## Features

- Utilizes the AVR ${ }^{\circledR}$ RISC Architecture
- AVR - High-performance and Low-power RISC Architecture
- 118 Powerful Instructions - Most Single Clock Cycle Execution
- 32 x 8 General Purpose Working Registers
- Up to 10 MIPS Throughput at 10 MHz
- Data and Nonvolatile Program Memory
- 2K Bytes of In-System Programmable Flash

Endurance 1,000 Write/Erase Cycles

- 128 Bytes of SRAM
- 128 Bytes of In-System Programmable EEPROM Endurance: 100,000 Write/Erase Cycles
- Programming Lock for Flash Program and EEPROM Data Security
- Peripheral Features
- One 8-bit Timer/Counter with Separate Prescaler
- One 16-bit Timer/Counter with Separate Prescaler, Compare, Capture Modes and 8-, 9- or 10-bit PWM
- On-chip Analog Comparator
- Programmable Watchdog Timer with On-chip Oscillator
- SPI Serial Interface for In-System Programming
- Full Duplex UART
- . Special Microcontroller Features
- Low-power Idle and Power Down Modes
- External and Internal Interrupt Sources
- . Specifications
- Low-power, High-speed CMOS Process Technology
- Fully Static Operation
- Power Consumption at $4 \mathrm{MHz}, 3 \mathrm{~V}, 25^{\circ} \mathrm{C}$
- Active: 2.8 mA
- Idle Mode: 0.8 mA
- Power Down Mode: <1 $\mu \mathrm{A}$
- I/O and Packages
- 15 Programmable I/O Lines
- 20-pin PDIP and SOIC
- Operating Voltages
- 2.7-6.0V (AT90S2313-4)
- 4.0-6.0V (AT90S2313-10)
- Speed Grades
- 0-4 MHz (AT90S2313-4)
- 0-10 MHz (AT90S2313-10)


## Description

The AT90S2313 is a low-power CMOS 8-bit microcontroller based on the AVR RISC architecture. By executing powerful instructions in a single clock cycle, the
(continued)

## Pin Configuration

PDIP/SOIC


AT90S2313 achieves throughputs approaching 1 MIPS per MHz allowing the system designer to optimize power consumption versus processing speed.
The AVR core combines a rich instruction set with 32 general purpose working registers. All the 32 registers are directly connected to the Arithmetic Logic Unit (ALU), allowing two independent registers to be accessed in one single instruction executed in one clock cycle. The resulting architecture is more code efficient while achieving throughputs up to ten times faster than conventional CISC microcontrollers.

## Block Diagram

Figure 1. The AT90S2313 Block Diagram


The AT90S2313 provides the following features: 2K bytes of In-System Programmable Flash, 128 bytes EEPROM, 128 bytes SRAM, 15 general purpose I/O lines, 32 general purpose working registers, flexible timer/counters with compare modes, internal and external interrupts, a programmable serial UART, programmable Watchdog Timer with internal oscillator, an SPI serial port for Flash Memory downloading and two software selectable power saving modes. The Idle Mode stops the CPU while allowing the SRAM, timer/counters, SPI port and interrupt system to continue functioning. The power down mode saves the register contents but freezes the oscillator, disabling all other chip functions until the next external interrupt or hardware reset.
The device is manufactured using Atmel's high density nonvolatile memory technology. The on-chip In-System Programmable Flash allows the program memory to be reprogrammed in-system through an SPI serial interface or by a conventional nonvolatile memory programmer. By combining an enhanced RISC 8-bit CPU with In-System Programmable Flash on a monolithic chip, the Atmel AT90S2313 is a powerful microcontroller that provides a highly flexible and cost effective solution to many embedded control applications.
The AT90S2313 AVR is supported with a full suite of program and system development tools including: C compilers, macro assemblers, program debugger/simulators, in-circuit emulators, and evaluation kits.

## Pin Descriptions

## VCC

Supply voltage pin.

## GND

Ground pin.

