



Agenda

Wide Area Network (WAN)

- **Environment**
- **Requirements**
- **Technologies**
- **Interface Signaling Protocols**

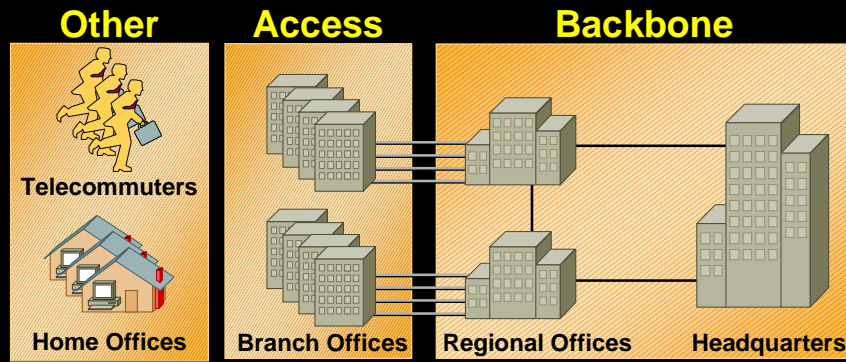
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Take-Away Message

- **By selecting the right WAN technology and hardware, you can:**
 - Build a uniform multiservice network for DVV consolidation**
 - Reduce recurring network operation costs (such as bandwidth)**

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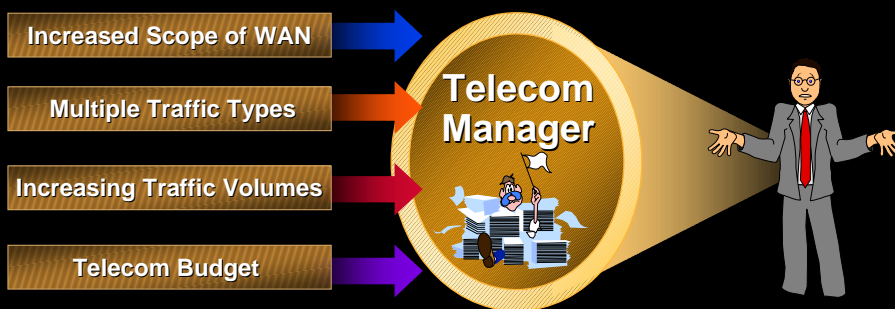
WAN Networking Environment



Multiple Networking Environments where All Users Require Ubiquitous Access to Corporate Applications

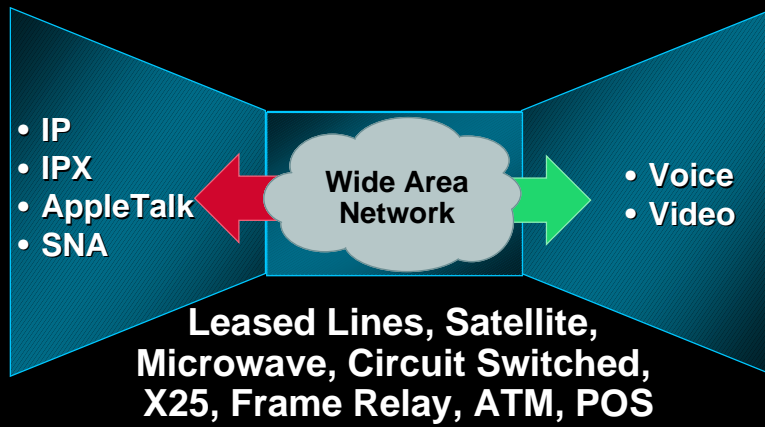
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WAN Networking Challenges



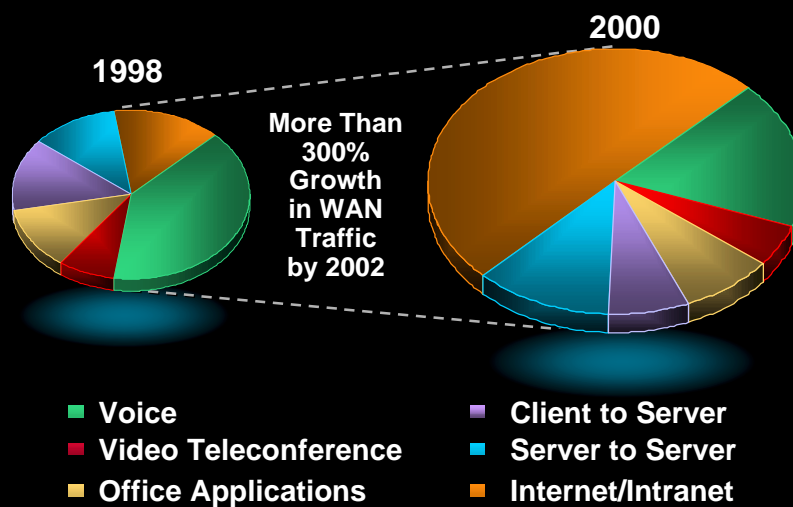
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Multiple Traffic Types



