## Features

- Utilizes the AVR ${ }^{\circledR}$ RISC Architecture
- AVR - High-performance and Low-power RISC Architecture
- 118 Powerful Instructions - Most Single Clock Cycle Execution
- 32 x 8 General Purpose Working Registers
- Up to 1.5 MIPS Throughput at 1.5 MHz
- Data and Nonvolatile Program Memory
- 8K Bytes Flash Program Memory

Endurance: 1,000 Write/Erase Cycles

- 256 Bytes Internal SRAM
- 512 Bytes EEPROM

Endurance: 100,000 Write/Erase Cycles

- Programming Lock for Flash Program and EEPROM Data Security
- Peripheral Features
- One 8-bit Timer/Counter with Separate Prescaler
- One 16-bit Timer/Counter with Separate Prescaler
- Special Microcontroller Features
- Low-power Idle and Power Down Modes
- External and Internal Interrupt Sources
- 6-channel, 10-bit ADC
- Specifications
- Low-power, High-speed CMOS Process Technology
- Fully Static Operation
- Power Consumption at $1.5 \mathrm{MHz}, 3.6 \mathrm{~V}, 25^{\circ} \mathrm{C}$
- Active: 1.2 mA
- Idle Mode: 0.2 mA
- Power Down Mode: <10 $\mu \mathrm{A}$
- I/O and Packages
- Seven General Output Lines
- Two External Interrupt Lines
- 48-pin LQFP/VQFP Package
- Operating Voltage
- 3.3-6.0V
- Speed Grade
- 0-1.5 MHz


## Description

The AT90C8534 is a low-power CMOS 8-bit microcontroller based on the AVR RISC architecture. By executing powerful instructions in a single clock cycle, the

## Pin Configuration

(continued)


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Note: This is a summary document. For the complete 56 page document, please visit our web site at www.atmel.com or e-mail at literature@atmel.com and request literature \#1229A.

AT90C8534 achieves throughputs approaching 1 MIPS per MHz allowing the system designer to optimize power consumption versus processing speed.

## Block Diagram

Figure 1. The AT90C8534 Block Diagram


The AVR core combines a rich instruction set with 32 general purpose working registers. All the 32 registers are directly connected to the Arithmetic Logic Unit (ALU), allowing two independent registers to be accessed in one single instruction executed in one clock cycle. The resulting architecture is more code efficient while achieving throughputs up to ten times faster than conventional CISC microcontrollers.
The AT90C8534 provides the following features: 8K bytes of Programmable Flash, 512 bytes EEPROM, 256 bytes SRAM, 7 general output lines, 2 external interrupt lines, 32 general purpose working registers, two flexible timer/counters, internal and external interrupts, 6 -channel, 10-bit ADC, and two software selectable power saving modes. The Idle mode stops the CPU while allowing the ADC, timer/counters and interrupt system to continue functioning. The Power Down mode saves the SRAM and register contents but freezes the oscillator, disabling all other chip functions until the next interrupt or hardware reset.
The device is manufactured using Atmel's high density nonvolatile memory technology. The on-chip Programmable Flash allows the program memory to be reprogrammed by a conventional nonvolatile memory programmer. By combining an 8 -bit RISC CPU with Programmable Flash on a monolithic chip, the Atmel AT90C8534 is a powerful microcontroller that provides a highly flexible and cost effective solution to many embedded control applications.
The AT90C8534 AVR is supported with a full suite of program and system development tools including: C compilers, macro assemblers, program debugger/simulators, in-circuit emulators, and evaluation kits.

## Pin Descriptions

## VCC

Digital supply voltage

## GND

Digital ground

## Port A (PA6..PAO)

Port A is a 7 -bit output port with tristate mode. The Port A output buffers can sink 20 mA and can drive LED displays directly. The port pins are tristated when a reset condition becomes active, even if the clock is not running.