## Port B (PB7..PB0)

Port B is an 8-bit bi-directional I/O port. Port pins can provide internal pull-up resistors (selected for each bit). PB0 and PB1 also serve as the positive input (AINO) and the negative input (AIN1), respectively, of the on-chip analog comparator. The Port B output buffers can sink 20 mA and can drive LED displays directly. When pins PB0 to PB7 are used as inputs and are externally pulled low, they will source current if the internal pull-up resistors are activated. The Port B pins are tri-stated when a reset condition becomes active, even if the clock is not active.

## Port D (PD6..PD0)

Port D has seven bi-directional I/O port with internal pull-up resistors, PD6..PD0. The Port D output buffers can sink 20 mA . As inputs, Port D pins that are externally pulled low will source current if the pull-up resistors are activated. The Port D pins are tri-stated when a reset condition becomes active, even if the clock is not active.

## RESET

Reset input. A low level on this pin for more than 50 ns will generate a reset, even if the clock is not running. Shorter pulses are not guaranteed to generate a reset.

## XTAL1

Input to the inverting oscillator amplifier and input to the internal clock operating circuit.

## XTAL2

Output from the inverting oscillator amplifier

## Architectural Overview

The fast-access register file concept contains $32 \times 8$-bit general purpose working registers with a single clock cycle access time. This means that during one single clock cycle, one ALU (Arithmetic Logic Unit) operation is executed. Two operands are output from the register file, the operation is executed, and the result is stored back in the register file - in one clock cycle.

Figure 2. The AT90S2313 AVR RISC Architecture
AVR AT90S2313 Architecture


Six of the 32 registers can be used as three 16 -bits indirect address register pointers for Data Space addressing - enabling efficient address calculations. One of the three address pointers is also used as the address pointer for the constant table look up function. These added function registers are the 16-bits X-register, Y-register and Z-register.
The ALU supports arithmetic and logic functions between registers or between a constant and a register. Single register operations are also executed in the ALU. Figure 2 shows the AT90S2313 AVR RISC microcontroller architecture.
In addition to the register operation, the conventional memory addressing modes can be used on the register file as well. This is enabled by the fact that the register file is assigned the 32 lowermost Data Space addresses ( $\$ 00-\$ 1 \mathrm{~F}$ ), allowing them to be accessed as though they were ordinary memory locations.

The I/O memory space contains 64 addresses for CPU peripheral functions as Control Registers, Timer/Counters, A/D-converters, and other I/O functions. The I/O memory can be accessed directly, or as the Data Space locations following those of the register file, \$20-\$5F.
The AVR has Harvard architecture - with separate memories and buses for program and data. The program memory is accessed with a two stage pipeline. While one instruction is being executed, the next instruction is pre-fetched from the program memory. This concept enables instructions to be executed in every clock cycle. The program memory is In -system Programmable Flash memory.
With the relative jump and call instructions, the whole 1 K address space is directly accessed. Most AVR instructions have a single 16 -bit word format. Every program memory address contains a 16 - or 32 -bit instruction.
During interrupts and subroutine calls, the return address program counter (PC) is stored on the stack. The stack is effectively allocated in the general data SRAM, and consequently the stack size is only limited by the total SRAM size and the usage of the SRAM. All user programs must initialize the SP in the reset routine (before subroutines or interrupts are executed). The 8 -bit stack pointer SP is read/write accessible in the I/O space.
The 128 bytes data SRAM + register file and I/O registers can be easily accessed through the five different addressing modes supported in the AVR architecture.
The memory spaces in the AVR architecture are all linear and regular memory maps.
Figure 3. Memory Mapss


A flexible interrupt module has its control registers in the I/O space with an additional global interrupt enable bit in the status register. All the different interrupts have a separate interrupt vector in the interrupt vector table at the beginning of the program memory. The different interrupts have priority in accordance with their interrupt vector position. The lower the interrupt vector address the higher the priority.