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Networked Application Growth

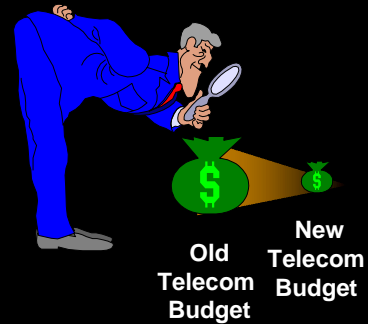


Source: Gartner Group

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Telecom Budgets

Although Traffic Volumes Are Increasing, Telecom Budgets Are Not— in Fact in Some Cases Decreasing



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Agenda

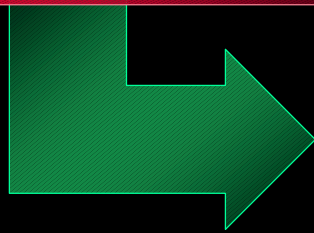
Wide Area Network (WAN)

- Environment
- **Requirements**
- Technologies
- Interface Protocols

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Wide Area Network Requirements

Minimize Bandwidth Costs
Maximize Efficiency
Maximize Performance
Support New/Emerging Applications
Maximize Availability
Minimize Management and Maintenance



Multiservice Consolidation
Bandwidth Efficiency
Performance and QoS Guarantees
Emerging IP Services
Carrier-Class Reliability
Ease of Operation and Management

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Wide Area Network Costs

- **Capital costs (hardware): non-recurring**

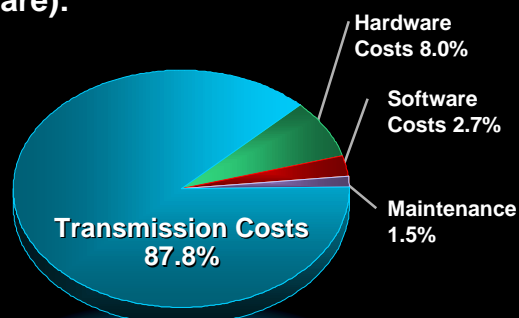
Written off over three to five years

- **Bandwidth rental costs: recurring**

One to five year contracts

Penalties for cancellation

Highest month recurring expense



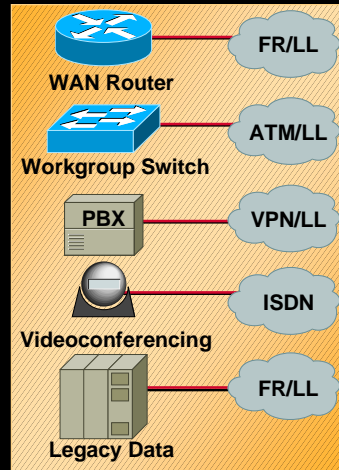
Source: Data Communications

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WAN Service Requirements

- **LAN interconnectivity**
Router and workgroup switches
- **Emerging IP services**
- **Voice toll bypass**
PBX tie trunking and voice switching
- **Videoconferencing**
Room-to-room, desktop
- **Legacy data**
SNA, X.25, async data
- **TDM migration**
Transport option for aging TDM systems

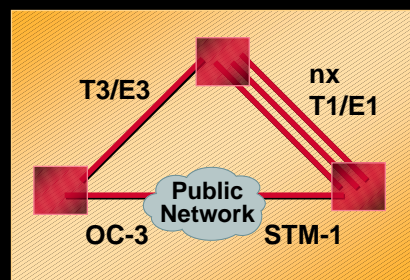
Multiservice Networking



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Wide Area Connectivity Requirements

- **Flexible trunking**
Sub-rate leased lines (64kbps to T1/E1)
 $n \times$ T1/E1 leased lines
T3/E3 leased lines
OC-3/STM-1 leased lines
Public/private services
- **Bandwidth efficiency**



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Agenda

Wide Area Network (WAN)