## INT1,0

External interrupt input pins. A falling or rising edge on either of these pins will generate an interrupt request. Interrupt pulses longer than 40 ns will generate an interrupt, even if the clock is not running.

ADIN5.. 0
ADC input pins. Any of these pins can be selected as the input to the ADC.

## RESET

Reset input. An external reset is generated by a low level on the $\overline{\operatorname{RESET}}$ pin. Reset pulses longer than 100 ns will generate a reset, even if the clock is not running. Shorter pulses are not guaranteed to generate a reset.

## XTAL1

Input to the inverting oscillator amplifier and input to the internal clock operating circuit.

## XTAL2

Output from the inverting oscillator amplifier

## AVCC

This is the supply voltage pin for the A/D Converter. It should be externally connected to VCC via a low-pass filter.

## AGND

Analog ground. If the board has a separate analog ground plane, this pin should be connected to this ground plane. Otherwise, connect to GND.

## Architectural Overview

The fast-access register file concept contains $32 \times 8$-bit general purpose working registers with a single clock cycle access time. This means that during one single clock cycle, one ALU (Arithmetic Logic Unit) operation is executed. Two operands are output from the register file, the operation is executed, and the result is stored back in the register file - in one clock cycle.
Six of the 32 registers can be used as three 16-bits indirect address register pointers for Data Space addressing - enabling efficient address calculations. One of the three address pointers is also used as the address pointer for the constant table look up function. These added function registers are the 16 -bits X-register, Y-register and Z-register.
The ALU supports arithmetic and logic functions between registers or between a constant and a register. Single register operations are also executed in the ALU. Figure 2 shows the AT90C8534 AVR RISC microcontroller architecture.
In addition to the register operation, the conventional memory addressing modes can be used on the register file as well. This is enabled by the fact that the register file is assigned the 32 lowermost Data Space addresses ( $\$ 00-\$ 1 F$ ), allowing them to be accessed as though they were ordinary memory locations.
The I/O memory space contains 64 addresses for CPU peripheral functions as Control Registers, Timer/Counters, A/Dconverters, and other I/O functions. The I/O Memory can be accessed directly, or as the Data Space locations following those of the register file, \$20-\$5F.
The AVR uses a Harvard architecture concept - with separate memories and buses for program and data. The program memory is executed with a single level pipelining. While one instruction is being executed, the next instruction is prefetched from the program memory. This concept enables instructions to be executed in every clock cycle. The program memory is programmable Flash memory.
With the relative jump and call instructions, the whole 4 K word ( 8 K bytes) address space is directly accessed. Most AVR instructions have a single 16-bit word format. Every program memory address contains a 16- or 32-bit instruction.
During interrupts and subroutine calls, the return address program counter (PC) is stored on the stack. The stack is effectively allocated in the general data SRAM, and consequently the stack size is only limited by the total SRAM size and the usage of the SRAM. All user programs must initialize the stack pointer (SP) in the reset routine (before subroutines or interrupts are executed). The 9-bit stack pointer is read/write accessible in the I/O space.
The 256 bytes data SRAM can be easily accessed through the five different addressing modes supported in the AVR architecture.
The memory spaces in the AVR architecture are all linear and regular memory maps.

Figure 2. The AT90C8534 AVR RISC Architecture

## AVR AT90C8534 Architecture



Figure 3. Memory Maps


A flexible interrupt module has its control registers in the I/O space with an additional global interrupt enable bit in the status register. All the different interrupts have a separate interrupt vector in the interrupt vector table at the beginning of the program memory. The different interrupts have priority in accordance with their interrupt vector position. The lower the interrupt vector address, the higher the priority.

## AT90C8534 Register Summary



Note: For compatibility with future devices, reserved bits should be written to zero if accessed. Reserved I/O memory addresses should never be written.

