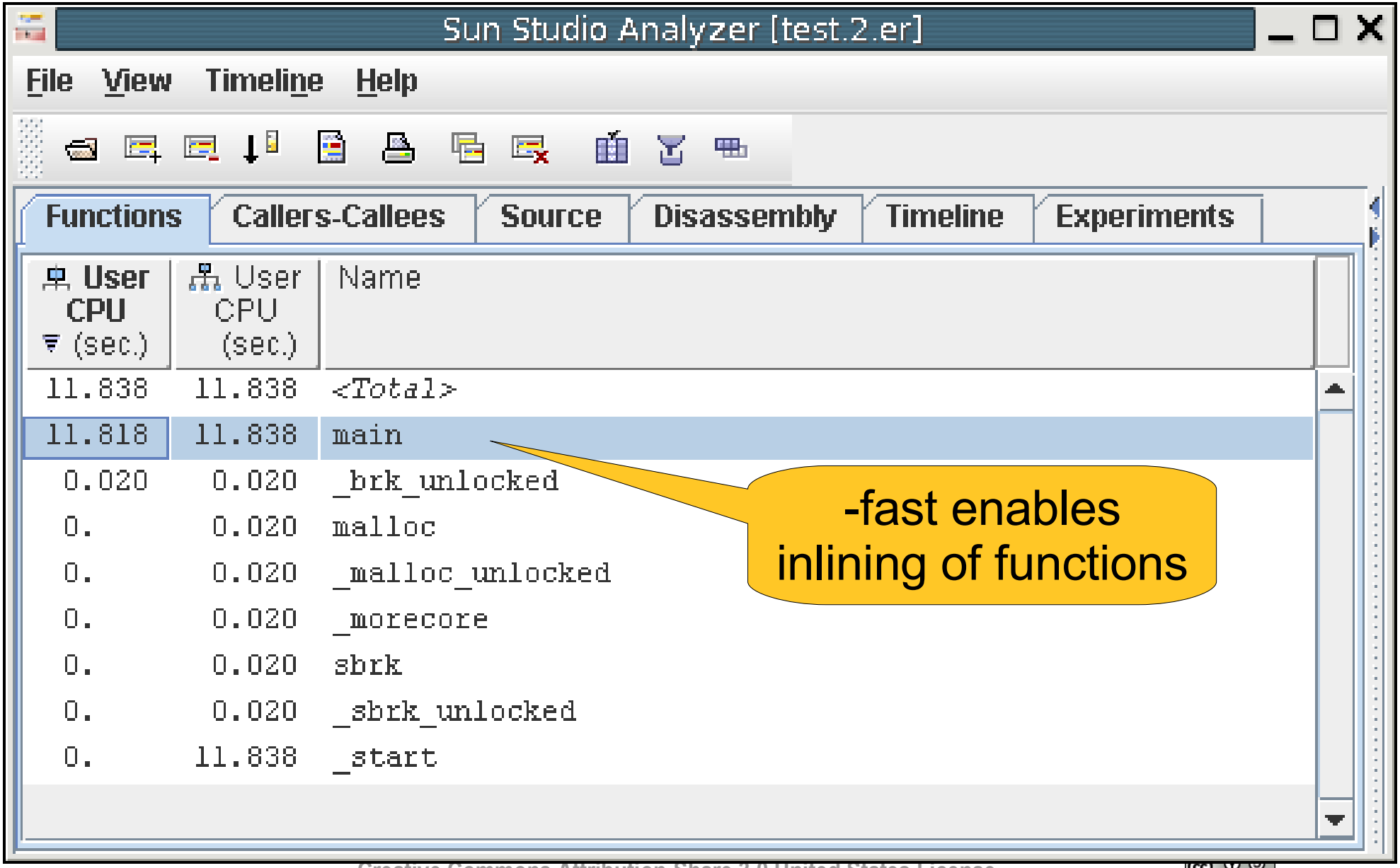
Runtime ~12s

Profiling the application




```
$ cc -g -fast -o m1 m1.c
$ collect m1
$ analyzer test.1.er
```

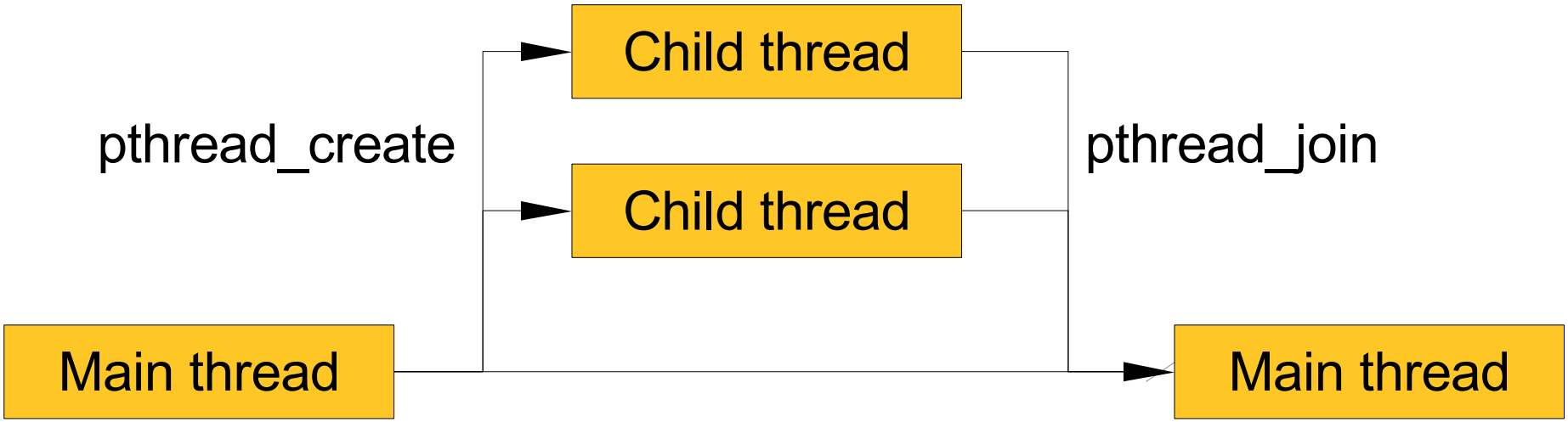
Application profile



POSIX threads (pthreads)

- Advantages:
  - High degree of control of parallelisation.
  - Flexibility in how threads cooperate.
- Disadvantages
  - Typically significant code changes
  - Increases program complexity
  - Can be hard to debug

Pthread model



Setting up pthreads

```
void main()
{
    pthread_t threads[2];
    int id[2];
    data = setup();
    for (int i=0; i<2; i++) {
      id[i]=i;
      pthread_create(&threads[i],0,
          calculate,(void*)&id[i]);
    }
    for (int i=0; i<2; i++) {
      pthread_join(threads[i],0);
    }
    validate();
}
```

Create threads

Wait for threads to finish

Parallelising using pthreads

```
void *calculate(void * arg)
{
    int x,y;
    double xv,yv;
    int id = *(int*)arg;
    int start = (int)(1.0*id/2*SIZE);
    int end =  (int)(1.0*(id+1)/2*SIZE);
    for (x=start; x<end; x++){
      for (y=0; y<SIZE; y++){
 ...
      }
    }
}
```