- Environment
- Requirements
- **Technologies**
- Interface Protocols

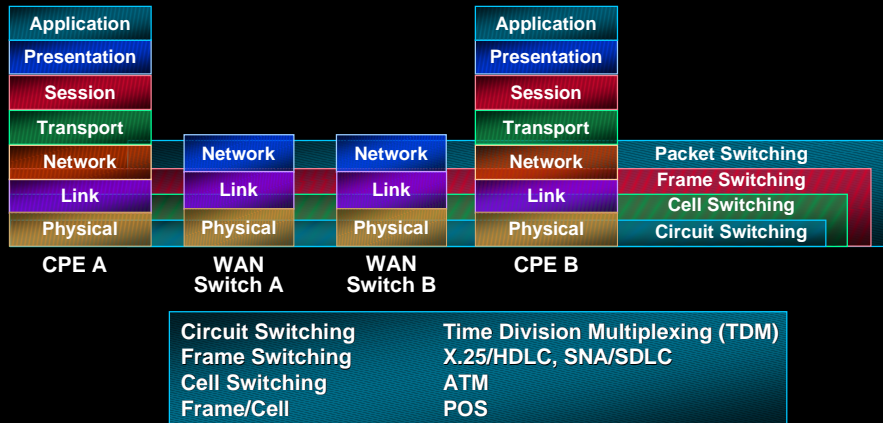
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Message

- **By selecting the right WAN technology and hardware, you can:**
 - Build a uniform multiservice network for DVV consolidation**
 - Reduce recurring network operation costs (such as bandwidth)**

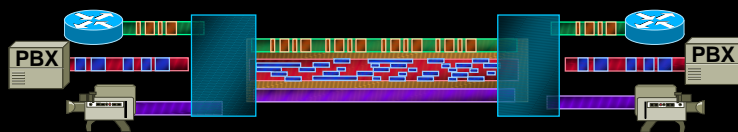
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WAN Technologies



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Circuit Switching (TDM) Overview



- Available bandwidth is statically partitioned amongst the applications
- Each application has access to the allocated bandwidth only; no single application can burst up to the entire bandwidth capacity
- Example: available bandwidth = 24 DS0s
 - Four channels for router data; eight for voice calls; six for video; six channels not allocated
- Circuit switching delay: application frame size/ allocated bandwidth
- Variability in delay: low

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Circuit Switching—Pros and Cons

Pros

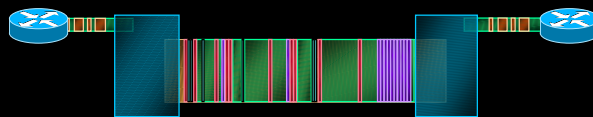
- **Predictable Performance**
As Bandwidth Is Allocated Statically, All Applications Get Consistent Performance and Predictable Delay
- **Protocol-Independent**
End Devices Can Be Using Any Protocol over the Wide Area Network
- **Can Be Used for Multiservice Applications Consolidation**

Cons

- **Inefficient Bandwidth Utilization**
No Dynamic Allocation of Bandwidth
No Bursting Capability—One Application Cannot Use Bandwidth Allocated to the Other Even if That Traffic Is Not Present
- **Bandwidth Limited**
Once All of the Available Bandwidth Is Allocated, Additional Bandwidth Must Be Procured

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X.25 Switching Overview



- Cells, frames, or packets within a switch but X.25 packets at egress
- Available bandwidth is dynamically allocated
- Only useful for data applications
- X.25 switching delay: can increase with number of switches in path
- Variability in delay: high
- Layer 3 switching—PVC, SVC-based

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Traffic Handling in an X.25 Network

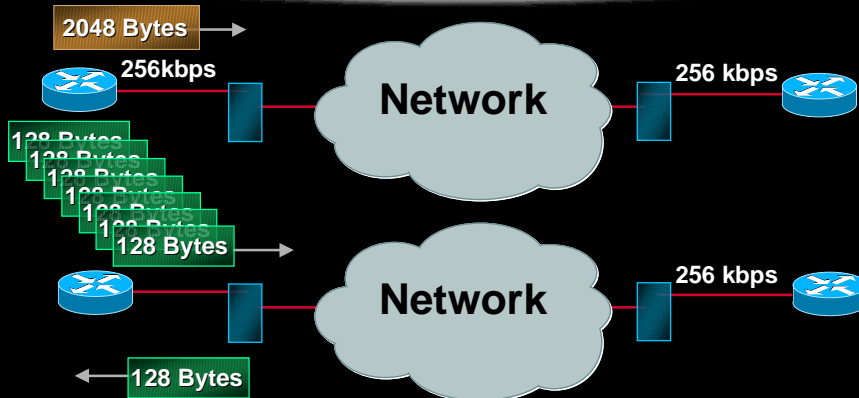


- Router sends X.25 packet to first X.25 switch

Assume that the call has been set up and a data size of 128 agreed by both ends with a window size of 7