Instruction Set Summary

| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| ARITHMETIC AND LOGIC INSTRUCTIONS |  |  |  |  |  |
| ADD | Rd, Rr | Add two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}$ | Z,C,N,V,H | 1 |
| ADC | Rd, Rr | Add with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}+\mathrm{Rr}+\mathrm{C}$ | Z,C,N,V,H | 1 |
| ADIW | Rdl, K | Add Immediate to Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl + K | Z,C,N,V,S | 2 |
| SUB | Rd, Rr | Subtract two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}$ | Z,C,N,V,H | 1 |
| SUBI | Rd, K | Subtract Constant from Register | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}$ | Z,C,N,V,H | 1 |
| SBC | Rd, Rr | Subtract with Carry two Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z,C,N,V,H | 1 |
| SBCI | Rd, K | Subtract with Carry Constant from Reg. | $\mathrm{Rd} \leftarrow \mathrm{Rd}-\mathrm{K}-\mathrm{C}$ | Z,C,N,V,H | 1 |
| SBIW | Rdl, K | Subtract Immediate from Word | Rdh:Rdl $\leftarrow$ Rdh:Rdl - K | Z,C,N,V,S | 2 |
| AND | Rd, Rr | Logical AND Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rr}$ | Z,N,V | 1 |
| ANDI | Rd, K | Logical AND Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{K}$ | Z,N,V | 1 |
| OR | Rd, Rr | Logical OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} v \mathrm{Rr}$ | Z,N,V | 1 |
| ORI | Rd, K | Logical OR Register and Constant | $\mathrm{Rd} \leftarrow \mathrm{Rd} v \mathrm{~K}$ | Z,N,V | 1 |
| EOR | Rd, Rr | Exclusive OR Registers | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rr}$ | Z,N,V | 1 |
| COM | Rd | One's Complement | $\mathrm{Rd} \leftarrow$ \$ FF - Rd | Z,C,N,V | 1 |
| NEG | Rd | Two's Complement | $\mathrm{Rd} \leftarrow \$ 00-\mathrm{Rd}$ | Z,C,N,V,H | 1 |
| SBR | Rd,K | Set Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} v \mathrm{~K}$ | Z,N,V | 1 |
| CBR | Rd,K | Clear Bit(s) in Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet(\$ \mathrm{FF}-\mathrm{K})$ | Z,N,V | 1 |
| INC | Rd | Increment | $\mathrm{Rd} \leftarrow \mathrm{Rd}+1$ | Z,N,V | 1 |
| DEC | Rd | Decrement | $\mathrm{Rd} \leftarrow \mathrm{Rd}-1$ | Z,N,V | 1 |
| TST | Rd | Test for Zero or Minus | $\mathrm{Rd} \leftarrow \mathrm{Rd} \bullet \mathrm{Rd}$ | Z,N,V | 1 |
| CLR | Rd | Clear Register | $\mathrm{Rd} \leftarrow \mathrm{Rd} \oplus \mathrm{Rd}$ | Z,N,V | 1 |
| SER | Rd | Set Register | $\mathrm{Rd} \leftarrow$ \$ FF | None | 1 |
| BRANCH INSTRUCTIONS |  |  |  |  |  |
| RJMP | k | Relative Jump | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 2 |
| IJMP |  | Indirect Jump to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 2 |
| RCALL | k | Relative Subroutine Call | $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 3 |
| ICALL |  | Indirect Call to (Z) | $\mathrm{PC} \leftarrow \mathrm{Z}$ | None | 3 |
| RET |  | Subroutine Return | $\mathrm{PC} \leftarrow$ STACK | None | 4 |
| RETI |  | Interrupt Return | $\mathrm{PC} \leftarrow$ STACK | I | 4 |
