Points missed: _____

Student's Name: ____

Total score: ____/100 points

East Tennessee State University Department of Computer and Information Sciences CSCI 4717 – Computer Architecture TEST 3 for Fall Semester, 2005

Section 201

Read this before starting!

- The total possible score for this test is 100 points.
- This test is closed book and closed notes
- You may use one sheet of scrap paper that you will turn in with your test.
- When possible, indicate final answers by drawing a box around them. This is to aid the grader (who might not be me!) Failure to do so might result in no credit for answer. Example:

- **1 point will be deducted** per answer for missing or incorrect units when required. **No** assumptions will be made for hexadecimal versus decimal, so you should always include the base in your answer.
- If you perform work on the back of a page in this test, indicate that you have done so in case the need arises for partial credit to be determined.

"Fine print"

Academic Misconduct: Section 5.7 "Academic Misconduct" of the East Tennessee State University Faculty Handbook, October 21, 2005:

"Academic misconduct will be subject to disciplinary action. Any act of dishonesty in academic work constitutes academic misconduct. This includes plagiarism, the changing of falsifying of any academic documents or materials, cheating, and the giving or receiving of unauthorized aid in tests, examinations, or other assigned school work. Penalties for academic misconduct will vary with the seriousness of the offense and may include, but are not limited to: a grade of 'F' on the work in question, a grade of 'F' of the course, reprimand, probation, suspension, and expulsion. For a second academic offense the penalty is permanent expulsion."

Memory Management

1. What does a "page fault" mean in terms of virtual memory? (3 points)

When a process requests code or data from a page that has not yet been loaded in real memory requiring the operating system to load the page.

2. How does virtual memory allow programs that are larger than real memory to be run? (2 points)

By only loading into memory the pages that a process is currently using. This means that unused pages can be left on the hard drive and therefore not take up as much memory as the full application.

3. What problem is typically caused by very small virtual memory page sizes? (2 points)

An increase in the number of page faults because of the large number of pages that will be required to support a process.

4. What problem is typically caused by large virtual memory page sizes? (2 points)

Large pages waste space because if only a small portion of the page is used, a large amount of code or data is loaded unnecessarily.



6. If 6 processes are running in paged memory with a page size of 512 bytes, what is the largest amount of memory that is being wasted? (2 points)

A process that has been divided into pages fills all of the pages except for the last one. The last one takes the remaining bytes after dividing the process by the page size, i.e., process size modulo the page size. The smallest value this can be is *one word* which means that for this case 512 - 1 = 511 bytes at most would be wasted per process. For six processes, this would be $511 \times 6 = 3066$.

7. True or false: The benefit of an inverted page table is that instead of maintaining an entry for every page in a process, it maintains an entry for every frame available in real memory. (2 points)

Assembly Language Details

8. Assume the address field of an instruction contains a decimal 105. The address field represents something different depending on the addressing mode used by the instruction. For each of the addressing modes below, describe where the processor is getting its data from. (2 points each)

a.) Immediate: The data is the operand, i.e., 105. No fetching of operands is needed.

b.) Indirect: The processor must go to memory address 105. At that location, it will find the data it needs to retrieve and use for the instruction.

9. Store the value ABCDEF₁₆ as individual bytes using big endian starting at address 100 in the figure to the right. Note that the figure will have blank locations after you are completed. (2 points) 0100_{16} AB 0101_{16} CD 0102_{16} EF

10. The code below uses a three-operand instruction. In the space below, write three short programs that do exactly the same thing, one with two-operand instructions, one with one-operand instructions, and one with zero-operand instructions. If it fits the need of the instruction, feel free to use register names R1, R2, etc. (5 points)

ADD C, A, B ; C = A + B

Two operand instructions	One operand instructions	Zero operand instructions		
MOV C, A ADD C, B	LOAD A ADD B STORE C	PUSH A PUSH B ADD POP C		

For the next 2 questions, use the following abbreviations for the stages of a 6-stage pipeline
Fetch instruction (FI)Fetch instruction (FI)Decode instruction (DI)Calculate operands (CO)Fetch operands (FO)Execute Instruction (EI)Write Operand (WO)

11. List all of the stages such that any 2 of them occurring at the same time could result in a bus resource conflict? (2 points)