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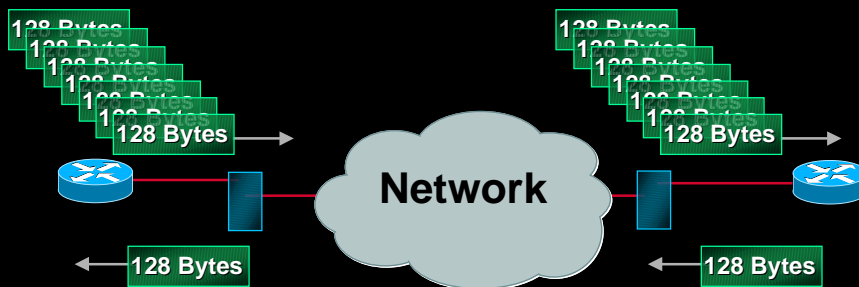
Traffic Handling in an X.25 Network



- Router sends X.25 a window (7) of (128) packets; the X.25 switch responds with "RR's" for correctly received packets or "RNR" for incorrect ones

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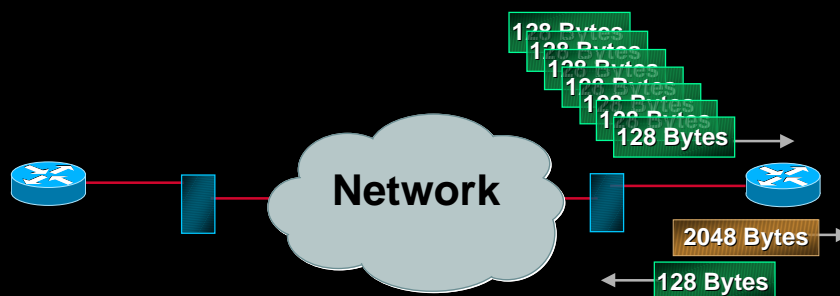
Traffic Handling in an X.25 Network



- The network transports the packets across the network; the packets are delivered to the far end using the same window mechanism as the input to the network

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Traffic Handling in an X.25 Network



- Last window sent to router and whole frame re-assembled and sent

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X.25 Switching—Pros and Cons

Pros

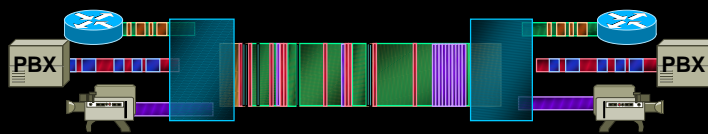
- **Excellent in Highly Errored Environments**
Protocol Goes Looking for an Error
- **Protocol-Independent**
End Devices Can Be Using any Protocol over the Wide Area Network

Cons

- **Unable to Guarantee Performance**
Due to Windowing Mechanism, Throughput and Performance Is Very Slow
- **High Delay and Variability in Delay**
Each Switch Has to Receive an Entire Window before Forwarding It to the Next Switch; Therefore Transit Delay Increases with Number of Switches in the Path
The FIFO Mode of Each Switch Causes a Variability at Each Switch
- **Typically Data Only**
As a Result of High Delay and Variability in Delay, X.25 Switches Are More Suited for Data-Only Environments that Have Little or No Delay-Sensitive Applications

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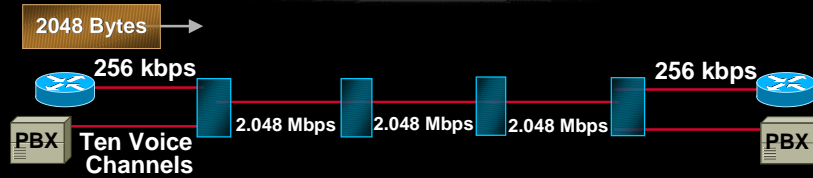
Frame Switching Overview



- **Either cells or frames within a switch but frames only at egress**
- **Available bandwidth is dynamically allocated**
- **Each application gets access to all available bandwidth; any application can burst up to the entire bandwidth capacity**
- **Frame switching delay: increases with number of switches in path (queuing and serialization delay)**
- **Variability in delay: can be high for FIFO queuing switches**

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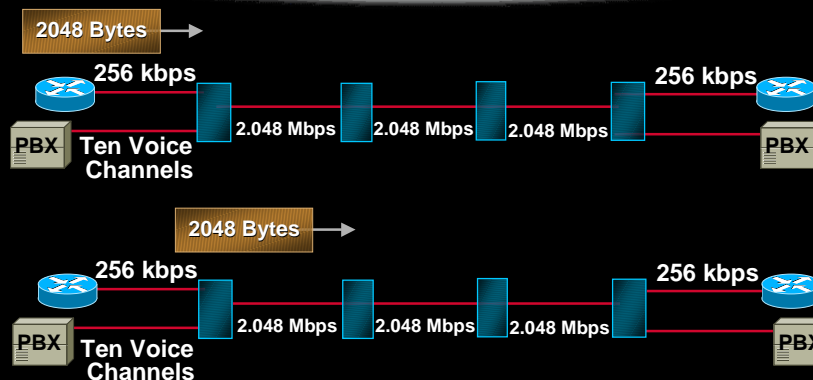
Traffic Handling in Frame-Switched Network