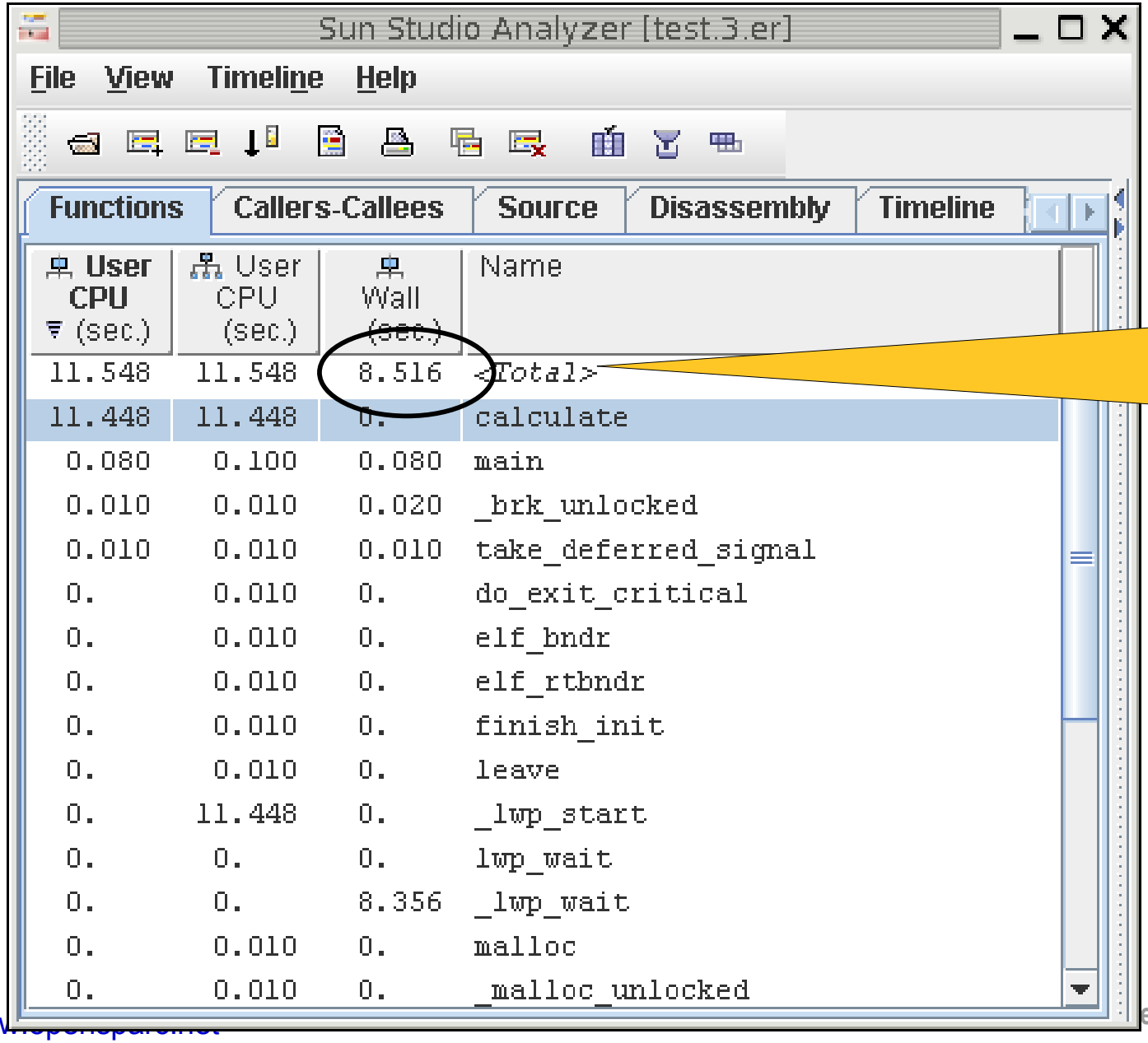
Divide work between threads

Running the application

compiler flags to support pthreads

```
$ cc -g -fast -o m2 m2.c -mt -lpthread
$ collect m2
$ analyzer test.2.er
```

Parallelising using pthreads

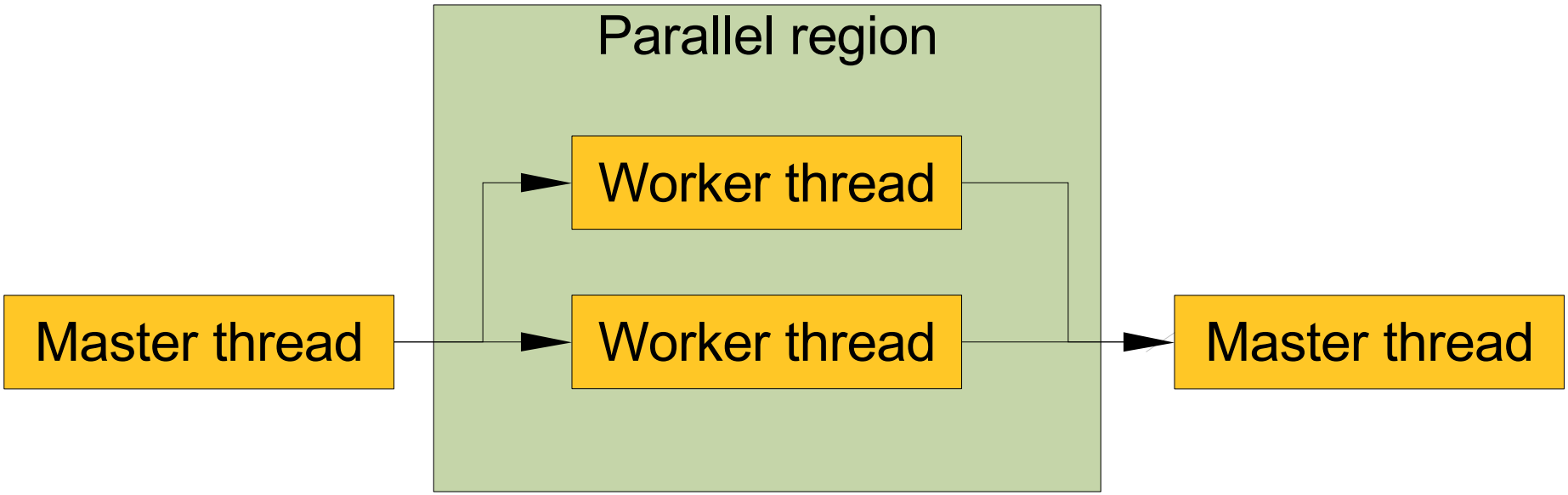


Wall time of 8.5 seconds. 1.35x faster than original 12s

OpenMP

- Advantages:
  - Easy to use
  - Minimal source changes
  - Incremental parallelisation
- Disadvantages
  - OpenMP 2.5 only supports parallelism through
    - Parallel loops
    - Parallel sections