FI, FO, WO

12. List all of the stages that are *always* used in the execution of an instruction. (Hint: Think of the simplest instruction and decide which stages are needed to execute it.) (2 points)

FI, DI, EI

Pipelines and Branch Prediction

- In an ideal implementation, what is the speed up of a processor with a k-stage pipeline over a nonpipelined processor if the duration of each stage is τ? Don't consider pipeline flushes or delays incurred between stages. (2 points)
 - a.) $k \cdot \tau$ b.) k-2 c.) k-1 d.) k e.) $(k+1) \cdot \tau$ f.) $k+\tau-1$ g.) none of the above
- 14. How many bits are required for each conditional branch in a branch history table in order to remember the past 3 branch outcomes for a specific instruction? (2 points)

a.) 2 (b.) 3 c.) 4 d.) 5 e.) 6 f.) 7 g.) 8

- 15. In what level of the memory hierarchy are the branch history bits typically stored? (2 points) The cache
- 16. Identify the two problems with having two pipelines (one for the branch stream and one for the non-branch stream) in order to minimize the delay caused by an incorrectly predicted branch. (3 points)
 - 1. If additional branches are encountered immediately after the first branch, we might need far more than two pipelines.
 - 2. Resource contention Additional active pipes would require multiple resources such as buses, ALUs, etc.
- 17. True or false: As far as the execution of a for-loop is concerned, the benefits of a loop buffer are only realized if the entire code executed for a for-loop can be contained in the buffer. (2 points)
- 18. How many cycles does prefetching the branch target save over not prefetching the branch target when a branch is incorrectly predicted? (2 points)
 - a.) 0 (b.) 1 c.) 2 d.) 3 e.) none of the above

For problems 20, 21, and 22, consider the following section of code.

19. Once compiled, how many conditional jumps would be contained in the machine code resulting from the above section of code? (2 points)

There would be one conditional jump to check if j reached 5 and a second to check if i reached 10. Therefore, there would be *2 conditional jumps*.

20. After fully executing the above section of code, how many conditional jumps would the CPU have encountered? (2 points)

The processor would have encountered the conditional jump on the outer loop 10 times and on the inner loop 10×5 times for a total of *60 times*.

21. Using the static branch prediction algorithm "branch always," how many of the conditional jumps calculated in the previous problem would have been predicted *incorrectly*? (2 points)

In a for-loop, branch always incorrectly predicts the branch once per full execution of the loop. For example, the conditional branch on the outer loop would have been incorrectly predicted only after i had been incremented to 10. Since the inner loop is executed 10 times and the outer loop is executed once, there will be *11 incorrect predictions*.

- 22. What are the three items that should be stored in a branch history table? (3 points)
 - 1. the address of the branch instruction itself
 - 2. the address of the branch target, i.e., where we are going if a branch occurs
 - 3. bits indicating the past branch history of this branch instruction

RISC Processors

- 23. (True) or false: The purpose behind the delayed branch is to avoid flushing the pipeline. (2 points)
- 24. Place a check mark next to each statement below that *tends to be true* for a RISC processor. (1 point each)
 - **RISC** processors have more specialized registers rather than general purpose.
 - **RISC** processors have a very limited number of addressing modes.
 - **RISC** processors using 2-way pipelined timing (only two instructions can be in the pipe at any on time) do so in order to avoid two simultaneous memory accesses.
 - **RISC** processors use a fixed instruction length.

25.	In order to av need to modif optimized del (2 points)	oid flushing the fy the code to ayed branch, v	rocessor, we o create an yed after L4?	L1: L2: L3: L4:	mov add cmp jne	al,bl ch,23 bl,ch label_	a	
(a.) L1 only	b.) L2 only	c.) Either L1 or L2	d.) Neither	۰СЦ	THC	II	

26. Find the absolute minimum number of registers required to execute the code below assuming that every variable in the code will need to be placed in a register. (In other words, don't try to optimize the code, just the use of registers.) (2 points)



The arrows identify the periods during which each of the variables is active. By aligning arrows that do not overlap in time, we can see that a minimum of 5 columns of arrows can be created. This means that a minimum of 5 registers will be needed.