- Router sends frame to first frame switch

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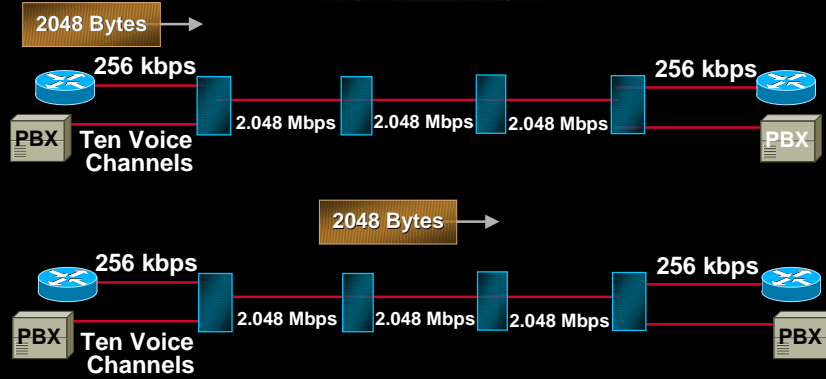
Traffic Handling in Frame-Switched Network



- First frame switch waits for the entire frame, adds header and forwards to next switch

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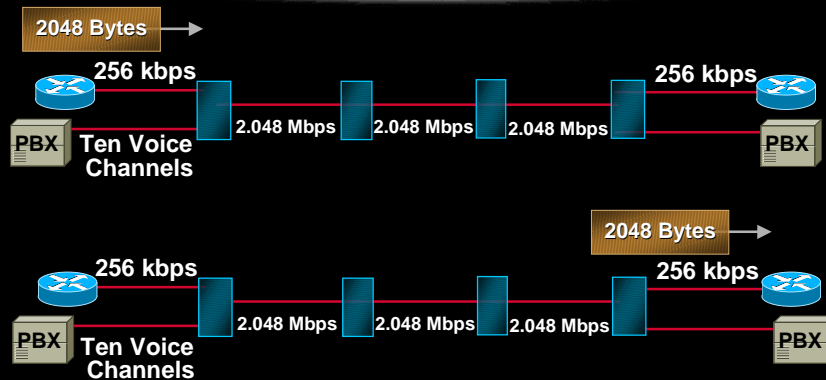
Traffic Handling in Frame-Switched Network



- Second frame switch waits for the entire frame, and forwards to next switch

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Traffic Handling in Frame-Switched Network



- Last frame switch waits for the entire frame, adds header and forwards it to router

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Frame Switching—Pros and Cons

Pros

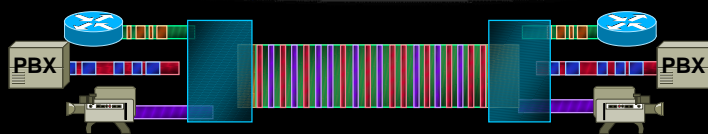
- **Dynamic Allocation of Bandwidth**
Available Bandwidth Is Allocated Dynamically to Any Application That Needs It
One Application Can Use Bandwidth Allocated to the Other If That Traffic Is Not Present
- **Can Be Used for Multiservice Applications**
Frame Switches Are Used for Multiservice Applications (DVV) (Less over Subscription and Reasonable Speed Links)

Cons

- **Unable to Guarantee Performance (in FIFO Mode)**
Frame Switches 'Typically' Operate in FIFO (First in-First out) Mode, so One Application Can Impact the Performance of Others
- **Medium Delay and Variability in Delay**
Each Switch Has to Receive an Entire Frame before Forwarding It to the Next Switch; Therefore Transit Delay Increases with Number of Switches in the Path
The FIFO Mode of Each Switch Causes a Variability at Each Switch

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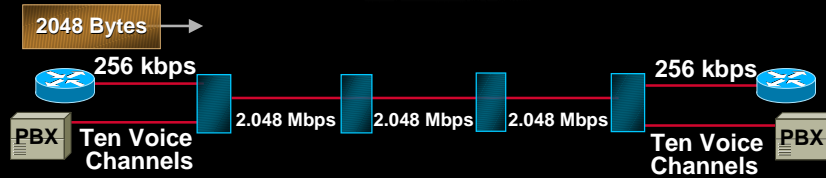
Cell Switching Overview



- **Cells only within a switch and cells only at the egress of a switch**
- **Available bandwidth is dynamically allocated**
- **Each application gets access to all available bandwidth; any application can burst up to the entire bandwidth capacity**
- **Cell switching delay: low delay**
- **Variability in delay: low**