OpenMP model

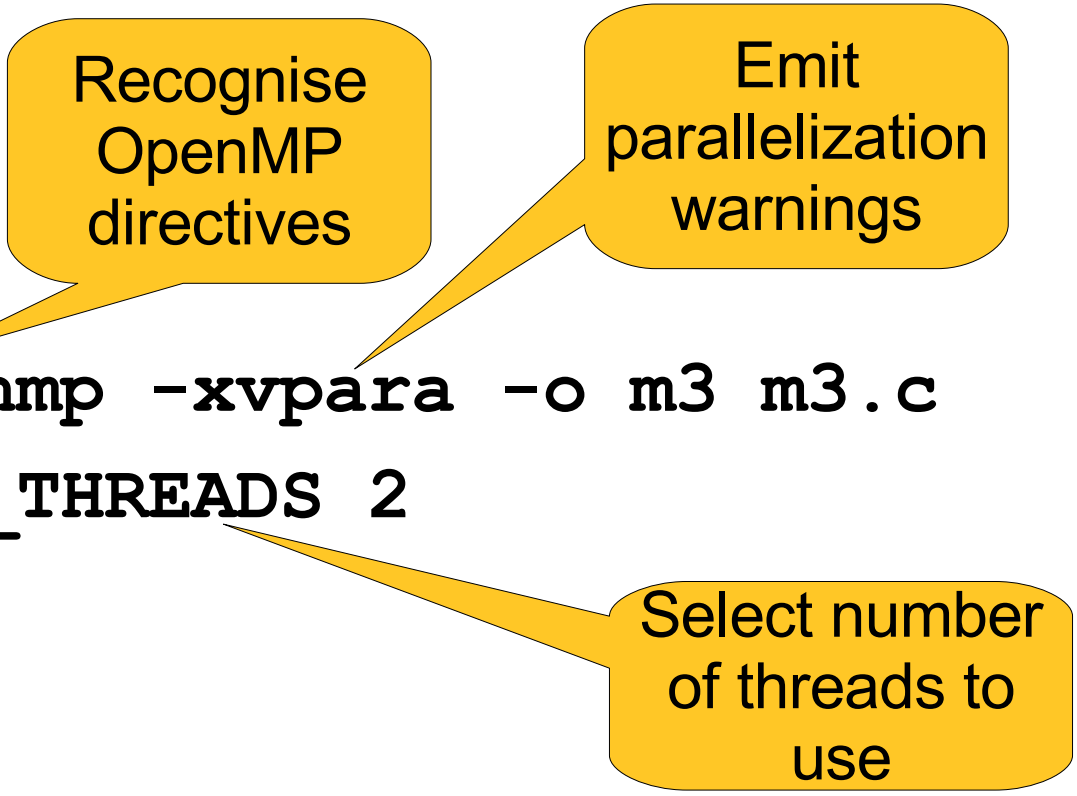


OpenMP

```
void calculate()
{
    int x,y;
    double xv,yv;
#pragma omp parallel for private(y,xv,yv)
    for (x=0; x<SIZE; x++){
      for (y=0; y<SIZE; y++){
        xv = ((double)(x-SIZE/2))
             /(double)(SIZE/4);
        yv = ((double)(y-SIZE/2))
             /(double)(SIZE/4);
        data[x][y]=inset(xv,yv);
      }
    }
}
```

Single source line change for parallelization

Using OpenMP



Same performance as Pthreads

Sun auto-scoping extension

```
void calculate()
{
    int x,y;
    double xv,yv;
#pragma omp parallel for default(__auto)
    for (x=0; x<SIZE; x++){
      for (y=0; y<SIZE; y++){
...
```

Compiler determines variable scoping

```
$ cc -g -fast -xopenmp -xvpara -o m3 m3.c
$ er_src -src calculate m3
...
   Source OpenMP region below has tag R1
   Variables autoscoped as SHARED in R1: data
   Variables autoscoped as PRIVATE in R1: xv, yv, y
...
```

Automatic parallelisation

- Advantages
  - Trivial to use
- Disadvantages
  - Limited ability to parallelise

Using autopar

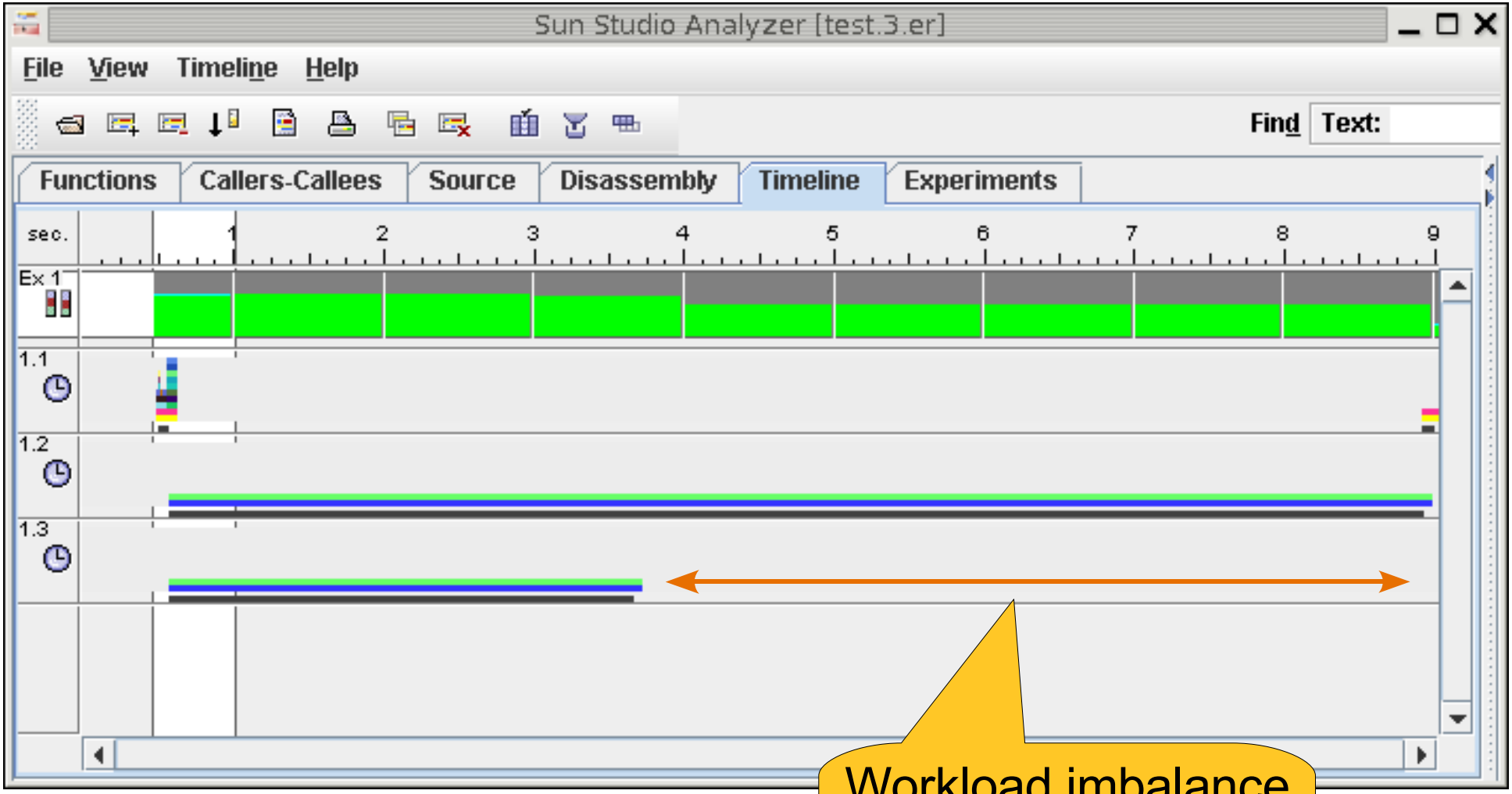
Enable auto-parallelisation

```
$ cc -fast -xautopar -xloopinfo -o m4 m1.c
...
"m1.c", line 40: PARALLELIZED, inl
  (inlined loop)
...

$ setenv OMP_NUM_THREADS 2
$ timex m4

real          6.00
user         11.69
sys           0.07
```