Superscalar

27. Which *two* of the following dependencies exhibit similar behavior? (2 points)

a.) Data dependency

b.) Procedural dependency

(c.) Resource conflict

28. Identify the write-read, write-write, and read-write dependencies in the instruction sequence below by entering each line pair with a dependency in the correct column of the table to the right. For example, if L1 and L4 had a write-write dependency (which they don't), you would enter L1-L4 in the column labeled "write-write". (4 points)

L1: R1 = R2 + R5 L2: R2 = R2 + 1 L3: R3 = R2 + R5 L4: R5 = R1 - R2 L5: R1 = R2 + 30	write-read L1 - L4 L2 - L3 L2 - L4 L2 - L5	write-write L1 – L5	read-write L1 - L2 L1 - L4 L3 - L4 L4 - L5
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29. In the code below, identify references to initial register values by adding the subscript 'a' to the register reference. Identify new allocations to registers with the next highest subscript and identify references to these new allocations using the same subscript. (3 points)

R1 _b	$= R2_a$	+	R5 _a
R2 <mark>b</mark>	= $R2_a$	+	1
R3 _b	= R2 _b	+	R5 <mark>a</mark>
r3 _c	= R3 _b	+	R2 <mark>b</mark>
R5 _b	= R1 _b	÷	R3 <mark>c</mark>
R1 _c	= R2 _b	×	R5 _b

Questions 32 through 35 are based on the "in-order issue/in-order completion" execution sequence shown in the figure to the right.

30. Why do I3 and I4 need to stay in the decode stage of the pipeline for two cycles? (2 points)

Decode		Execute		Write		Cycle		
I1	I2							1
I3	I4	I1		I2				2
I3	I4	I1						3
I5	I6		I3	I4		I1	I2	4
I5	I6		I3					5
	I6		I5			I3	I4	6
			I6					7
						I5	I6	8

Due to in-order issue, they must wait for I1 and I2 to fully execute before they can enter the execute stage. Because I1 takes two cycles to execute, they must wait in the decode stage until I1 and I2 *both* go to the write stage.

31. Why doesn't I5 get written during cycle 7 instead of cycle 8? (2 points)

Because of the in-order completion requirement, I6 must complete before both I5 and I6 can be written together. An out-of order completion would have allowed I5 to be written during cycle 7.

- 32. Why does I6 stay in the decode stage of the pipeline for one more cycle than I5? (2 points)I6 uses the same resource as I5, so it must wait until I5 is finished executing.
- 33. Which instructions could have been written earlier if this had been an "out-of-order completion" machine. (2 points)

I2, I4, and I5

SMP and Clusters

34. Which cache coherence protocol uses distributed control, the directory protocol or snoopy protocol? (2 points)

The snoopy protocol

- 35. Assume a multiprocessor system uses the MESI protocol. If the current state of a line in processor A's cache is exclusive and processor B loads the same line into its cache, what does the state of that line in processor A's cache change to? (2 points)
 - a.) modified b.) exclusive (c.) shared d.) invalid e.) cannot be determined
- 36. Assume a multiprocessor system uses the MESI protocol. If the current state of a line in processor A's cache is exclusive and processor A modifies that data, but *does not* write the new value to main memory, what does the state of that line in processor A's cache change to? (2 points)

a.) modified b.) exclusive c.) shared d.) invalid e.) cannot be determined

37. Assume a multiprocessor system uses the MESI protocol. If the current state of a line in processor A's cache is shared and processor B modifies that data without updating main memory, what does the state of that line in processor A's cache change to? (2 points)

a.) modified b.) exclusive c.) shared (d.) invalid e.) cannot be determined

38. Which SMP bus configuration has the most complex printed circuit board? (2 points)

a.) time-shared bus	b.) multiport memory	c.) central controller
39. The snoopy protocol is more symmetric multiprocessors	interconnection method for	
a.) time-shared bus	b.) multiport memory	c.) central controller

For the next 6 questions, identify whether the statement more closely identifies an SMP system or a cluster. (1 point each)

	SMP	Cluster	
40.		\checkmark	Processors are loosely coupled – communication usually at the file level
41.	$\overline{\mathbf{A}}$		Operation is closer to that of a single processor system
42.		\checkmark	The newer of the two architectures
43.		$\overline{\mathbf{A}}$	Easier to build, especially by end-user
44.	\checkmark		Takes less physical space
45.		$\overline{\mathbf{A}}$	Superior scalability