| CPSE | Rd,Rr | Compare, Skip if Equal | if ( $\mathrm{Rd}=\mathrm{Rr}$ ) $\mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| CP | Rd,Rr | Compare | Rd - Rr | Z, N,V,C,H | 1 |
| CPC | Rd,Rr | Compare with Carry | $\mathrm{Rd}-\mathrm{Rr}-\mathrm{C}$ | Z, N,V,C,H | 1 |
| CPI | Rd, K | Compare Register with Immediate | Rd-K | Z, N,V,C,H | 1 |
| SBRC | Rr, b | Skip if Bit in Register Cleared | if $(\operatorname{Rr}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| SBRS | Rr, b | Skip if Bit in Register is Set | if $(\operatorname{Rr}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| SBIC | P, b | Skip if Bit in I/O Register Cleared | if $(\mathrm{P}(\mathrm{b})=0) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| SBIS | P, b | Skip if Bit in I/O Register is Set | if $(\mathrm{P}(\mathrm{b})=1) \mathrm{PC} \leftarrow \mathrm{PC}+2$ or 3 | None | 1/2 |
| BRBS | s, k | Branch if Status Flag Set | if (SREG(s) $=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRBC | s, k | Branch if Status Flag Cleared | if (SREG(s) $=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BREQ | k | Branch if Equal | if $(Z=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRNE | k | Branch if Not Equal | if $(Z=0)$ then $P C \leftarrow P C+k+1$ | None | $1 / 2$ |
| BRCS | k | Branch if Carry Set | if $(C=1)$ then $P C \leftarrow P C+k+1$ | None | 1/2 |
| BRCC | k | Branch if Carry Cleared | if $(C=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRSH | k | Branch if Same or Higher | if $(C=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRLO | k | Branch if Lower | if $(C=1)$ then $P C \leftarrow P C+k+1$ | None | 1/2 |
| BRMI | k | Branch if Minus | if $(\mathrm{N}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRPL | k | Branch if Plus | if $(\mathrm{N}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRGE | k | Branch if Greater or Equal, Signed | if $(\mathrm{N} \oplus \mathrm{V}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRLT | k | Branch if Less Than Zero, Signed | if ( $\mathrm{N} \oplus \mathrm{V}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHS | k | Branch if Half Carry Flag Set | if ( $\mathrm{H}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRHC | k | Branch if Half Carry Flag Cleared | if ( $\mathrm{H}=0$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRTS | k | Branch if T Flag Set | if ( $\mathrm{T}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRTC | k | Branch if T Flag Cleared | if $(\mathrm{T}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRVS | k | Branch if Overflow Flag is Set | if $(\mathrm{V}=1)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | $1 / 2$ |
| BRVC | k | Branch if Overflow Flag is Cleared | if $(\mathrm{V}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRIE | k | Branch if Interrupt Enabled | if ( $\mathrm{I}=1$ ) then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |
| BRID | k | Branch if Interrupt Disabled | if $(\mathrm{I}=0)$ then $\mathrm{PC} \leftarrow \mathrm{PC}+\mathrm{k}+1$ | None | 1/2 |