Emit auto-parallelisation information

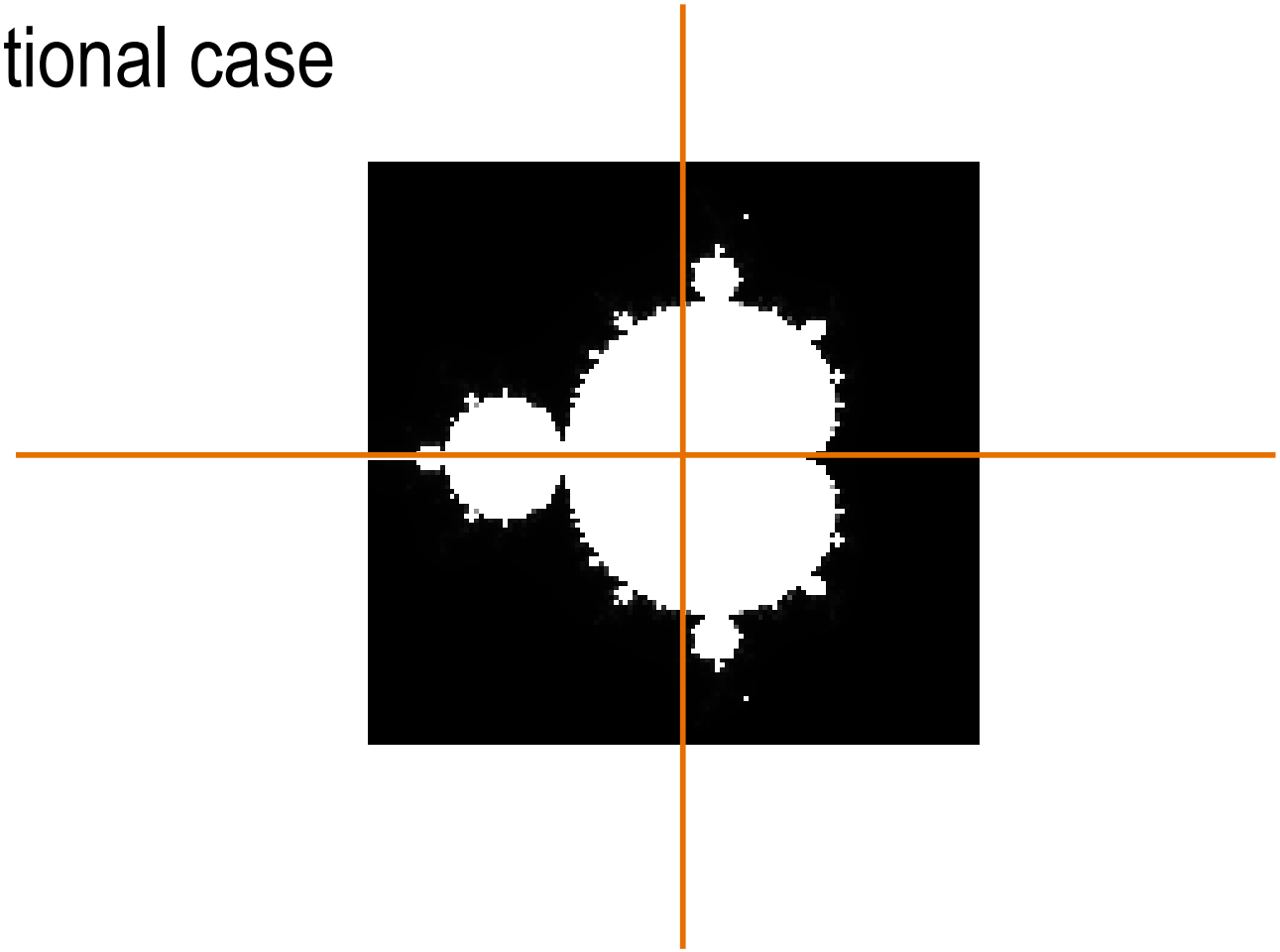
Wall time of 6 seconds
1.95x faster than original

Why is OpenMP/pthread slower?



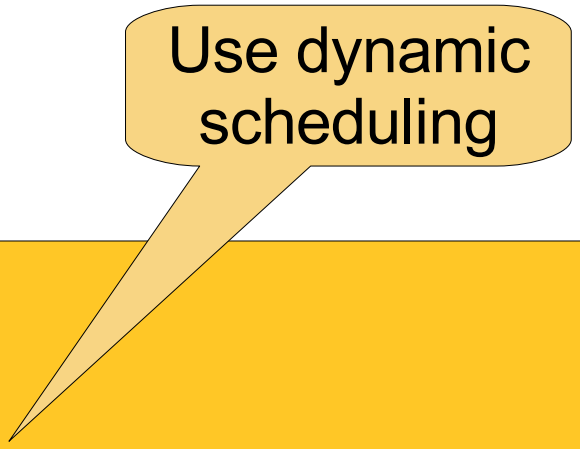
Load imbalance

- Horizontal symmetry, not vertical
- Exceptional case



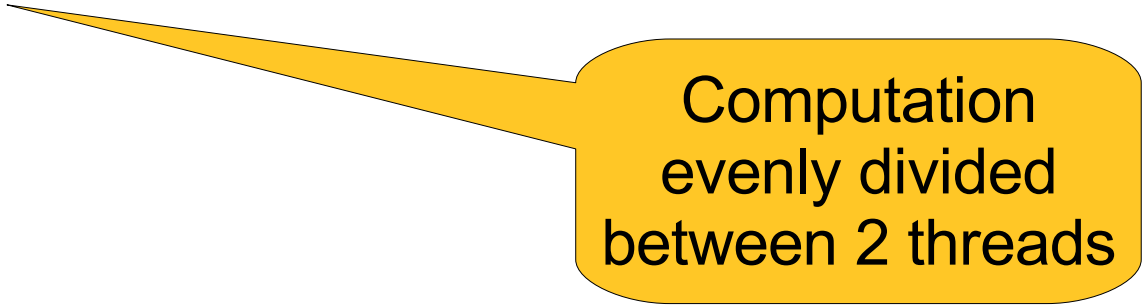
Dynamic scheduling OpenMP

```
void calculate()
{
    int x,y;
    double xv,yv;
#pragma omp parallel for \
        private (y,xv,yv)\
        schedule (dynamic)
    for (x=0; x<SIZE; x++) {
      for (y=0; y<SIZE; y++) {
        xv = ((double)(x-SIZE/2))
             /(double)(SIZE/4);
        yv = ((double)(y-SIZE/2))
             /(double)(SIZE/4);
        data[x][y]=inset(xv,yv);
      } } }
```



Using OpenMP

```
$ cc -fast -xopenmp -xvpara -o m3 m3.c
$ setenv OMP_NUM_THREADS 2
$ timex m3
real          5.84
user         11.38
sys           0.06
```



Notes on OpenMP and Autopar

What stops autopar?

- Function calls (could modify or read data)
  - Inlining with `-xipo` or `-xinline=<func>`
- Pointer aliasing
  - Use `restrict` keyword or `-xrestrict`
  - Use `-xalias_level=<level>`

Reductions

- Multiple threads cooperating to produce a single value.

```
double sum(double *a, int n)
{
  double t=0;
  for (int i=0; i<n; i++)
  { t+=a[i]; }
  return t;
}
% cc -xautopar -xloopinfo -c -O red.c
line 4: not parallelized, unsafe dependence(t)
```

Reductions

- Order of calculation will be different to serial case

```
double sum(double *a, int n)
{
  double t=0;
  for (int i=0; i<n; i++)
  { t+=a[i]; }
  return t;
}
% cc -xautopar -xreduction -xloopinfo -c -O red.c
line 4: PARALLELIZED, reduction, and serial version
  generated
```

Reductions

- OpenMP

```
double sum(double *a, int n)
{
  double t=0;
#pragma omp parallel for reduction(+:t)
  for (int i=0; i<n; i++)
  { t+=a[i]; }
  return t;
}
```

OpenMP 2.5 – parallel sections

```
#pragma omp parallel sections
{
   #pragma omp section
   {
      /*Region 1*/
   }
   #pragma omp section
   {
      /*Region 2*/
   }
}
```

OpenMP 3.0 - tasks

- Spread tasks over multiple threads

```
node * p = head;
while (p)
{
   #pragma omp task
   {
     process(p);
   }
   p = p->next;
}
```

Sharing data between threads

Sharing data between threads

- Multiple threads updating same data (variable, array etc.)
- Data needs to be `volatile`
  - To avoid being held in register
- Only one thread should update at a time