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Traffic Handling in Cell-Switched Network



- Router sends frame to first cell switch

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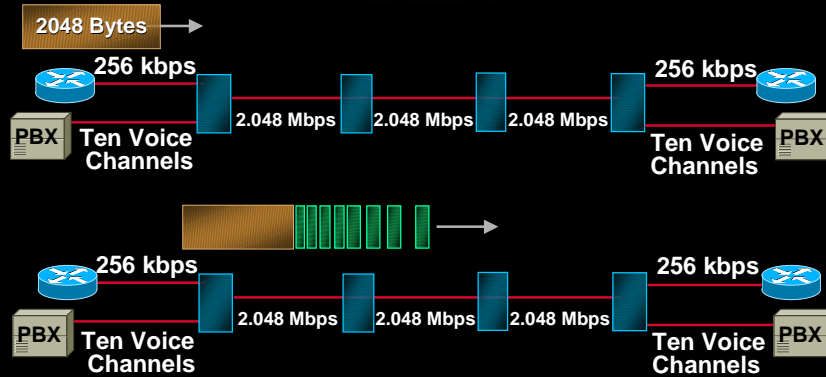
Traffic Handling in Cell-Switched Network



- First cell switch starts to receive the frame; as it receives enough data to fill up the payload of one cell, it adds header and forwards to next switch

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Traffic Handling in Cell-Switched Network



- Second cell switch receives the cell and switches it to the next switch

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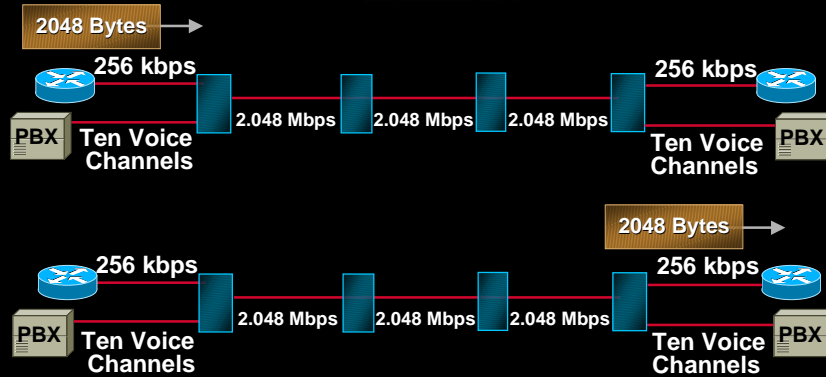
Traffic Handling in Cell-Switched Network



- Next cell switch receives the cells and forwards to next switch till the last switch where all cells are collected until the last cell is received

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Traffic Handling in Cell-Switched Network



- Last switch assembles the frame and forwards to the attached router

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Cell Switching—Pros and Cons

Pros

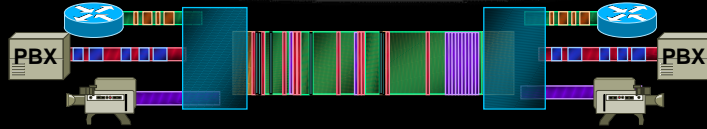
- **Dynamic Allocation of Bandwidth**
Available Bandwidth Is Allocated Dynamically to Any Application That Needs It
One Application Can Use Bandwidth Allocated to the Other if That Traffic Is Not Present
- **Guaranteed Performance**
Cell Switches with Efficient Traffic and Bandwidth Management Schemes Can Ensure That Each Application Receives Guaranteed Performance (TM, QoS Queuing, CAC, PNNI/UNI, Etc.)
- **Low Delay (Controlled and Bounded) and Low Variability in Delay**
Using Fixed Length Cells Ensures That Network Transit Delay and Variability in Delay Is Minimized
Switches Use QoS-Based Queuing and Scheduling Such as CBR, VBR, ABR
- **Typically Multiservice**
As a Result of Low Delay, Low Variability in Delay and the Ability to Guarantee Performance, Cell Switches Are Ideally Suited to Support Multiple Services Concurrently

Cons

- **Overhead**
However the Bandwidth Efficiency and Ability to Provide Low Delay and Low Variability in Delay in Cell Switching Easily Overcomes the Small Incremental Overhead

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Frame/Cell Switching Overview



- Either cells or frames within a switch and both frames and cells at egress of switch
- Frame switching is used for data and cells are injected for delay-sensitive traffic
- Available bandwidth is dynamically allocated
- Each application gets access to all available bandwidth; any application can burst up to the entire bandwidth capacity
- Frame/cell switching delay: increases with number of switches in path
- Variability in delay: high

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Frame/Cell Switching— Pros and Cons