Register Summary


Notes: 1. For compatibility with future devices, reserved bits should be written to zero if accessed. Reserved I/O memory addresses should never be written.
2. Some of the status flags are cleared by writing a logical one to them. Note that the CBI and SBI instructions will operate on all bits in the I/O register, writing a one back into any flag read as set, thus clearing the flag. The CBI and SBI instructions work with registers $\$ 00$ to $\$ 1 F$ only.

Instruction Set Summary

| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ARITHMETIC AND LOGIC INSTRUCTIONS |  |  |  |  |  |
| ADD | Rd, Rr | Add two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}$ | Z,C,N,V,H | 1 |
| ADC | Rd , Rr | Add with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}+\mathrm{C}$ | Z,C,N,V,H | 1 |
| ADIW | Rdl, K | Add Immediate to Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl + K | Z,C,N,V,S | 2 |
| SUB | Rd , Rr | Subtract two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}$ | Z,C,N,V,H | 1 |
| SUBI | Rd, K | Subtract Constant from Register | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}$ | Z,C,N,V,H | 1 |
| SBIW | Rdl, K | Subtract Immediate from Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl - K | Z,C,N,V,S | 2 |
| SBC | Rd , Rr | Subtract with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z,C,N,V,H | 1 |
| SBCI | Rd, K | Subtract with Carry Constant from Reg. | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}-\mathrm{C}$ | Z,C,N,V,H | 1 |
| AND | Rd , Rr | Logical AND Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rr}$ | Z,N,V | 1 |
| ANDI | Rd, K | Logical AND Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{K}$ | Z,N,V | 1 |
| OR | Rd, Rr | Logical OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} v \mathrm{Rr}$ | Z,N,V | 1 |
| ORI | Rd, K | Logical OR Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rdv} \mathrm{K}$ | Z,N,V | 1 |
| EOR | Rd , Rr | Exclusive OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rr}$ | Z,N,V | 1 |
| COM | Rd | One's Complement | $\mathrm{Rd} \leftarrow$ \$FF - Rd | Z,C,N,V | 1 |
| NEG | Rd | Two's Complement | $\mathrm{Rd} \leftarrow \$ 00-\mathrm{Rd}$ | Z,C,N,V,H | 1 |
| SBR | Rd, K | Set Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \mathrm{v}$ K | Z,N,V | 1 |
| CBR | Rd,K | Clear Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet(\$ \mathrm{FF}-\mathrm{K})$ | Z,N,V | 1 |
| INC | Rd | Increment | $\mathrm{Rd} \leftarrow \mathrm{Rd}+1$ | Z,N,V | 1 |
| DEC | Rd | Decrement | $\mathrm{Rd} \leftarrow \mathrm{Rd}-1$ | Z,N,V | 1 |
| TST | Rd | Test for Zero or Minus | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rd}$ | Z,N,V | 1 |
| CLR | Rd | Clear Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rd}$ | Z,N,V | 1 |
| SER | Rd | Set Register | $\mathrm{Rd} \leftarrow$ \$FF | None | 1 |
| BRANCH INSTRUCTIONS |  |  |  |  |  |
| RJMP | k | Relative Jump | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 2 |
| IJMP |  | Indirect Jump to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 2 |
| RCALL | k | Relative Subroutine Call | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 3 |
| ICALL |  | Indirect Call to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 3 |
| RET |  | Subroutine Return | $\mathrm{PC} \leftarrow$ STACK | None | 4 |
| RETI |  | Interrupt Return | $\mathrm{PC} \leftarrow$ STACK | I | 4 |
| CPSE | Rd,Rr | Compare, Skip if Equal | if ( $\mathrm{Rd}=\mathrm{Rr}$ ) $\mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | $1 / 2$ |
| CP | Rd, Rr | Compare | Rd - Rr | Z, N,V,C,H | 1 |
| CPC | Rd, Rr | Compare with Carry | $\mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z, N,V,C,H | 1 |
| CPI | Rd, K | Compare Register with Immediate | Rd-K | Z, N,V,C,H | 1 |
| SBRC | Rr, b | Skip if Bit in Register Cleared | if $(\operatorname{Rr}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| SBRS | Rr, b | Skip if Bit in Register is Set | if $(\operatorname{Rr}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| SBIC | P, b | Skip if Bit in I/O Register Cleared | if $(\mathrm{P}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| SBIS | P, b | Skip if Bit in I/O Register is Set | if ( $\mathrm{R}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| BRBS | s, k | Branch if Status Flag Set | if (SREG(s) $=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRBC | s, k | Branch if Status Flag Cleared | if (SREG(s) $=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BREQ | k | Branch if Equal | if $(\mathrm{Z}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRNE | k | Branch if Not Equal | if $(\mathrm{Z}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRCS | k | Branch if Carry Set | if ( $\mathrm{C}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRCC | k | Branch if Carry Cleared | if ( $\mathrm{C}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRSH | k | Branch if Same or Higher | if ( $\mathrm{C}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRLO | k | Branch if Lower | if ( $\mathrm{C}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRMI | k | Branch if Minus | if $(\mathrm{N}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRPL | k | Branch if Plus | if $(\mathrm{N}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRGE | k | Branch if Greater or Equal, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRLT | k | Branch if Less Than Zero, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHS | k | Branch if Half Carry Flag Set | if ( $\mathrm{H}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHC | k | Branch if Half Carry Flag Cleared | if ( $\mathrm{H}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRTS | k | Branch if T Flag Set | if ( $\mathrm{T}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRTC | k | Branch if T Flag Cleared | if ( $\mathrm{T}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRVS | k | Branch if Overflow Flag is Set | if ( $\mathrm{V}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRVC | k | Branch if Overflow Flag is Cleared | if $(\mathrm{V}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRIE | k | Branch if Interrupt Enabled | if $(\mathrm{I}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRID | k | Branch if Interrupt Disabled | if $(\mathrm{I}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |