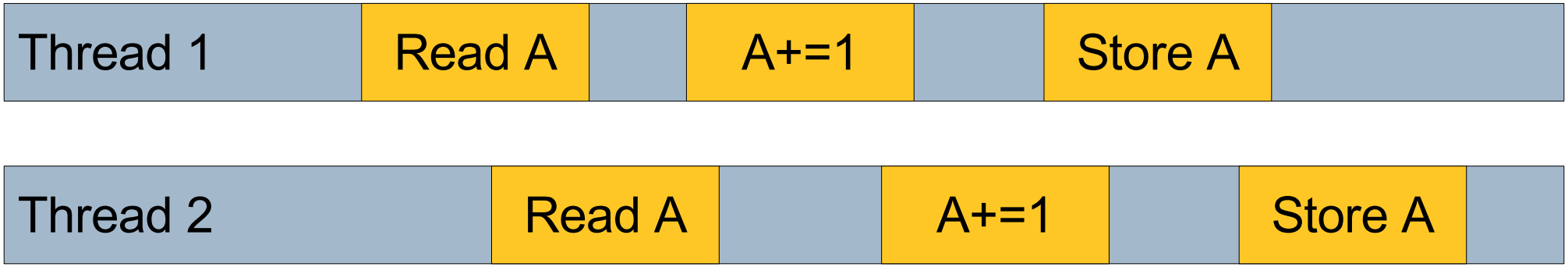
Sharing between Pthreads

```
volatile int sum=0;
...
void *calculate(void * arg) {
    int x,y;
    double xv,yv;
    int id = *(int*)arg;
    int start = (int)(1.0*id/2*SIZE);
    int end =  (int)(1.0*(id+1)/2*SIZE);
    for (x=start; x<end; x++) {
      for (y=0; y<SIZE; y++) {
        ...
        sum+=inset(xv,yv);
      }
    }
}
```

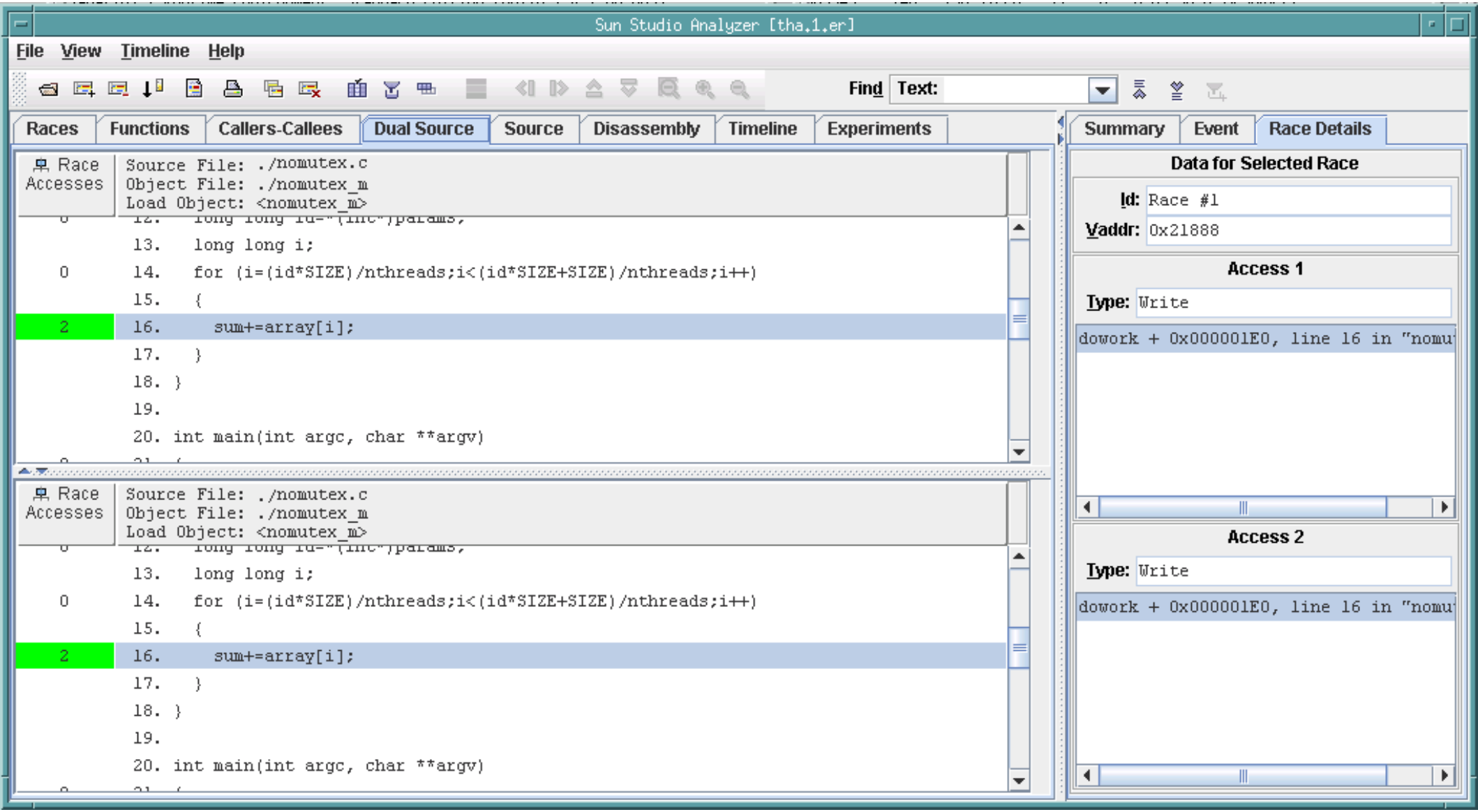
Variable `sum` is `volatile` but shared between multiple threads