Pros

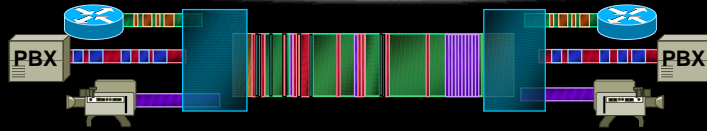
- **Dynamic Allocation of Bandwidth**
Available Bandwidth Is Allocated Dynamically to Any Application That Needs It
One Application Can Use Bandwidth Allocated to the Other if That Traffic Is Not Present
- **Multiple Traffic Types**
Cell Injection Process Is Used to 'Insert' Delay Sensitive Traffic (Like Voice) During the Transfer of a Data Frame

Cons

- **Proprietary**
There Are No Standards for Frame/Cell Switching
- **Unable to Guarantee Performance**
Cell Injection Causes Frames to Be Partially Transmitted and Possibly Buffered Along the Path; Frame-Based Traffic (Data) Suffers at the Expense of Cell Traffic; Increased Volumes of Cell Traffic Can Impact Quality Throughput and Performance of Frame-Based Traffic
- **Bandwidth Inefficiency**
Increased Cell Traffic Can Cause Frame Traffic to Be Discarded in the Network, Thereby Reducing the 'Quality Throughput' or 'Goodput' of the Network
- **High Delay and Variability in Delay**
Frame/Cell Switches Operate as Frame Switches for Data and Have Similar Delay Characteristics; Cell Injection Causes Additional Delays for Frame Traffic

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Packet over SONET Overview



- Either cells or frames within a switch and SONET/SDH payload at egress of switch
- POS is serial transmission of multiprotocol packets over SONET/SDH frames—efficient high-speed transport with low overhead
- Available bandwidth is dynamically allocated
- Any application can burst up to the entire available bandwidth capacity-packet based infrastructure
- Switching delay: low
- Variability in delay: low

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Packet over SONET— Pros and Cons

Pros

- **Dynamic Allocation of Bandwidth**
Bandwidth Is Available to Any Application That Needs It
- **Scalable**
Underlying Technology Can Scale in Speeds with Increase in Traffic Volume
- **Efficiency-Overhead**
Direct Mapping of Packets over SONET Payload Increases Efficiency by Removing Overhead
- **Fault-Tolerant, Reliability Carrier Class Redundancy with APS**
With Automatic Protection Switching, Very Fast Switchovers in Case of Failure
- **Underlying Transport Flexibility**
Can Be Used with Different Underlying Transport-SONET/SDH ADMs, DWDM Equipment

Cons

- **More Suited Towards Packet-Based Applications**
Direct Mapping of Packets over Transport SONET/SDH Media
- **QoS Evolving from Best-Effort to Guaranteed Approach**
Efforts Underway to Include Differentiated Service Mechanisms
- **Not Suitable for Low Speed Links**
'Typically' Used over High-Speed Fiber-Based Infrastructures

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Summary of Technologies

	Bandwidth Efficiency	Multiservice	Traffic Mgmt./QoS	Appl. Performance	Scalability
Circuit Switching (TDM)	*	***	*	**	*
X.25	**	*	*	*	*
Frame Switching	***	***	***	***	***
Cell Switching (ATM)	****	****	****	****	****
Frame/Cell Switching	***	***	***	***	***
Packet over SONET	****	**	**	***	***

* = Weak ** = Fair *** = Good **** = Excellent

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Agenda

Wide Area Network (WAN)

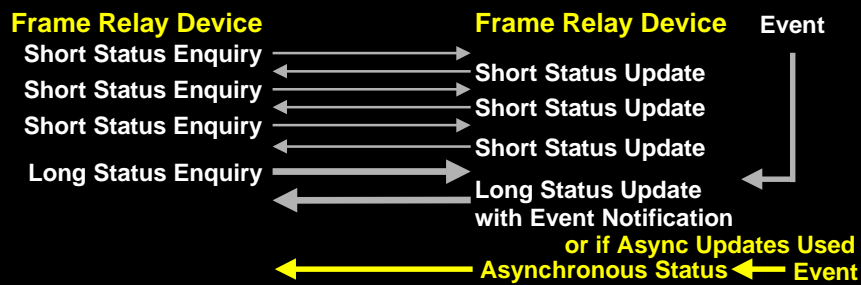
- Environment
- Requirements
- Technologies
- Interface Protocols

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Frame Relay UNI LMI

- UNI between end device and network
- Event notification
- Unidirectional update requests
- E-LMI allows one config. of switches and router
- Same frame format as NNI

ANSI	T1.617 Annex D
ITU-T	Q.933 Annex A Signaling
LMI	"Gang of Four" Cisco, NT, DEC