## Instruction Set Summary (Continued)

| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| DATA TRANSFER INSTRUCTIONS |  |  |  |  |  |
| MOV | Rd, Rr | Move Between Registers | $\mathrm{Rd} \leftarrow \mathrm{Rr}$ | None | 1 |
| LDI | Rd, K | Load Immediate | $\mathrm{Rd} \leftarrow \mathrm{K}$ | None | 1 |
| LD | Rd, X | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, $\mathrm{X}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{X}), \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| LD | Rd, - X | Load Indirect and Pre-Dec. | $X \leftarrow X-1, \mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, Y | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LD | Rd, $\mathrm{Y}+$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Y}), \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| LD | Rd, - Y | Load Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1, \mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LDD | Rd, $\mathrm{Y}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Y}+\mathrm{q})$ | None | 2 |
| LD | Rd, Z | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LD | Rd, $\mathrm{Z}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Z}), \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| LD | Rd, -Z | Load Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1, \mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LDD | Rd, $\mathrm{Z}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Z}+\mathrm{q})$ | None | 2 |
| LDS | Rd, k | Load Direct from SRAM | $\mathrm{Rd} \leftarrow \mathrm{L}$ ( $)$ | None | 2 |
| ST | X, Rr | Store Indirect | $(\mathrm{X}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{X}+$, Rr | Store Indirect and Post-Inc. | $(\mathrm{X}) \leftarrow \mathrm{Rr}, \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| ST | - X, Rr | Store Indirect and Pre-Dec. | $X \leftarrow X-1,(X) \leftarrow \operatorname{Rr}$ | None | 2 |
| ST | Y, Rr | Store Indirect | $(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{Y}+$, Rr | Store Indirect and Post-Inc. | $(\mathrm{Y}) \leftarrow \mathrm{Rr}, \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| ST | - Y, Rr | Store Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1,(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | $\mathrm{Y}+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(Y+q) \leftarrow R \mathrm{Rr}$ | None | 2 |
| ST | Z, Rr | Store Indirect | $(Z) \leftarrow \operatorname{Rr}$ | None | 2 |
| ST | $\mathrm{Z}_{+}, \mathrm{Rr}$ | Store Indirect and Post-Inc. | $(Z) \leftarrow \operatorname{Rr}, \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| ST | -Z, Rr | Store Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1,(\mathrm{Z}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | $\mathrm{Z}+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(Z+q) \leftarrow \operatorname{Rr}$ | None | 2 |
| STS | k, Rr | Store Direct to SRAM | $(\mathrm{k}) \leftarrow \mathrm{Rr}$ | None | 2 |
| LPM |  | Load Program Memory | $\mathrm{R} 0 \leftarrow(\mathrm{Z})$ | None | 3 |
| IN | Rd, P | In Port | $\mathrm{Rd} \leftarrow \mathrm{P}$ | None | 1 |