Data races

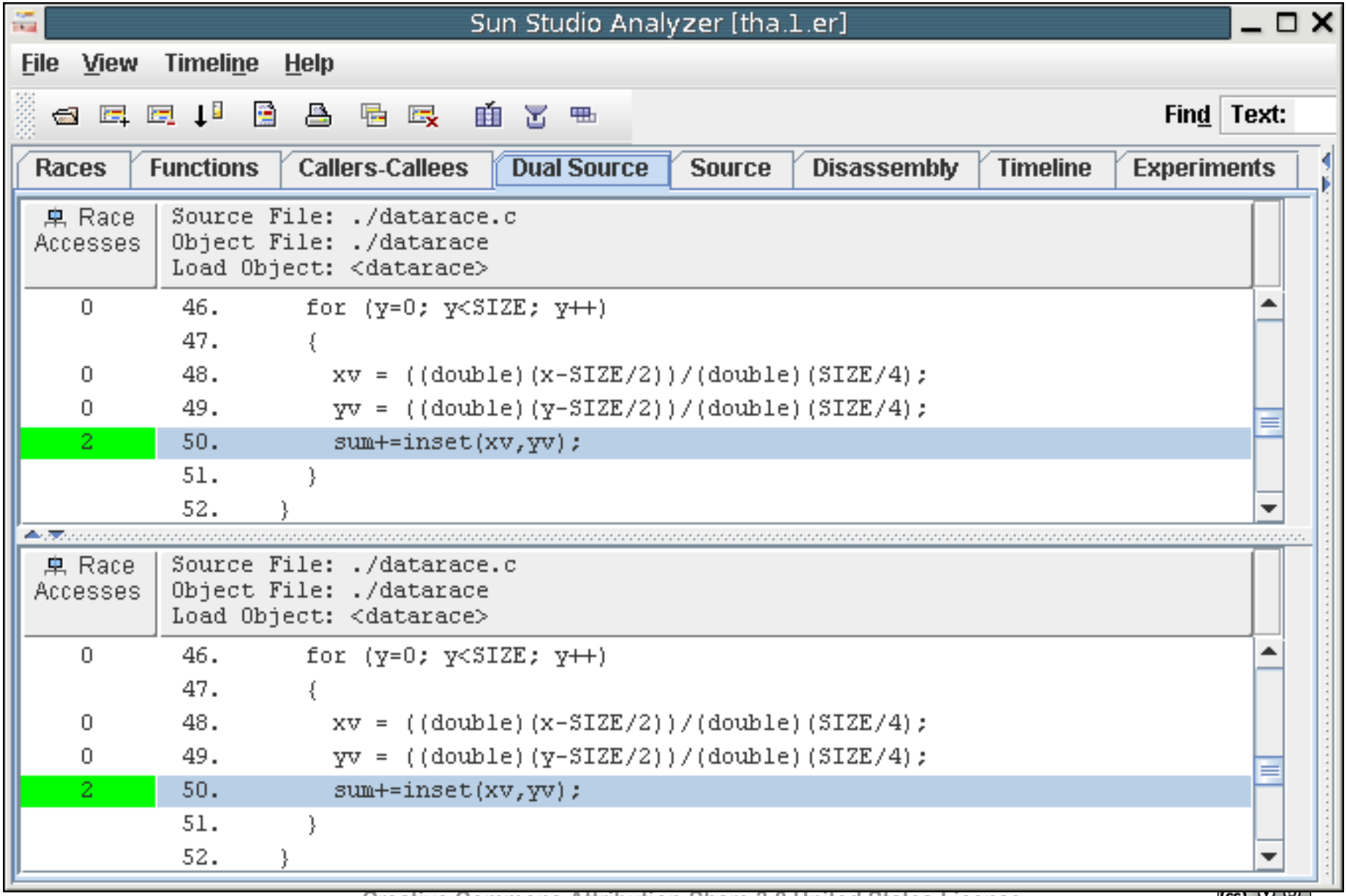
- When:
  - > several threads access same variable
  - > without synchronisation
  - > one or more accesses are writes
- Results in hard to debug non-deterministic behaviour



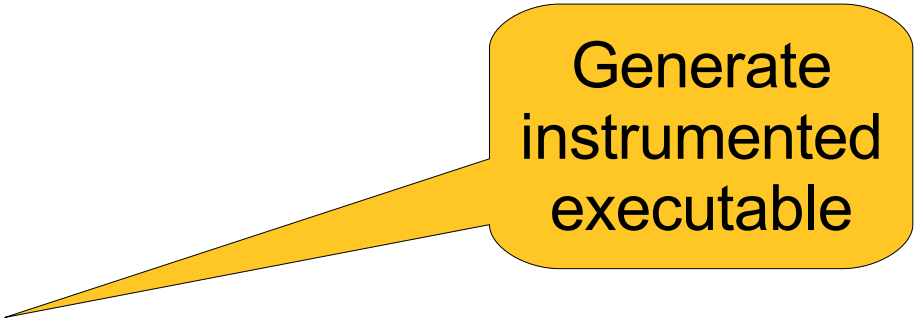
Thread analyzer



View source code



Notes: Data races




```
$ cc -g -xinstrument=datarace -mt \
  -lpthread -o race race.c
$ collect -r on datarace
$ analyzer tha.1.er
```

Fixing data accesses

- Single thread access:
  - Mutex locks
  - Critical regions (OpenMP)
- Lock-less:
  - Atomic operations (**`man atomic_ops`**)
- Data sharing:
  - Thread local storage
  - OpenMP reduction directive

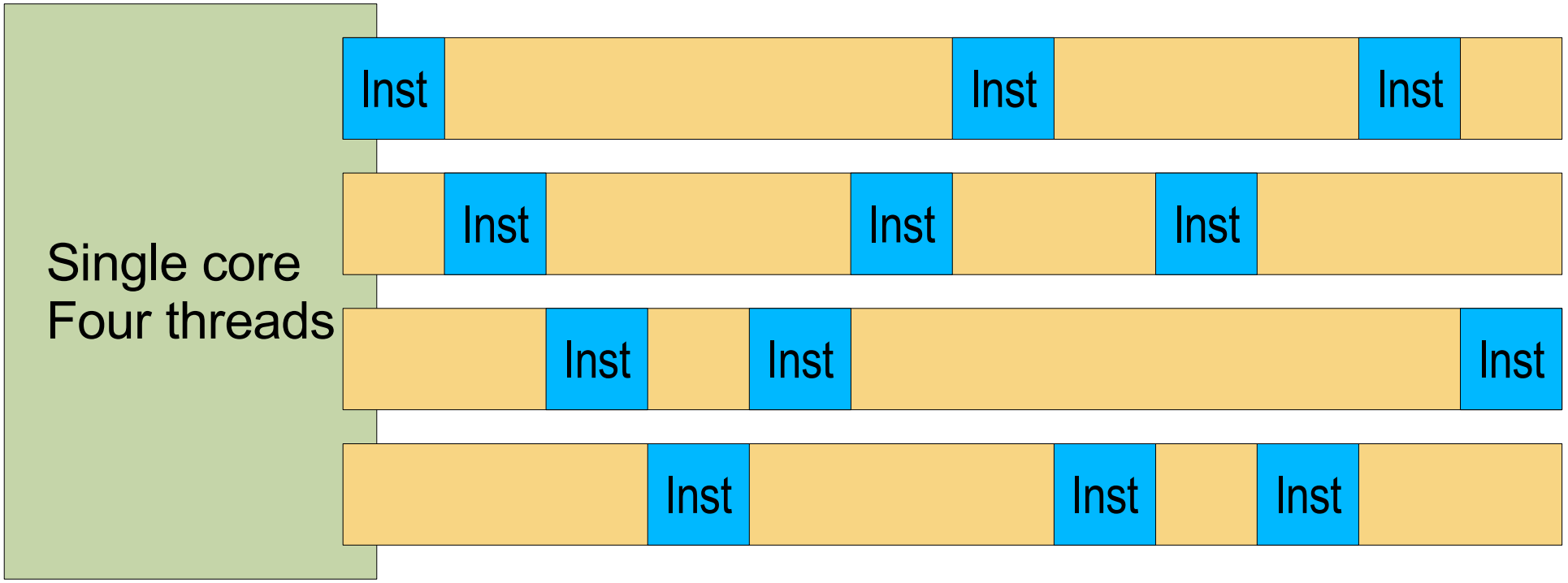
Agenda

- Compiler and tools options
- Compiling applications
- Profiling applications
- Writing parallel applications
- System utilisation
- Other resources