## Instruction Set Summary (Continued)

| Mnemonics | Operands | Description | Operation | Flags | \#Clocks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| DATA TRANSFER INSTRUCTIONS |  |  |  |  |  |
| MOV | Rd, Rr | Move Between Registers | $\mathrm{Rd} \leftarrow \mathrm{Rr}$ | None | 1 |
| LDI | Rd, K | Load Immediate | $\mathrm{Rd} \leftarrow \mathrm{K}$ | None | 1 |
| LD | Rd, X | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, $\mathrm{X}^{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{X}), \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| LD | Rd, - X | Load Indirect and Pre-Dec. | $\mathrm{X} \leftarrow \mathrm{X}-1, \mathrm{Rd} \leftarrow(\mathrm{X})$ | None | 2 |
| LD | Rd, Y | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LD | Rd, Y+ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Y}), \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| LD | Rd, - Y | Load Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1, \mathrm{Rd} \leftarrow(\mathrm{Y})$ | None | 2 |
| LDD | Rd, $\mathrm{Y}+\mathrm{q}$ | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Y}+\mathrm{q})$ | None | 2 |
| LD | Rd, Z | Load Indirect | $\mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LD | Rd, $\mathrm{Z}_{+}$ | Load Indirect and Post-Inc. | $\mathrm{Rd} \leftarrow(\mathrm{Z}), \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| LD | Rd, -Z | Load Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1, \mathrm{Rd} \leftarrow(\mathrm{Z})$ | None | 2 |
| LDD | Rd, Z+q | Load Indirect with Displacement | $\mathrm{Rd} \leftarrow(\mathrm{Z}+\mathrm{q})$ | None | 2 |
| LDS | Rd, k | Load Direct from SRAM | $\mathrm{Rd} \leftarrow(\mathrm{k})$ | None | 2 |
| ST | X, Rr | Store Indirect | $(\mathrm{X}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{X}+$, Rr | Store Indirect and Post-Inc. | $(\mathrm{X}) \leftarrow \mathrm{Rr}, \mathrm{X} \leftarrow \mathrm{X}+1$ | None | 2 |
| ST | - X, Rr | Store Indirect and Pre-Dec. | $\mathrm{X} \leftarrow \mathrm{X}-1,(\mathrm{X}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | Y, Rr | Store Indirect | $(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| ST | $\mathrm{Y}+$, Rr | Store Indirect and Post-Inc. | $(\mathrm{Y}) \leftarrow \mathrm{Rr}, \mathrm{Y} \leftarrow \mathrm{Y}+1$ | None | 2 |
| ST | - Y, Rr | Store Indirect and Pre-Dec. | $\mathrm{Y} \leftarrow \mathrm{Y}-1,(\mathrm{Y}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | $\mathrm{Y}+\mathrm{q}, \mathrm{Rr}$ | Store Indirect with Displacement | $(Y+q) \leftarrow R r$ | None | 2 |
| ST | Z, Rr | Store Indirect | $(Z) \leftarrow R r$ | None | 2 |
| ST | Z+, Rr | Store Indirect and Post-Inc. | $(\mathrm{Z}) \leftarrow \mathrm{Rr}, \mathrm{Z} \leftarrow \mathrm{Z}+1$ | None | 2 |
| ST | -Z, Rr | Store Indirect and Pre-Dec. | $\mathrm{Z} \leftarrow \mathrm{Z}-1,(\mathrm{Z}) \leftarrow \mathrm{Rr}$ | None | 2 |
| STD | Z+q,Rr | Store Indirect with Displacement | $(Z+q) \leftarrow \operatorname{Rr}$ | None | 2 |
| STS | k, Rr | Store Direct to SRAM | $(\mathrm{k}) \leftarrow \mathrm{Rr}$ | None | 2 |
| LPM |  | Load Program Memory | $\mathrm{R} 0 \leftarrow(\mathrm{Z})$ | None | 3 |
| IN | Rd, P | In Port | $\mathrm{Rd} \leftarrow \mathrm{P}$ | None | 1 |
| OUT | P, Rr | Out Port | $\mathrm{P} \leftarrow \mathrm{Rr}$ | None | 1 |
| PUSH | Rr | Push Register on Stack | STACK $\leftarrow \mathrm{Rr}$ | None | 2 |