| OUT | $\mathrm{P}, \mathrm{Rr}$ | Out Port | $\mathrm{P} \leftarrow \mathrm{Rr}$ | None | 1 |
| PUSH | Rr | Push Register on Stack | STACK $\leftarrow \mathrm{Rr}$ | None | 2 |
| POP | Rd | Pop Register from Stack | $\mathrm{Rd} \leftarrow$ STACK | None | 2 |
| BIT AND BIT-TEST INSTRUCTIONS |  |  |  |  |  |
| SBI | P, b | Set Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 1$ | None | 2 |
| CBI | P, b | Clear Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 0$ | None | 2 |
| LSL | Rd | Logical Shift Left | $\mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{Rd}(0) \leftarrow 0$ | Z,C,N,V | 1 |
| LSR | Rd | Logical Shift Right | $\mathrm{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \operatorname{Rd}(7) \leftarrow 0$ | Z,C,N,V | 1 |
| ROL | Rd | Rotate Left Through Carry | $\mathrm{Rd}(0) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{C} \leftarrow \operatorname{Rd}(7)$ | Z,C,N,V | 1 |
| ROR | Rd | Rotate Right Through Carry | $\operatorname{Rd}(7) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \mathrm{C} \leftarrow \operatorname{Rd}(0)$ | Z,C,N,V | 1 |
| ASR | Rd | Arithmetic Shift Right | $\mathrm{Rd}(\mathrm{n}) \leftarrow \mathrm{Rd}(\mathrm{n}+1), \mathrm{n}=0 . .6$ | Z,C,N,V | 1 |
| SWAP | Rd | Swap Nibbles | $\operatorname{Rd}(3.0) \leftarrow \operatorname{Rd}(7 . .4), \operatorname{Rd}(7 . .4) \leftarrow \operatorname{Rd}(3 . .0)$ | None | 1 |
| BSET | s | Flag Set | SREG(s) $\leftarrow 1$ | SREG(s) | 1 |
| BCLR | s | Flag Clear | SREG(s) $\leftarrow 0$ | SREG(s) | 1 |
| BST | $\mathrm{Rr}, \mathrm{b}$ | Bit Store from Register to T | $\mathrm{T} \leftarrow \operatorname{Rr}(\mathrm{b})$ | T | 1 |
| BLD | Rd, b | Bit load from T to Register | $\mathrm{Rd}(\mathrm{b}) \leftarrow \mathrm{T}$ | None | 1 |
| SEC |  | Set Carry | $\mathrm{C} \leftarrow 1$ | C | 1 |
| CLC |  | Clear Carry | $\mathrm{C} \leftarrow 0$ | C | 1 |
| SEN |  | Set Negative Flag | $\mathrm{N} \leftarrow 1$ | N | 1 |
| CLN |  | Clear Negative Flag | $\mathrm{N} \leftarrow 0$ | N | 1 |
| SEZ |  | Set Zero Flag | $\mathrm{Z} \leftarrow 1$ | Z | 1 |
| CLZ |  | Clear Zero Flag | $\mathrm{Z} \leftarrow 0$ | Z | 1 |
| SEI |  | Global Interrupt Enable | $1 \leftarrow 1$ | 1 | 1 |
| CLI |  | Global Interrupt Disable | $1 \leftarrow 0$ | , | 1 |
| SES |  | Set Signed Test Flag | $\mathrm{S} \leftarrow 1$ | S | 1 |
| CLS |  | Clear Signed Test Flag | $\mathrm{S} \leftarrow 0$ | S | 1 |
| SEV |  | Set Twos Complement Overflow | $\mathrm{V} \leftarrow 1$ | V | 1 |
| CLV |  | Clear Twos Complement Overflow | $\mathrm{V} \leftarrow 0$ | V | 1 |
| SET |  | Set T in SREG | $\mathrm{T} \leftarrow 1$ | T | 1 |
| CLT |  | Clear T in SREG | $\mathrm{T} \leftarrow 0$ | T | 1 |
| SEH |  | Set Half Carry Flag in SREG | $\mathrm{H} \leftarrow 1$ | H | 1 |
| CLH |  | Clear Half Carry Flag in SREG | $\mathrm{H} \leftarrow 0$ | H | 1 |
| NOP |  | No Operation |  | None | 1 |
| SLEEP |  | Sleep | (see specific descr. for Sleep function) | None | 3 |
| WDR |  | Watchdog Reset | (see specific descr. for WDR/timer) | None | 1 |