System utilisation

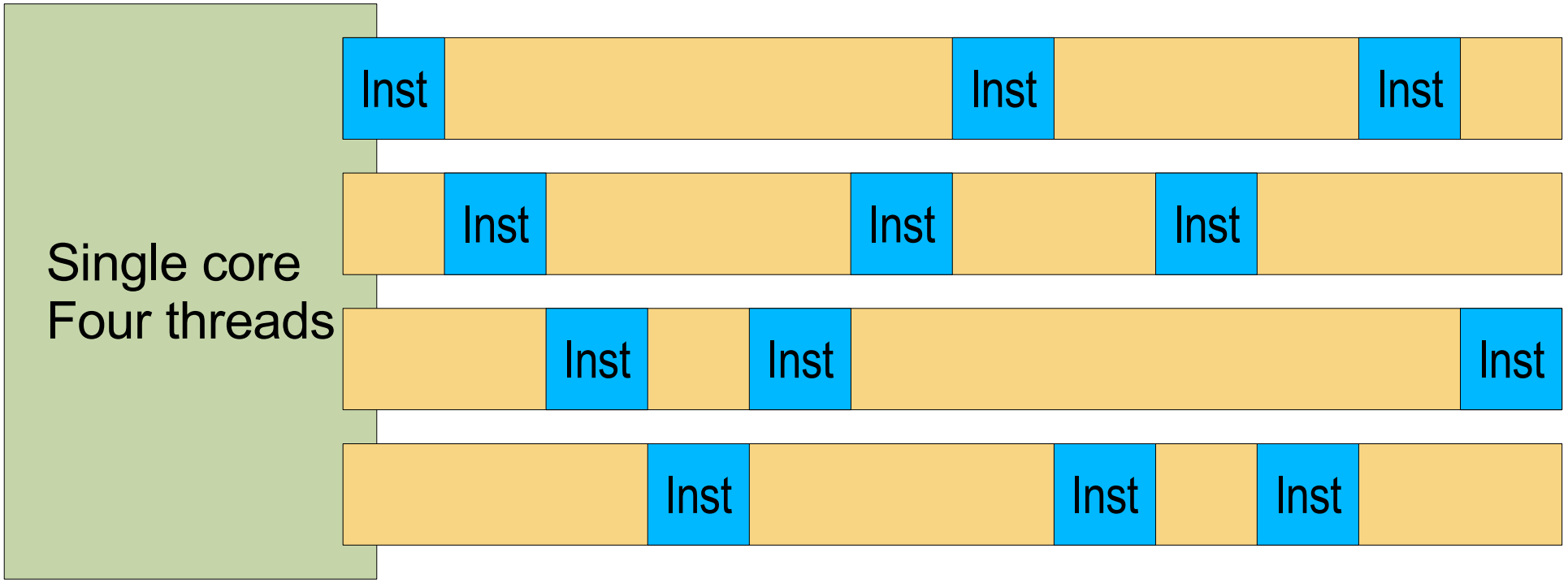
- Tools:
  - System aggregate: **`vmstat`**
  - Processes: **`prstat`**
  - Processors: **`mpstat`**
- Indicate that processors are busy
- But not the utilisation of the cores

Core utilisation



- Multiple threads share core
- Instruction count indicates core utilisation
- Utilisation = Instructions issued / Max. Instruction Issue rate

Stall and instruction budgets

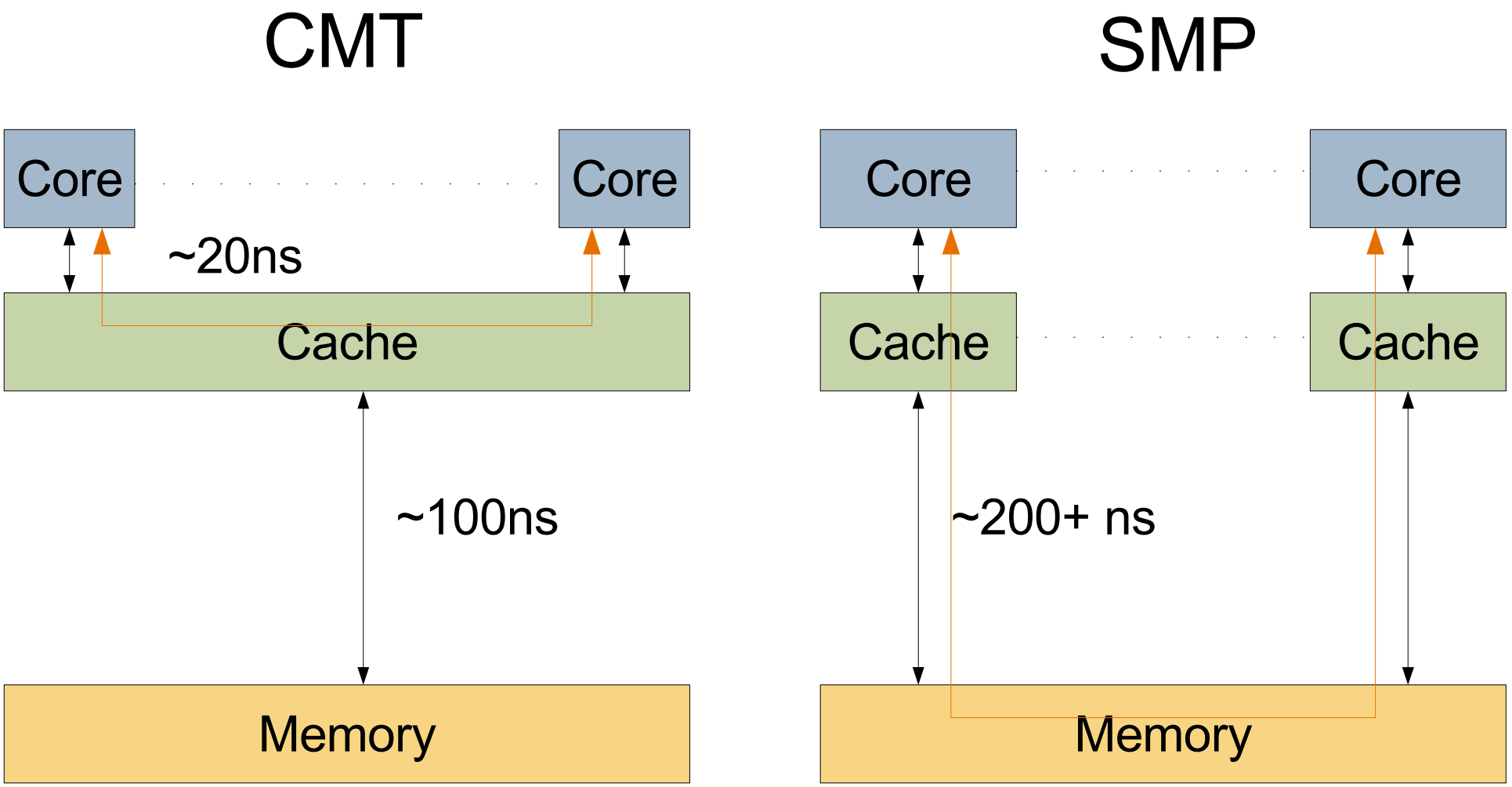


- One instruction issued from one of four threads
- Three threads do not issue – can be stalled or ready
- Stall budget of 3x instruction issue
- Instruction issue rate not impacted until stall budget exceeded

Processor utilisation

- Utilisation = percentage of peak instruction issue rate
- Raw data: **`cpustat`**
- Formatted: **`corestat`**
- Per process:
  - **`cputrack`**
  - **`ripc`** (from spot)

CMT vs SMP

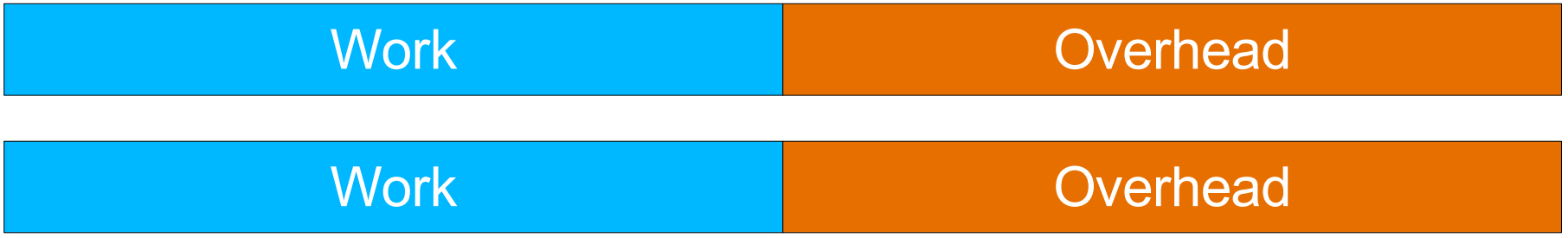


Limits of parallelisation

- Limit to parallelisation is:
  - > Costs of synchronisation
  - > Is greater than gain from more threads