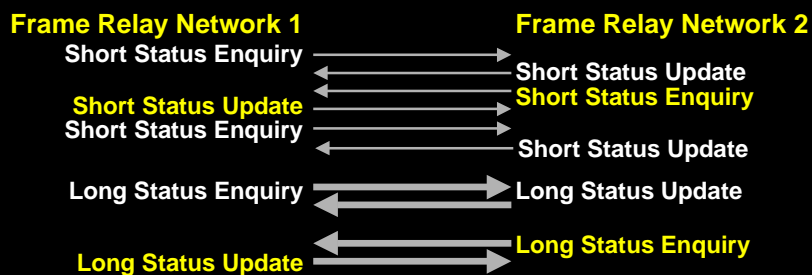


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Frame Relay NNI

- NNI between networks
- Event notification
- Bi-directional update requests
- Allows foresight updates between two Cisco networks using CLLM
- Same frame format as UNI

ANSI	T1.618 Annex D
ITU-T	Q.933 Annex A



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ATM UNI and NNI

Generic Flow Control	Virtual Path Identifier	
Virtual Path Identifier	Virtual Channel Identifier	
Virtual Channel Identifier		
Virtual Channel Identifier	Payload Type	CLP
Header Error Control (HEC)		

UNI Header Format

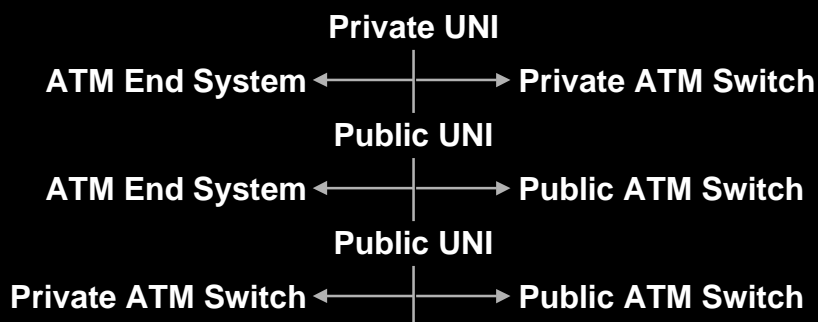
Virtual Path Identifier	
Virtual Path Identifier	Virtual Channel Identifier
Virtual Channel Identifier	
Virtual Channel Identifier	Payload Type
CLP	Header Error Control (HEC)

NNI Header Format

- Unlike Frame Relay ATM UNI and NNI formats are different
- UNI VPI range is 256 and VCI range is 65,535
- NNI VPI range is 4,096 and VCI range is 65,535

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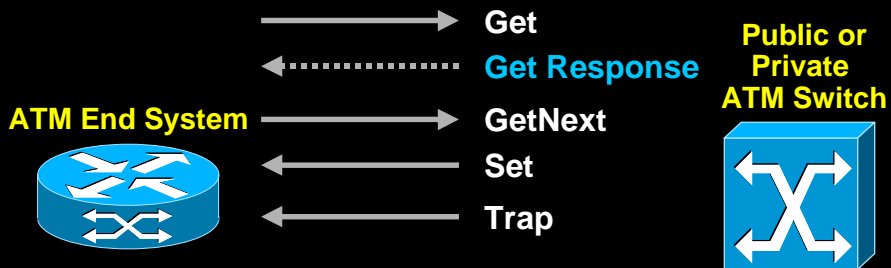
ATM I—LMI



- Integrated local management interface-ilmi
- Use SNMP across UNI and NNI for ILMI MIB
- Uses AAL 5 encapsulation
- Used for ATM end system address (AESAs) formerly NSAP addressing for services
- Automatic recognition of UNI or NNI interface protocol

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ATM I—LMI (Cont.)



- **Get**—retrieve specified management information
- **GetNext**—retrieve via traversal of MIB, management information
- **Set**—alter management information
- **Trap**—report extraordinary information

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Summary

- **By selecting the right WAN technology and hardware, you can:**
 - Build a uniform multiservice network for DVV consolidation**
 - Reduce recurring network operation costs (such as bandwidth)**

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Follow-on Sessions

- **WAN**

Session #102 (Level 2)—Deploying WAN Technologies

Session #103 (Level 3)—Advanced WAN Concepts and Troubleshooting

Session #1201 (Product)—IGX/BPX/MGX Product Update

Session #1208 (Product)—26xx/36xx/3810 Product Update

- **Optical**

Session #604 (Level 1)—Introduction to Optical Carrier Services

Session #605 (Level 2)—Deploying Optical Infrastructure

Session #606 (Level 3)—Advanced Optical Technology Concepts

Session #1202 (Product)—GSR Product Update

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Session 101

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