| POP | Rd | Pop Register from Stack | $\mathrm{Rd} \leftarrow$ STACK | None | 2 |
| BIT AND BIT-TEST INSTRUCTIONS |  |  |  |  |  |
| SBI | P,b | Set Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 1$ | None | 2 |
| CBI | P,b | Clear Bit in I/O Register | $\mathrm{l} / \mathrm{O}(\mathrm{P}, \mathrm{b}) \leftarrow 0$ | None | 2 |
| LSL | Rd | Logical Shift Left | $\mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{Rd}(0) \leftarrow 0$ | Z,C,N,V | 1 |
| LSR | Rd | Logical Shift Right | $\operatorname{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \operatorname{Rd}(7) \leftarrow 0$ | Z,C,N,V | 1 |
| ROL | Rd | Rotate Left Through Carry | $\mathrm{Rd}(0) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}+1) \leftarrow \operatorname{Rd}(\mathrm{n}), \mathrm{C} \leftarrow \mathrm{Rd}(7)$ | Z,C,N,V | 1 |
| ROR | Rd | Rotate Right Through Carry | $\mathrm{Rd}(7) \leftarrow \mathrm{C}, \mathrm{Rd}(\mathrm{n}) \leftarrow \mathrm{Rd}(\mathrm{n}+1), \mathrm{C} \leftarrow \operatorname{Rd}(0)$ | Z,C,N,V | 1 |
| ASR | Rd | Arithmetic Shift Right | $\operatorname{Rd}(\mathrm{n}) \leftarrow \operatorname{Rd}(\mathrm{n}+1), \mathrm{n}=0 . .6$ | Z,C,N,V | 1 |
| SWAP | Rd | Swap Nibbles | $\operatorname{Rd}(3 . .0) \leftarrow \operatorname{Rd}(7 . .4), \operatorname{Rd}(7 . .4) \leftarrow \operatorname{Rd}(3 . .0)$ | None | 1 |
| BSET | s | Flag Set | SREG(s) $\leftarrow 1$ | SREG(s) | 1 |
| BCLR | s | Flag Clear | SREG(s) $\leftarrow 0$ | SREG(s) | 1 |
| BST | $\mathrm{Rr}, \mathrm{b}$ | Bit Store from Register to T | $\mathrm{T} \leftarrow \operatorname{Rr}(\mathrm{b})$ | T | 1 |
| BLD | Rd, b | Bit load from T to Register | $\mathrm{Rd}(\mathrm{b}) \leftarrow \mathrm{T}$ | None | 1 |
| SEC |  | Set Carry | $C \leftarrow 1$ | C | 1 |
| CLC |  | Clear Carry | $\mathrm{C} \leftarrow 0$ | C | 1 |
| SEN |  | Set Negative Flag | $\mathrm{N} \leftarrow 1$ | N | 1 |
| CLN |  | Clear Negative Flag | $\mathrm{N} \leftarrow 0$ | N | 1 |
| SEZ |  | Set Zero Flag | $\mathrm{Z} \leftarrow 1$ | Z | 1 |
| CLZ |  | Clear Zero Flag | $\mathrm{Z} \leftarrow 0$ | Z | 1 |
| SEI |  | Global Interrupt Enable | $1 \leftarrow 1$ | I | 1 |
| CLI |  | Global Interrupt Disable | $1 \leftarrow 0$ | 1 | 1 |
| SES |  | Set Signed Test Flag | $\mathrm{S} \leftarrow 1$ | S | 1 |
| CLS |  | Clear Signed Test Flag | $\mathrm{S} \leftarrow 0$ | S | 1 |
| SEV |  | Set Twos Complement Overflow. | $\mathrm{V} \leftarrow 1$ | V | 1 |
| CLV |  | Clear Twos Complement Overflow | $\mathrm{V} \leftarrow 0$ | V | 1 |
| SET |  | Set T in SREG | $\mathrm{T} \leftarrow 1$ | T | 1 |
| CLT |  | Clear T in SREG | $\mathrm{T} \leftarrow 0$ | T | 1 |
| SEH |  | Set Half Carry Flag in SREG | $\mathrm{H} \leftarrow 1$ | H | 1 |
| CLH |  | Clear Half Carry Flag in SREG | $\mathrm{H} \leftarrow 0$ | H | 1 |
| NOP |  | No Operation |  | None | 1 |
| SLEEP |  | Sleep | (see specific description. for Sleep func- | None | 3 |
| WDR |  | Watchdog Reset | this command has no effect | None | 1 |

Ordering Information

| Speed (MHz) | Power Supply | Ordering Code | Package | Operation Range |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $3.3-6.0 \mathrm{~V}$ | AT90VC8534-1AC | 48 A | Commercial |
|  |  |  | $\left(0^{\circ} \mathrm{C}\right.$ to $\left.70^{\circ} \mathrm{C}\right)$ |  |

## Packaging Information



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