Single thread

Two threads



Microparallelism

- Many cores
- Low communication latency between cores
- => Synchronisation cost is low
- => Profitable to parallelise smaller regions

Microparallelism outline

- Get next quanta of work
- Check that it is safe to start work
  - No other thread touching writing to read data
  - Avoids data races
  - Ensures correct ordering
- Perform work
- Update status to indicate work completed

Optimising for CMT

- Traditional approach
  - First reduce latency
  - Then parallelise
- CMT approach
  - First parallelise
  - Check instruction issue rate
    - If at max, need to reduce instruction count
    - If not, improve instruction latency

Compiling for CMT

- Generic compiler flags ok:
  - **`-xtarget=generic`**
- Usual optimisations are good:
  - Profile feedback: **`-xprofile=[collect:|use:]`**
  - Crossfile optimisation: **`-xipo`**
  - At least **`-O`**
  - Perhaps **`-fast`**

Summary

- Always profile your application
- Always use optimisation
- Use multiple threads
  - Autopar
  - OpenMP
  - PThreads
- On CMT synchronisation costs are lower
- Check instruction count
  - corestat
  - BIT

Agenda

- Compiler and tools options
- Compiling applications
- Profiling applications
- Writing parallel applications
- Calculating system utilisation
- Other resources

Web-based resources

- Developer portal:
  http://developers.sun.com/
- Documentation:
  http://developers.sun.com/sunstudio/documentation/index.jsp
- Forums:
  http://developers.sun.com/sunstudio/community/forums/index.jsp
- Blogs:
  http://developers.sun.com/sunstudio/community/blogs/index.jsp

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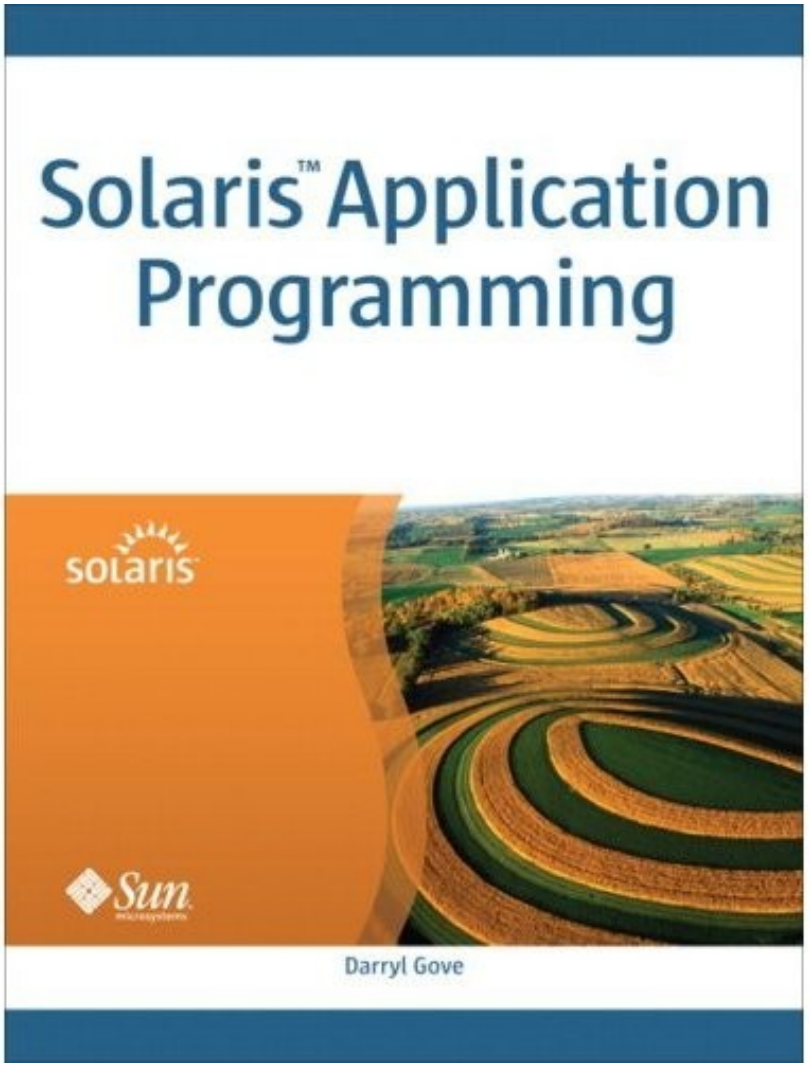
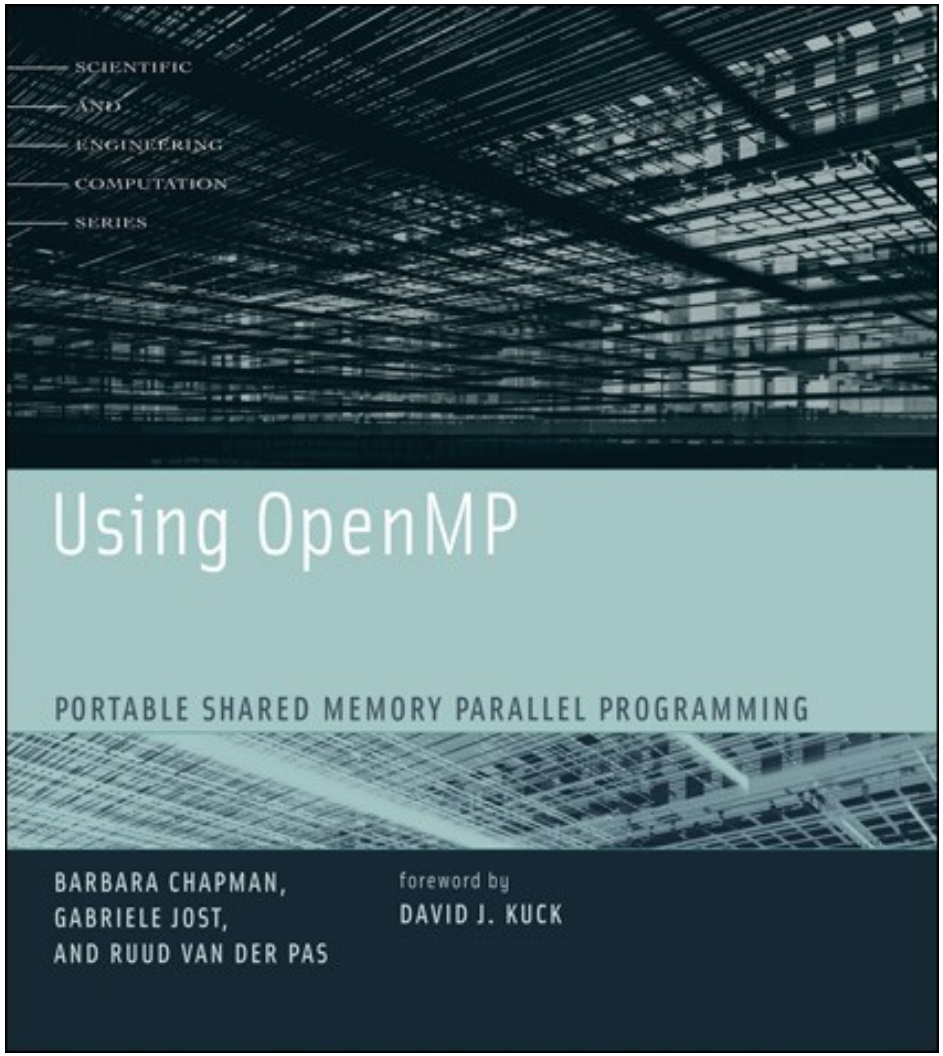


Using OpenMP



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