## Ordering Information

| Speed (MHz) | Power Supply | Ordering Code | Package | Operation Range |
| :---: | :---: | :--- | :--- | :---: |
| 4 | $2.7-6.0 \mathrm{~V}$ | AT90S2313-4PC | 20 P 3 | Commercial |
|  |  | AT90S2313-4SC | 20 S | $\left(0^{\circ} \mathrm{C}\right.$ to $\left.70^{\circ} \mathrm{C}\right)$ |
|  |  | AT90S2313-4PI | 20 P 3 | Industrial |
|  |  | AT90S2313-4SI | 20 S | $\left(-40^{\circ} \mathrm{C}\right.$ to $\left.85^{\circ} \mathrm{C}\right)$ |
| 10 |  | AT90S2313-10PC | 20 P 3 | Commercial |
|  |  |  | AT90S2313-10SC | 20 S |
|  |  |  |  |  |
|  |  | AT90S2313-10PI | 20 P 3 | Industrial |
|  |  | AT90S2313-10SI | 20 S | $\left(-40^{\circ} \mathrm{C}\right.$ to $\left.85^{\circ} \mathrm{C}\right)$ |


| Package Type |  |
| :--- | :--- |
| 20P3 | 20-lead, 0.300" Wide, Plastic Dual In-Line Package (PDIP) |
| $20 S$ | 20-lead, 0.300" Wide, Plastic Gull-Wing Small Outline (SOIC) |

## Packaging Information



## Atmel Headquarters

Corporate Headquarters 2325 Orchard Parkway San Jose, CA 95131 TEL (408) 441-0311 FAX (408) 487-2600

## Europe

Atmel U.K., Ltd.
Coliseum Business Centre Riverside Way Camberley, Surrey GU15 3YL England
TEL (44) 1276-686-677
FAX (44) 1276-686-697

## Asia

Atmel Asia, Ltd.
Room 1219
Chinachem Golden Plaza
77 Mody Road Tsimhatsui
East Kowloon
Hong Kong
TEL (852) 2721-9778
FAX (852) 2722-1369

## Japan

Atmel Japan K.K.
9F, Tonetsu Shinkawa BIdg. 1-24-8 Shinkawa
Chuo-ku, Tokyo 104-0033
Japan
TEL (81) 3-3523-3551
FAX (81) 3-3523-7581

## Atmel Operations

Atmel Colorado Springs

1150 E. Cheyenne Mtn. Blvd.
Colorado Springs, CO 80906
TEL (719) 576-3300
FAX (719) 540-1759
Atmel Rousset
Zone Industrielle
13106 Rousset Cedex
France
TEL (33) 4-4253-6000
FAX (33) 4-4253-6001

Fax-on-Demand<br>North America:<br>1-(800) 292-8635<br>International:<br>1-(408) 441-0732<br>e-mail<br>literature@atmel.com<br>Web Site<br>http://www.atmel.com<br>\section*{BBS}<br>1-(408) 436-4309

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